

LINEAR PROGRAMMING, ZERO-FORCING, AND MAXIMALLY DIVERSE OPTIMA

(PROGRESS DOCUMENTATION)

1 Diversity Model/Integer Program

Denote the diversity model as $\text{DZF}(G, s^*)$, where $G \in \mathbb{G}$ and s^* is an optimal solution to $\text{ZIP}(G)$.

$$\text{minimize} \quad \sum_{v \in V} z_v \tag{1a}$$

$$\text{subject to} \quad s_v + \sum_{e \in \delta^-(v)} y_e = 1, \quad \forall v \in V, \tag{1b}$$

$$s'_v + \sum_{e \in \delta^-(v)} y'_e = 1, \quad \forall v \in V, \tag{1c}$$

$$x_u - x_v + (T+1)y_e \leq T, \quad \forall e = (u, v) \in E, \tag{1d}$$

$$x'_u - x'_v + (T+1)y'_e \leq T, \quad \forall e = (u, v) \in E, \tag{1e}$$

$$x_w - x_v + (T+1)y_e \leq T, \quad \forall e = (u, v) \in E, \forall w \in N(u) \setminus \{v\}, \tag{1f}$$

$$x'_w - x'_v + (T+1)y'_e \leq T, \quad \forall e = (u, v) \in E, \forall w \in N(u) \setminus \{v\}, \tag{1g}$$

$$\sum_{v \in V} s_v = s^*, \tag{1h}$$

$$\sum_{v \in V} s'_v = s^*, \tag{1i}$$

$$s_v + s'_v - z_v \leq T, \quad \forall v \in V, \tag{1j}$$

$$s, s' \in \{0, 1\}^n, \quad x, x' \in \{0, \dots, T\}^n, \quad y, y' \in \{0, 1\}^m \tag{1k}$$

Proposition 1.1. *Given an optimal solution to $\text{DZF}(G, s^*)$, $z_v = 1$ for some $v \in V$ if and only if $s_v = s'_v = 1$.*

Proof. Suppose $z_v = 1$ for some $v \in V$. By way of contradiction, also suppose either $s_v = 0$ or $s'_v = 0$. Then constraint (1j) will be true if $z_v = 0$. This contradicts the fact that we have an optimal solution to $\text{DZF}(G, s^*)$, i.e., the objective function (1a) would not be minimized. Thus, our assumption that $s_v = 0$ or $s'_v = 0$ must be false. Hence, $s_v = s'_v = 1$.

Conversely, let $s_v = s'_v = 1$ for some $v \in V$. Then constraint (1j) holds only when $z_v = 1$. Therefore, given an optimal solution to $\text{DZF}(G, s^*)$, $z_v = 1$ for some $v \in V$ if and only if $s_v = s'_v = 1$. \square