

I. INTRODUCTION

Vulnerabilities can be present and triggered by any layer of a computer system (Figure 1). However, each vulnerability has a certain risk level depending on what can be compromised.

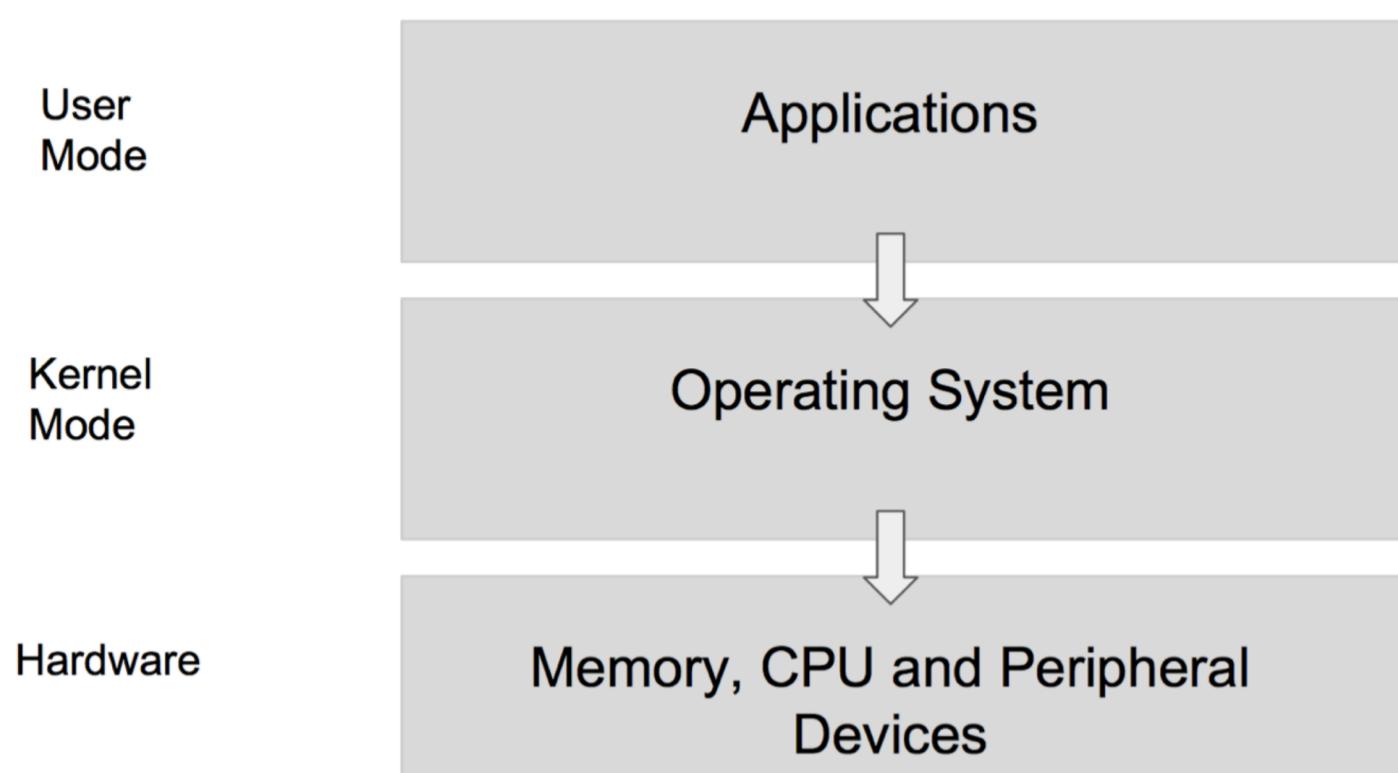
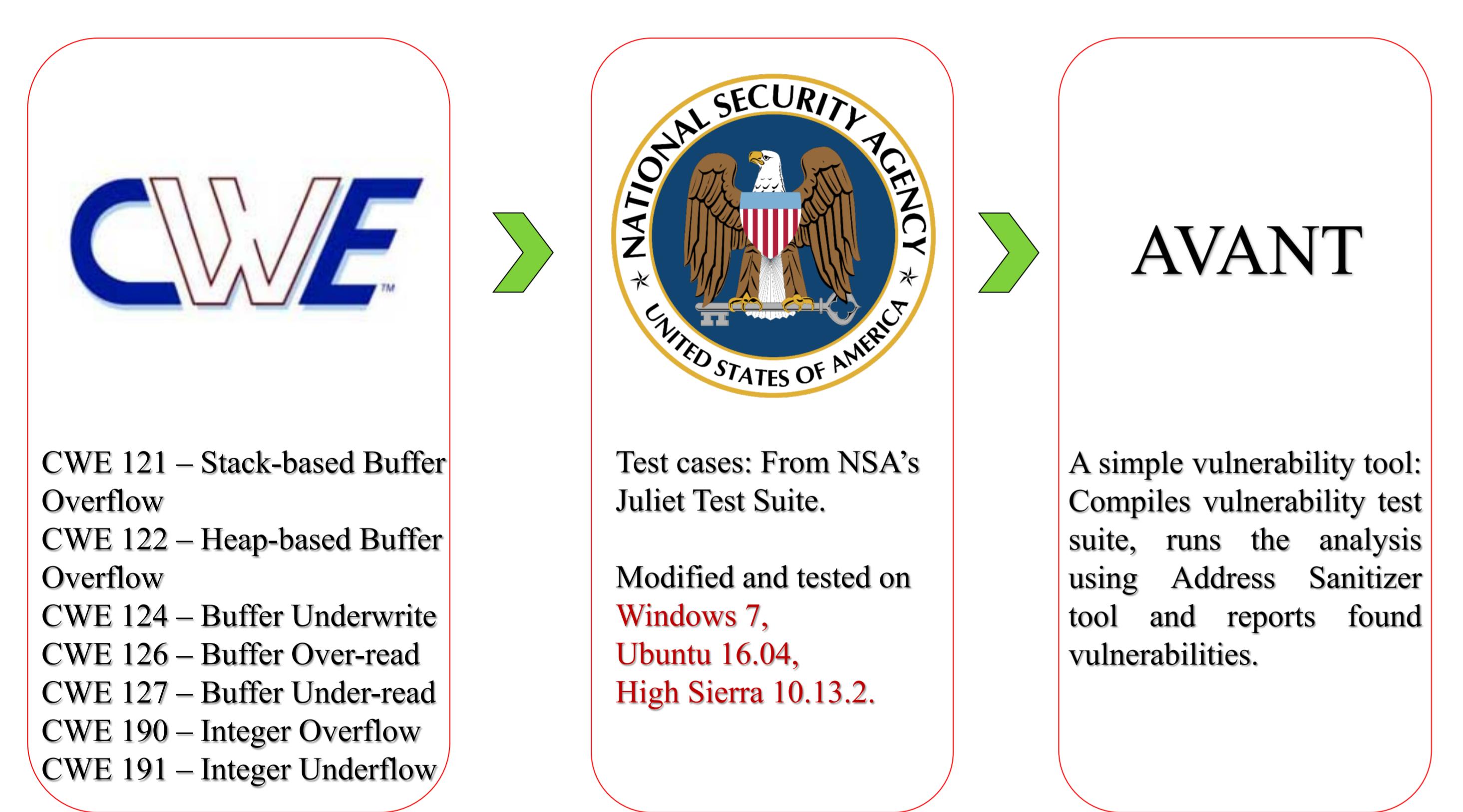


Figure 1: Typical representation of layers in computer system.

We present Automated Vulnerability Analysis Tool (AVANT), a vulnerability checking tool that is architecture-agnostic and reports vulnerabilities found for binaries in our test suite.

III. DESIGN MODEL



V. PRELIMINARY RESULTS

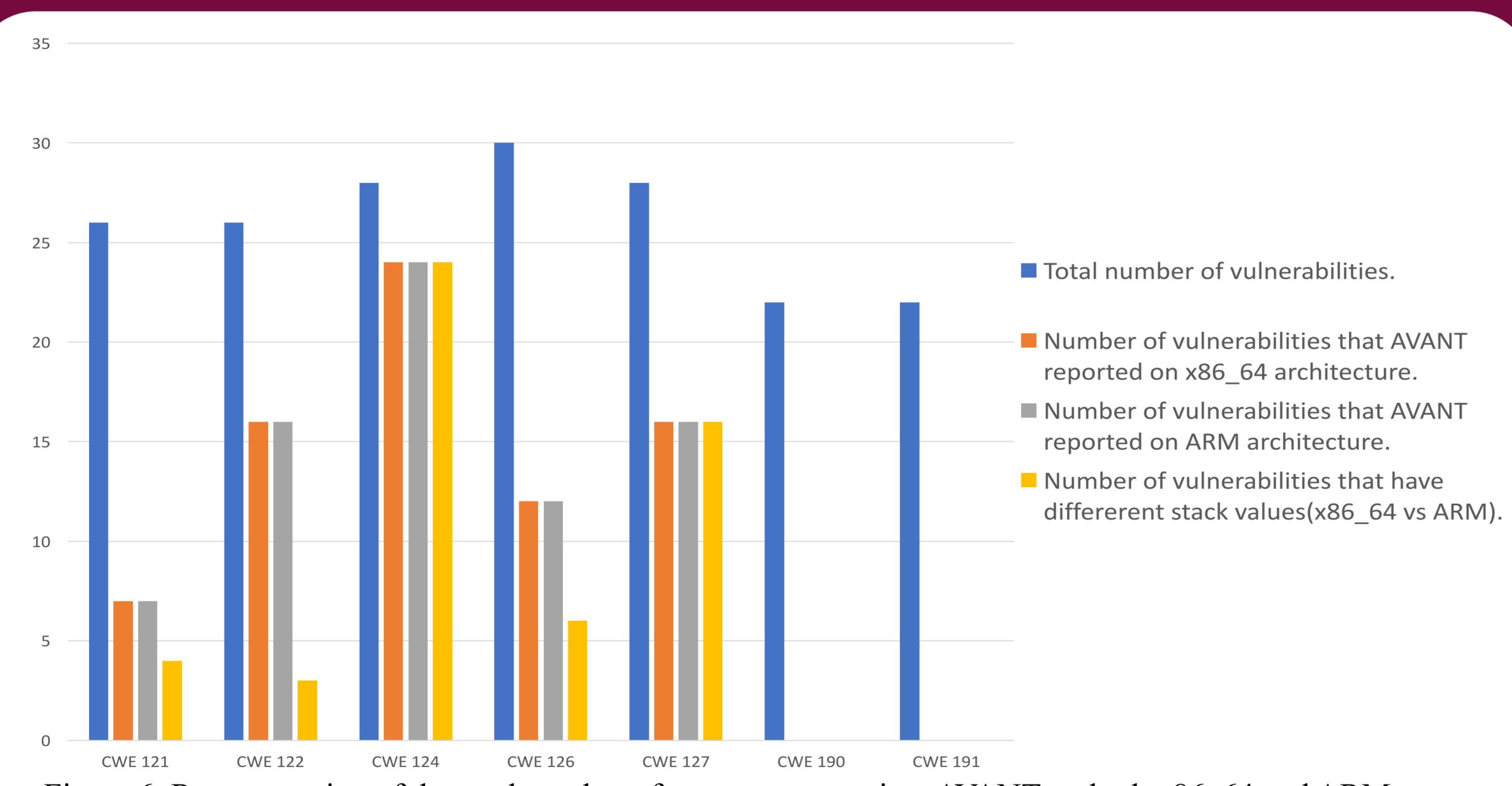


Figure 6: Representation of the total number of test cases ran using AVANT on both x86_64 and ARM architectures and the number of reported vulnerabilities.

SUMMARY: AddressSanitizer: stack-buffer-underflow (/home/pi/OSFA-compiled/OSFA-

Benchmarks/testcases/CWE124_Buffer_Underwrite/s01/CWE124_Buffer_Underwrite_char_declare_memcpy_01.out+0x4a2b44)
Shadow bytes around the buggy address:
0x10002e003f50: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x10002e003f60: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x10002e003f70: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x10002e003f80: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x10002e003f90: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
=>0x10002e003fa0: 00 00 00 f1 f1 f1 [f1]00 00 00 00 00 00 00 00 00
0x10002e003fb0: 00 00 00 00 04 f3 f3 f3 f3 f3 00 00 00
0x10002e003fc0: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x10002e003fd0: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x10002e003fe0: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x10002e003ff0: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00

Figure 7: Stack representation of the CWE124 Buffer Underwrite test case ran on x86_64.

SUMMARY: AddressSanitizer: stack-buffer-underflow (/home/pi/OSFA-compiled/OSFA-
Benchmarks/testcases/CWE124_Buffer_Underwrite/s01/CWE124_Buffer_Underwrite_char_declare_memcpy_01.out+0xab6db) in _asan_memset
Shadow bytes around the buggy address:
0x2fd21e40: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x2fd21e50: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x2fd21e60: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x2fd21e70: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x2fd21e80: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
=>0x2fd21e90: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x2fd21ea0: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x2fd21eb0: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x2fd21ec0: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x2fd21ed0: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
0x2fd21ee0: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00

Figure 8: Stack representation of the CWE124 Buffer Underwrite test case ran on Raspberry Pi 3.

II. MOTIVATION

Technology innovations and improvements

Number of smart devices increases

Larger attack surface

Below is an example of a buffer overflow vulnerability:

```
char pass[] = "abcd";
int validate_user() {
    char buff[5];
    printf("Enter your password:\n-> ");
    gets(buff);
    return !strcmp(buff, pass);
}
```

Figure 2: CWE-120 Buffer Overflow based example.

```
0x7fffffff9b0:0x00 0x00 0x00 0x00 0x00 0x00
0x00 0x00
0x7fffffff9b8:0x00 0x00 0x00 0x00 0x00 0x00
0x00 0x00
0x7fffffff9c0:0xe0 0xd9 0xff 0xff 0xff 0x7f
0x00 0x00
0x7fffffff9c8:0xb7 0x06 0x40 0x00 0x00 0x00
0x00 0x00
```

```
0x7fffffff9d0:0xe0 0xd9 0xff 0xff 0x7f
0x00 0x00
0x7fffffff9d8:0xb7 0x06 0x40 0x00 0x00 0x00
0x00 0x00
```

Figure 3: Representation of the stack pointer before inputted password.

Figure 4: Representation of the stack pointer after inputted "AAAAAA" as a password.

In this buffer overflow the part of the input that does not fit in the buffer keeps being written in adjacent memory address.

Meltdown and Spectre attacks showed us that a bug in Intel chips allows access to higher parts of computer's memory.

IV. EXPERIMENTAL SETUP

100+ Test cases

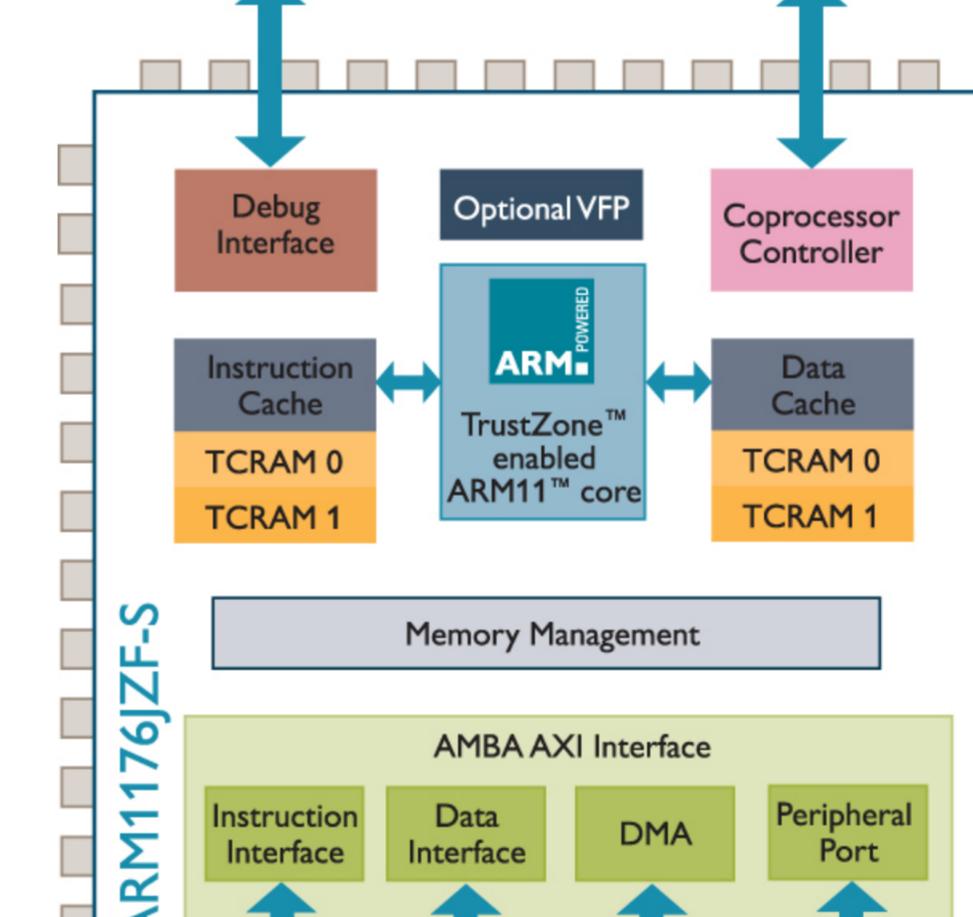


Figure 5: ARM architecture.

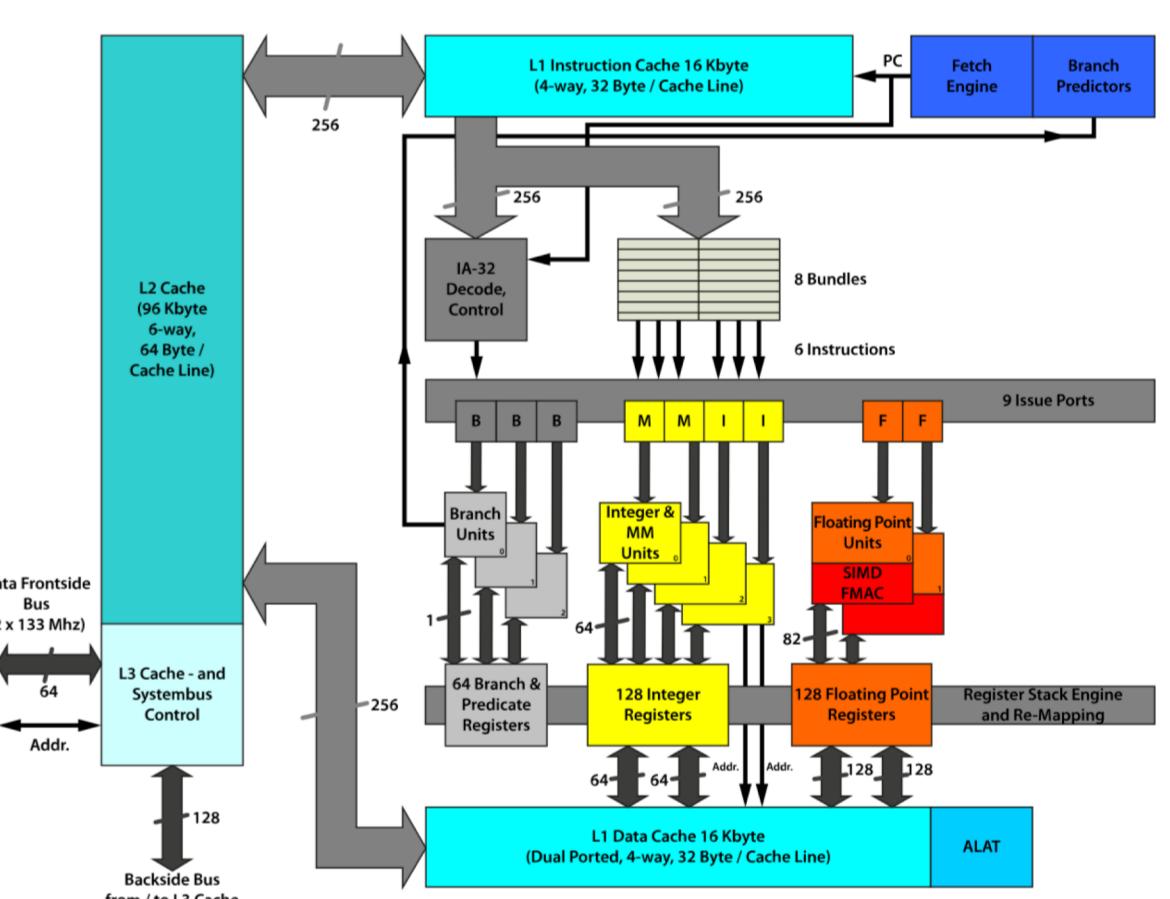


Figure 6: Intel x86_64 architecture.

	ARM Architecture	Intel x86_64
Architecture	armv7l	x86_64
Byte Order:	Little Endian	Little Endian
CPU(s):	4	8
On-line CPU(s) list:	0-3	0-7
Thread(s) per core:	1	2
Core(s) per socket:	4	4
Socket(s):	1	1
Model:	4	94
Model name:	ARMv7 Processor rev 4 (v7l)	Intel(R) Core(TM) i7-6700K CPU @ 4.00GHz
CPU max MHz:	1200.0000	4200.0000
CPU min MHz:	600.0000	800.0000
BogoMIPS:	38.40	8015.91

VI. CONCLUSION & FUTURE WORK

- These results show us that both Intel x86_64 and ARM architectures report same vulnerabilities; this is good from the security standpoint. The interesting part is that same reported vulnerabilities on two architectures differ in their stack values/addresses which can lead to some other attacks on the security of those architectures.
- The research done in this project will be extended to look more deeply into this problem by expanding the test set in order to identify more vulnerabilities. We are developing a test suite that includes binaries with annotated and categorized vulnerabilities.

VII. REFERENCE

- [1] Trecakov, S., Tran, C., Badawy, H., Siddique, N., Acosta, J., & Misra, S. (2017, October). Can Architecture Design Help Eliminate Some Common Vulnerabilities?. In 2017 IEEE 14th International Conference on Mobile Ad Hoc and Sensor Systems (MASS) (pp.590-593). IEEE.