02-0: **Algorithm Analysis**

When is algorithm A better than algorithm B?

02-1: **Algorithm Analysis**

When is algorithm A better than algorithm B?

- Algorithm A runs faster

02-2: **Algorithm Analysis**

When is algorithm A better than algorithm B?

- Algorithm A runs faster
- Algorithm A requires less space to run

02-3: **Algorithm Analysis**

When is algorithm A better than algorithm B?

- Algorithm A runs faster
- Algorithm A requires less space to run

Space / Time Trade-off

- Can often create an algorithm that runs faster, by using more space

For now, we will concentrate on time efficiency

02-4: **Best Case vs. Worst Case**

How long does the following function take to run:

```
boolean find(int A[], int element) {
  for (i=0; i<A.length; i++) {
     if (A[i] == elem)
       return true;
  }
  return false;
}
```

02-5: **Best Case vs. Worst Case**

How long does the following function take to run:

```
boolean find(int A[], int element) {
  for (i=0; i<A.length; i++) {
     if (A[i] == elem)
       return true;
  }
  return false;
}
```

It depends on if – and where – the element is in the list

02-6: **Best Case vs. Worst Case**

- Best Case – What is the fastest that the algorithm can run
- Worst Case – What is the slowest that the algorithm can run
- Average Case – How long, on average, does the algorithm take to run

Worst Case performance is almost always important.
Usually, Best Case performance is unimportant (why?)
Usually, Average Case = Worst Case (but not always!)

02-7: **Measuring Time Efficiency**

How long does an algorithm take to run?

02-8: **Measuring Time Efficiency**

How long does an algorithm take to run?

- Implement on a computer, time using a stopwatch.

02-9: **Measuring Time Efficiency**

How long does an algorithm take to run?

- Implement on a computer, time using a stopwatch.
  Problems:
  - Not just testing algorithm – testing implementation of algorithm
  - Implementation details (cache performance, other programs running in the background, etc) can affect results
  - Hard to compare algorithms that are not tested under *exactly the same conditions*

02-10: **Measuring Time Efficiency**

How long does an algorithm take to run?

- Implement on a computer, time using a stopwatch.
  Problems:
  - Not just testing algorithm – testing implementation of algorithm
  - Implementation details (cache performance, other programs running in the background, etc) can affect results
  - Hard to compare algorithms that are not tested under *exactly the same conditions*
- Better Method: Build a mathematical model of the running time, use model to compare algorithms

02-11: **Competing Algorithms**

- Linear Search

```
for (i=low; i <= high; i++)
  if (A[i] == elem) return true;
return false;
```

- Binary Search

```
int BinarySearch(int low, int high, elem) {
  if (low > high) return false;
  mid = (high + low) / 2;
  if (A[mid] == elem) return true;
  if (A[mid] < elem)
    return BinarySearch(mid+1, high, elem);
  else
    return BinarySearch(low, mid-1, elem);
}
```

02-12: **Linear vs Binary**

- Linear Search

```
for (i=low; i <= high; i++)
  if (A[i] == elem) return true;
return false;
```

Time Required, for a problem of size n (worst case):

02-13: **Linear vs Binary**

- Linear Search

```
for (i=low; i <= high; i++)
  if (A[i] == elem) return true;
return false;
```

Time Required, for a problem of size n (worst case):

$c_1 * n$ for some constant c_1

02-14: **Linear vs Binary**

- Binary Search

```
int BinarySearch(int low, int high, elem) {
  if (low > high) return false;
  mid = (high + low) / 2;
  if (A[mid] == elem) return true;
  if (A[mid] < elem)
    return BinarySearch(mid+1, high, elem);
  else
    return BinarySearch(low, mid-1, elem);
}
```

Time Required, for a problem of size n (worst case):

02-15: **Linear vs Binary**

- Binary Search

```
int BinarySearch(int low, int high, elem) {
  if (low > high) return false;
  mid = (high + low) / 2;
  if (A[mid] == elem) return true;
  if (A[mid] < elem)
    return BinarySearch(mid+1, high, elem);
  else
    return BinarySearch(low, mid-1, elem);
}
```

Time Required, for a problem of size n (worst case): $c_2 * lg(n)$ for some constant c_2

02-16: **Do Constants Matter?**

- Linear Search requires time $c_1 * n$, for some c_1
- Binary Search requires time $c_2 * lg(n)$, for some c_2

What if there is a *very* high overhead cost for function calls?

What if c_2 is *1000 times larger* than c_1?

02-17: **Constants *Do Not* Matter!**

Length of list	Time Required for Linear Search	Time Required for Binary Search
10	0.001 seconds	0.3 seconds
100	0.01 seconds	0.66 seconds
1000	0.1 seconds	1.0 seconds
10000	1 second	1.3 seconds
100000	10 seconds	1.7 seconds
1000000	2 minutes	2.0 seconds
10000000	17 minutes	2.3 seconds
10^{10}	11 days	3.3 seconds
10^{15}	30 centuries	5.0 seconds
10^{20}	300 million years	6.6 seconds

02-18: **Growth Rate**

We care about the *Growth Rate* of a function – how much more we can do if we add more processing power

Faster Computers $\neq$ Solving Problems Faster
Faster Computers $=$ Solving Larger Problems

- Modeling more variables
- Handling bigger databases
- Pushing more polygons

02-19: **Growth Rate Examples**

	Size of problem that can be solved					
time	$10n$	$5n$	$n \lg n$	n^2	n^3	2^n
1 s	1000	2000	1003	100	21	13
2 s	2000	4000	1843	141	27	14
20 s	20000	40000	14470	447	58	17
1 m	60000	120000	39311	774	84	19
1 hr	3600000	7200000	1736782	18973	331	25

02-20: **Constants and Running Times**

- When calculating a formula for the running time of an algorithm:
  - Constants aren't as important as the growth rate of the function
  - Lower order terms don't have much of an impact on the growth rate
    - $x^3 + x$ vs x^3
- We'd like a formal method for describing what is important when analyzing running time, and what is not.

02-21: **Big-Oh Notation**

$O(f(n))$ is the set of all functions that are bound from above by $f(n)$

$T(n) \in O(f(n))$ if

$\exists c, n_0$ such that $T(n) \leq c * f(n)$ when $n > n_0$

02-22: **Big-Oh Examples**

$$\begin{aligned}
n &\in O(n) \text{ ?} \\
10n &\in O(n) \text{ ?} \\
n &\in O(10n) \text{ ?} \\
n &\in O(n^2) \text{ ?} \\
n^2 &\in O(n) \text{ ?} \\
10n^2 &\in O(n^2) \text{ ?} \\
n \lg n &\in O(n^2) \text{ ?} \\
\ln n &\in O(2n) \text{ ?} \\
\lg n &\in O(n) \text{ ?} \\
3n+4 &\in O(n) \text{ ?} \\
5n^2+10n-2 &\in O(n^3)\text{? } O(n^2) \text{ ? } O(n) \text{ ?}
\end{aligned}$$

02-23: **Big-Oh Examples**

$$\begin{aligned}
n &\in O(n) \\
10n &\in O(n) \\
n &\in O(10n) \\
n &\in O(n^2) \\
n^2 &\not\in O(n) \\
10n^2 &\in O(n^2) \\
n \lg n &\in O(n^2) \\
\ln n &\in O(2n) \\
\lg n &\in O(n) \\
3n+4 &\in O(n) \\
5n^2+10n-2 &\in O(n^3), \in O(n^2), \not\in O(n) \text{ ?}
\end{aligned}$$

02-24: **Big-Oh Examples II**

$$\begin{aligned}
\sqrt{n} &\in O(n) \text{ ?} \\
\lg n &\in O(2^n) \text{ ?} \\
\lg n &\in O(n) \text{ ?} \\
n \lg n &\in O(n) \text{ ?} \\
n \lg n &\in O(n^2) \text{ ?} \\
\sqrt{n} &\in O(\lg n) \text{ ?} \\
\lg n &\in O(\sqrt{n}) \text{ ?} \\
n \lg n &\in O(n^{\frac{3}{2}}) \text{ ?} \\
n^3 + n \lg n + n\sqrt{n} &\in O(n \lg n) \text{ ?} \\
n^3 + n \lg n + n\sqrt{n} &\in O(n^3) \text{ ?} \\
n^3 + n \lg n + n\sqrt{n} &\in O(n^4) \text{ ?}
\end{aligned}$$

02-25: **Big-Oh Examples II**

$$\begin{array}{rcl}
\sqrt{n} & \in & O(n) \\
\lg n & \in & O(2^n) \\
\lg n & \in & O(n) \\
n \lg n & \not\in & O(n) \\
n \lg n & \in & O(n^2) \\
\sqrt{n} & \not\in & O(\lg n) \\
\lg n & \in & O(\sqrt{n}) \\
n \lg n & \in & O(n^{\frac{3}{2}}) \\
n^3 + n \lg n + n\sqrt{n} & \not\in & O(n \lg n) \\
n^3 + n \lg n + n\sqrt{n} & \in & O(n^3) \\
n^3 + n \lg n + n\sqrt{n} & \in & O(n^4)
\end{array}$$

02-26: **Big-Oh Examples III**

$f(n) = \begin{cases} n & \text{for } n \text{ odd} \\ n^3 & \text{for } n \text{ even} \end{cases}$

$g(n) = n^2$

$$\begin{array}{rcl}
f(n) & \in & O(g(n)) \text{ ?} \\
g(n) & \in & O(f(n)) \text{ ?} \\
n & \in & O(f(n)) \text{ ?} \\
f(n) & \in & O(n^3) \text{ ?}
\end{array}$$

02-27: **Big-Oh Examples III**

$f(n) = \begin{cases} n & \text{for } n \text{ odd} \\ n^3 & \text{for } n \text{ even} \end{cases}$

$g(n) = n^2$

$$\begin{array}{rcl}
f(n) & \not\in & O(g(n)) \\
g(n) & \not\in & O(f(n)) \\
n & \in & O(f(n)) \\
f(n) & \in & O(n^3)
\end{array}$$

02-28: **Big-Ω Notation** $\Omega(f(n))$ is the set of all functions that are bound from *below* by $f(n)$

$T(n) \in \Omega(f(n))$ if

$\exists c, n_0$ such that $T(n) \geq c * f(n)$ when $n > n_0$

02-29: **Big-Ω Notation** $\Omega(f(n))$ is the set of all functions that are bound from *below* by $f(n)$

$T(n) \in \Omega(f(n))$ if

$\exists c, n_0$ such that $T(n) \geq c * f(n)$ when $n > n_0$

$$f(n) \in O(g(n)) \Rightarrow g(n) \in \Omega(f(n))$$

02-30: **Big-Θ Notation** $\Theta(f(n))$ is the set of all functions that are bound *both* above *and* below by $f(n)$. Θ is a *tight bound*

$T(n) \in \Theta(f(n))$ if

$T(n) \in O(f(n))$ and $T(n) \in \Omega(f(n))$

02-31: **Big-Oh Rules**

1. If $f(n) \in O(g(n))$ and $g(n) \in O(h(n))$, then $f(n) \in O(h(n))$
2. If $f(n) \in O(kg(n)$ for any constant $k > 0$, then $f(n) \in O(g(n))$
3. If $f_1(n) \in O(g_1(n))$ and $f_2(n) \in O(g_2(n))$, then $f_1(n) + f_2(n) \in O(\max(g_1(n), g_2(n)))$
4. If $f_1(n) \in O(g_1(n))$ and $f_2(n) \in O(g_2(n))$, then $f_1(n) * f_2(n) \in O(g_1(n) * g_2(n))$

(Also work for Ω, and hence Θ)

02-32: **Big-Oh Guidelines**

- Don't include constants/low order terms in Big-Oh
- Simple statements: $\Theta(1)$
- Loops: Θ(inside) * # of iterations
  - Nested loops work the same way
- Consecutive statements: Longest Statement
- Conditional (if) statements:
  O(Test + longest branch)

02-33: **Calculating Big-Oh**

```
for (i=1; i<n; i++)
  sum++;
```

02-34: **Calculating Big-Oh**

```
for (i=1; i<n; i++)      Executed n times
  sum++;                  O(1)
```

Running time: $O(n), \Omega(n), \Theta(n)$

02-35: **Calculating Big-Oh**

```
for (i=1; i<n; i=i+2)
  sum++;
```

02-36: **Calculating Big-Oh**

```
for (i=1; i<n; i=i+2)      Executed n/2 times
  sum++;                    O(1)
```

Running time: $O(n), \Omega(n), \Theta(n)$

02-37: **Calculating Big-Oh**

```
for (i=1; i<n; i++)
  for (j=1; j < n/2; j++)
    sum++;
```

02-38: **Calculating Big-Oh**

```
for (i=1; i<n; i++)        Executed n times
  for (j=1; j < n/2; j++)   Executed n/2 times
    sum++;                   O(1)
```

Running time: $O(n^2), \Omega(n^2), \Theta(n^2)$

02-39: **Calculating Big-Oh**

```
for (i=1; i<n; i=i*2)
  sum++;
```

02-40: **Calculating Big-Oh**

```
for (i=1; i<n; i=i*2)   Executed lg n times
  sum++;                 O(1)
```

Running Time: $O(\lg n), \Omega(\lg n), \Theta(\lg n)$

02-41: **Calculating Big-Oh**

```
for (i=0; i<n; i++)
  for (j = 0; j<i; j++)
    sum++;
```

02-42: **Calculating Big-Oh**

```
for (i=0; i<n; i++)      Executed n times
  for (j = 0; j<i; j++)   Executed <= n times
    sum++;                 O(1)
```

Running Time: $O(n^2)$. Also $\Omega(n^2)$?

02-43: **Calculating Big-Oh**

```
for (i=0; i<n; i++)
  for (j = 0; j<i; j++)
    sum++;
```

Exact # of times `sum++` is executed:

$$\begin{aligned} \sum_{i=1}^{n} i &= \frac{n(n+1)}{2} \\ &\in \Theta(n^2) \end{aligned}$$

02-44: **Calculating Big-Oh**

```
sum = 0;
for (i=0; i<n; i++)
  sum++;
for (i=1; i<n; i=i*2)
  sum++;
```

02-45: **Calculating Big-Oh**

```
sum = 0;                O(1)
for (i=0; i<n; i++)     Executed n times
  sum++;                 O(1)
for (i=1; i<n; i=i*2)   Executed lg n times
  sum++;                 O(1)
```

Running Time: $O(n), \Omega(n), \Theta(n)$

02-46: **Calculating Big-Oh**

```
sum = 0;
for (i=0; i<n; i=i+2)
  sum++;
for (i=0; i<n/2; i=i+5)
  sum++;
```

02-47: **Calculating Big-Oh**

```
sum = 0;                   O(1)
for (i=0; i<n; i=i+2)      Executed n/2 times
  sum++;                    O(1)
for (i=0; i<n/2; i=i+5)    Executed n/10 times
  sum++;                    O(1)
```

Running Time: $O(n), \Omega(n), \Theta(n)$

02-48: **Calculating Big-Oh**

```
for (i=0; i<n;i++)
  for (j=1; j<n; j=j*2)
    for (k=1; k<n; k=k+2)
      sum++;
```

02-49: **Calculating Big-Oh**

```
for (i=0; i<n;i++)         Executed n times
  for (j=1; j<n; j=j*2)     Executed lg n times
    for (k=1; k<n; k=k+2)    Executed n/2 times
      sum++;                  O(1)
```

Running Time: $O(n^2 \lg n), \Omega(n^2 \lg n), \Theta(n^2 \lg n)$

02-50: **Calculating Big-Oh**

```
sum = 0;
for (i=1; i<n; i=i*2)
  for (j=0; j<n; j++)
    sum++;
```

02-51: **Calculating Big-Oh**

```
sum = 0;                  O(1)
for (i=1; i<n; i=i*2)     Executed lg n times
  for (j=0; j<n; j++)      Executed n times
    sum++;                  O(1)
```

Running Time: $O(n \lg n), \Omega(n \lg n), \Theta(n \lg n)$

02-52: **Calculating Big-Oh**

```
sum = 0;
for (i=1; i<n; i=i*2)
  for (j=0; j<i; j++)
    sum++;
```

02-53: **Calculating Big-Oh**

```
sum = 0;                   O(1)
for (i=1; i<n; i=i*2)      Executed lg n times
  for (j=0; j<i; j++)       Executed <= n times
    sum++;                   O(1)
```

Running Time: $O(n \lg n)$. Also $\Omega(n \lg n)$?

02-54: **Calculating Big-Oh**

```
sum = 0;
for (i=1; i<n; i=i*2)
  for (j=0; j<i; j++)
    sum++;
```

# of times `sum++` is executed:

$$\begin{aligned} \sum_{i=0}^{\lg n} 2^i &= 2^{\lg n + 1} - 1 \\ &= 2n - 1 \\ &\in \Theta(n) \end{aligned}$$

02-55: **Calculating Big-Oh**

Of course, a little change can mess things up a bit ...

```
sum = 0;
for (i=1; i<=n; i=i+1)
  for (j=1; j<=i; j=j*2)
    sum++;
```

02-56: **Calculating Big-Oh**

Of course, a little change can mess things up a bit ...

```
sum = 0;
for (i=1; i<=n; i=i+1)     Executed n times
  for (j=1; j<=i; j=j*2)   Executed <= lg n times
    sum++;                 O(1)
```

So, this is code is $O(n \lg n)$ – but is it also $\Omega(n \lg n)$?

02-57: **Calculating Big-Oh**

Of course, a little change can mess things up a bit ...

```
sum = 0;
for (i=1; i<=n; i=i+1)     Executed n times
  for (j=1; j<=i; j=j*2)   Executed <= lg n times
    sum++;                 O(1)
```

Total time `sum++` is executed:

$$\sum_{i=1}^{n} \lg i$$

This can be tricky to evaluate, but we only need a bound ...

02-58: **Calculating Big-Oh**

Total # of times `sum++` is executed:

$$\begin{aligned}
\sum_{i=1}^{n} \lg i &= \sum_{i=1}^{n/2-1} \lg i + \sum_{i=n/2}^{n} \lg i \\
&\geq \sum_{i=n/2}^{n} \lg i \\
&\geq \sum_{i=n/2}^{n} \lg n/2 \\
&= n/2 \lg n/2 \\
&\in \Omega(n \lg n)
\end{aligned}$$