# The Two Set Relations Generating Geometry

### George. M. Van Treeck

ABSTRACT. A ruler (measuring stick) partitions both domain and range intervals approximately into sets of same-sized subintervals. As the subinterval size converges to zero: 1) The union of range sets of subintervals, where each domain set corresponds to a less or equal-sized range set, generates the properties of metric space and all Lp norms, in particular, Manhattan and Euclidean distance. 2) The domain and range set constraints that generates Manhattan distance also generates Euclidean volume. The proofs allow simpler derivations of Coulomb's charge force, Newton's gravity force, and spacetime equations without using other laws of physics or Gauss's divergence theorem. Time limits physical distance and volume to 3 dimensions. All proofs are verified in Coq.

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### 1. Introduction

Metric space, the Euclidean distance metric/vector norm, and Euclidean area/volume in the Lebesgue, Borel, and Hausdorff measures have been definitions [Gol76] [Rud76] rather than derived from more fundamental set definitions. In this article, both distance and volume are defined as the cardinals of set operations. A ruler (measuring stick) measures interval lengths by partitioning all domain and range intervals approximately into sets of size c subintervals and summing the sizes. The distance and volume set operations on the sets of subintervals converge to the properties of metric space, the  $L^p$  norms, and Euclidean volume as  $c \to 0$ .

<sup>2010</sup> Mathematics Subject Classification. Primary 28A75, 28E15. Secondary 03E75, 51M99. Copyright © 2020 George M. Van Treeck. Creative Commons Attribution License.

The derivations of metric space, distance equations, and volume from the set operations provide some insights into geometry and physics, for example: the constraint on the union operation generating the properties of metric space; the mapping between sets that makes Euclidean distance the smallest possible distance between two distinct points in  $\mathbb{R}^n$ ; the domain and range set constraints that generate both Manhattan distance and Euclidean volume; how the Euclidean distance and volume proofs allow deriving Coulomb's charge force, Newton's gravity force, and space time equations without using other laws of physics or Gauss's divergence theorem; the proportionate interval principle generates the inverse square law; how time places an additional constraint on physical sets, which limits physical distance and volume to 3 dimensions.

All the proofs in this article have been formally verified using using the Coq proof verification system [Coq15]. The formal proofs are located in the Coq files, "euclidrelations.v" and "threed.v," at: https://github.com/treeck/RASRGeometry.

### 2. Ruler measure and convergence

DEFINITION 2.1. Ruler measure: A ruler measures the size, M, of a closed, open, or semi-open interval as the sum of the sizes of the nearest integer number of whole subintervals, p, each subinterval having the same size, c. Notionally:

(2.1) 
$$\forall c, s \in \mathbb{R}, [a, b] \subset \mathbb{R}, s = |a - b| \land c > 0 \land$$

$$(p = floor(s/c) \lor p = ceiling(s/c)) \land M = \sum_{i=1}^{p} c = pc.$$

Theorem 2.2. Ruler convergence:  $\forall [a,b] \subset \mathbb{R}, \ s = |a-b| \Rightarrow M = \lim_{c \to 0} pc = s.$ 

The theorem, "limit\_c\_0\_M\_eq\_exact\_size," and formal proof is in the Coq file,

euclidrelations.v.

Proof. (epsilon-delta proof)

By definition of the floor function,  $floor(x) = max(\{y : y \le x, y \in \mathbb{Z}, x \in \mathbb{R}\})$ :

$$(2.2) \hspace{0.1in} \forall \hspace{0.1in} c>0, \hspace{0.1in} p=floor(s/c) \hspace{0.1in} \wedge \hspace{0.1in} 0 \leq |floor(s/c)-s/c| < 1 \hspace{0.1in} \Rightarrow \hspace{0.1in} 0 \leq |p-s/c| < 1.$$

Multiply all sides of inequality 2.2 by |c|:

$$(2.3) \hspace{1cm} \forall \hspace{0.1cm} c>0, \quad 0 \leq |p-s/c| < 1 \quad \Rightarrow \quad 0 \leq |pc-s| < |c|.$$

$$\begin{array}{lll} (2.4) & \forall \ \delta \ : \ |pc-s| < |c| = |c-0| < \delta \\ & \Rightarrow & \forall \ \epsilon = \delta : \ |c-0| < \delta \ \land \ |pc-s| < \epsilon \ := \ M = \lim_{c \to 0} pc = s. \end{array} \ \Box$$

The proof steps using the ceiling function (the outer measure) are the same as the steps in the previous proof using the floor function (the inner measure). The following is an example of ruler convergence for the interval,  $[0,\pi]$ :  $s=|0-\pi|$ , and  $p=floor(s/c) \Rightarrow p \cdot c = 3.1_{c=10^{-1}}, \ 3.14_{c=10^{-2}}, \ 3.141_{c=10^{-3}}, ..., \pi_{\lim_{c\to 0}}$ .

### 3. Distance

**Notation convention:** Vertical bars around a set or list,  $|\cdots|$ , indicates the cardinal (number of members in the set or list).

**3.1. Countable distance space.** An concrete example of a countable distance is the number of steps walked in a range (distance) set, y, which equals the number of pieces of land in a corresponding domain set, x: |x| = |y|. Generalizing, each disjoint domain set,  $x_i$ , has a corresponding range (distance) set,  $y_i$ . Therefore, the countable distance spanning the disjoint domain sets is the number of members,  $d_c$ , in the union range set:

Definition 3.1. Countable distance space,  $d_c$ :

$$\bigcap_{i=1}^{n} x_i = \emptyset \quad \land \quad d_c = \left| \bigcup_{i=1}^{n} y_i \right| \quad \land \quad |x_i| = |y_i|^q, \ q \ge 1.$$

THEOREM 3.2. Inclusion-exclusion Inequality:  $|\bigcup_{i=1}^n y_i| \leq \sum_{i=1}^n |y_i|$ .

The inclusion-exclusion inequality follows from the inclusion-exclusion principle [CG15]. But, a more intuitive and simple proof follows from the associative law of addition where the sum of set sizes is equal to the size of all the set members appended into a list and the commutative law of addition that allows sorting that list into a list of unique members (the *union* set) and a list of duplicates. The duplicates being  $\geq 0$  implies the union size is always  $\leq$  the sum of set sizes.

A formal proof, inclusion\_exclusion\_inequality, using sorting into a set of unique members (*union* set) and a list of duplicates, is in the file euclidrelations.v.

Proof.

(3.1) 
$$\sum_{i=1}^{n} |y_i| = |append_{i=1}^n y_i| = |sort(append_{i=1}^n y_i)|$$
$$= |\bigcup_{i=1}^{n} y_i| + |duplicates_{i=1}^n y_i|.$$

(3.2) 
$$|\bigcup_{i=1}^{n} y_i| + |duplicates_{i=1}^{n} y_i| = \sum_{i=1}^{n} |y_i| \wedge |duplicates_{i=1}^{n} y_i| \ge 0$$
  
 $\Rightarrow |\bigcup_{i=1}^{n} y_i| \le \sum_{i=1}^{n} |y_i|. \square$ 

**3.2.** Metric Space. All function range intervals, d(u, w), satisfying the countable distance space definition,  $d_c = |\bigcup_{i=1}^n y_i|$ , where the ruler is applied, generates three of the four metric space properties: triangle inequality, non-negativity, and identity of indiscernibles. The set-based reason for the fourth property of metric space, symmetry [d(u, v) = d(v, u)], will be identified in the last section of this article. The formal proofs: triangle\_inequality, non\_negativity, and identity\_of\_indiscernibles are in the Coq file, euclidrelations.v.

THEOREM 3.3. Triangle Inequality: 
$$d_c = |y_1 \cup y_2| \Rightarrow d(u, w) \leq d(u, v) + d(v, w)$$
.

PROOF. Apply the ruler measure (2.1), the countable distance space condition (3.1), inclusion-exclusion inequality (3.2), and then ruler convergence (2.2).

$$(3.3) \quad \forall c > 0, \ d(u, w), \ d(u, v), \ d(v, w) :$$

$$|y_1| = floor(d(u, v)/c) \quad \land \quad |y_2| = floor(d(v, w)/c) \quad \land$$

$$d_c = floor(d(u, w)/c) \quad \land \quad d_c = |y_1 \cup y_2| \le |y_1| + |y_2|$$

$$\Rightarrow floor(d(u, w)/c) \le floor(d(u, v)/c) + floor(d(v, w)/c)$$

$$\Rightarrow floor(d(u, w)/c) \cdot c \le floor(d(u, v)/c) \cdot c + floor(d(v, w)/c) \cdot c$$

$$\Rightarrow \lim_{c \to 0} floor(d(u, w)/c) \cdot c \le \lim_{c \to 0} floor(d(u, v)/c) \cdot c + \lim_{c \to 0} floor(d(v, w)/c) \cdot c$$

$$\Rightarrow d(u, w) \le d(u, v) + d(v, w). \quad \Box$$

THEOREM 3.4. Non-negativity:  $d_c = |y_1 \cup y_2| \Rightarrow d(u, w) \geq 0$ .

PROOF. By definition, a set always has a size (cardinal)  $\geq 0$ :

$$(3.4) \quad \forall \ c > 0, \ d(u, w) : \quad floor(d(u, w)/c) = d_c \quad \land \quad d_c = |y_1 \cup y_2| \ge 0$$

$$\Rightarrow \quad floor(d(u, w)/c) = d_c \ge 0 \quad \Rightarrow \quad d(u, w) = \lim_{c \to 0} d_c \cdot c \ge 0. \quad \Box$$

Theorem 3.5. Identity of Indiscernibles: d(w, w) = 0.

PROOF. Apply the triangle inequality property (3.3):

$$(3.5) \quad \forall \ d(u,v) = d(v,w) = 0 \ \land \ d(u,w) \le d(u,v) + d(v,w) \ \Rightarrow \ d(u,w) \le 0.$$

Combine the non-negativity property (3.4) and the previous inequality (3.5):

$$(3.6) d(u, w) \ge 0 \wedge d(u, w) \le 0 \Leftrightarrow 0 \le d(u, w) \le 0 \Rightarrow d(u, w) = 0.$$

Combine the result of step 3.6 and the condition, d(u, v) = 0, in step 3.5.

(3.7) 
$$d(u, w) = 0 \land d(u, v) = 0 \Rightarrow w = v.$$

Combine the condition, d(v, w) = 0, in step 3.5 and the result of step 3.7.

$$(3.8) d(v,w) = 0 \wedge w = v \Rightarrow d(w,w) = 0.$$

**3.3. Distance space range.** From the countable distance space definition,  $d_c = |\bigcup_{i=1}^n y_i|$ , as the amount of intersection increases, more domain set members can map to a single range set member. Therefore, the number of domain-to-range set member mappings is a function of the amount of range set intersection.

From the countable distance space property (3.1), where  $|x_i| = |y_i| = p_i$ , the range of possible domain-to-range set member mappings is the number of one-to-one (bijective),  $1 \cdot |y_i| = p_i$ , mappings, to the number of many-to-many,  $|x_i| \cdot |y_i| = p_i^2$ , mappings. Therefore, the total number of domain-to-range set mappings varies from  $\sum_{i=1}^{n} (1 \cdot |y_i|) = \sum_{i=1}^{n} p_i$  to  $\sum_{i=1}^{n} (|y_i| \cdot |x_i|) = \sum_{i=1}^{n} p_i^2$  mappings. Applying the ruler (2.1) and ruler convergence theorem (2.2) to the smallest and largest total number of domain-to-range set mapping cases converges to the real-valued Manhattan and Euclidean distance relations.

#### 3.4. Manhattan distance.

Theorem 3.6. Manhattan (largest) distance, d, is the size of the range interval,  $[d_0, d_m]$ , corresponding to a set of disjoint domain intervals,

 $\{[a_1,b_1],[a_2,b_2],\ldots,[a_n,b_n]\}, where:$ 

$$d = \sum_{i=1}^{n} s_i$$
,  $d = |d_0 - d_m|$ ,  $s_i = |a_i - b_i|$ .

The theorem, "taxicab\_distance," and formal proof is in the Coq file, euclidrelations.v.

Proof.

From the countable distance space definition (3.1) and the inclusion-exclusion inequality (3.2), the largest possible countable distance,  $d_c$ , is the equality case:

(3.9) 
$$d_c = |\bigcup_{i=1}^n y_i| \le \sum_{i=1}^n |y_i| \wedge |x_i| = |y_i| = p_i$$
  
 $\Rightarrow d_c \le \sum_{i=1}^n |y_i| = \sum_{i=1}^n p_i \Rightarrow \exists p_i, d_c : d_c = \sum_{i=1}^n p_i.$ 

Multiply both sides of equation 3.11 by c and take the limit:

(3.10) 
$$d_c = \sum_{i=1}^n p_i \Rightarrow d_c \cdot c = \sum_{i=1}^n (p_i \cdot c) \Rightarrow \lim_{c \to 0} d_c \cdot c = \sum_{i=1}^n \lim_{c \to 0} (p_i \cdot c).$$

Apply the ruler (2.1) and ruler convergence theorem (2.2) to the definition of d:

$$(3.11) d = |d_0 - d_m| \Rightarrow \exists c d: floor(d/c) = d_c \Rightarrow d = \lim_{c \to 0} d_c \cdot c.$$

Apply the ruler (2.1) and ruler convergence theorem (2.2) to each domain interval:

Combine equations 3.11, 3.10, 3.12:

(3.13) 
$$d = \lim_{c \to 0} d_c \cdot c \quad \wedge \quad \lim_{c \to 0} d_c \cdot c = \sum_{i=1}^n \lim_{c \to 0} (p_i \cdot c) \quad \wedge \\ \lim_{c \to 0} (p_i \cdot c) = s_i \quad \Rightarrow \quad d = \lim_{c \to 0} d_c \cdot c = \sum_{i=1}^n \lim_{c \to 0} (p_i \cdot c) = \sum_{i=1}^n s_i. \quad \Box$$

### 3.5. Euclidean distance.

THEOREM 3.7. Euclidean (smallest) distance, d, is the size of the range interval,  $[d_0, d_m]$ , corresponding to a set of disjoint domain intervals,

$$\{[a_1,b_1],[a_2,b_2],\ldots,[a_n,b_n]\}, where:$$

$$d^2 = \sum_{i=1}^n s_i^2$$
,  $d = |d_0 - d_m|$ ,  $s_i = |a_i - b_i|$ .

The theorem, "Euclidean\_distance," and formal proof is in the Coq file, euclidrelations.v.  $\,$ 

### Proof.

Apply the rule of product to the largest number of domain-to-range set mappings, where all  $p_i$  number of range set members,  $y_i$ , map to each of the  $p_i$  number of members in the domain set,  $x_i$ , which is the Cartesian product,  $|y_i| \cdot |x_i|$ :

$$(3.14) |x_i| = |y_i| = p_i \quad \Rightarrow \quad \sum_{i=1}^n |y_i| \cdot |x_i| = \sum_{i=1}^n p_i^2.$$

From the countable distance space definition (3.1) and the inclusion-exclusion inequality (3.2), choose the equality case:

Square both sides of equation 3.15  $(x = y \Leftrightarrow f(x) = f(y))$ :

$$(3.16) \exists p_i, d_c : d_c = \sum_{i=1}^n p_i \quad \Leftrightarrow \quad \exists p_i, d_c : d_c^2 = (\sum_{i=1}^n p_i)^2.$$

Apply the square of sum inequality,  $(\sum_{i=1}^{n} p_i)^2 \ge \sum_{i=1}^{n} p_i^2$ , to equation 3.16 and select the smallest area (the equality) case:

(3.17) 
$$d_c^2 = \left(\sum_{i=1}^n p_i\right)^2 = \sum_{i=1}^n p_i \sum_{j=1}^n p_j$$

$$= \sum_{i=1}^n p_i^2 + \sum_{i=1}^n p_i \sum_{j=1, j \neq i}^n p_j \ge \sum_{i=1}^n p_i^2 \quad \Rightarrow \quad \exists \ p_i : d_c^2 = \sum_{i=1}^n p_i^2.$$

Multiply both sides of equation 3.17 by  $c^2$ , simplify, and take the limit.

(3.18) 
$$d_c^2 = \sum_{i=1}^n p_i^2 \implies d_c^2 \cdot c^2 = \sum_{i=1}^n p_i^2 \cdot c^2 \iff (d_c \cdot c)^2 = \sum_{i=1}^n (p_i \cdot c)^2$$
  
 $\implies \lim_{c \to 0} (d_c \cdot c)^2 = \sum_{i=1}^n \lim_{c \to 0} (p_i \cdot c)^2.$ 

Apply the ruler (2.1) and ruler convergence theorem (2.2) and square both sides:

$$(3.19) \quad \exists \ c \ d \in \mathbb{R}: \ floor(d/c) = d_c \quad \Rightarrow \quad d = \lim_{c \to 0} d_c \cdot c \quad \Rightarrow \quad d^2 = \lim_{c \to 0} (d_c \cdot c)^2.$$

Apply the ruler (2.1) and ruler convergence theorem (2.2) to each domain interval:

$$(3.20) \ \ s_i = |a_i - b_i| \quad \land \quad floor(s_i/c) = |x_i| = |y_i| = p_i \quad \Rightarrow \quad \lim_{c \to 0} p_i \cdot c = s_i.$$

Combine equations 3.19, 3.18, 3.20:

(3.21) 
$$d^2 = \lim_{c \to 0} (d_c \cdot c)^2 \wedge \lim_{c \to 0} (d_c \cdot c)^2 = \sum_{i=1}^n \lim_{c \to 0} (p_i \cdot c)^2 \wedge \lim_{c \to 0} (p_i \cdot c) = s_i \Rightarrow d^2 = \lim_{c \to 0} (d_c \cdot c)^2 = \sum_{i=1}^n \lim_{c \to 0} (p_i \cdot c)^2 = \sum_{i=1}^n s_i^2. \square$$

### 4. Euclidean Volume

The case of disjoint range sets and a 1-1 correspondence of domain-to-range set mappings,  $\bigcap_{i=1}^n x_i = \emptyset \quad \land \quad \bigcap_{i=1}^n y_i = \emptyset \quad \land \quad d_c = |\bigcup_{i=1}^n y_i| \quad \land \quad |x_i| = |y_i|$ , generates Manhattan distance (3.6). Each sub-Manhattan distance is the sum of a unique n-tuple of subintervals, one size c subinterval from each range set,  $y_i$ . There are multiple possible unique sub-Manhattan distance n-tuples (paths) between the same two distinct points, A and B, that sum to the same Manhattan distance as  $c \to 0$ . The countable volume,  $v_c$ , is the number all possible unique sub-Manhattan distance n-tuples, which is the number of n-tuples generated by Cartesian product of the range sets:

Definition 4.1. Countable Volume,  $v_c$ :

$$\bigcap_{i=1}^{n} x_i = \emptyset \quad \land \quad \bigcap_{i=1}^{n} y_i = \emptyset \quad \land \quad v_c = |\times_{i=1}^{n} y_i| \quad \land \quad |x_i| = |y_i|.$$

THEOREM 4.2. Euclidean volume, v, is size of the range interval,  $[v_0, v_m]$ , corresponding to the Cartesian product of all the members of the domain intervals,  $\{[a_1, b_1], [a_2, b_2], \ldots, [a_n, b_n]\}$ . Notionally:

$$v = \prod_{i=1}^{n} s_i, \ v = |v_0 - v_m|, \ s_i = |a_i - b_i|.$$

The theorem, "Euclidean\_volume," and formal proof is in the Coq file, euclidrelations.v.

Proof.

Use the ruler (2.1) to partition each of the domain intervals,  $[a_i, b_i]$ , into a set,  $x_i$ , containing  $p_i$  number of subintervals.

$$(4.1) \forall i \ n \in \mathbb{N}, \quad i \in [1, n], \quad c > 0 \quad \land \quad floor(s_i/c) = p_i = |x_i| = |y_i|.$$

Apply the ruler convergence theorem (2.2) to equation 4.1:

$$(4.2) floor(s_i/c) = p_i \Rightarrow \lim_{c \to 0} (p_i \cdot c) = s_i.$$

Apply the associative law of multiplication to derive the countable volume (4.1) in terms of  $p_i$ :

$$(4.3) v_c = |\times_{i=1}^n y_i| \wedge |y_i| = p_i \Rightarrow v_c = |\times_{i=1}^n y_i| = \prod_{i=1}^n |y_i| = \prod_{i=1}^n p_i.$$

Multiply both sides of equation 4.3 by  $c^n$ :

$$(4.4) v_c \cdot c^n = (\prod_{i=1}^n p_i) \cdot c^n = \prod_{i=1}^n (p_i \cdot c).$$

Use those cases, where  $v_c$  has an integer  $n^{th}$  root.

$$(4.5) \quad \forall n, p, v_c \in \mathbb{N} : p^n = v_c \Rightarrow v_c \cdot c^n = p^n \cdot c^n = (p \cdot c)^n = \prod_{i=1}^n (p_i \cdot c).$$

Apply the ruler (2.1) and ruler convergence (2.2) to the range interval,  $[v_0, v_m]$  (where  $v = |v_0 - v_m|$ ), and then combine with equations 4.5 and 4.2:

(4.6) 
$$floor(v/c^n) = p^n \Rightarrow v = \lim_{c \to 0} (p \cdot c)^n = \prod_{i=1}^n \lim_{c \to 0} (p_i \cdot c) = \prod_{i=1}^n s_i.$$

## 5. Applications to physics

**5.1.** Coulomb's charge force. The sizes,  $q_1$  and  $q_2$ , of two charges are independent domain variables, where each size, c, component of a charge exerts a force on each same size, c, component of the other charge. The total force, F, is proportionate to the total number of forces (the Cartesian product of the number of same-sized, infinitesimal components) multiplied times a quantum charge force,  $m_{C}a_{C}$ . From the volume proof (4.2), the Cartesian product converges to  $q_1q_2$ :

(5.1) 
$$F \propto (m_C a_C) (\lim_{c \to 0} floor(q_1/c) \cdot c) (\lim_{c \to 0} floor(q_2/c) \cdot c) =$$

$$(m_C a_C) \iint dq_1 dq_2 = (m_C a_C) (q_1 q_2).$$

From equation 5.1, an increase in charge, q, causes a proportionate increase in force, F. Solving for the constant force,  $F = m_C a_C$ , requires a proportionate variable, r, to offset the effect of a change in q:  $r \propto q \Rightarrow \exists q_C/r_C \in \mathbb{R} : r(q_C/r_C) = q$ :

$$(5.2) \quad \forall \ q_1, q_2 \ge 0 \ \exists \ q \in \mathbb{R} \ : q^2 = q_1 q_2 \quad \land \quad r(q_C/r_C) = q \quad \Rightarrow \quad (r(q_C/r_C))^2 = q_1 q_2.$$

$$(5.3) \quad (r(q_C/r_C))^2 = q_1 q_2 \quad \land \quad F \propto (m_C a_C)(q_1 q_2)$$

$$\Rightarrow \quad F \propto (m_C a_C)(r(q_C/r_C))^2 = (m_C a_C)(q_1 q_2)$$

$$\Rightarrow \quad F = m_C a_C = (m_C a_C r_C^2/q_C^2)q_1 q_2/r^2 = k_c q_1 q_2/r^2.$$

where  $k_C = m_C a_C r_C^2/q_C^2$  corresponds to the SI units:  $Nm^2C^{-2}$ .

**5.2.** Newton's gravity force equation. The sizes,  $m_1$  and  $m_2$ , of two masses are independent domain variables, where each size, c, component of a mass exerts a force on each same size, c, component of the other mass. The total force, F, is proportionate to the total number of forces (the Cartesian product of the number of same-sized, infinitesimal components) multiplied times a quantum gravity force,  $m_{G}a_{G}$ . From the volume proof (4.2), the Cartesian product converges to  $m_{1}m_{2}$ :

(5.4) 
$$F \propto (m_G a_G) (\lim_{c \to 0} floor(m_1/c) \cdot c) (\lim_{c \to 0} floor(m_2/c) \cdot c) =$$

$$(m_C a_C) \iint dm_1 dm_2 = (m_G a_G) (m_1 m_2).$$

From equation 5.4, an increase in mass, m, causes a proportionate increase in force, F. Solving for the constant force,  $F = m_G a_G$ , requires a proportionate variable, r, to offset the effect of a change in m:  $r \propto m \Rightarrow \exists m_G/r_G \in \mathbb{R} : r(m_G/r_G) = m$ :

(5.5) 
$$\forall m_1, m_2 \ge 0 \exists m \in \mathbb{R} : m^2 = m_1 m_2 \land r(m_G/r_G) = m$$
  
 $\Rightarrow (r(m_G/r_G))^2 = m_1 m_2.$ 

(5.6) 
$$(r(m_G/r_G))^2 = m_1 m_2 \wedge F \propto (m_G a_G)(m_1 m_2)$$
  
 $\Rightarrow F \propto (m_G a_G)(r(m_G/r_G))^2 = (m_G a_G)(m_1 m_2)$   
 $\Rightarrow F = m_G a_G = (m_G a_G r_G^2/q_G^2)m_1 m_2/r^2.$ 

(5.7) 
$$\exists t_G \in \mathbb{R} : r_G/t_G^2 = a_G \land F = m_G a_G = (m_G a_G r_G^2/m_G^2) m_1 m_2/r^2$$
  
 $\Rightarrow F = m_G a_G = (r_G^3/m_G t_G^2) m_1 m_2/r^2 = G m_1 m_2/r^2,$ 

where  $G = r_G^3/m_G t_G^2$  corresponds to the SI units:  $m^3 kg^{-1}s^{-2}$ .

**5.3. Spacetime equations.** Applying the ruler measure, if sequencing across each same-sized subinterval of a *physical*, Euclidean distance (range) interval, [0, r], corresponds to a proportionate number of same-sized subintervals of a time interval, [0, t], then, as the subinterval size converges to zero, the interval, [0, t], is proportionate to the range interval, [0, r], where there is a conversion constant, c, that is the ratio of some value,  $r_c$  to some value,  $t_c$ , such that  $r = (r_c/t_c)t = ct$ .

Applying the ruler, to two intervals,  $[0, d_1]$  and  $[0, d_2]$ , in two inertial (independent, non-accelerating) frames of reference, the distance and time to sequence over the subintervals the two intervals converges to a range of distances (and times) from Manhattan (3.6) to Euclidean distance (3.7).

(5.8) 
$$r^2 = d_1^2 + d_2^2 \quad \land \quad r = (r_c/t_c)t = ct$$
  

$$\Rightarrow \quad (ct)^2 = d_1^2 + d_2^2 \quad \Rightarrow \quad d_2 = \sqrt{(ct)^2 - d_1^2}.$$

(5.9) 
$$d_2 = \sqrt{(ct)^2 - d_1^2} \quad \land \quad d = d_2 \quad \land \quad d_1 = vt$$
  
 $\Rightarrow \quad d = \sqrt{(ct)^2 - (vt)^2} = ct\sqrt{1 - (v^2/c^2)},$ 

which is the spacetime dilation equation. [Bru17].

(5.10) 
$$d_2^2 = (ct)^2 - d_1^2 \quad \land \quad s = d_2 \quad \land \quad d_1^2 = x^2 + y^2 + z^2$$
  

$$\Rightarrow \quad s^2 = (ct)^2 - (x^2 + y^2 + z^2),$$

which is one form of the spacetime interval equation [Bru17].

**5.4. 3 dimensions of physical geometry.** The set and arithmetic operations used to calculate distance and volume requires sequencing through a totally ordered set of dimensions, for example, the countable distance space:  $d_c = |\bigcup_{i=1}^n y_i|$ , Euclidean distance:  $d^2 = \sum_{i=1}^n s_i^2$ , countable volume:  $v_c = |\times_{i=1}^n x_i|$ , and Euclidean volume:  $v = \prod_{i=1}^n s_i$ . The commutative property of the union, addition, and multiplication operations also allows sequencing through a set of n number of dimensions in all n! number of possible orders.

From the derivation of the spacetime equations (5.3), physical sets have the additional constraint that sequencing across the members of a non-empty set takes some greater than zero amount of time. Determining that a physical sequencer sequenced a physical set in the order,  $[x_5, x_4, \dots, x_1]$ , and next sequenced in the order,  $[x_1, x_2, \dots, x_5]$ , requires the total order, at most one successor and at most one predecessor per set member, to not change during the time of the sequencing.

Note that to "jump" to any member of an ordered, *physical* set requires calculating offsets, which is implicitly counting (sequencing through) successor/predecessor links to reach the indexed member, where the calculations take some greater than zero amount of time. Deterministic sequencing of a totally ordered, *physical* set via successor/predecessor links in each possible order requires each set member to be either a successor or predecessor to every other set member (sequentially adjacent) during the *time* of sequencing, herein referred to as a symmetric geometry.

It will now be proved that a set satisfying the constraints of a single total order and also symmetric defines a cyclic set containing at most 3 members, in this case, 3 dimensions of physical space.

Definition 5.1. Ordered geometry:

$$\forall i \ n \in \mathbb{N}, \ i \in [1, n-1], \ \forall x_i \in \{x_1, \dots, x_n\},$$
$$successor \ x_i = x_{i+1} \ \land \ predecessor \ x_{i+1} = x_i.$$

DEFINITION 5.2. Symmetric geometry (every set member is sequentially adjacent to every other member):

$$\forall i \ j \ n \in \mathbb{N}, \ \forall \ x_i \ x_j \in \{x_1, \dots, x_n\}, \ successor \ x_i = x_j \ \Leftrightarrow \ predecessor \ x_j = x_i.$$

Theorem 5.3. An ordered and symmetric set is a cyclic set.

$$successor x_n = x_1 \land predecessor x_1 = x_n.$$

The theorem, "ordered\_symmetric\_is\_cyclic," and formal proof is in the Coq file, threed.v.

PROOF. A total order (5.1) defines unique successors and predecessors for all set members except for the successor of  $x_n$  and the predecessor of  $x_1$ . Therefore, the only member that can be a successor of  $x_n$ , without creating a contradiction, is  $x_1$ . And the only member that can be a predecessor of  $x_1$ , without creating a contradiction, is  $x_n$ . From the properties of a symmetric geometry (5.2):

$$(5.11) i = n \land j = 1 \land successor x_i = x_j \Rightarrow successor x_n = x_1.$$

Applying the definition of a symmetric geometry (5.2) to conclusion 5.11:

(5.12) 
$$i = n \land j = 1 \land predecessor x_j = x_i \Rightarrow predecessor x_1 = x_n.$$

Theorem 5.4. An ordered and symmetric set is limited to at most 3 members.

The lemmas and formal proofs in the Coq file threed.v are:

Lemmas: adj111, adj122, adj212, adj123, adj133, adj233, adj213, adj313, adj323, and not\_all\_mutually\_adjacent\_gt\_3.

The following proof uses Horn clauses (a subset of first-order logic) that uses unification and resolution. Horn clauses make it clear which facts satisfy a proof goal.

PROOF.

It was proved that an ordered and symmetric set is a cyclic set (5.3). In other words, the successors and predecessors of an ordered and symmetric set are cyclic:

Definition 5.5. Cyclic successor of m is n:

$$(5.13) \ Successor(m, n, setsize) \leftarrow (m = setsize \land n = 1) \lor (n = m + 1 \le setsize).$$

Definition 5.6. Cyclic predecessor of m is n:

$$(5.14) \quad Predecessor(m, n, setsize) \leftarrow (m = 1 \land n = setsize) \lor (n = m - q \ge 1).$$

DEFINITION 5.7. Adjacent: member m is sequentially adjacent to member n if the cyclic successor of m is n or the cyclic predecessor of m is n. Notionally: (5.15)

 $Adjacent(m, n, setsize) \leftarrow Successor(m, n, setsize) \lor Predecessor(m, n, setsize).$ 

Prove that every member is adjacent to every other member, where  $setsize \in \{1, 2, 3\}$ :

$$(5.16) Adjacent(1,1,1) \leftarrow Successor(1,1,1) \leftarrow (m = setsize \land n = 1).$$

$$(5.17) Adjacent(1,2,2) \leftarrow Successor(1,2,2) \leftarrow (n=m+1 \leq setsize).$$

$$(5.18) \hspace{1cm} Adjacent(2,1,2) \leftarrow Successor(2,1,2) \leftarrow (n = setsize \land m = 1).$$

$$(5.19) \qquad Adjacent(1,2,3) \leftarrow Successor(1,2,3) \leftarrow (n=m+1 \leq setsize).$$

$$(5.20) \qquad Adjacent(2,1,3) \leftarrow Predecessor(2,1,3) \leftarrow (n=m-q \geq 1).$$

$$(5.21) \qquad Adjacent(3,1,3) \leftarrow Successor(3,1,3) \leftarrow (n = setsize \land m = 1).$$

$$(5.22) Adjacent(1,3,3) \leftarrow Predecessor(1,3,3) \leftarrow (m=1 \land n=setsize).$$

$$(5.23) Adjacent(2,3,3) \leftarrow Successor(2,3,3) \leftarrow (n=m+1 \leq setsize).$$

$$(5.24) Adjacent(3,2,3) \leftarrow Predecessor(3,2,3) \leftarrow (n=m-q \ge 1).$$

Must prove that for all setsize > 3, there exist non-adjacent members. For example, the first and third members are not  $(\neg)$  adjacent:

(5.25) 
$$\forall setsize > 3: \neg Successor(1, 3, setsize > 3) \\ \leftarrow Successor(1, 2, setsize > 3) \leftarrow (n = m + 1 < setsize).$$

That is, member 2 is the only successor of member 1 for all setsize > 3, which implies member 3 is not a successor of member 1 for all setsize > 3.

(5.26) 
$$\forall \ set size > 3: \neg Predecessor(1, 3, set size > 3) \\ \leftarrow Predecessor(1, set size, set size > 3) \leftarrow (m = 1 \land n = set size > 3).$$

That is, member n = set size > 3 is the only predecessor of member 1, which implies member 3 is not a predecessor of member 1 for all set size > 3.

(5.27) 
$$\forall setsize > 3: \neg Adjacent(1, 3, setsize > 3)$$
  
 $\leftarrow \neg Successor(1, 3, setsize > 3) \land \neg Predecessor(1, 3, setsize > 3). \square$ 

That is, for all setsize > 3, some elements are not sequentially adjacent to every other element (not symmetric).

## 6. Insights and implications

Applying the ruler measure (2.1) and ruler convergence (2.2) to the set relations, countable distance space (3.1) and countable volume (4.1) yields the following insights and implications:

- (1) The properties of metric space, Euclidean distance and area/volume are derived from two set relations without using the notions of Euclidean geometry like plane, side, angle, perpendicular, congruence, intersection, etc. [Joy98], which makes analysis a more self-contained discipline and also provides a formal (set-based) understanding of geometry.
- (2) Making the measures of countable distance (3.1) and volume (4.1) independent of the measure of intervals (2.1) allows deriving relations, equations, and insights into geometry and physics that metric space, the Lebesgue, Borel, and Hausdorff measures are incapable of providing.

(3) The case of countable distance (3.1), where  $\bigcap_{i=1}^{n} y_i = \emptyset$  and  $|x_i| = |y_i|$ , converges to the Manhattan distance. Countable volume (4.1) as the cardinal of the set of all possible unique sub-Manhattan distance n-tuples is a set-based pattern relating volume to distance:

 $\bigcap_{i=1}^{n} x_{i} = \emptyset \quad \land \quad \bigcap_{i=1}^{n} y_{i} = \emptyset \quad \land \quad d_{c} = |\bigcup_{i=1}^{n} y_{i}| \quad \land \quad |x_{i}| = |y_{i}| \Rightarrow \bigcap_{i=1}^{n} x_{i} = \emptyset \quad \land \quad \bigcap_{i=1}^{n} y_{i} = \emptyset \quad \land \quad v_{c} = |\times_{i=1}^{n} y_{i}| \quad \land \quad |x_{i}| = |y_{i}|,$  that metric space, the Lebesgue, Borel, and Hausdorff measures have not shown.

- (4) The largest intersection (smallest distance),  $d_c = |\bigcup_{i=1}^n y_i|$ , has the most domain set members mapping to each range set member. The equality case,  $|x_i| = |y_i| = p_i$ , of the countable distance space constraint,  $|x_i| = |y_i|^q$ ,  $q \ge 1$ , (3.1) limits the largest total number of domain-to-range set mappings to  $\sum_{i=1}^n |x_i| \cdot |y_i| = \sum_{i=1}^n p_i^2$ , which is the set-based reason Euclidean distance (3.7) is the smallest possible distance between two distinct points in  $\mathbb{R}^n$  (flat space, where  $|x_i| = |y_i|$ ).
- (5)  $|x_i| = |y_i|^q$ ,  $q \ge 1$ , (3.1) generates all the  $L^p$  norms,  $||L||_p = (\sum_{i=1}^n s_i^p)^{1/p}$ . For example, using the same proof pattern as for Euclidean distance (3.7):  $p_i = |y_i| \Rightarrow |x_i| = p_i^q \Rightarrow \sum_{i=1}^n |x_i| \cdot |y_i| = \sum_{i=1}^n p_i^{q+1} \le d_c^{q+1} \dots$
- (6) It was shown that the countable distance set operation,  $d_c = |\bigcup_{i=1}^n y_i|$ , (3.1) generates three of the metric space properties. The countable distance constraint,  $|x_i| = |y_i|^q$ ,  $q \ge 1$ , (3.1) is the reason for the fourth property of metric space, symmetry, d(u,v) = d(v,u), where the value, q, causes the same combinatorial domain-to-range set mapping for every domain-range set pair that manifests as the same value, p, in each term of the  $L^p$  norm.
- (7) The Euclidean volume proof was used to derive the Coulomb's charge force (5.1) and Newton's gravity force (5.4) without using laws of physics or Gauss's divergence theorem. The Euclidean distance proof was used to derive the spacetime equations (5.8) without a constant speed of light assumption or even the notion of light.
- (8) The Proportionate Interval Principle: The derivations of the charge force, gravity force, and spacetime equations shows that all Euclidean distance range intervals having a size, r, have proportionately sized intervals of other types, for example:  $r = (r_C/q_C)q = (r_G/m_G)m = (r_c/t_c)t$ .
  - (a) From the derivations of the charge and gravity forces, the proportionate interval principle generates the inverse square law, which is a mapping of Cartesian product (rectangular) geometric area ( $r^2$ ) to rectangular charge area ( $q_1q_2$ ) and rectangular gravity area ( $m_1m_2$ ).
    - (i) Contrary to the common belief, for fields, the size of the vector normal to the surface of a sphere is a function of the rectangular area  $(r^2)$  not "geometric dilution" (flux/divergence) on the surface area of a sphere  $(4\pi r^2)$ . Flux applies to particle flow in a fluid but does not apply to field forces.
    - (ii) Therefore, the charge constant,  $k_c = 1/(4\pi\epsilon_0)$ , is invalid. The correct constant should be,  $k_c = 1/\epsilon_0$ , where the permeability in a vacuum,  $\epsilon_0$ , is  $4\pi$  times larger.
  - (b) If there are quantum values of charge,  $q_C$ , and mass,  $m_G$ , then there are quantum distances,  $r_C$  and  $r_G$ , where the forces do not exist (not

- defined) at smaller distances, which might have implications for the strong force, particle collisions, and the density of black holes.
- (c) Discrete states (like spin) do not have proportionate continuous distance intervals. Therefore, discrete state changes with respect to time are independent of distance (for example, the change in spin of two quantum coupled particles).
- (9) Relativity theory assumes that only 3 dimensions of physical space exist [Bru17]. The proof in this article shows that time constrains physical distance and volume to at most three dimensions (5.4). Higher dimensional Hilbert spaces (inner product vector spaces) are valid if no more than three of the dimensions represent physical space. Theories of higher dimensions of physical distance and volume would would have contradictions.
- (10) The proof of at most 3 members in any ordered and symmetric set (5.4), implies that each *physical* infinitesimal volume (ball) can have at most 3 ordered and symmetric dimensions of discrete *physical* states of the same type. And each dimension of discrete physical states can have at most 3 ordered and symmetric discrete state values, which allows  $3 \cdot 3 \cdot 3 = 27$  possible combinations of discrete values corresponding to 27 possible "types" of infinitesimal balls.
- (11) If each of the three possible ordered and symmetric dimensions of discrete physical states contained unordered sets of discrete state values, for example, unordered binary values, then there would be  $2 \cdot 2 \cdot 2 = 8$  possible combinations of values. Unordered sets (states) are non-deterministic. For example, every time an unordered state is physically measured, there is a 50 percent chance of having one of the binary values.
- (12) Where infinitesimal balls intersect, an algebra of the interactions of the discrete states with respect to time needs to be developed. The interaction of the discrete states associated with overlapping infinitesimal balls with respect to time might result in what we perceive as waves, motion, mass, charge, etc.

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George Van Treeck, 668 Westline Dr., Alameda, CA 94501