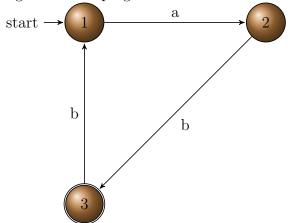
CS375 Week 6

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Figure 1: Pumping Lemma Contradiction



0.1

$$L = \{a^nb^{2n} \mid n \geq 1\} = \{abb, aabbbb, aaabbbbbb, ...\}$$

Proof. For any regular language L, there exists a number p such that for any string w in L of length at least p there are strings x,y,z such that

- \bullet w = xyz
- $\bullet \mid xy \mid \leq p$
- $\bullet \mid y \mid \geq 1$
- Then $x = a^n, y = a^n, z = b^{p+1}$
- $\bullet \ xy^2z\ni L$
- This is a contradiction so the shown language is nonregular.

0.2

1. Palindromes

Proof. For any regular language L, there exists a number p such that for any string w in L of length at least p there are strings x,y,z such that

1

- (a) If w = xyz.
- (b) And if x = a, y = b, z = a.
- (c) or if x = b, y = a, z = b.
- (d) Then $w = a^{90+1}ba^{90}$
- (e) and $w = b^{90+1}ab^{90}$
- (f) Therefore Palindromes are nonregular

2. Equal

Proof. For any regular language L, there exists a number p such that for any string w in L of length at least p there are strings x,y,z such that

- (a) $EQUAL = \{ \Lambda \ ab \ ba \ aabb \ abab \ abba \ baba \ baba \ baba \ baba \ baba \ aaabbb... \}$
- (b) $\{a^nb^n\} = \mathbf{a}^*\mathbf{b}^* \cap EQUAL$
- (c) If a^*b^* is regular then so is the result of $\mathbf{a}^*\mathbf{b}^* \cap EQUAL$
- (d) w = xyz
- (e) $|xy| \le p$
- (f) $|y| \ge 1$
- (g) Then $x = a^n, y = a^2, z = b^n$
- (h) So then xyyz would allow aaab which is not in this language.
- (i) Therefore this language is nonregular

0.3

Prove that the below generates the language defined by the regular expression: $\mathbf{a}^*\mathbf{b}\mathbf{b}$

Prod1
$$S \rightarrow aS \mid bb$$

$$S \Longrightarrow aS$$

$$\Longrightarrow aaS$$

$$\Longrightarrow aaaS$$

$$\Longrightarrow aaaaS$$

$$\Longrightarrow aaaaaS$$

$$\Longrightarrow aaaaabb$$
(1)

This derivation could continue infinitely until terminal bb is appended.

0.4

To generate aabbab using: a(a + b)*

Prod1
$$S \rightarrow aX$$

Prod2 $X \rightarrow aX \mid bX \mid \Lambda$
 $S \Longrightarrow aX$ (by Prod 1)
 $\Longrightarrow aaX$ (by Prod 2)
 $\Longrightarrow aabX$ (by Prod 2)
 $\Longrightarrow aabbX$ (by Prod 2)
 $\Longrightarrow aabbaX$ (by Prod 2)
 $\Longrightarrow aabbabX$ (by Prod 2)
 $\Longrightarrow aabbabA$ (by Prod 2)
 $\Longrightarrow aabbab\Lambda$ (by Prod 2)

0.5

To generate bbabaaa using: $(a + b)^*a(a + b)^*a(a + b)^*$

Prod1
$$S \rightarrow XaXaX$$

Prod2 $X \rightarrow aX \mid bX \mid \Lambda$
 $S \Longrightarrow XaXaX$ (by Prod 1)
 $\Longrightarrow XaXaaX$ (by Prod 2)
 $\Longrightarrow bXaXaaX$ (by Prod 2)
 $\Longrightarrow bXaXaa$ (by Prod 2)
 $\Longrightarrow bXabXaa$ (by Prod 2)
 $\Longrightarrow bXabaXaa$ (by Prod 2)
 $\Longrightarrow bbXabaXaa$ (by Prod 2)
 $\Longrightarrow bbAabaXaa$ (by Prod 2)
 $\Longrightarrow bbabaXaa$ (by Prod 2)
 $\Longrightarrow bbabaAaa$ (by Prod 2)
 $\Longrightarrow bbabaAaa$ (by Prod 2)

0.6

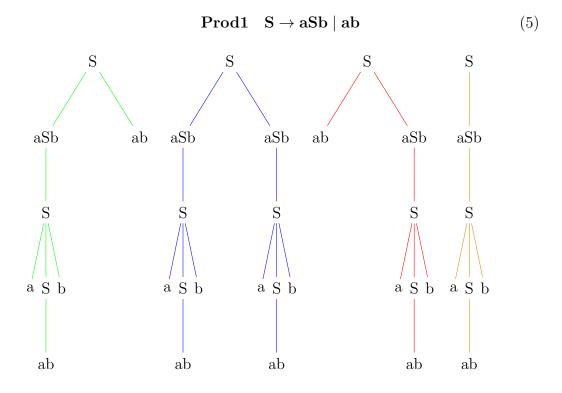
Prod1
$$S \rightarrow Xa$$

Prod2 $X \rightarrow bbX \mid bbS \mid bb \mid \Lambda$ (4)

0.7

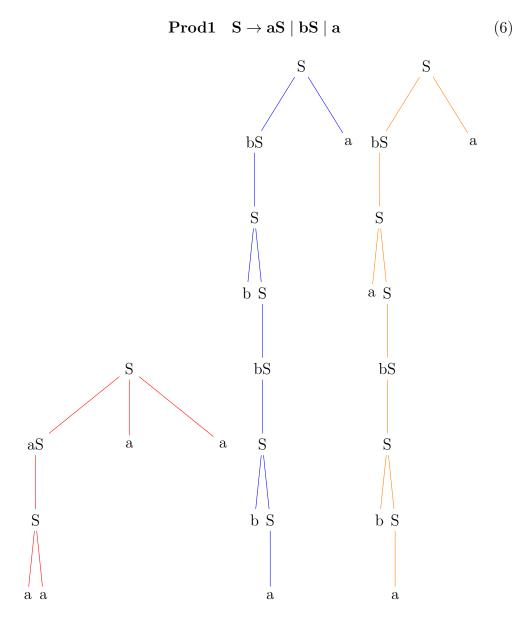
0.7.1

This CFG cannot match aaaa, abaa, or bbaa. The only string of length 4 is aabb:



0.7.2

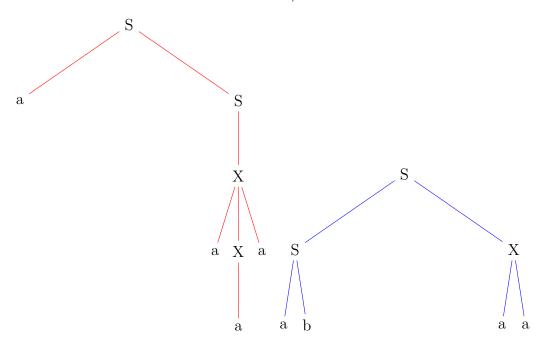
For strings aaaa, abaa, or bbaa. This CFG can create the string aaaa, bbaa:



0.7.3

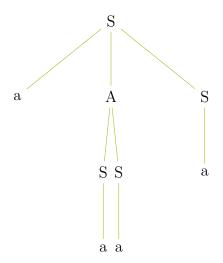
For strings aaaa, abaa, or bbaa. This CFG can create the string aaaa, abaa:

$$\begin{array}{ll} \mathbf{Prod1} & \mathbf{S} \rightarrow \mathbf{aS} \mid \mathbf{aSb} \mid \mathbf{X} \\ \mathbf{Prod2} & \mathbf{X} \rightarrow \mathbf{aXa} \mid \mathbf{a} \end{array}$$



0.7.4

$$\begin{array}{ll} \mathbf{Prod1} & \mathbf{S} \rightarrow \mathbf{aAS} \mid \mathbf{a} \\ \mathbf{Prod2} & \mathbf{A} \rightarrow \mathbf{SbA} \mid \mathbf{SS} \mid \mathbf{ba} \end{array}$$



0.7.5

$$\begin{array}{ll} \mathbf{Prod1} & \mathbf{S} \rightarrow \mathbf{aB} \mid \mathbf{bA} \\ \mathbf{Prod2} & \mathbf{A} \rightarrow \mathbf{a} \mid \mathbf{aS} \mid \mathbf{bAA} \\ \mathbf{Prod3} & \mathbf{B} \rightarrow \mathbf{b} \mid \mathbf{bS} \mid \mathbf{aBB} \end{array} \tag{9}$$

S

- 0.8
- 0.9
- 0.10
- 0.11