CS375 Week 6

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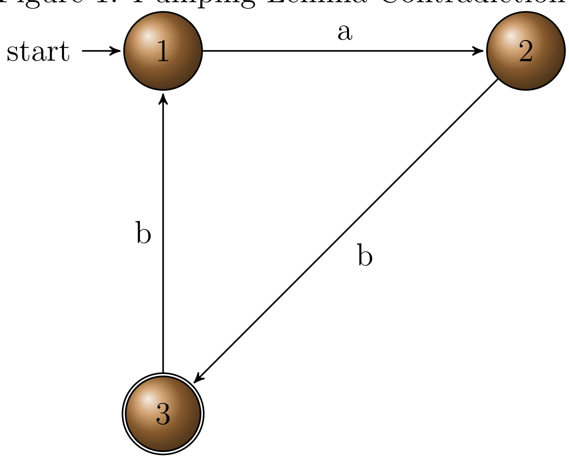
Figure 1: Pumping Lemma Contradiction



0.1

$$\mathbf{L = \{a^n b^{2n} \mid n \geq 1\} = \{abb, aabbbb, aaabbbbbb, ...\}}$$

Proof. For any regular language L, there exists a number p such that for any string w in L of length at least p there are strings x,y,z such that

- $w = xyz$
- $\mid xy \mid \leq p$
- $\mid y \mid \geq 1$
- Then $x = a^n, y = a^n, z = b^{p+1}$
- $xy^2z \ni L$
- This is a contradiction so the shown language is nonregular.

□

0.2

1. Palindromes

   Proof. For any regular language L, there exists a number p such that for any string w in L of length at least p there are strings x,y,z such that

(a) If w = xyz.
(b) And if x = a, y = b, z = a.
(c) or if x = b, y = a, z = b.
(d) Then $w = a^{90+1}ba^{90}$
(e) and $w = b^{90+1}ab^{90}$
(f) Therefore Palindromes are nonregular

□

2. Equal

Proof. For any regular language L, there exists a number p such that for any string w in L of length at least p there are strings x,y,z such that

(a) EQUAL = $\{\Lambda \quad ab \quad ba \quad aabb \quad abab \quad abba \quad baab \quad baba \quad bbaa \quad aaabbb...\}$
(b) $\{a^n b^n\} = \mathbf{a}^*\mathbf{b}^* \cap EQUAL$
(c) If a^*b^* is regular then so is the result of $\mathbf{a}^*\mathbf{b}^* \cap EQUAL$
(d) $w = xyz$
(e) $\mid xy \mid \leq p$
(f) $\mid y \mid \geq 1$
(g) Then $x = a^n, y = a^2, z = b^n$
(h) So then xyyz would allow aaab which is not in this language.
(i) Therefore this language is nonregular

□

0.3

Prove that the below generates the language defined by the regular expression: $\mathbf{a}^*\mathbf{bb}$

Prod1 S → aS | bb

$$\begin{aligned} S &\Longrightarrow aS \\ &\Longrightarrow aaS \\ &\Longrightarrow aaaS \\ &\Longrightarrow aaaaS \\ &\Longrightarrow aaaaaS \\ &\Longrightarrow aaaaabb \end{aligned} \tag{1}$$

This derivation could continue infinitely until terminal bb is appended.

0.4

To generate aabbab using: $a(a+b)*$

$$\begin{aligned}
&\textbf{Prod1} \quad \mathbf{S \to aX} \\
&\textbf{Prod2} \quad \mathbf{X \to aX \mid bX \mid \Lambda} \\
&S \Longrightarrow aX \quad \text{(by Prod 1)} \\
&\Longrightarrow aaX \quad \text{(by Prod 2)} \\
&\Longrightarrow aabX \quad \text{(by Prod 2)} \\
&\Longrightarrow aabbX \quad \text{(by Prod 2)} \\
&\Longrightarrow aabbaX \quad \text{(by Prod 2)} \\
&\Longrightarrow aabbabX \quad \text{(by Prod 2)} \\
&\Longrightarrow aabbab\Lambda \quad \text{(by Prod 2)}
\end{aligned} \tag{2}$$

0.5

To generate bbabaaa using: $(a+b)^*a(a+b)^*a(a+b)^*$

$$\begin{aligned}
&\textbf{Prod1} \quad \mathbf{S \to XaXaX} \\
&\textbf{Prod2} \quad \mathbf{X \to aX \mid bX \mid \Lambda} \\
&S \Longrightarrow XaXaX \quad \text{(by Prod 1)} \\
&\Longrightarrow XaXaaX \quad \text{(by Prod 2)} \\
&\Longrightarrow bXaXaaX \quad \text{(by Prod 2)} \\
&\Longrightarrow bXaXaa \quad \text{(by Prod 2)} \\
&\Longrightarrow bXabXaa \quad \text{(by Prod 2)} \\
&\Longrightarrow bXabaXaa \quad \text{(by Prod 2)} \\
&\Longrightarrow bbXabaXaa \quad \text{(by Prod 2)} \\
&\Longrightarrow bbabaXaa \quad \text{(by Prod 2)} \\
&\Longrightarrow bbabaaa \quad \text{(by Prod 2)}
\end{aligned} \tag{3}$$

0.6

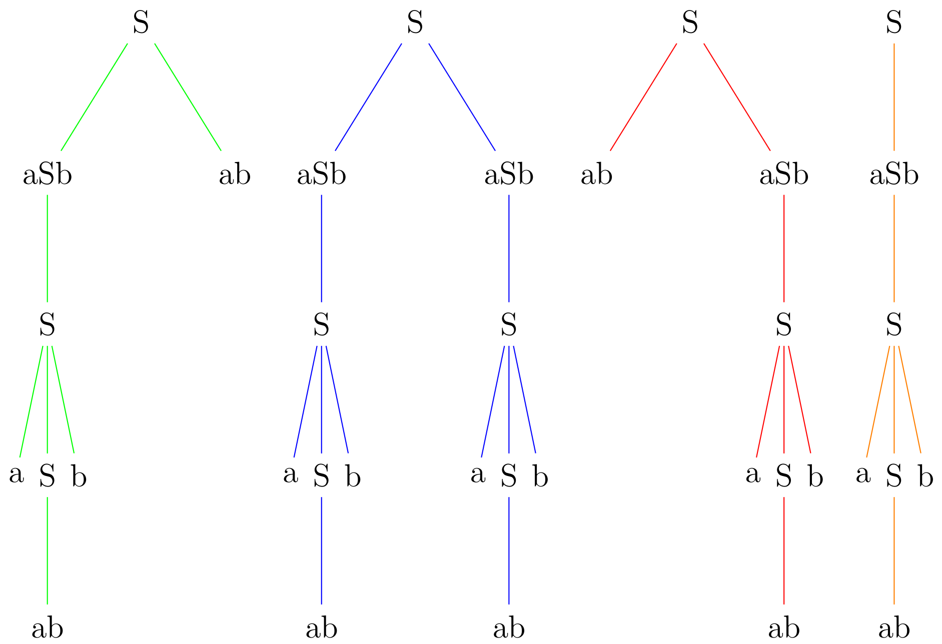
$$\begin{aligned}
&\textbf{Prod1} \quad \mathbf{S \to Xa} \\
&\textbf{Prod2} \quad \mathbf{X \to bbX \mid bbS \mid bb \mid \Lambda}
\end{aligned} \tag{4}$$

0.7

0.7.1

This CFG cannot match aaaa, abaa, or bbaa. The only string of length 4 is aabb:

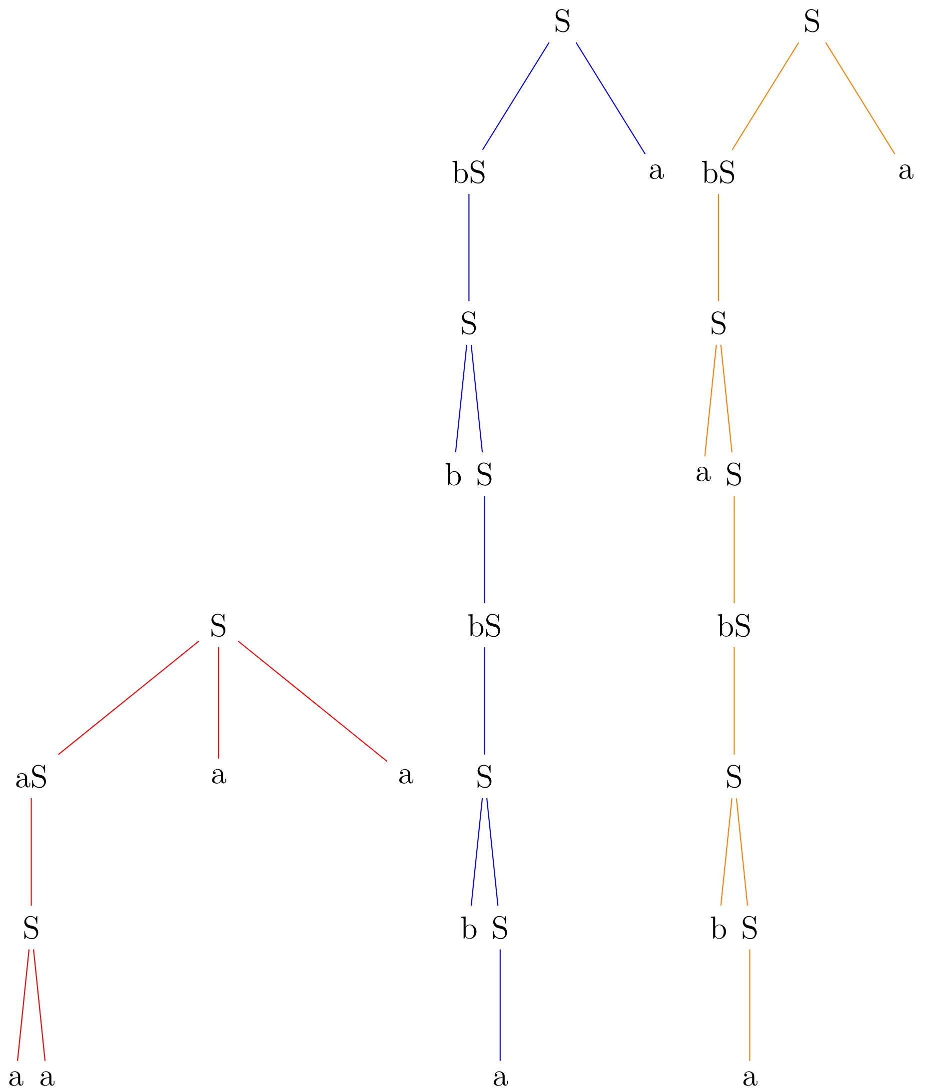
$$\textbf{Prod1} \quad \mathbf{S \rightarrow aSb \mid ab} \tag{5}$$



0.7.2

For strings aaaa, abaa, or bbaa. This CFG can create the string aaaa, bbaa:

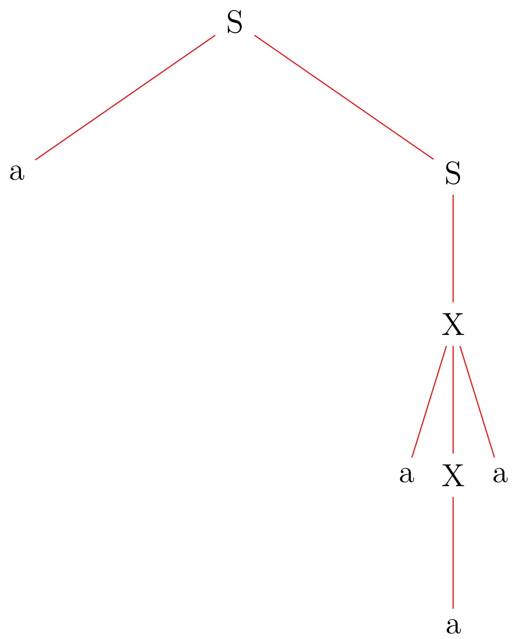
$$\textbf{Prod1} \quad \mathbf{S \rightarrow aS \mid bS \mid a} \tag{6}$$

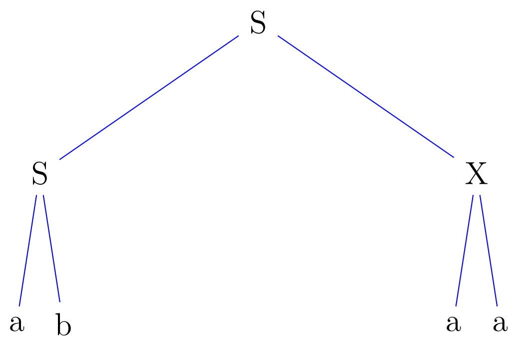


0.7.3

For strings aaaa, abaa, or bbaa. This CFG can create the string aaaa, abaa:

$$\begin{array}{ll} \textbf{Prod1} & \mathbf{S \rightarrow aS \mid aSb \mid X} \\ \textbf{Prod2} & \mathbf{X \rightarrow aXa \mid a} \end{array} \tag{7}$$

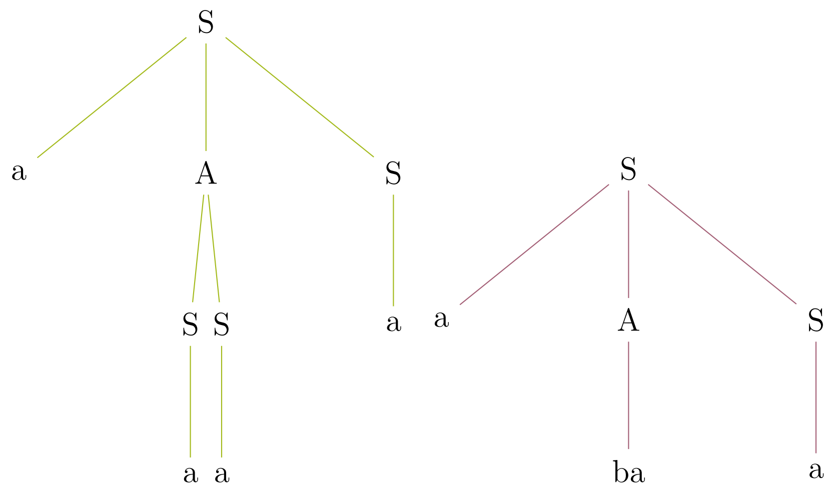




0.7.4

For strings aaaa, abaa, or bbaa. This CFG can create the string aaaa, abaa:

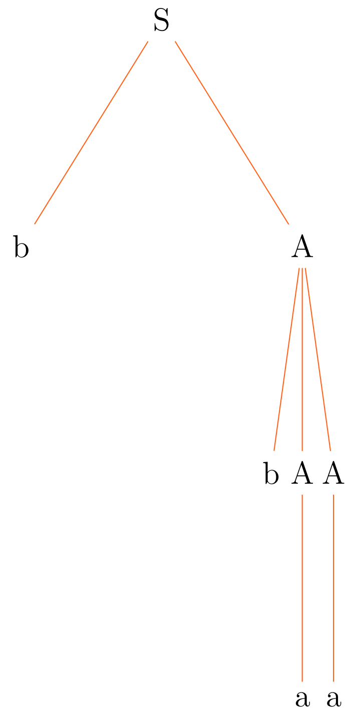
$$\begin{aligned} &\textbf{Prod1} \quad \mathbf{S \rightarrow aAS \mid a} \\ &\textbf{Prod2} \quad \mathbf{A \rightarrow SbA \mid SS \mid ba} \end{aligned} \tag{8}$$



0.7.5

For strings aaaa, abaa, or bbaa. This CFG can create the string bbaa:

$$\begin{aligned} &\textbf{Prod1} \quad \mathbf{S \rightarrow aB \mid bA} \\ &\textbf{Prod2} \quad \mathbf{A \rightarrow a \mid aS \mid bAA} \\ &\textbf{Prod3} \quad \mathbf{B \rightarrow b \mid bS \mid aBB} \end{aligned} \tag{9}$$

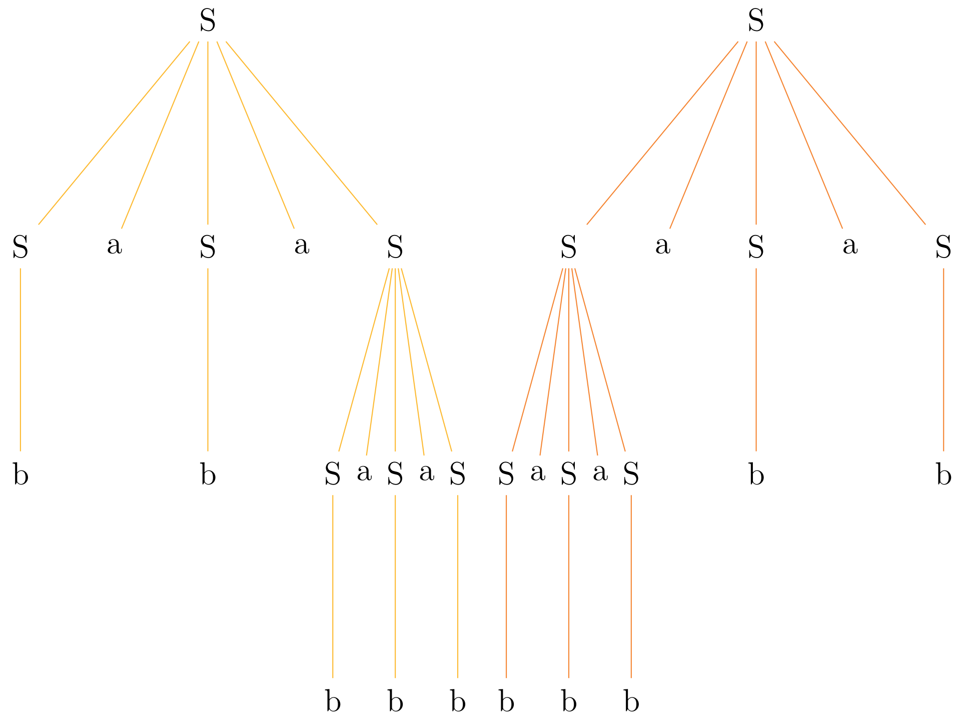


0.8

0.8.1

Ambiguous CFG with the string: babababab

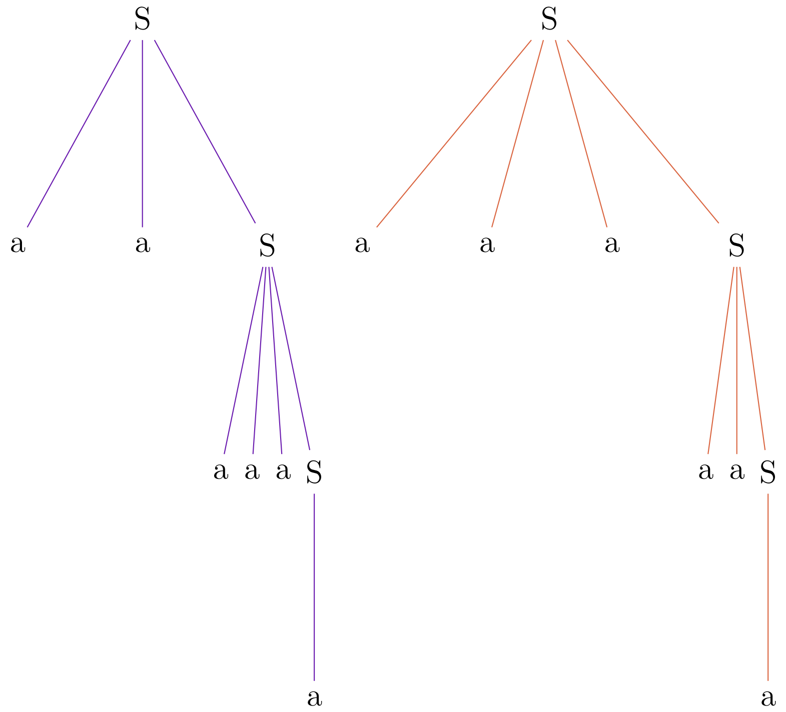
$$\textbf{Prod1} \quad \mathbf{S} \rightarrow \mathbf{SaSaS} \mid \mathbf{b} \tag{10}$$



0.8.2

Ambiguous CFG with the string: aaaaaa

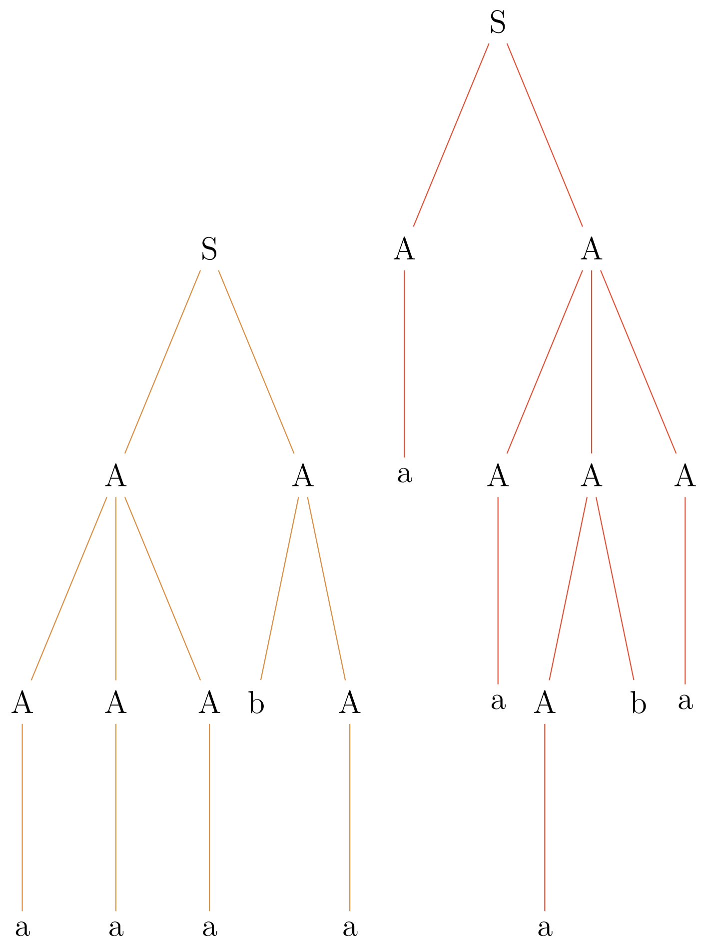
$$\textbf{Prod1} \quad \mathbf{S \rightarrow aaS \mid aaaS \mid a} \tag{11}$$



0.8.3

Ambiguous CFG with the string: aaaba

$$\begin{aligned} &\textbf{Prod1} \quad \mathbf{S \rightarrow AA} \\ &\textbf{Prod2} \quad \mathbf{A \rightarrow AAA \mid a \mid bA \mid Ab} \end{aligned} \tag{12}$$

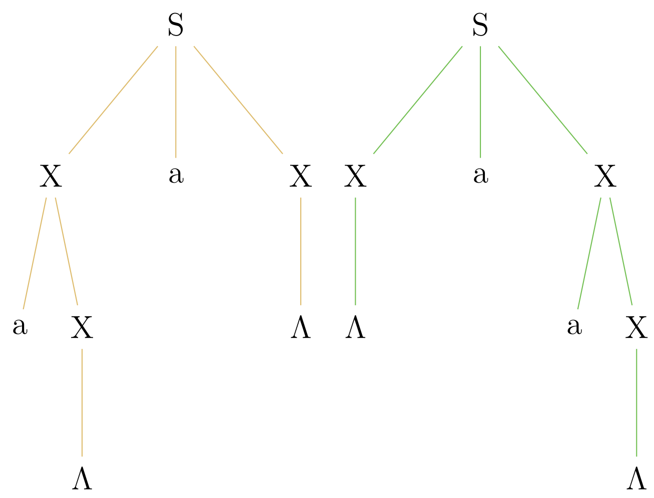


0.9

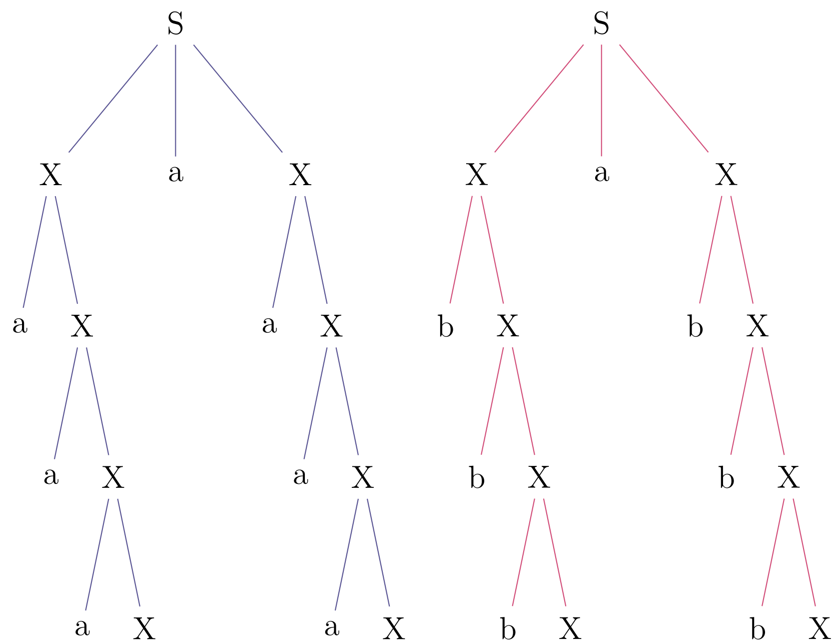
0.9.1

$$\begin{aligned} &\textbf{Prod1} \quad \mathbf{S} \rightarrow \mathbf{XaX} \\ &\textbf{Prod2} \quad \mathbf{X} \rightarrow \mathbf{aX} \mid \mathbf{bX} \mid \boldsymbol{\Lambda} \end{aligned} \tag{13}$$

Ambiguous

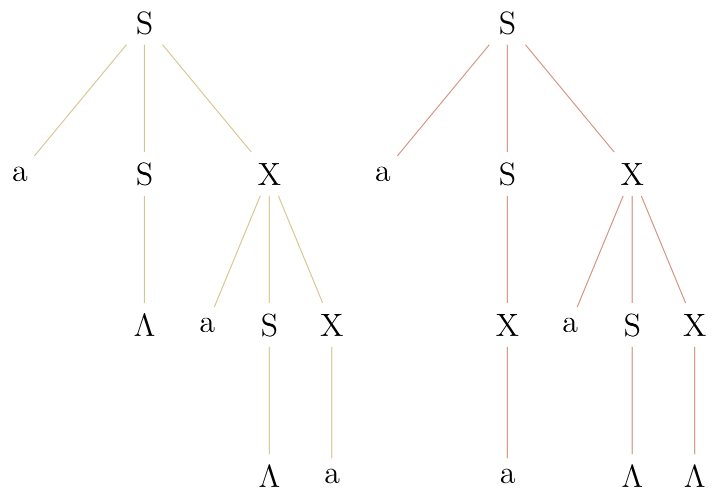


Non-Ambiguous



0.9.2

$$\begin{aligned} &\textbf{Prod1} \quad \mathbf{S \rightarrow aSX \mid \Lambda} \\ &\textbf{Prod2} \quad \mathbf{X \rightarrow aX \mid a} \end{aligned} \tag{14}$$



0.9.3

$$\textbf{Prod1} \quad \mathbf{S \rightarrow aS \mid bS \mid aaS \mid \Lambda} \tag{15}$$

0.10

0.10.1

$$\textbf{Prod1} \quad \mathbf{S \rightarrow aS \mid bS \mid a} \tag{16}$$

0.10.2

$$\begin{aligned} &\textbf{Prod1} \quad \mathbf{S \rightarrow aSb \mid bX} \\ &\textbf{Prod2} \quad \mathbf{X \rightarrow bX \mid b} \end{aligned} \tag{17}$$

0.11

The CFG's generated by the Regular Languages over the alphabet $\Sigma = \{ab\}$:

0.11.1

The language defined by $(aaa + b)^*$:

0.11.2

The language defined by $(a+b)^*(bbb+aaa)(a+b)^*$:

0.11.3

All the strings that end in b and have an even number of b's in total:

0.11.4

The set of all strings of odd length:

0.11.5

All strings with exactly one a or exactly one b:

0.11.6

All strings with an odd number of a's or and even number of b's: