Combinatorics in matroid Kazhdan-Lusztig polynomials

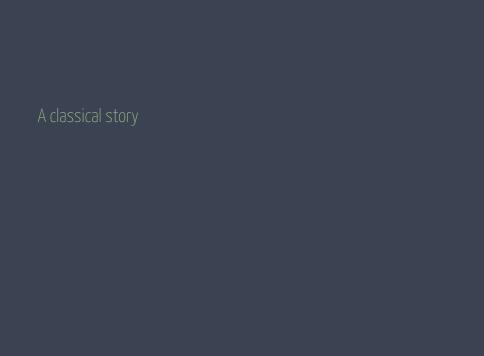
by Trevor K. Karn (U. Minnesota - Twin Cities) (joint with George Nasr, Nick Proudfoot, and Lorenzo Vecchi) on Thursday, March 23, 2023

A classical story

Matroids

Combinatorial formulas





A <u>partition</u> $\lambda \vdash n$ is a weakly decreasing sequence of nonnegative integers $\lambda_1 \geq \lambda_2 \geq \cdots$ summing to n.

Example
$$\lambda = (5,4,4,2,1) \vdash 16 \text{ has Ferrers diagram}$$

A <u>Young tableau</u> T is a filling of a Ferrers diagram by positive integers. T is <u>standard</u> if it is filled by $\{1,2,\ldots,n\}$ and increasing in rows and columns. Define f^{λ} as the number of standard tableaux of shape λ .

Example

One of $f^{(5,4,4,2,1)} = \overline{549120}$ standard Young tableaux:

1	6	10	13	16
2	7	11	14	
3	8	12	15	
4	9			
5				

Fix *n*. Then

$$\sum_{\lambda} (f^{\lambda})^2 = r$$

Fix n. Then

$$\sum_{\lambda \vdash n} (f^{\lambda})^2 = n!$$

Proof 1:

The Robinson-Schensted bijection:

pairs of standard tableaux of same shape \longleftrightarrow permutation in \mathfrak{S}_n

A representation of a group G is a homomorphism

$$\rho: G \to \operatorname{GL}_n(\mathbb{C}).$$

A representation is <u>irreducible</u> if there is no *G*-stable subspace $W \subseteq \mathbb{C}^n$.

There is a representation of \mathfrak{S}_3 defined by sending

$$(12) \mapsto \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$(23) \mapsto \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}.$$

The action: $\pi \in \mathfrak{S}_3$ permutes coordinates. So

$$(12) \cdot \langle 1, 2, 3 \rangle = \langle 2, 1, 3 \rangle$$

and

$$\pi \cdot \langle 1, 1, 1 \rangle = \langle 1, 1, 1 \rangle$$

Irreducible \mathfrak{S}_n representations are indexed by $\lambda \vdash n$. Denote them by S^{λ} . Then

$$\dim \mathcal{S}^{\lambda} = f^{\lambda}.$$

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Example

The representation on the last slide contains $S^{(3)}$ and $S^{(2,1)}$.

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The representation on the last slide contains $S^{(3)}$ and $S^{(2,1)}$.

Fact

Let d_1, d_2, \ldots, d_r be the dimensions of all irreducible representations of a finite group. Then

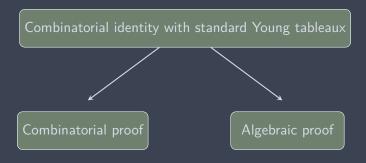
$$\sum_{i}d_{i}^{2}=|G|.$$

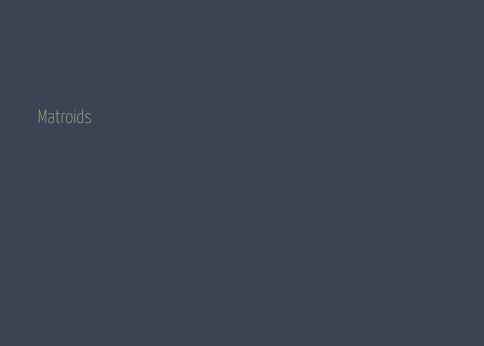
Fix n. Then

$$\sum_{\lambda \vdash n} (f^{\lambda})^2 = n$$

Proof 2:

$$\sum_{\lambda} (f^{\lambda})^2 = \sum_{i} d_i^2 = |G| = |\mathfrak{S}_n| = n!$$





Let v_1, \ldots, v_n be vectors in a vector space V (not all 0). Then any two bases A, B for the span of $v_1, \ldots v_n$ satisfy the following requirements:

- 1) There must be at least one basis
- 2) If $a \in A B$ then there is a $b \in B$ with $(A a) \cup b$ a basis.



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A **matroid** $M = (E, \mathcal{B})$ is a finite set E with $\emptyset \neq \mathcal{B} \subseteq 2^E$ such that if $A, B \in \mathcal{B}$ and $a \in A$, there exists $b \in B$ such that

$$(A-a)\cup b\in \mathcal{B}$$

Call ${\cal B}$ the bases of the matroid.

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Matroids have the combinatorics of vectors without the "zeroth postulate".

A matroid M = (E, C) is a set E with $C \subseteq 2^E$ such that

 $\emptyset \not\in \mathcal{C}$

If $C_1, C_2 \in \mathcal{C}$ with $C_1 \subseteq C_2$, then $C_1 = C_2$.

If $C_1, C_2 \in \mathcal{C}$ are distinct, and $e \in C_1 \cap C_2$, then there is a $C_3 \in \mathcal{C}$ such that $C_3 \subseteq (C_1 \cup C_2) - e$

Call ${\mathcal C}$ the circuits.

The **uniform matroid** $U_{k,n}$ models n-many k-dimensional vectors in general position

Bases \longleftrightarrow any set of k-many vectors

Circuits \longleftrightarrow any set of k+1-many vectors

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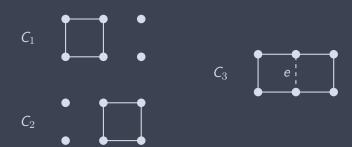
Example of the example

 $U_{3,12}$ corresponds to 12 generic vectors in \mathbb{R}^3 . One choice of basis is $\{e_1,e_2,e_3\}$. On the other hand $\{e_1,e_2,e_3,v\}$ is dependent for any $v\in\mathbb{R}^3$.

A graph Γ with edges E forms a matroid:

Bases ←→ spanning trees

Circuits ←→ cycles



The columns of a matrix form a matroid:

$$\begin{bmatrix} 1 & 2 & 0 & 0 \\ 2 & 4 & 2 & 4 \end{bmatrix}$$



The projective plane is the set of lines through the origin in \mathbb{F}^3 .

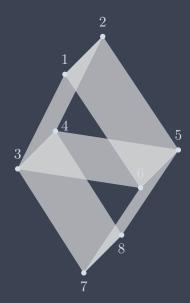
The projective plane is the set of lines through the origin in \mathbb{F}^3 .

If
$$\mathbb{F} = \mathbb{F}_2$$
,

$$egin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 & 1 \ 0 & 1 & 0 & 1 & 0 & 1 & 1 \ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 \end{bmatrix}$$

Almost all matroids cannot be written as a matrix.

» Example: Vámos matroid





Finite set of vectors	Matroids	
Maximally independent sets	Bases ${\cal B}$	
Minimally dependent sets	Circuits ${\cal C}$	
Dimension of span	Rank	
Subspaces	Flats ${\mathcal F}$	
Codimension 1 subspaces	Hyperplanes ${\cal H}$	

> Vocab

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Definition

 $\mathcal{CH} = \mathcal{C} \cap \mathcal{H}$ is the set of circuit hyperplanes.

A hyperplane H is <u>stressed</u> if every subset of H of size rk(E) is in C. Denote the set of (nontrivial) stressed hyperplanes by SH.

A paving matroid is sparse if $\binom{E}{k} = \mathcal{CH} \sqcup \mathcal{B}$

M is a paving matroid if all circuits have size at least $k = \operatorname{rk}(E)$

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Conjecture (Mayhew, Newman, Welsh, Whittle '11)

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Theorem (Pendavingh, van der Pol '15)

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 $\mathbb{PF}_3 \leftrightarrow \mathsf{paving}$

 $V \leftrightarrow \text{sparse paving}$

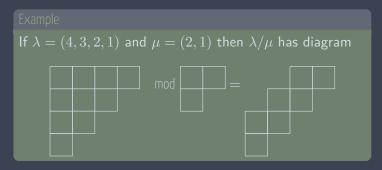
Fact

The matroid Kazhdan–Lusztig polynomial $P_M(t)$ is an interesting polynomial invariant of a matroid M, introduced by Elias, Proudfoot, and Wakefield in 2016.

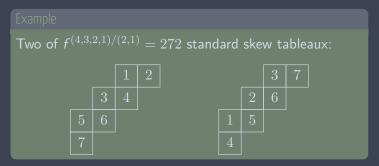
It is defined in terms of \mathcal{F} .



A skew partition λ/μ is a pair of partitions where the diagram of μ is contained in the diagram of λ



A skew tableau T is a filling of a skew diagram by positive integers. T is standard if it is filled by $\{1,2,\ldots,|\lambda|-|\mu|\}$ and increasing in rows and columns. Define $f^{\lambda/\mu}$ as the number of standard skew tableaux of shape λ/μ .



Theorem (Lee, Nasr, Radcliffe '21)

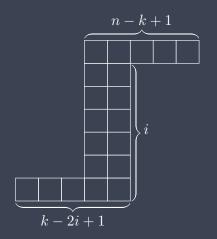
Let M be a rank-k, sparse paving matroid with E = [n] and circuit hyperplanes \mathcal{CH} . The t^i coefficient in $P_M(t)$ is

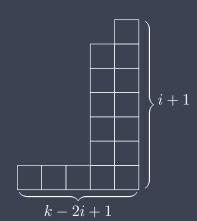
$$f^{\lambda/\mu} - |\mathcal{CH}| f^{\lambda'/\mu'}$$

where

$$\lambda = [n - 2i, (k - 2i + 1)^{i}], \mu = [(k - 2i - 1)^{i}]$$

$$\lambda' = [(k-2i+1)^{i+1}], \mu' = [k-2i, (k-2i-1)^{i-1}]$$





Theorem (Lee, Nasr, Radcliffe '21)

For a sparse paving matroid M, the t^i coefficient in $P_M(t)$ is

 $f^{\lambda/\mu} - |\mathcal{CH}| f^{\lambda'/\mu'}$

Proof 1 (LNR '21): Combinatorial argument with recursion.

Fact

There is a (reducible) \mathfrak{S}_n representation $S^{\lambda/\mu}$ of dimension $f^{\lambda/\mu}$.

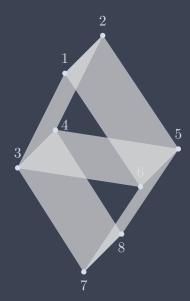
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For a sparse paving matroid M, the t^i coefficient in $P_M(t)$ is

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Proof 1 (LNR '21): Combinatorial argument with recursion. Proof 2 (KNPV '22): $\dim(\text{some } S^{\lambda/\mu} \text{ coming from } M)$.

» Example: Vámos matroid



Some general facts:

Know $P_M(t)$ always has constant term 1.

Know deg
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.

rk
$$V = 4$$
 so $P_V(t) = 1 + ?t$.

Only need to compute linear coefficient!

$$\lambda = [6, 3], \ \mu = [1] \longrightarrow$$

$$\lambda' = [3, 3], \ \mu' = [2] \longrightarrow$$

$$|\mathcal{CH}| = 5$$

$$f^{\lambda/\mu} - 5f^{\lambda'/\mu'} = 48 - 15 = 33$$

$$P_{V}(t) = 1 + 33t$$

» Example: Projective plane over **F**₃

$$\lambda/\mu = \boxed{ }$$

$$\mathcal{C}\mathcal{H} = \emptyset$$

$$f^{\lambda/\mu} = 65 \qquad \neq 0$$

From Elias, Proudfoot, and Wakefield, we know

$$P_M(t)=1$$

» Example: Projective plane over \mathbb{F}_3

$$\lambda/\mu =$$

$$|\mathcal{SH}| = 13$$

$$\lambda'/\mu' =$$

$$f^{\lambda/\mu} - 13f^{\lambda'/\mu'} = 65 - 13 * 5 = 0$$

From Elias, Proudfoot, and Wakefield, we know

$$P_{M}(t)=1$$

Theorem

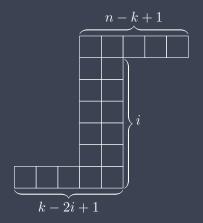
For a (arbitrary!) paving matroid M, the t^i coefficient in $P_M(t)$ is

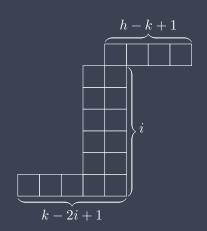
$$f^{\lambda/\mu} - \sum_{H \in \mathcal{SH}} f^{\lambda'(|H|)/\mu'}$$

where

$$\lambda = [n - 2i, (k - 2i + 1)^{i}], \mu = [(k - 2i - 1)^{i}]$$
$$= [h - 2i + 1, (k - 2i + 1)^{i}], \mu' = [h - 2i, (k - 2i - 1)^{i-1}]$$

Proof: Our proof of LNR's theorem implies this more general result



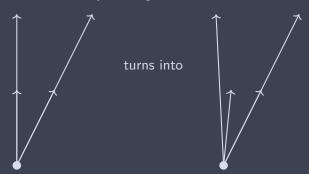


» Proofidea

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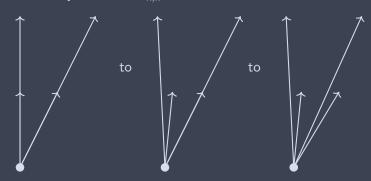


» Proof idea

Do this until you obtain $U_{k,n}$

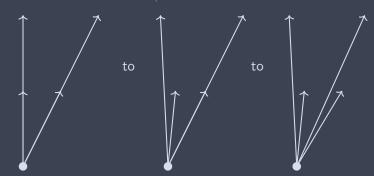
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Proof idea

Do this until you obtain $U_{k,n}$



Each step accounts for $S^{\lambda'(h)/\mu'}$.

Certain $P_{M}(t)$ in terms of standard skew Young tableaux

Combinatorial proof [LNR21]

Algebraic proof [KNPV22] + extension

THANK YOU!

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