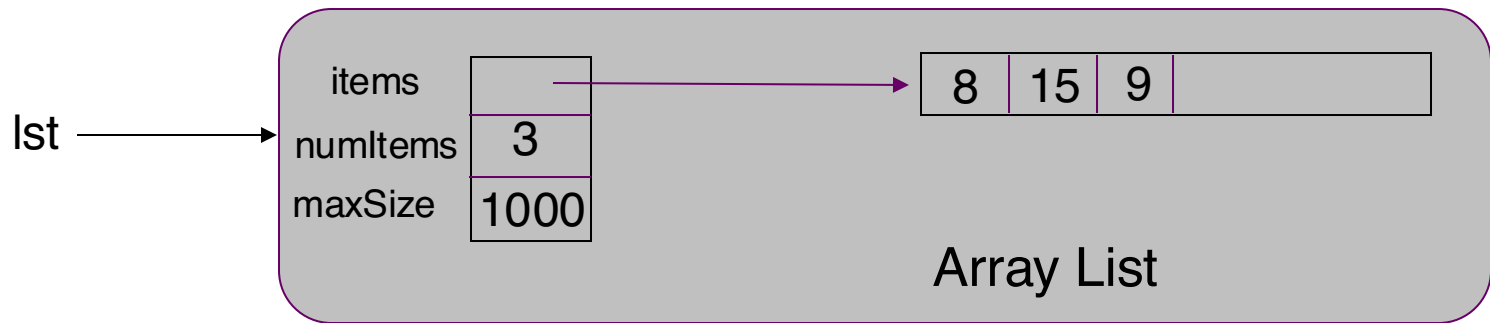


Lists

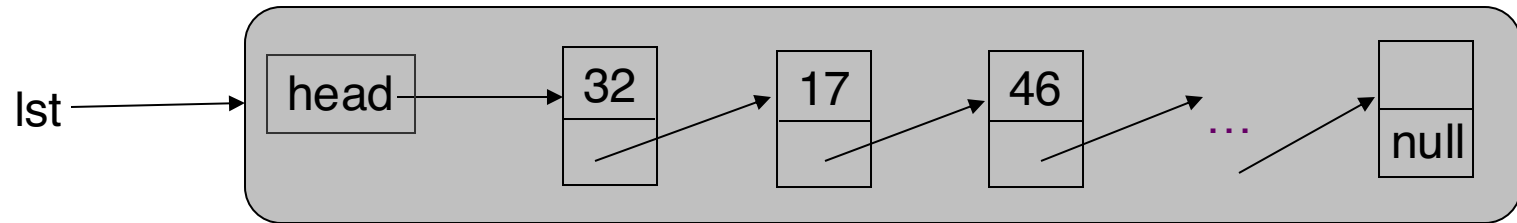
EECS 233

Array Representation



- n Store list elements wall-to-wall in memory in an array
- n Keep a variable recording the current number of elements
- n Advantages
 - Easy and efficient access to *any* item in the sequence
 - 4 `items[i]` gives you the item at position `i`
 - 4 **Random access**
 - Every item can be accessed in constant time given its index
 - Very compact: no auxiliary fields are required
- n Disadvantages of using an array:
 - The need to specify an initial array size and resize as required (how?)
 - Difficult to insert/delete items at arbitrary positions (running time?)
 - May have many empty positions

Linked List Representation

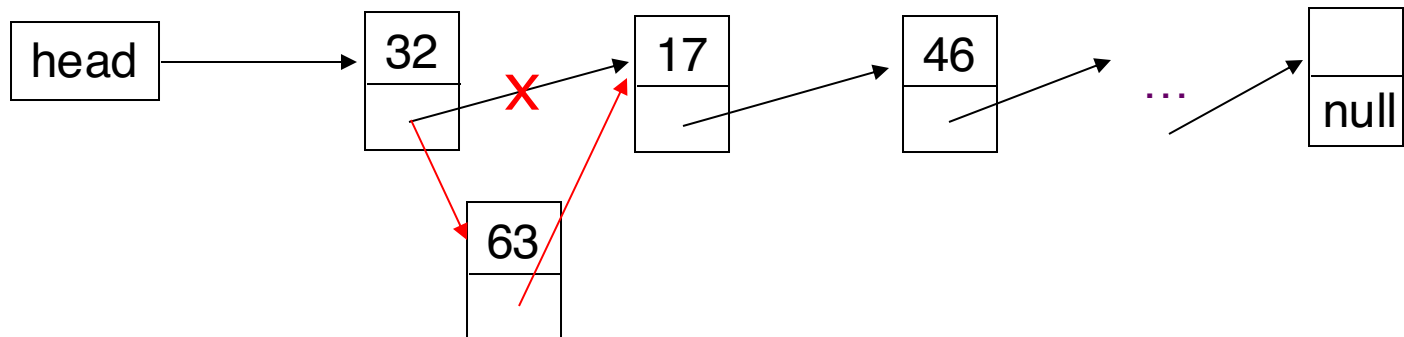


- n More efficient data structure for a dynamic sequence (with frequent insert/delete operations)
 - n A linked list stores a sequence of items in separate *nodes*. Each node contains:
 - a single item, and
 - a “link” (i.e., a reference/pointer) to the node containing the next item
- The last node in the linked list has a link value of **NULL or null**.
- n The linked list starts with a variable that holds a reference to the first node – **head of the list**

Advantages/Disadvantages of Linked List

n Advantages:

- No capacity limit (provided there is enough memory).
- Easy to insert/delete an item – no need to “shift over” other items.
 - 4 Done in constant time.



n Disadvantages:

- No random access
 - 4 “walk down” the list to access an item
- Memory overhead for the links

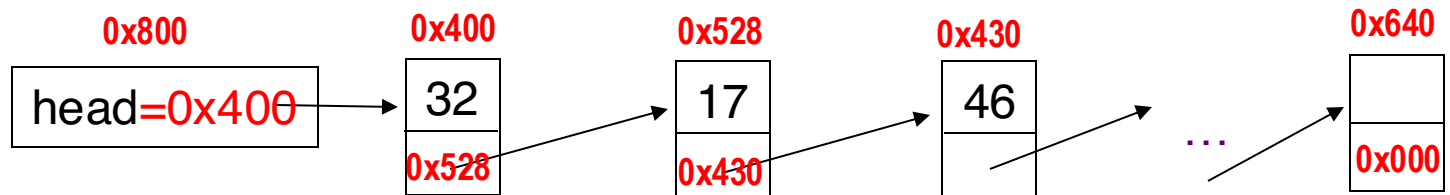
Memory Management for Linked Lists

- n In an array, the elements occupy consecutive memory locations in the heap:

➤ Address in red

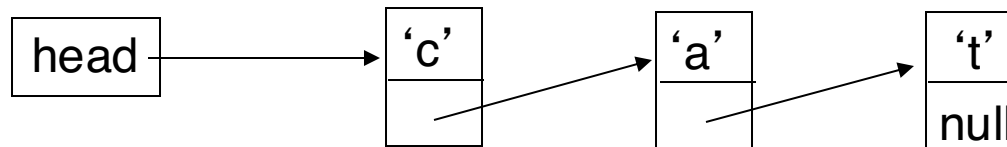
0x800	0x804	0x808		
32	17	46	...	

- n In a linked list, each node is a distinct object in the heap. The nodes do *not* have to be next to each other in memory.



An Example Linked List

- n A string represented using a linked list: **LLString**. Each node in the linked list represents one character.

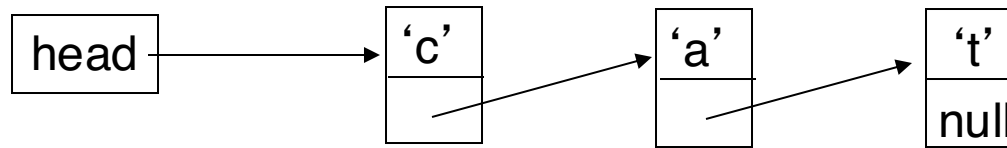


- n Java classes

```
public class StringNode {  
    private char ch;  
    private StringNode next;  
    ...  
}
```

```
public class LLString {  
    private StringNode head;  
    private int theSize;  
    ...  
}
```

Under the Hood of the String Linked List



- n The string as a whole will be represented (*internally*) by a variable that holds a reference to the node containing the first character.

```
StringNode str1;
```

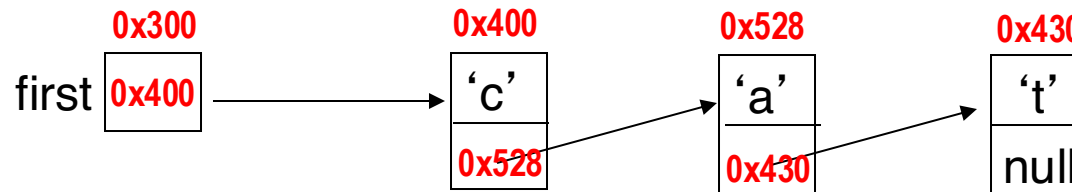
- n An empty string will be represented by a NULL value.

```
StringNode str2 = null;
```

- n We will use helper methods that take the first node of the string as a parameter.
 - We will have `length(str1)` instead of `str1.length()`
 - This is necessary so that the methods can handle empty strings.
 - 4 if `str1 == NULL`, `length(str1)` will work, but `str1.length()` will produce a runtime error

More on References

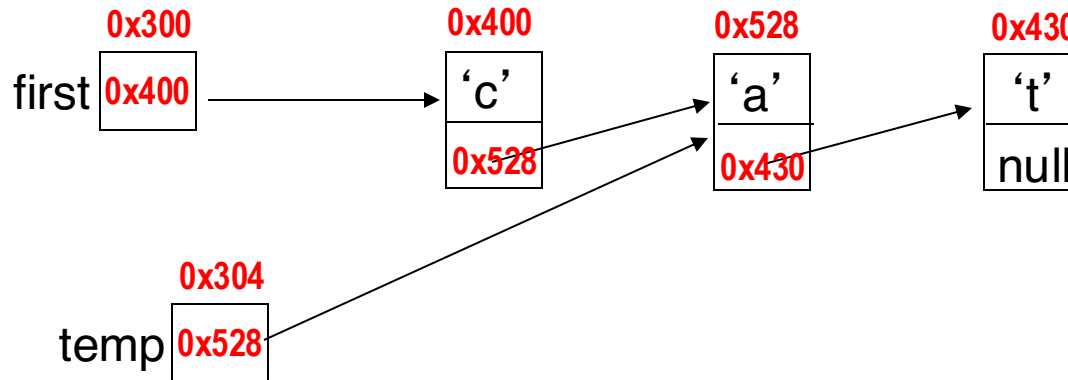
- n A reference is also a variable
 - that has its location in the memory, and
 - whose value is the address (i.e., location) of data



- e.g., first: address=0x300, value=0x400
- How about first.next.next ?
- How about first.next.ch ?

More on References

n Example: temp.next.ch



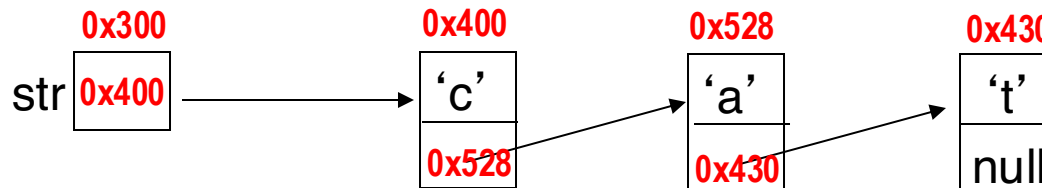
- Start with expression: “**temp.**” It represents the node to which **temp** refers.
- Now consider “**temp.next**”. It represents the field “next” of the node to which temp refers.
 - 4 address = ?
 - 4 value = ?
- Next, consider “**temp.next.**”. It represents the node with address 0x430.
- Finally, **temp.next.ch** represents ‘t’.

Recursion on Linked Lists

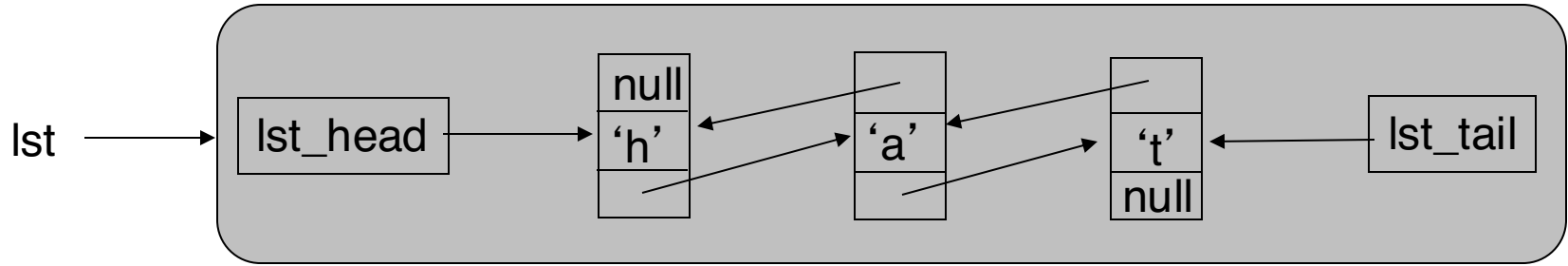
- n Recursive definition of a linked list: a linked list is either
 - empty or
 - a single node, followed by a linked list
- n Recursive definition lends itself to recursive methods.

- n Example: length of a string
 - length of “cat” = 1 + length of “at”
 - length of “at” = 1 + length of “t”
 - length of “t” = 1 + length of the empty string (which is 0)

```
private static int length(StringNode str) {  
    if (str == null)  
        return 0;  
    else  
        return 1 + length(str.next);  
}
```



Doubly Linked List



- n Both next and prev are defined in StringNode
- n Why needed?

```
public class LLString {  
    private StringNode lst_head;  
    private StringNode lst_tail;  
    private int theSize;  
    ...  
}
```

```
public class StringNode {  
    private char ch;  
    private StringNode next;  
    private StringNode prev;  
    ...  
}
```

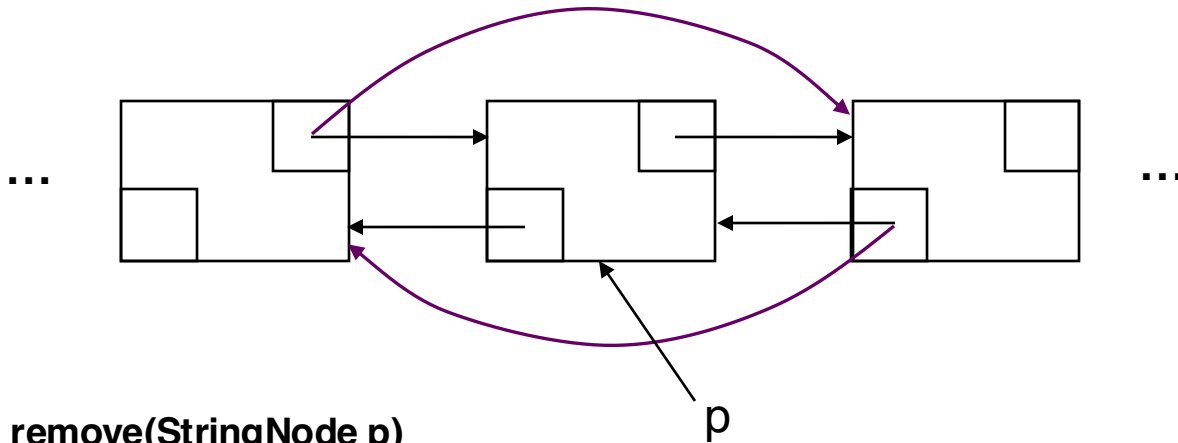
Example: Traversing Linked List

- n Access the node at position i in a doubly linked list

```
public StringNode getNode(int i) {  
    if (i < 0 || i >= theSize) throw new Exception("Index out of bounds");  
    StringNode ptr;  
    if (i < theSize/2) {  
        ptr = lst_head;  
        for (j = 0; j != i; j++) ptr = ptr.next;  
    } else {  
        ptr = lst_tail;  
        for (j = theSize-1; j != i; j--) ptr = ptr.prev;  
    }  
    return ptr;  
}
```

What is the running time?

Example: Removing a Node



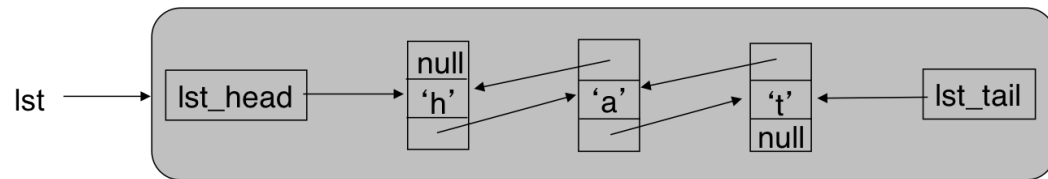
```
public char remove(StringNode p)
{
    if (p == lst_head || p == lst_tail)
```

?

```
    p.next.prev = p.prev;
    p.prev.next = p.next;
    theSize--;
```

```
    return p.ch;
```

```
}
```

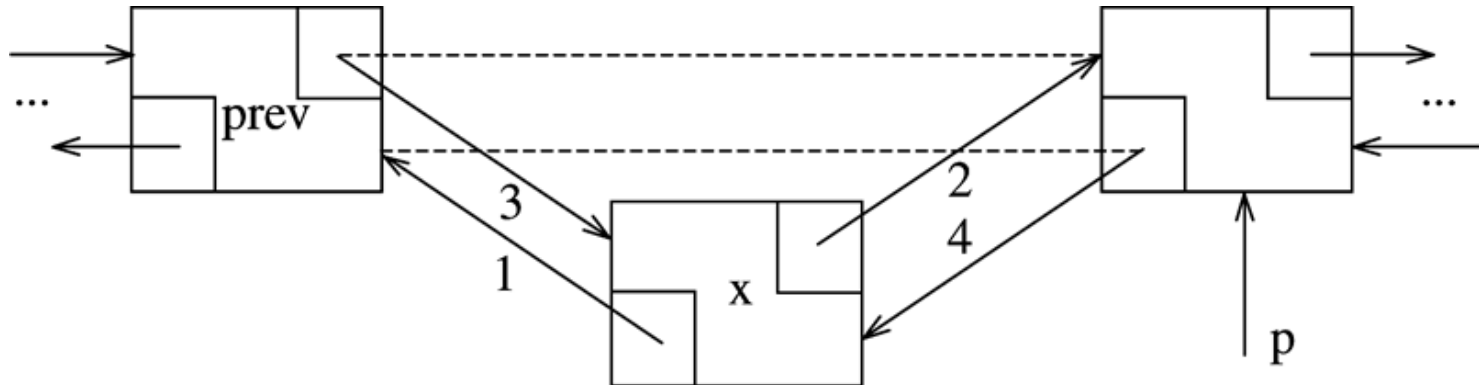


Do we need to explicitly de-allocate p?

Example: Inserting a Node

- Insert a new node before p.

What if p is the first element? Last element?
What if p == null?
What if p is not part of a list?



Other Operations

Either simple linked list or doubly linked list

- n Count the occurrences of an item in the linked list
- n Remove all occurrences of an item
- n Reverse a linked list (trivial for doubly linked list)
- n Duplicate a linked list