

# Graph Traversal

EECS 233

# Graph Traversals

- Traversing a graph involves starting at some vertex and visiting all of the vertices that can be reached from that vertex.
  - visiting a vertex = processing its data in some way
  - if the graph is connected, all of the vertices will be visited

- Example:
  - A Web crawler (vertices = pages, edges = links)

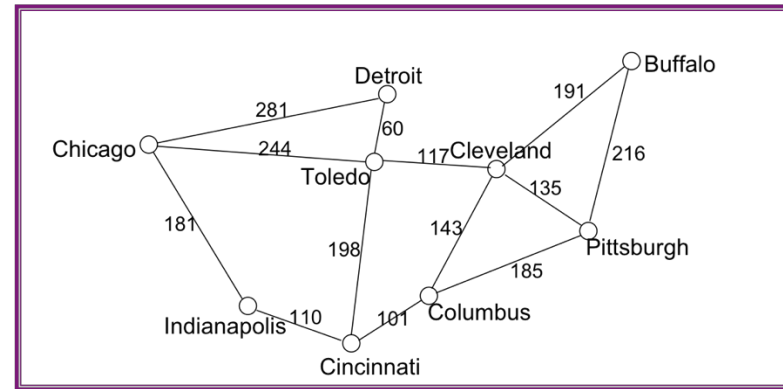
- We will consider two types of traversals:

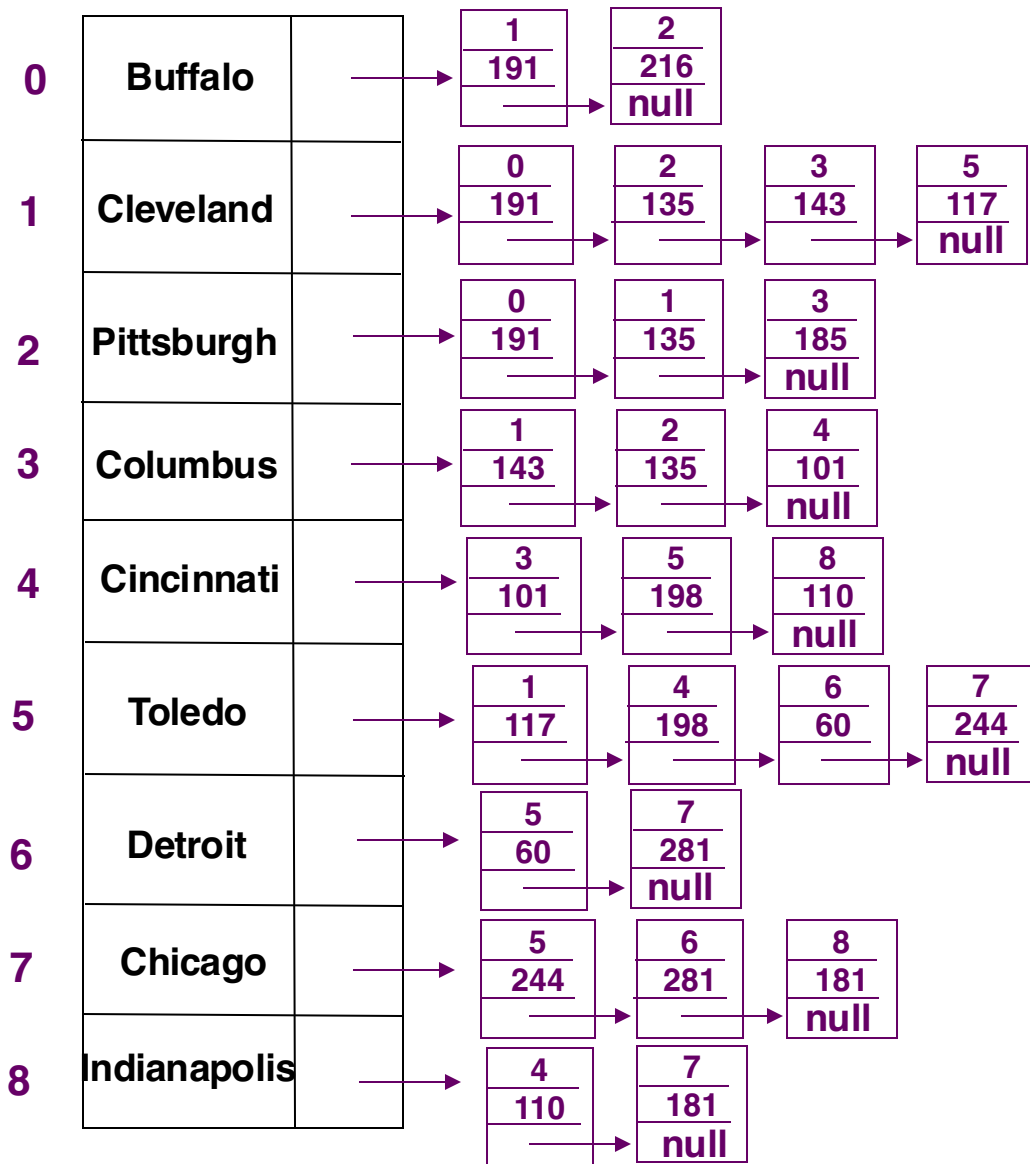
- **depth-first:**

- Visit the starting vertex
- Proceed as far as possible along a given path (via a neighbor) before “backtracking” and going along the next path

- **breadth-first:**

- visit the starting vertex
- visit all of its neighbors
- visit all unvisited vertices 2 edges away
- visit all unvisited vertices 3 edges away, etc.



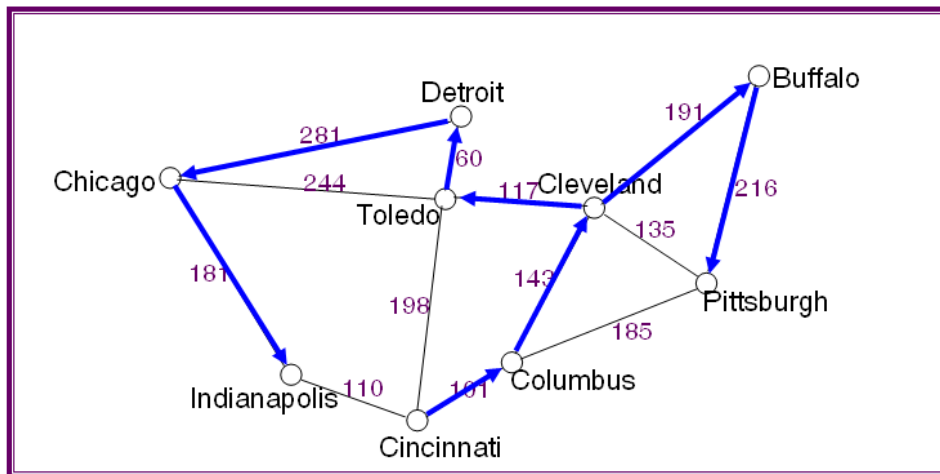


# Depth-first Traversal

- Visit a vertex, then make recursive calls on all of its yet-to-be-visited neighbors:

*depthFirstTrav(v)*  
*myDepthFirstTrav(v, NULL)*

*myDepthFirstTrav(node, parent)*  
*visit node and mark it as visited*  
*node.parent = parent*  
*for each vertex w in node's neighbors*  
*if (w has not been visited)*  
*myDepthFirstTrav(w, node)*



```
public class Graph {
    class Vertex {
        private String id;
        private linkedList <Edge> edges;
        // adjacency list
        private boolean encountered;
        private boolean done;
        private Vertex parent;
        private double cost;
        ...
    }
    class Edge {
        private int endNode;
        private double cost;
        ...
    }
    private Vertex[] vertices;
    private int numVertices;
    private int maxNum;
    ...
}
```

# Depth-first Traversal

- Visit a vertex, then make recursive calls on all of its yet-to-be-visited neighbors:

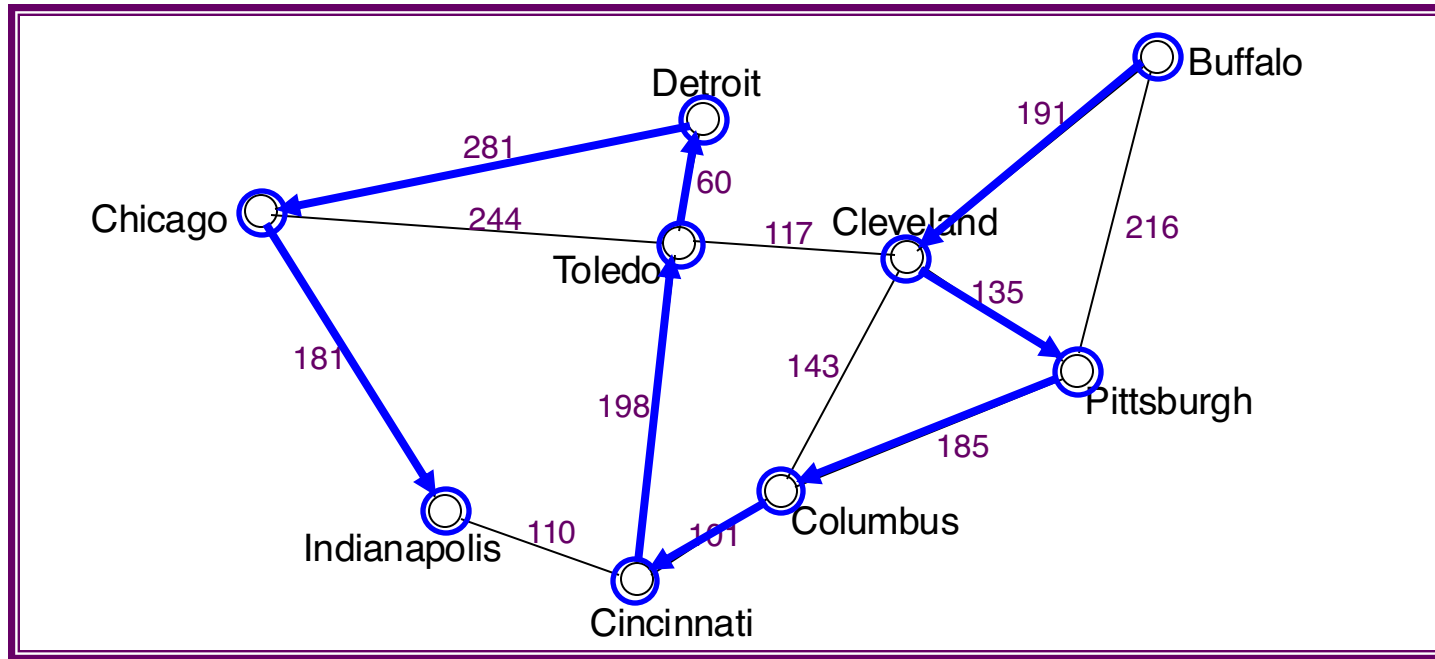
```
myDepthFirstTrav(v, parent)  
visit v and mark it as visited  
v.parent = parent  
for each vertex w in v's neighbors  
if (w has not been visited)  
myDepthFirstTrav(w, v)
```

- Implementation:

```
public void myDepthFirstTrav(int i, int parent) {  
    System.output.println(vertices[i].id);  
    vertices[i].encountered = true;  
    vertices[i].parent = parent;  
    Iterator<Edge> edgeIter = edges.iterator();  
    while (edgeIter.hasNext()) {  
        Edge curEdge = edgeIter.next();  
        j = curEdge.endNode;  
        if (vertices[j].encountered == false)  
            myDepthFirstTrav(j, i);  
    }  
}
```

```
public class Graph {  
    class Vertex {  
        private String id;  
        private linkedList <Edge> edges;  
        // adjacency list  
        private boolean encountered;  
        private boolean done;  
        private Vertex parent;  
        private double cost;  
        ...  
    }  
    class Edge {  
        private int endNode;  
        private double cost;  
        ...  
    }  
    private Vertex[] vertices;  
    private int numVertices;  
    private int maxNum;  
    ...  
}
```

# Example: From Buffalo



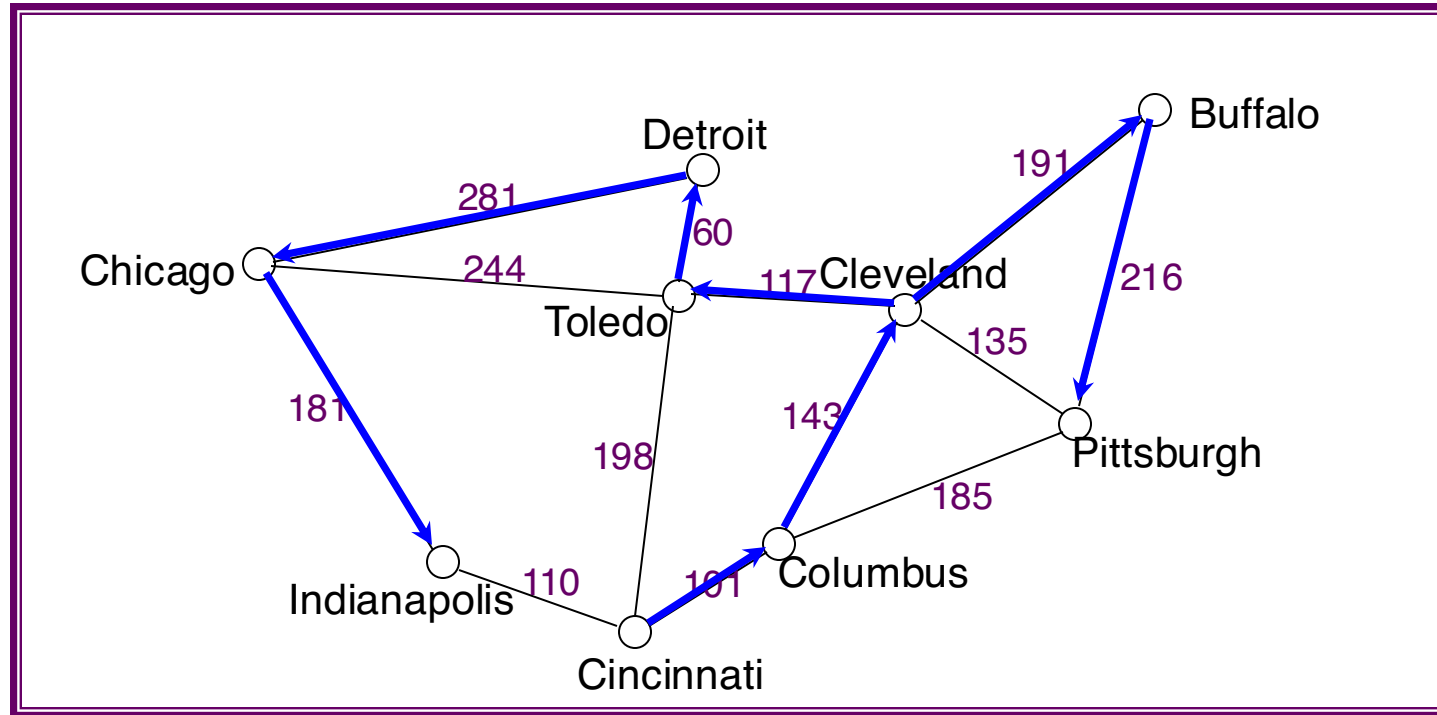
```

myDepthFirstTrav("Buffalo", null)
  visit Buffalo; Buffalo.parent = NULL
  w = "Cleveland"
  myDepthFirstTrav("Cleveland", "Buffalo")
    visit Cleveland; Cleveland.parent = Buffalo
    w = "Pittsburgh"
    myDepthFirstTrav("Pittsburgh", "Cleveland")
      visit Pittsburgh; Pittsburgh.parent = Cleveland
      w = "Buffalo"; Buffalo has been visited
      w = "Columbus"
      myDepthFirstTrav("Columbus", "Pittsburgh")
        visit Columbus; Columbus.parent = Pittsburgh;
        w = "Cincinnati"
        myDepthFirstTrav("Cincinnati", "Columbus")
          ...
  
```

```

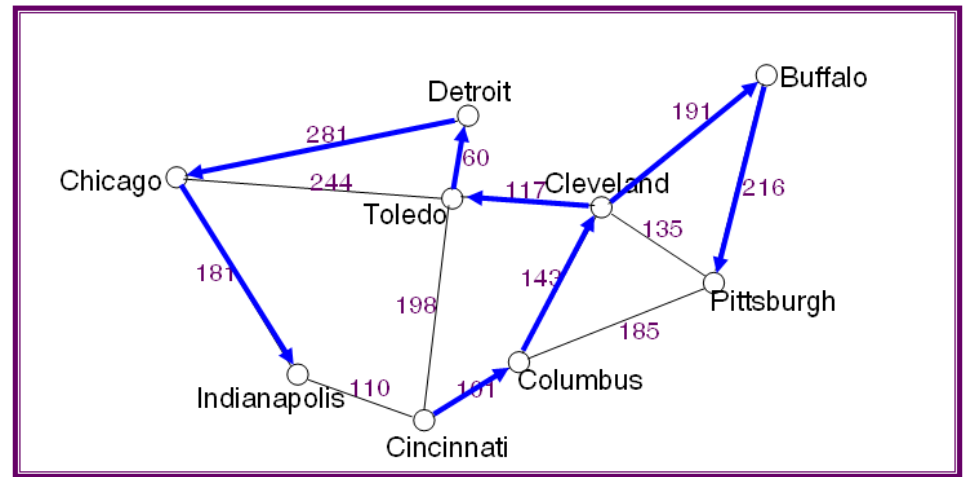
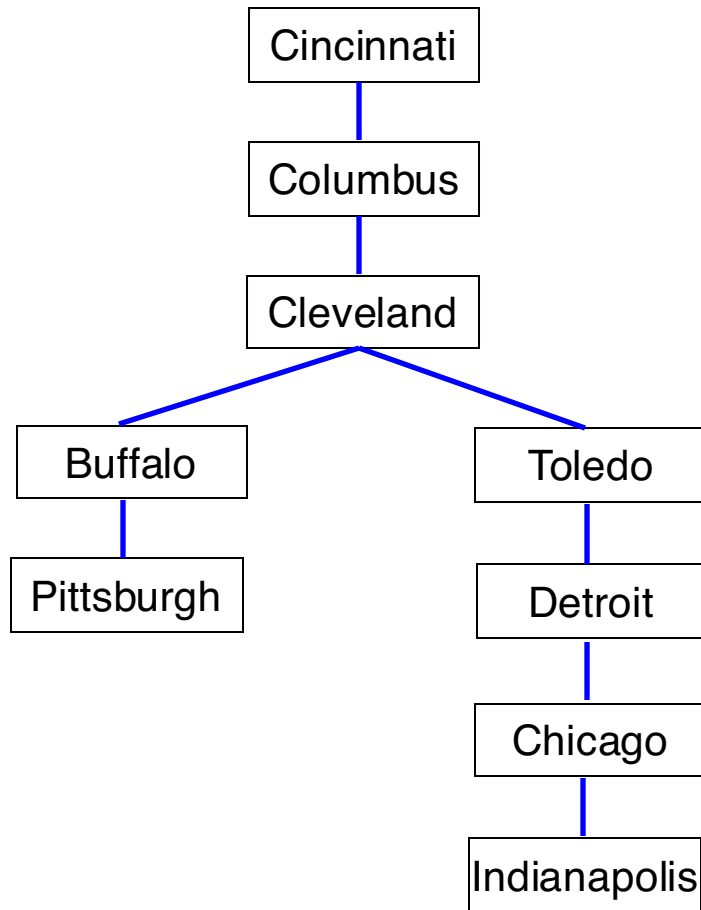
myDepthFirstTrav(v, parent)
  visit v and mark it as visited
  v.parent = parent
  for each vertex w in v's neighbors
    if (w has not been visited)
      myDepthFirstTrav(w, v)
  
```

# Example: From Cincinnati



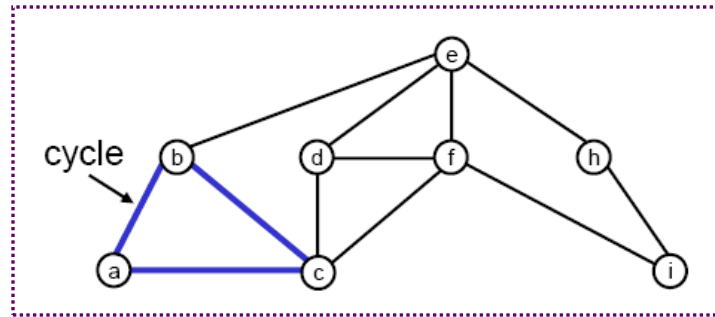
**Order:** Cincinnati, Columbus, Cleveland, Buffalo, Pittsburgh, Toledo, Detroit, Chicago, Indianapolis

# Depth-First Spanning Trees





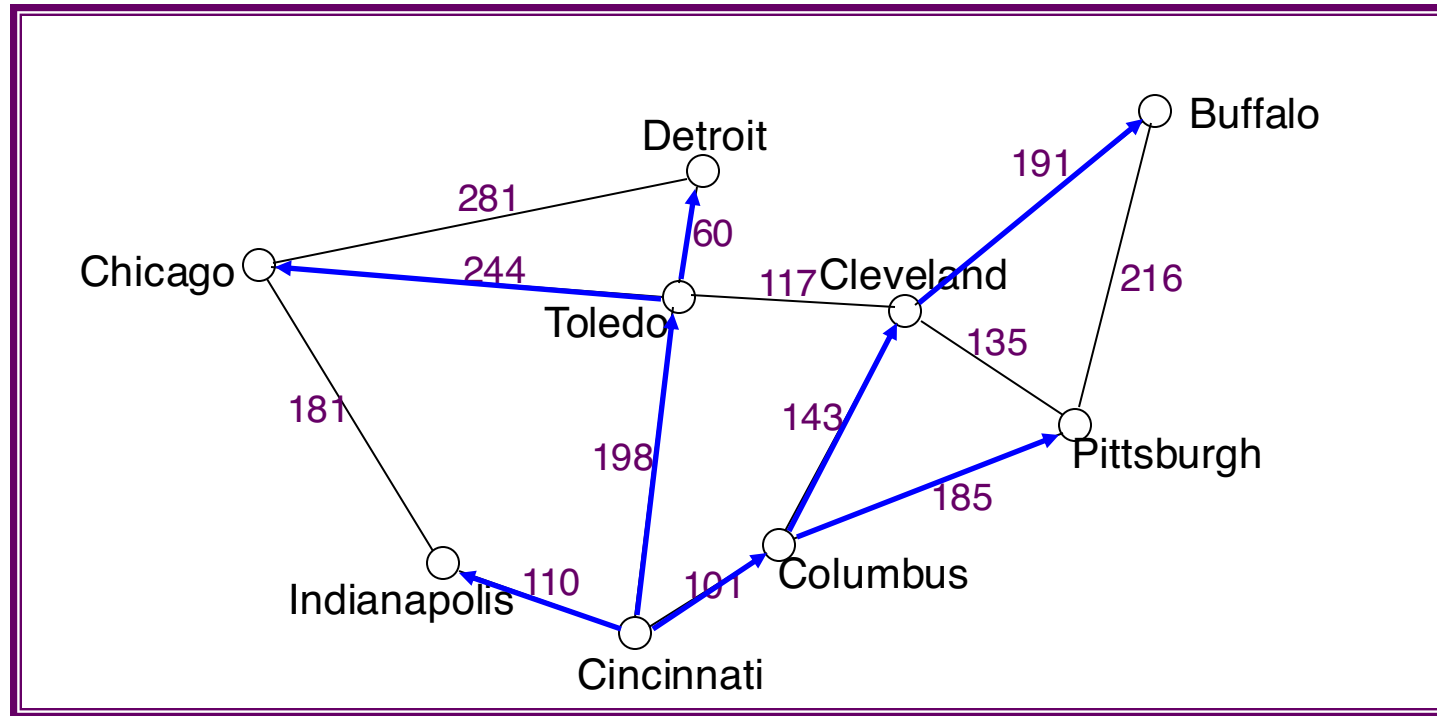
# An Application: Checking for Cycles



To discover a cycle:

- perform a depth-first traversal
  - when considering the non-parent neighbors of a current vertex, if we discover one that is already marked as visited, there is a cycle
- If we discover no cycles during the course of the traversal, the graph is acyclic.

# Breadth-First Traversal



**Order:** Cincinnati, Columbus, Toledo, Indianapolis, Pittsburgh, Cleveland, Detroit, Chicago, Buffalo

- Starting from Cincinnati, what would be the order of visits?
- **breadth-first:**
  - visit a vertex
  - visit all of its neighbors
  - visit all unvisited vertices 2 edges away
  - visit all unvisited vertices 3 edges away, etc.

# Breadth-First Traversal Method (pseudo code)

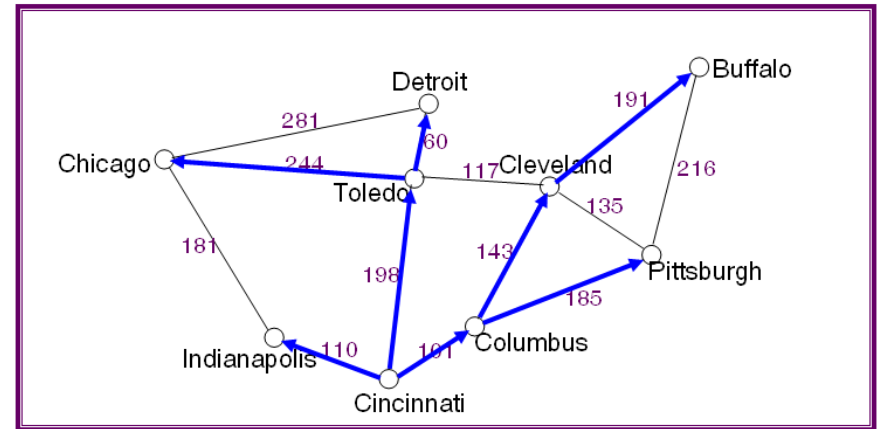
- Use a queue, as we did for level-order tree traversal:

**bfTrav(origin)**

**origin.parent = null**  
**create a new queue q**  
**q.insert(origin)**

**while (!q.isEmpty())**  
    **v = q.remove()**  
    **visit v**

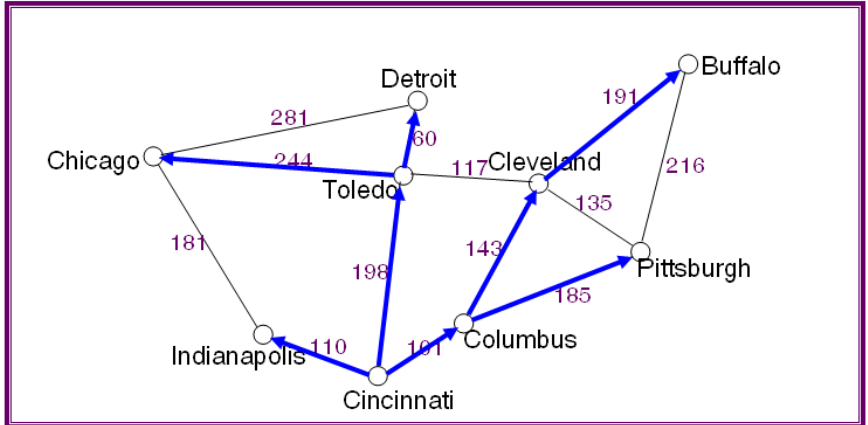
**for each vertex w in v's neighbors**  
        **if w is not encountered**  
            **mark w as encountered**  
            **w.parent = v;**  
            **q.insert(w);**



# Example: From Cincinnati

## □ Tracing the queue operations

```
insert("Cincinnati");  
remove("Cincinnati"); // and print it  
insert("Columbus");  
insert("Toledo");  
insert("Indianapolis");  
remove("Columbus"); // and print it  
insert("Cleveland");  
insert("Pittsburgh");  
remove("Toledo"); // and print it  
insert("Detroit");  
insert("Chicago");  
remove("Indianapolis"); // and print it  
remove("Cleveland"); // and print it  
insert("Buffalo");  
remove("Pittsburgh"); // and print it  
remove("Detroit"); // and print it  
remove("Chicago"); // and print it  
remove("Buffalo"); // and print it
```



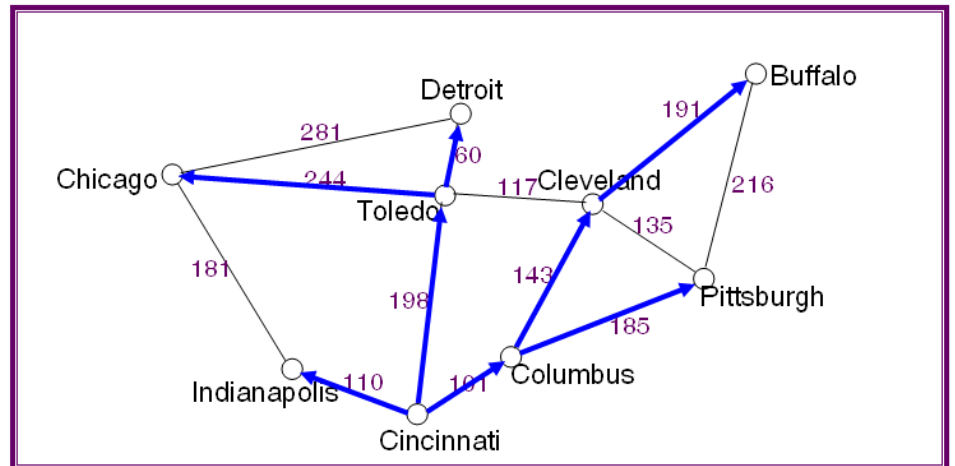
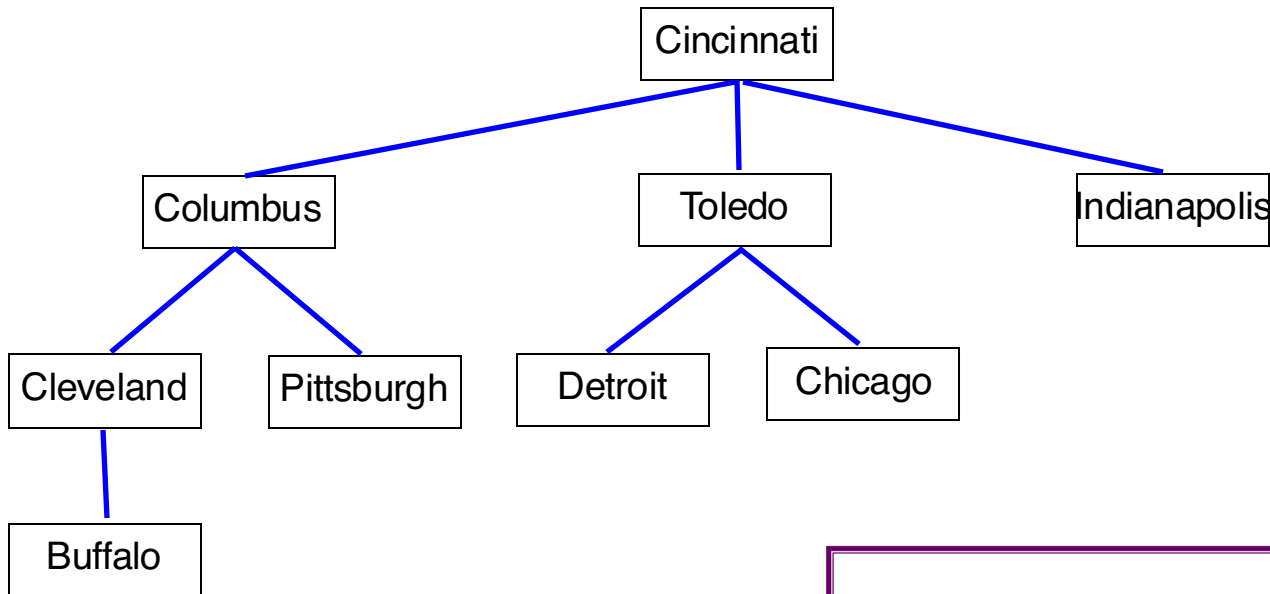
### bfTrav(origin)

```
origin.parent = null  
create a new queue q  
q.insert(origin)
```

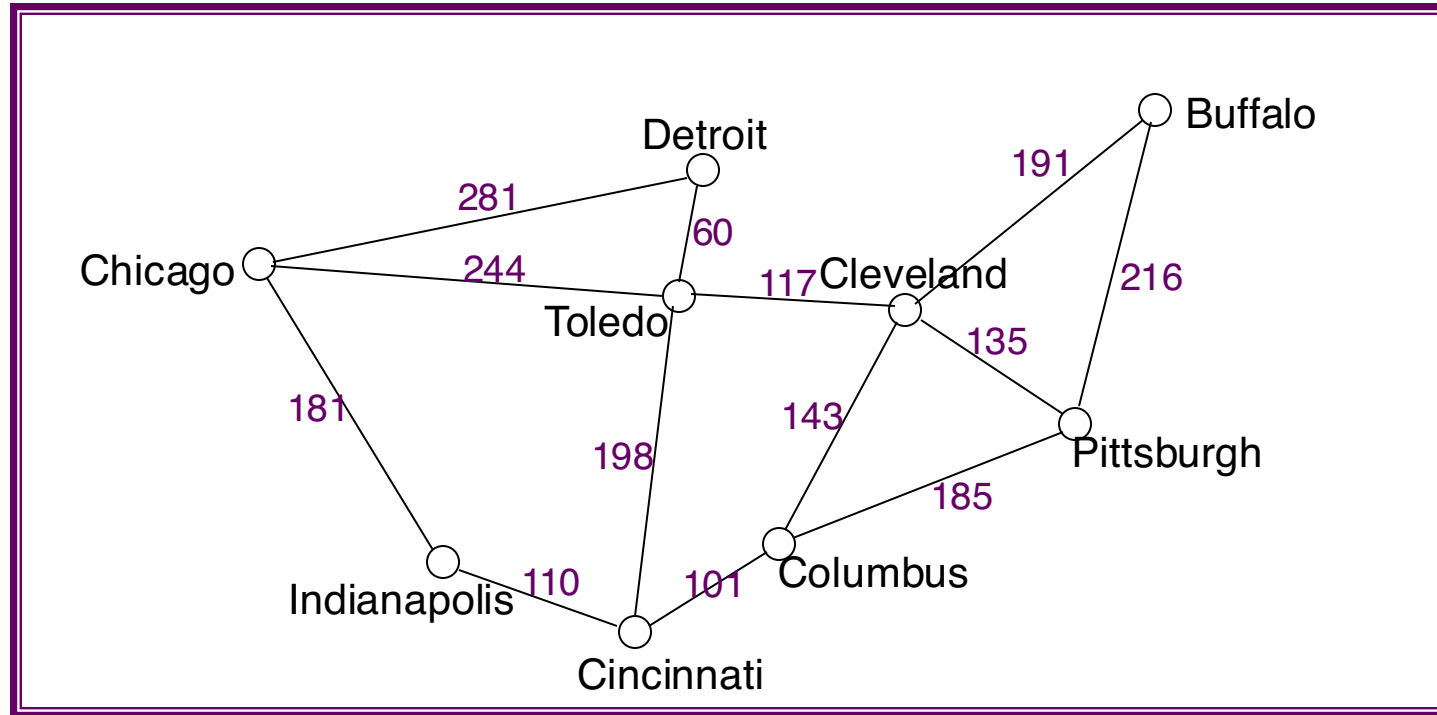
```
while (!q.isEmpty())  
    v = q.remove()  
    visit v
```

```
for each vertex w in v's neighbors  
    if w is not encountered  
        mark w as encountered  
        w.parent = v;  
        q.insert(w);
```

# Breadth-First Spanning Trees



# Example: From Cleveland



# Running Time of Graph Traversal

- Let  $V$  = number of vertices in the graph, and  $E$  = number of edges
- Assume we use an adjacency list as the data structure
- A traversal requires  $O(V + E)$  steps.
  - visit each vertex once
  - traverse each vertex's adjacency list at most once
    - the total length of the adjacency lists is  $2E = O(E)$
  - for a dense graph,  $E \rightarrow V^2$ , so the worst-case bound is  $O(V^2)$
  - for a sparse graph,  $O(V + E) \ll O(V^2)$