

Topics and Sample Questions for EECS 302 Exam 1:

1. **Topics:** Sets, Cartesian Products, Relations and Functions

Sample Problem: Given $A = \{1, 2, 4, 7\}$ and $B = \{2, 4, 5\}$, check ALL boxes that are applicable in describing the following sets and briefly justify your answers

The set in each row represents.....	Relation	Function	Neither	Justification
V: $\{(1,2), (4,2), (7,5), (4,2)\}$	✓			$V \subset A \times B$ (relation) but $2 \in A$ is not mapped to B (not function)
W: $\{(1,2), (4,2), (2,5), (7,5)\}$	✓	✓		$W \subset A \times B$ (relation) and each member of A has exactly one map to B (function)
X: $\{(1,2), (4,2), (2,3), (7,5)\}$			✓	$X \not\subset A \times B$ (not a relation hence not a function)
Y: $\{(1,2), (4,2), (2,4), (7,5), (4,4)\}$	✓			$Y \subset A \times B$ (relation) but $4 \in A$ has two mappings in B (not a function)
Z: $\{(1,2), (4,2), (2,4), (7,5), (4,2)\}$	✓	✓		$Z \subset A \times B$ (relation) and each member of A has exactly one map to B (a function)

2. **Topics:** Statements (or propositions) and logical operators ($\sim, \wedge, \vee, \rightarrow, \leftrightarrow$); Using logical operators to create negation of a statement, compound statements, and conditional statements; Truth values, truth tables and logical equivalences; Informal language and formal Language

Sample Problems:

2.1 Use logical identities (logical equivalence laws such as Associative laws, Commutative laws, Distributive laws, De Morgan's laws, Double negative laws, Idempotent laws, Absorption laws, and logical identities involving \rightarrow and \leftrightarrow) to show logical equivalences of the following conditional statements: Let p, q , and r be proposition variables. Then

$$p \vee \sim q \rightarrow r \vee q \equiv (\sim p \wedge q) \vee (r \vee q) \equiv \text{negation of } \sim(\sim p \wedge q) \wedge (\sim r \wedge \sim q)$$

$$\begin{aligned} p \vee \sim q \rightarrow r \vee q &\equiv \sim(p \vee \sim q) \vee (r \vee q) \\ &\equiv (\sim p \wedge q) \vee (r \vee q) \end{aligned}$$

Identity involving " \rightarrow "
De Morgan's and Double negative laws

$$\begin{aligned} \sim(\sim(\sim p \wedge q) \wedge (\sim r \wedge \sim q)) &\equiv \sim((p \vee \sim q) \wedge (\sim r \wedge \sim q)) \\ (\text{negation of } \sim(\sim p \wedge q) \wedge (\sim r \wedge \sim q)) &= (\sim(p \vee \sim q)) \vee \sim(\sim r \wedge \sim q) \\ &\equiv (\sim p \wedge q) \vee (r \vee q) \end{aligned}$$

De Morgan's and Double negative laws
De Morgan's law

De Morgan's and Double negative laws

QED

2.2 Use a truth table to reconfirm the logical equivalence of the first two statements above

p	q	r	$\sim q$	$p \vee \sim q$	$r \vee q$	$p \vee \sim q \rightarrow r \vee q$	$\sim p$	$\sim p \wedge q$	$(\sim p \wedge q) \vee (r \vee q)$
T	T	T	F	T	T	T	F	F	T
T	T	F	F	T	F	F	F	F	F
T	F	T	T	T	T	T	F	F	T
T	F	F	T	T	F	F	F	F	F
F	T	T	F	F	T	T	T	T	T
F	T	F	F	F	F	T	T	T	T
F	F	T	T	T	T	T	T	F	T
F	F	F	T	T	F	F	T	F	F

$$\therefore p \vee \sim q \rightarrow r \vee q \equiv (\sim p \wedge q) \vee (r \vee q)$$

Same

3. **Topics:** Arguments: Proof of validity using (1) Truth Tables or (2) Rules of Inference/Valid Argument Forms (namely Modus Ponens, Modus Tollens, Generalization, Specialization, Elimination, Transitivity, Proof by Division into Cases, and Contradiction Rule)

Sample problems:

Given the following argument in informal language (English)

Sam is a CS major or an economics major.

If Sam is a CS major, Sam is required to take EECS 302

\therefore Sam is an economics major or Sam is not required to take EECS 302

- 3.1 Write the above argument using formal language (using appropriate logical operators).

Let $p \equiv$ "Sam is a CS major"; $q \equiv$ "Sam is an Econ major"; $r \equiv$ "Sam takes EECS 302".

\therefore Argument: $p \vee q$
 $p \rightarrow r$
 $\therefore q \vee \sim r$

- 3.2 Use appropriate rules of inference to deduce whether the argument is valid. (If you think it is invalid you can provide a case to show invalidity).

The argument is invalid since the case of $(p, \sim q, r)$ makes both premises true but the conclusion $q \vee \sim r$ is false.

Although finding a case to show invalidity of the argument suffices, we will show how to use the rules of inference to do this.

We will consider the case $\sim q$
 We have $p \vee q$

$\sim q$
 $\therefore p$ (by Elimination)

$p \rightarrow r$

$\therefore r$ (Modus Ponens)

and $\sim q \vee \sim r$
 $\sim q$ $\therefore \sim r$ (Elimination)

$\therefore p \vee q$
 $p \rightarrow r$
 $q \vee \sim r$
 $\sim q$
 $\therefore r \wedge \sim r = C$

This shows the invalidity of the argument

- 3.3 Use a truth table to confirm the conclusion found in 3.2.

p	q	r	$p \vee q$	$p \rightarrow r$	$\sim r$	$q \vee \sim r$
T	T	T	T	T	F	T
T	T	F	T	F	T	T
T	F	T	T	T	F	(F) ←
T	F	F	T	F	T	T
F	T	T	T	T	F	T
F	T	F	T	T	T	T
F	F	T	F	T	F	
F	F	F	F	T	T	

This critical line leads to false conclusion
 Hence the argument is invalid.

4. **Topics:** Quantified statements (including quantified conditional statements and multiply-quantified statements) and negation of quantified statements

Sample problem:

Consider the following (multiply quantified) statement S :

$$S: \forall e \in \mathbb{Z}^+, \exists m \in \mathbb{N} \text{ such that } \forall n \in \mathbb{N}, \text{ if } n \geq m, \text{ then } \forall x \in \mathbb{R}, |f_n(x) - f(x)| < e.$$

- 4.1. What is the negation of the statement S ?

$$\sim S = \exists e \in \mathbb{Z}^+, \forall m \in \mathbb{N}, \exists n \in \mathbb{N} \text{ such that } n \geq m \text{ and } (\exists x \in \mathbb{R} \text{ such that } |f_n(x) - f(x)| \geq e)$$

- 4.2. What is the contrapositive form of the statement S ? (Hint: Just concentrate on the if-then part)

Contrapositive form of S

$$= \forall e \in \mathbb{Z}^+, \exists m \in \mathbb{N}, \forall n \in \mathbb{N} \text{ such that if } \exists x \in \mathbb{R} \text{ such that } |f_n(x) - f(x)| \geq e, \text{ then } n < m$$

5. **Topics:** Relationships among: i) logical conditioning quantifiers \rightarrow and \leftrightarrow ; ii) conditional informal language operators ("if", "only if" and "if and only if"); iii) implicit quantifiers \Rightarrow and \Leftrightarrow ; and (iv) necessary conditions, sufficient conditions and necessary and sufficient conditions

Sample problems:

Let $f(x)$ be a twice continuously differentiable real-valued function of a real variable x (i.e. the domain of f is \mathbb{R} and the range (co-domain) of f is also \mathbb{R}). The following results are well known:

- A necessary condition for a real x^* to be a minimizer of f is that $f'(x^*) = 0$ and $f''(x^*) \geq 0$
- A sufficient condition for a real x^* to be a strict local minimizer of f is that $f'(x^*) = 0$ and $f''(x^*) > 0$
- If f is a convex function, then a necessary and sufficient condition for a real x^* to be a minimizer of f is that $f'(x^*) = 0$

Write each of the above statements in a formal form using first (\rightarrow or \leftrightarrow) then (\Rightarrow or \Leftrightarrow):

Answer:

- a) $\forall x \in \mathbb{R}, x \text{ is a minimizer of } f(x) \rightarrow f'(x) = 0 \text{ and } f''(x) \geq 0$
 or
 A real x is a minimizer of $f(x) \Rightarrow f'(x) = 0$ and $f''(x) \geq 0$

Now write similar formal statements for b) and c).

- b) $\forall x \in \mathbb{R}, f'(x) = 0 \text{ and } f''(x) > 0 \rightarrow x \text{ is a strict local minimizer of } f.$
 or For a real x such that $f'(x) = 0$ and $f''(x) > 0 \Rightarrow x \text{ is a strict local minimizer of } f$
- c) $\forall x \in \mathbb{R}, x \text{ is a minimizer of a convex function } f \leftrightarrow f'(x) = 0$
 or A real x is a minimizer of a convex function $f \Leftrightarrow f'(x) = 0$

It is often easier to prove the above results by a contrapositive proof.

What is the contrapositive form of the necessary condition in part a) above?

Contrapositive form of the necessary condition in part a) is
such that either $f'(x) \neq 0$ or $f''(x) < 0$, x is not a minimizer of $f(x)$

5. **Topics:** Direct proof, counterexample, proof by contradiction, contrapositive proof

Sample problems:

5.1 Prove or disprove the following statements/claims formally and briefly justify each step in the body of your proof. Also clearly specify which method of proof you are using.

- a) Claim: There exists an integer $m \geq 4$ such that $2m^2 - 5m + 2$ is prime

We will disprove the claim by showing a counterexample i.e. will show

\forall integer $m \geq 4$, $2m^2 - 5m + 2$ is not a prime.

Since $2m^2 - 5m + 2 = (2m-1)(m-2)$ and both $(2m-1)$ and $(m-2)$ are integers greater than 1 for all $m \geq 4$.

$\therefore 2m^2 - 5m + 2$ is a composite. QED

- b) Claim: The difference between any odd integer and an even integer is odd

We will use a direct proof.

Let n be an odd number and m be an even number.

$\therefore n = 2k_1 + 1$ and $m = 2k_2$ for some integers k_1 and k_2

$\therefore n - m = 2k_1 + 1 - 2k_2 = 2(k_1 - k_2) + 1 = 2r + 1$

integer r (axiom of algebra)

$\therefore n - m$ is odd by definition. QED

- c) Claim: If a sum of two real numbers is less than 70, then at least one of the numbers is less than 35. We will prove by a contrapositive proof

Suppose x and y are real numbers greater than 35 i.e. $x \geq 35$ and $y \geq 35$.
(or equal to)

$\therefore x + y \geq 35 + 35 = 70$

\therefore The supposition is wrong

\therefore At least one of x and y must be less than 35. QED

- d) Claim: If x and y are rational numbers with $y \neq 0$, and if z is an irrational number, then $x + yz$ is irrational. We will prove this by contradiction.

Let x and y be rational numbers with $y \neq 0$

$\therefore x = \frac{m}{n}$ and $y = \frac{k}{l}$ for some integers m, n, k, l with $n \neq 0, k \neq 0, l \neq 0$

Assume that $x + yz$ is rational (where z is irrational)

$\therefore x + yz = \frac{r}{s}$ for some integers r and s with $s \neq 0$

$\therefore \frac{m}{n} + \frac{k}{l}z = \frac{r}{s} \therefore z = \left(\frac{r}{s} - \frac{m}{n}\right) \frac{l}{k} = \frac{l(rk - sm)}{snk}$ integer M

$\therefore z$ is rational

A contradiction since z is irrational $\therefore x + yz$ is irrational QED

5.2 Provide a brief justification for each claim made in the following body of proof

Theorem: There are infinitely many primes.

Proof:

Suppose there are finitely many prime with p being the largest prime.

Let $M = p! + 1$. Clearly every prime divides $p!$.

But none divides M (Due to a theorem which says if $p|a$, $p \nmid a+1$).

However, since M is an integer greater than 1, M is divisible by a prime.

Let q be a prime that divides M . So $q > p$.

So the supposition that p is the greatest prime is false.

Thus, there must be infinitely many primes. Hence the theorem is proved

Reason

Contrapositive argument

Since every prime $\leq p$ and

is a factor of $p!$

Due to a theorem which says

every integer > 1 is divisible by a prime

p cannot be the largest prime

Contrapositive argument

6. **Topics:** Key results and concepts in Number Theory: the quotient-remainder theorem, div and mod, floor and ceiling, the uniqueness factorization theorem

Sample Problems:

6.1 For an odd integer n , show that $\left\lfloor \frac{n^2}{4} \right\rfloor = \left(\frac{n-1}{2} \right) \left(\frac{n+1}{2} \right)$

Let n be an odd integer

$\therefore n = 2k+1$ for some integer k

$\therefore \frac{n^2}{4} = \frac{4k^2 + 4k + 1}{4} = k^2 + k + \frac{1}{4} \Rightarrow \left\lfloor \frac{n^2}{4} \right\rfloor = k^2 + k$. Since $k^2 + k \leq k^2 + k + \frac{1}{4} < k^2 + k + 1$

But $k^2 + k = k(k+1) = \frac{2k}{2} \cdot \frac{2k+2}{2} = \frac{2k+1-1}{2} \cdot \frac{2k+1+1}{2} = \frac{n-1}{2} \cdot \frac{n+1}{2}$

$\therefore \left\lfloor \frac{n^2}{4} \right\rfloor = \left(\frac{n-1}{2} \right) \left(\frac{n+1}{2} \right)$

6.2 Show that if $n \bmod 15 = 3$, then $10n \bmod 15 = 0$.

Since $n \bmod 15 = 3$, $n = 15k + 3$ for some integer k (By QR Theorem)

$\therefore 10n = 150k + 30 = (10k+2)15 + 0$

$\therefore 10n \bmod 15 = 0$

6.3 If today is Tuesday, use the quotient-remainder theorem to find out the day of the week 127 days from now.

Let Tuesday = 0, Wednesday = 1, ..., Monday = 6

Let d be the day of the week 127 days from now (which is Tuesday)

$\therefore d = 127 \bmod 7 = 1$

Thus 127 days from now will be Wednesday.

6.4 Find the smallest value of k such that $2^5 \cdot 5^1 \cdot 7^1 \cdot 11^3 \cdot k$ is a perfect cube.

Let $k = 2^1 \cdot 5^2 \cdot 7^2 \cdot 11^0 = 2450$

$\therefore 2^5 \cdot 5^1 \cdot 7^1 \cdot 11^3 \cdot k = 2^6 \cdot 5^3 \cdot 7^3 \cdot 11^3 = (2^2 \cdot 5^1 \cdot 7^1 \cdot 11^1)^3 = (75460)^3$

Note any other number k that will make the given number a perfect cube will have to be larger than 2450 (by the unique factorization theorem)

$\therefore 2450$ is the smallest required number.