

# Notes

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## 1 The type theory

The type theory is a two-level type theory, meaning it contains two notions of equality.

## 2 Algorithmic Definitions

We present an algorithmic procedure for deciding the judgments of the type theory.

**Definition 2.1.** Head reduction is defined as follows:

$$\begin{aligned} \Gamma \triangleright (\lambda x. e) f &\rightarrow e[f/x] \\ \Gamma \triangleright ef &\rightarrow e'f \quad \text{if } \Gamma \triangleright e \rightarrow e' \\ \Gamma \triangleright [\alpha \rightarrow e, \beta \rightarrow f] &\rightarrow e \quad \text{if } \Gamma \vdash \alpha \\ \Gamma \triangleright [\alpha \rightarrow e, \beta \rightarrow e] &\rightarrow f \quad \text{if } \Gamma \vdash \beta \\ \Gamma \triangleright \text{out}_\psi(a) &\rightarrow t \quad \text{if } \Gamma \triangleright a \uparrow A[\psi \rightarrow t] \text{ and } \Gamma \vdash \psi \\ \Gamma \triangleright \text{out}_\psi(\text{in}(a)) &\rightarrow a \\ \Gamma \triangleright \text{out}_\psi(a) &\rightarrow a' \quad \text{if } \Gamma \triangleright a \rightarrow a' \\ \Gamma \triangleright \text{out}_{[\alpha]}(\text{in}(a)) &\rightarrow a \\ \Gamma \triangleright \text{out}_{[\alpha]}(a) &\rightarrow a' \quad \text{if } \Gamma \triangleright a \rightarrow a' \end{aligned}$$

**Definition 2.2.** Term equivalence is defined as follows:

$$\begin{aligned} \Gamma \triangleright M \Leftrightarrow N : A[\psi \rightarrow t] &\quad \text{if } \Gamma \triangleright \text{out}_\psi(M) \Leftrightarrow \text{out}_\psi(N) : A \\ \Gamma \triangleright M \Leftrightarrow N : [\alpha] \rightarrow A &\quad \text{if } \Gamma, \alpha \triangleright \text{out}_{[\alpha]}(M) \Leftrightarrow \text{out}_{[\alpha]}(N) : A \\ \Gamma \triangleright M \Leftrightarrow N : (x : \mathbb{I}) \rightarrow A &\quad \text{if } \Gamma, x : \mathbb{I} \triangleright Mx \Leftrightarrow Nx : A(x) \\ \Gamma \triangleright M \Leftrightarrow N : (x : \mathbf{cof}) \rightarrow A &\quad \text{if } \Gamma, x : \mathbf{cof} \triangleright Mx \Leftrightarrow Nx : A(x) \end{aligned}$$

**Definition 2.3.** Neutral equivalence is defined as follows:

$$\begin{array}{ll}
\Gamma \triangleright [\alpha \rightarrow p_1, \beta \rightarrow p_2] \leftrightarrow q \uparrow A & \text{if } \Gamma, \alpha \triangleright p_1 \Leftrightarrow q : A \text{ and } \Gamma, \beta \triangleright p_2 \Leftrightarrow q : A \\
\Gamma \triangleright p \leftrightarrow [\alpha \rightarrow q_1, \beta \rightarrow q_2] \uparrow A & \text{if } \Gamma, \alpha \triangleright p \Leftrightarrow q_1 : A \text{ and } \Gamma, \beta \triangleright p \Leftrightarrow q_2 : A \\
\Gamma \triangleright \text{out}_\psi(p) \Leftrightarrow \text{out}_\psi(q) \uparrow A & \text{if } \Gamma \triangleright p \leftrightarrow q \uparrow A[\psi \rightarrow t] \\
\Gamma \triangleright \text{out}_{[\alpha]}(p) \Leftrightarrow \text{out}_\psi(q) \uparrow A & \text{if } \Gamma \triangleright p \leftrightarrow q \uparrow [\alpha] \rightarrow A
\end{array}$$