A   Proof of Theorem 3

Proof. The proof is presented here in two cases, first for the user v_i whose edges' privacy levels are not all 2:

$$
\begin{aligned}
\mathbb{E}[\hat{w}_i] &= \mathbb{E}[w_i + Lap(\frac{K_i/(2p_1 - 1)}{(1 - \beta)(1 - \alpha)\epsilon_1})] \\
&= \mathbb{E}[w_i^{(1)} + w_i^{(2)}] \\
&= \mathbb{E}[\frac{t_i^{(1)} - (1 - p_1)s_i^{(1)}}{2p_1 - 1} + \frac{t_i^{(2)} - (1 - p_2)s_i^{(2)}}{2p_2 - 1}] \\
&= \frac{1}{2p_1 - 1}\mathbb{E}[t_i^{(1)} - (1 - p_1)s_i^{(1)}] + \frac{1}{2p_2 - 1}\mathbb{E}[t_i^{(2)} - (1 - p_2)s_i^{(2)}].
\end{aligned}
\tag{8}
$$

Let $s_i^{*(g)}(g \in \{1, 2\})$ be the count of the 2-star counts of user v_i excluding the triangle part, i.e., $s_i^{*(g)} = |\{(v_i, v_j, v_k) : i < j < k, a_{i,j} = a_{i,k} = 1, (v_j, v_k) \notin E\}|$. For simplicity, it is assumed that the user does not perform an edge clipping operation. T_i denotes the real triangle counts in the local graph for user v_i. $t_i^{(g)}$ is the noisy triangle counts, i.e., $t_i = |\{(v_i, v_j, v_k) : i < j < k, a_{i,j} = a_{i,k} = 1, (v_j, v_k) \in E'\}|$. $s_i^{(g)}$ is the true 2-star counts for user v_i, then we have $s_i^{(g)} = s_i^{*(g)} + T_i^{(g)}$. By the properties of RR:

$$
\mathbb{E}[t_i^{(g)}] = T_i^{(g)}p + s_i^{*(g)}(1 - p),
\tag{9}
$$

so,

$$
\begin{aligned}
\mathbb{E}[\hat{w}_i] &= \frac{1}{2p_1 - 1}\mathbb{E}[t_i^{(1)} - (1 - p_1)s_i^{(1)}] + \frac{1}{2p_2 - 1}\mathbb{E}[t_i^{(2)} - (1 - p_2)s_i^{(2)}] \\
&= \frac{1}{2p_1 - 1}\mathbb{E}[T_i^{(1)}p_1 + s_i^{*(1)}(1 - p_1) - (1 - p_1)s^{(1)}] \\
&\quad + \frac{1}{2p_2 - 1}\mathbb{E}[T_i^{(2)}p_2 + s_i^{*(2)}(1 - p_2) - (1 - p_2)s^{(2)}] \\
&= \frac{(2p_1 - 1)T_i^{(1)}}{2p_1 - 1} + \frac{(2p_2 - 1)T_i^{(2)}}{2p_2 - 1} \\
&= T_i^{(1)} + T_i^{(2)} \\
&= T_i.
\end{aligned}
\tag{10}
$$

For the user v_i whose edges all have privacy level 2:

$$
\begin{aligned}
\mathbb{E}[\hat{w}_i] &= \mathbb{E}[w_i + Lap(\frac{K_i/(2p_2 - 1)}{(1 - \beta)(1 - \alpha)\epsilon_1})] \\
&= \mathbb{E}[w_i^{(2)}] \\
&= \frac{1}{2p_2 - 1}\mathbb{E}[t_i^{(2)} - (1 - p_2)s_i^{(2)}] \\
&= \frac{1}{2p_2 - 1}\mathbb{E}[T_i^{(2)}p_2 + s_i^{*(2)}(1 - p_2) - (1 - p_2)s^{(2)}] \\
&= \frac{(2p_2 - 1)T_i^{(2)}}{2p_2 - 1} \\
&= T_i^{(2)} \\
&= T_i.
\end{aligned}
\tag{11}
$$

B Proof of Theorem 4

For a user v_i whose associated edges all have a privacy level of 2:

Proof.

$$
\begin{aligned}
Var[\hat{w}_i] &= Var[w_i^{(2)} + Lap(\frac{K_i/(2p_2 - 1)}{(1 - \beta)(1 - \alpha)\epsilon_1})] \\
&= Var[\frac{t_i^{(2)} - (1 - p_2)s_i^{(2)}}{2p_2 - 1} + Lap(\frac{K_i/(2p_2 - 1)}{(1 - \beta)(1 - \alpha)\epsilon_1})] \\
&= \frac{1}{(2p_2 - 1)^2}Var[t_i^{(2)}] + 2(\frac{K_i/(2p_2 - 1)}{(1 - \beta)(1 - \alpha)\epsilon_1})^2 \\
&= \frac{p_2q_2(T_i^{(2)} + s_i^{*(2)})}{(2p_2 - 1)^2} + 2(\frac{K_i/(2p_2 - 1)}{(1 - \beta)(1 - \alpha)\epsilon_1})^2 \\
&= \frac{p_2q_2s_i^{(2)}}{(2p_2 - 1)^2} + 2(\frac{K_i/(2p_2 - 1)}{(1 - \beta)(1 - \alpha)\epsilon_1})^2 \\
&\leq \frac{p_2q_2}{(2p_2 - 1)^2}\frac{K_i(K_i - 1)}{2} + 2(\frac{K_i/(2p_2 - 1)}{(1 - \beta)(1 - \alpha)\epsilon_1})^2.
\end{aligned}
\tag{12}
$$

For a user v_i whose associated edges' privacy levels are not all 2:

Proof.

$$Var[\hat{w}_i] = Var[w_i^{(1)} + w_i^{(2)} + Lap(\frac{K_i/(2p_1-1)}{(1-\beta)(1-\alpha)\epsilon_1})]$$

$$= Var[\frac{t_i^{(1)} - (1-p_1)s_i^{(1)}}{2p_1-1} + \frac{t_i^{(2)} - (1-p_2)s_i^{(2)}}{2p_2-1} + Lap(\frac{K_i/(2p_1-1)}{(1-\beta)(1-\alpha)\epsilon_1})]$$

$$= \frac{1}{(2p_1-1)^2}Var[t_i^{(1)}] + \frac{1}{(2p_2-1)^2}Var[t_i^{(2)}] + 2(\frac{K_i/(2p_1-1)}{(1-\beta)(1-\alpha)\epsilon_1})^2$$

$$= \frac{p_1q_1(T_i^{(1)} + s_i^{*(1)})}{(2p_1-1)^2} + \frac{p_2q_2(T_i^{(2)} + s_i^{*(2)})}{(2p_2-1)^2} + 2(\frac{K_i/(2p_1-1)}{(1-\beta)(1-\alpha)\epsilon_1})^2$$

$$= \frac{p_1q_1s_i^{(1)}}{(2p_1-1)^2} + \frac{p_2q_2s_i^{(2)}}{(2p_2-1)^2} + 2(\frac{K_i/(2p_1-1)}{(1-\beta)(1-\alpha)\epsilon_1})^2$$

$$\leq \frac{p_1q_1(s_i^{(1)} + s_i^{(2)})}{(2p_1-1)^2} + 2(\frac{K_i/(2p_1-1)}{(1-\beta)(1-\alpha)\epsilon_1})^2$$

$$\leq \frac{p_1q_1}{(2p_1-1)^2}\frac{K_i(K_i-1)}{2} + 2(\frac{K_i/(2p_1-1)}{(1-\beta)(1-\alpha)\epsilon_1})^2.$$

$$(13)$$