Lecture Outline

Reminders to self:

- ☐ Turn on lecture recording to Cloud
- ☐ Turn on Zoom microphone

Last Lecture

- Finished ROM
- Programmable Logic Arrays
- Started Latches & Flip-Flops (groundwork for S-R Latch)

Today's Lecture

- Latches
- Gated Latches
- Flip-Flops



Handouts and Announcements

Announcements

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- Homework Problem 9-4
 - Posted on Carmen yesterday (2/21)
 - Due: 11:25am Monday 2/27
- Homework Reminder:
 - HW 9-1 past due
 - HW 9-2 & 9-3 due: 11:59pm Thursday 2/23
- Read for Friday: pages 352-358, 370, 375-380

Set-Reset Latch

Set-Reset Latch (SR Latch

- Can be made from two NOR gates with feedback in Cross-coupled form
- Term present state [Q(t)] used to denote state of Q output of the latch or flip-flop at time any input signal changes
- Term rext state $[Q(t + \varepsilon)]$ denotes state of Q output after Pres
 - latch or flip-flop has reacted to input change and stabilized
- Normally we will write next state with + superscript, without explicitly showing $t + \varepsilon$

hange	S			_)		Q
)] den	otes			R - L	_	
	Presen	t State		Next	State	
S	R	Q	P	Q^+	P^+	
0	0	0	1	0	1	} Hold state
0	0	1	0	1	0) HOLD 1141 E
0	1	0	1	D	1	} Reset state
0	1	1	0	D	J	
1	0	0	1	- 1	0	201511
1	0	1	0		0	} Set State

NOR = D IF

any is

Set-Reset Latch

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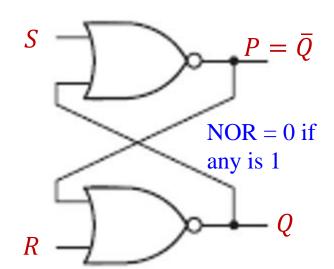
Set-Reset Latch (SR Latch)

$$S = \{1\}$$
 BAD!

- Makes both P = Q = 0, but then $P \neq \overline{Q}$
- If *S* and *R* are simultaneously changed to 0, both *P* and *Q* may both change to 1
- The 1s propagate through the NOR gates changing *P* and *Q* to 0
- Latch may continue to oscillate if the gate delays are equal
- Unstable : S = R = 1 must be avoided

The <u>designer</u> of the circuit using the SR Latch is responsible for avoiding it!



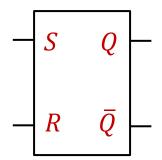


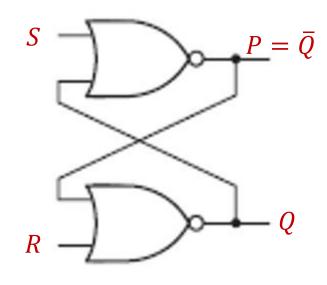
Set-Reset Latch

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Set-Reset Latch (SR Latch)

SR Latch Component





- S = 0, R = 0 "Hold" $Q^* = Q$
- S = 0, R = 1 "Reset" $Q^{+} = 0$
- S = 1, R = 0 "Set" $Q^{*} = 1$
- S = 1, R = 1 Do Not Allow Very Bad

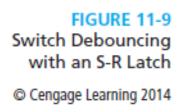
- Note: No clock input in component figure
- We will later use SR flip-flops, which will have a clock input
- That is how you can tell them apart

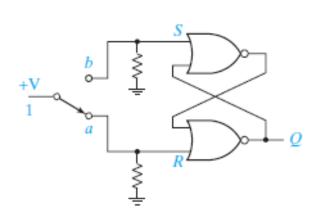
Set-Reset Latch

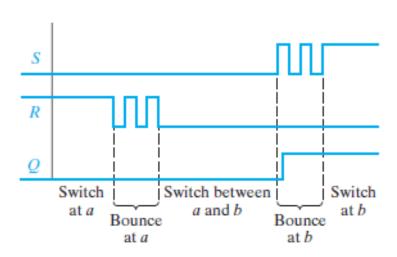
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Set-Reset Latch (SR Latch)

- A useful application of the SR Latch is switch debouncing
- Switch contacts tend to vibrate open and closed several times before settling to their final position
- Produces an electronically noisy transition
- This noise can interfere with proper operation of a logic circuit







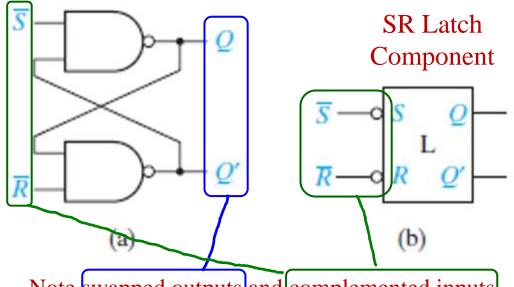
- This debouncing circuit works for double-throw switches (position a or position b)
- Debouncing a single-throw switch requires a different circuit



\bar{S} - \bar{R} Latch

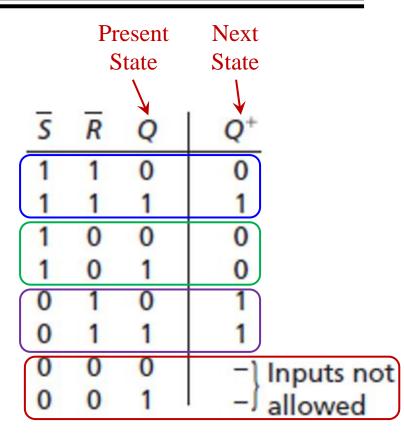
A NAND version is also possible

\bar{S} – \bar{R} Latch Circuit



Note swapped outputs and complemented in	ıputs
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- $\bar{S} = 1, \bar{R} = 1$ Hold
- $\bar{S} = 0, \bar{R} = 1$ Set
- $\bar{S}=1, \bar{R}=0$ Reset
- $\bar{S} = 0$, $\bar{R} = 0$ Unstable: Do Not Allow



Gated Latches

Gated Latches

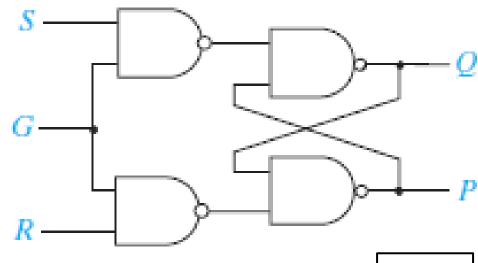
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- Gated latches have an additional input called the gate or enable input
- When the gate input is inactive
 - Which may be the high or low value
 - State of the latch cannot change
- When the gate input is active
 - Active-high or active-low
 - Latch is controlled by the other inputs
 - Operates as previously described

Gated S-R Latch

FIGURE 11-11 NAND-Gate Gated S-R Latch

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$$G=0 \implies X=Y=1$$
 Hold $G=1 \implies \begin{cases} X=ar{S} \\ Y=ar{R} \end{cases}$ Regular operation

Avoid
$$X = Y = 0$$
: $G = 1, S = 1, R = 1$

Avoid changing S or R when G = 1

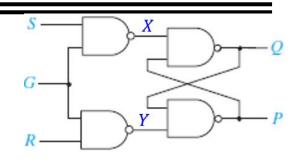
- Let S and R stabilize before changing G from 0 to 1
- Else glitches from static or dynamic hazards in circuits that generate S or R can change latch to undesired state if they arrive after G = 1

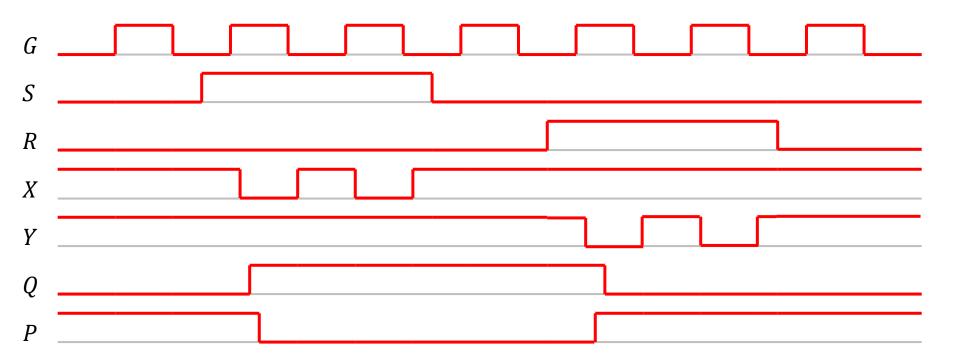


Gated S-R Latch

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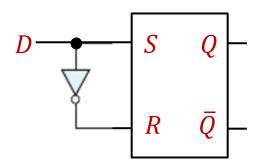
Timing diagram with G serving as a "clock" \rightarrow | \leftarrow Gate delay





D Latch ECE2060

D-Latch: Data Latch



$$D=1 \implies S=1, R=0, \quad Q^+=1$$
 $D=0 \implies S=0, R=1, \quad Q^+=0$ Reset

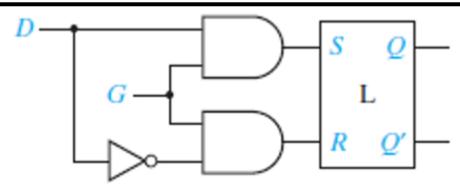
Inverter results in: No Hold State (00); No Bad State (11)

Make it gated But without Hold state it does not really have memory ⇒

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Gated D Latch

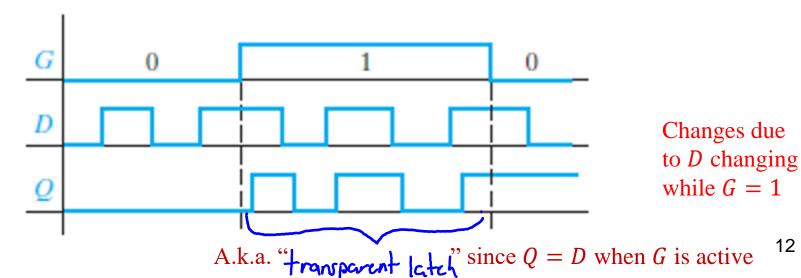


$$G = 1 \implies S = D, R = \overline{D}$$

Set or Reset based on value of *D*

$$G = 0 \implies \begin{cases} S = 0 \\ R = 0 \end{cases}$$
 Hold

Change due to G = 1 while D is already stable



Gated Latches vs Edge-Triggered Flip-Flops

- With Gated Latches (Gated Memory)
 - Wait until input (S and R, O) is stable before setting G=1 (or 0 if actively) to write to Q
 - This is " Level-Sensitive " memory
 - Q updated based on G level of 0 or 1
 - Must avoid changing input(s) while G active
 - Using *G* as a clock for synchronizing writing to memory might be doable for one or a few latches
 - But ensuring all that timing for stable inputs is problematic for large numbers of latches
- Edge-Triggered Flip-Flops
 - Better if we could set up latches that would update at a specific moment in time
 - Use "Edge-triggered" Flip-Flops, designed to change states based on either
 - Rising edge of clock only, or
 - Falling edge of clock only.
 - Rather than whenever clock is at 1 (or at 0)

Edge-Triggered Flip-Flops

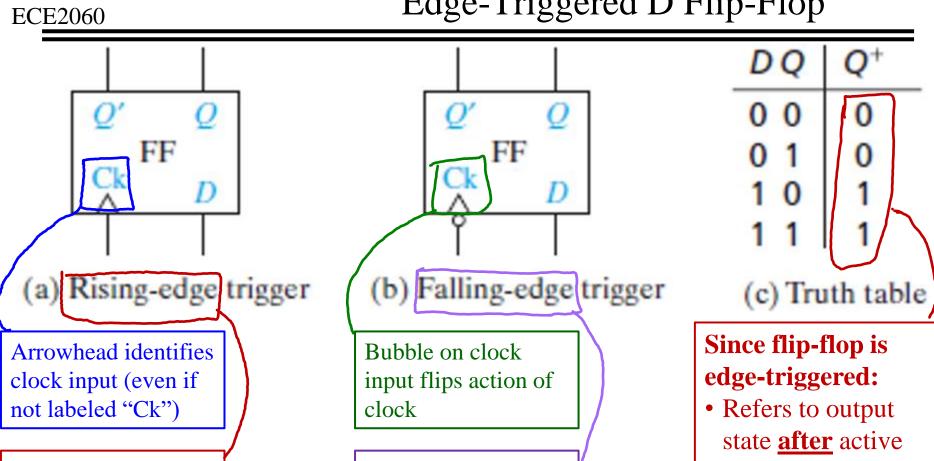
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Edge-Triggered and Leader-Follower Flip-Flops

- Edge-Triggered
 - If inputs to flip-flop only need to be stable for a short period of time around the clock edge,
 - Then we refer to the flip-flop as edge-triggered
- Leader-Follower Edge-Triggered
 - Refers to a particular implementation
 - Uses two gated latches in such a way that the flip-flop outputs only change on a clock edge
- Alert!
 - Not all leader-follower flip-flops are edge-triggered
 - Leader-follower can be used to make an edge-triggered flip-flop
 - But is not sufficient to guarantee it is edge-triggered (details matter)



Edge-Triggered D Flip-Flop



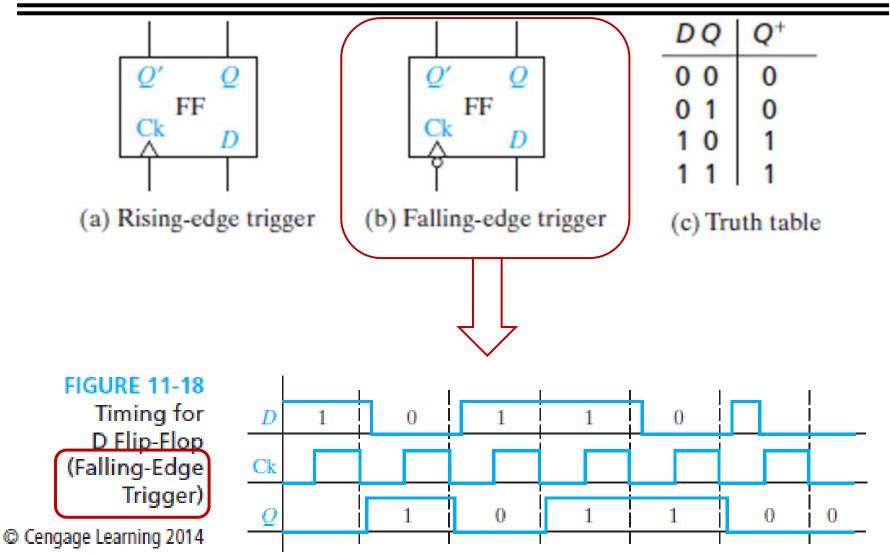
Output can change in response to a $0 \rightarrow 1$ transition on the clock (positive sloped edge

Output can change in response to a $1 \rightarrow 0$ transition on the clock (regative

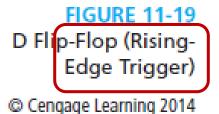
sloped edge

- edge of clock hits
- Output does not change on change of D by itself

Edge-Triggered D Flip-Flop

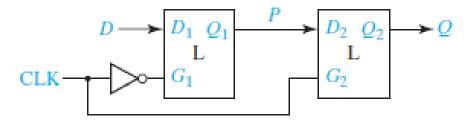


Edge-Triggered D Flip-Flop

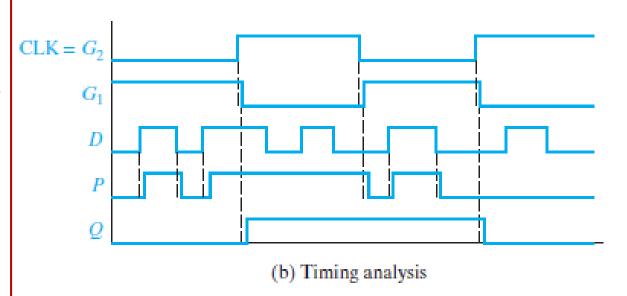


The two *D* latches will not be transparent at the same time

- There is a small interval due to inverter gate delay when it looks like both might be transparent
- But there is also a delay through the latches during that interval
- Careful attention to internal timing needed

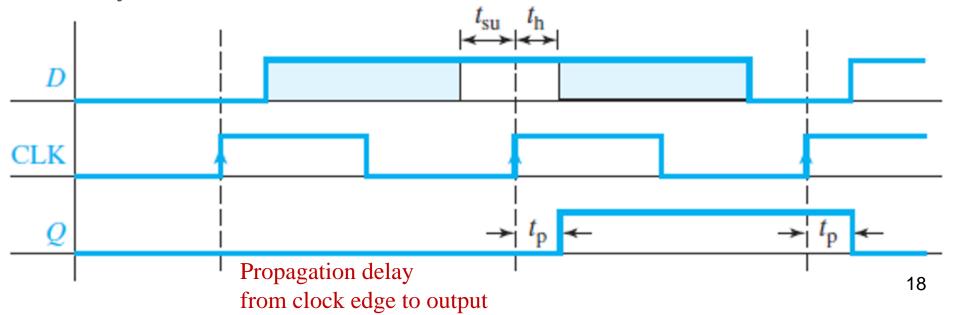


(a) Construction from two gated D latches



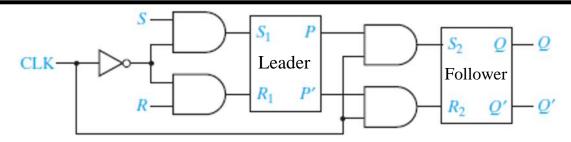
Edge-Triggered D Flip-Flop

- To function properly, *D* input must be stable for a period of time before and after the active edge of clock
 - Set up time $\equiv t_{su}$ is time D must be stable before active edge of clock
 - Allows *D* to propagate through first gated latch
 - Hold time $\equiv t_h$ is time D must be stable after active edge of clock
 - Needed so that current D is stored in first latch, before any change
- Shown here for rising-edge triggered case
- Usually found on data sheets





S-R Flip-Flop



Operation Details

- Leader-Follower configuration
- CLK = 0

S and R set or reset outputs P and P' of Leader gated latch

• *CLK* changes to 1

Leader gated latch holds values of P and P' and

• CLK = 1

Leader gated latch continues to hold values of P and P'

• *CLK* changes to 0

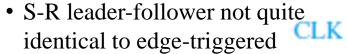
Values of Q and Q' latched into the Follower gated latch

S-R Flip-Flop

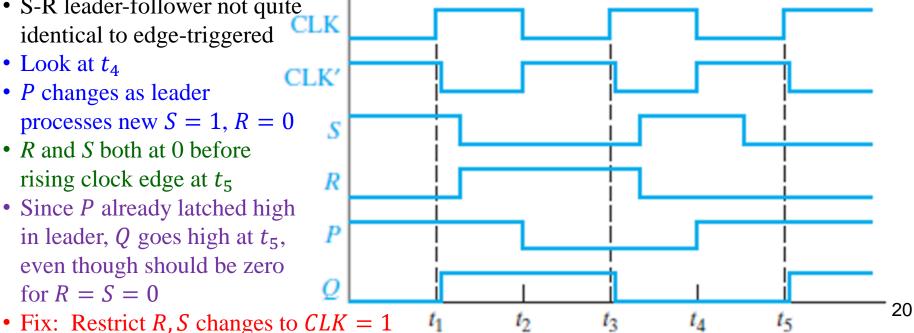
Operation Details

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- CLK = 0 S and $R \rightarrow P$, Follower holds previous |Q|
- $CLK \rightarrow 1$ Leader holds P, transfers it to Follower
- CLK = 1 Leader continues to hold P, Q does not change even if S or R change
- $CLK \rightarrow 0$ Q latched into Follower, Leader processes new inputs



- Look at t_4
- *P* changes as leader processes new S = 1, R = 0
- R and S both at 0 before rising clock edge at t_5
- Since *P* already latched high in leader, Q goes high at t_5 , even though should be zero for R = S = 0



Leader

Follower

J-K Flip-Flop

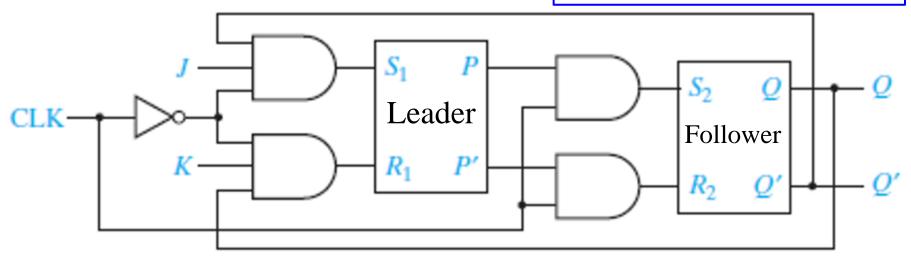
- Uses present values of Q and Q' to make [S=1,R=1] into a "Toggle" condition
- One way to implement with rising-edge *CLK*:

Gate for follower enabled by *CLK*

Note: While I = K = 1 is OK,

feeding present Q and Q' through

ANDs ensures S_1 and R_1 not = 1 at



Gate for leader, enabled by \overline{CLK} , with of Q and Q' also fed back in cross-connected manner

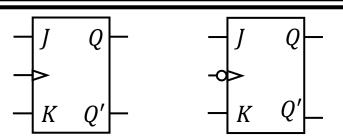
• CLK = 0, J = 1, K = 1 $S_1 = Q'$ and $R_1 = Q$ same time If $Q = 1 \rightarrow S_1 = 0, R_1 = 1 \Rightarrow$ If $Q = 0 \rightarrow S_1 = 1, R_1 = 0 \Rightarrow$

• Toggle mode:

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J-K Flip-Flop



J-K Flip-Flop

General purpose do everything flip-flop

•
$$J = K = 0$$

Hold

•
$$J = 0, K = 1$$

Reset

•
$$J = 1, K = 0$$

Set

•
$$J = K = 1$$

Toggle

Present	After active-	
state	edge of clock	
<u> </u>	<u> </u>	
JKQ	Q	
000	0	
001	1	
010	0	
0 1 1	0	
100	1	
101	1	
110	1	
111	0	

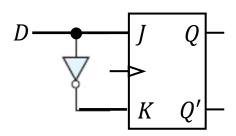
J-K Flip-Flop

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J-K Flip-Flop

General purpose, do everything flip-flop

• For example, D Flip-Flop from J-K



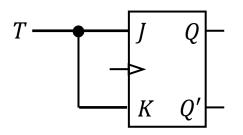
D	Q	Q^+
0	X	0
1	X	1

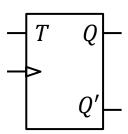
T Flip-Flop

J-K Flip-Flop

General purpose, do everything flip-flop

• For example, T Flip-Flop from J-K





T	Q	Q^+
0	X	Q
1	X	$ar{Q}$