

## CSE 2321 Homework 7

### Problem 1

$$\begin{aligned} T(n) &= \sum_{k=0}^{n-1} ar^k && \text{Geometric series is really useful!} \\ &= a \left( \frac{1-r^n}{1-r} \right) && \text{If } r \neq 1 \end{aligned}$$

#### 1A

Recurrence relation:

$$T(n) = 2T\left(\frac{n}{2}\right) + 1$$

Running time:

$$T(n) = \Theta(n)$$

Generic formula after k substitutions:

$$T(n) = 2^{k+1}T\left(\frac{n}{2^{k+1}}\right) + \sum_{i=0}^k 2^i \quad (1)$$

$$\text{Stop when } \frac{n}{2^{k+1}} = 1$$

$$k = \log_2(n) - 1$$

$$T(n) = 2^{(\log_2(n)-1+1)}T\left(\frac{n}{2^{(\log_2(n)-1+1)}}\right) + \sum_{i=0}^{(\log_2(n)-1)} 2^i \quad (2)$$

$$n \cdot 1 + n - 1$$

$$T(n) = \Theta(n)$$

## 1B

Recurrence relation:

$$T(n) = 3T\left(\frac{n}{3}\right) + 1$$

Running time:

$$T(n) = \Theta(n)$$

Generic formula after k substitutions:

$$T(n) = 3^{k+1}T\left(\frac{n}{3^{k+1}}\right) + \sum_{i=0}^k 3^i \quad (3)$$

$$\text{Stop when } \frac{n}{3^{k+1}} = 1$$

$$k = \log_3(n) - 1$$

$$T(n) = 3^{(\log_3(n)-1+1)}T\left(\frac{n}{3^{(\log_3(n)-1+1)}}\right) + \sum_{i=0}^{(\log_3(n)-1)} 3^i \quad (4)$$

$$n \cdot 1 + \frac{n-1}{2}$$

$$T(n) = \Theta(n)$$

## 1C

Recurrence relation:

$$T(n) = T(n-1) + 1$$

Running time:

$$T(n) = \Theta(n)$$

Generic formula after k substitutions:

$$T(n) = T(n - (k+1)) + (k+1) \quad (5)$$

$$\text{Stop when } n - (k+1) = 1$$

$$k = n - 2$$

$$T(n - ((n-2) + 1)) + ((n-2) + 1) \quad (6)$$

$$1 + n - 1 = n$$

$$T(n) = \Theta(n)$$

## 1D

Recurrence relation:

$$T(n) = T(n - 2) + 1$$

Running time:

$$T(n) = \Theta(n)$$

Generic formula after k substitutions:

$$T(n) = T(n - 2(k + 1)) + (k + 1) \quad (7)$$

$$\text{Stop when } n - 2(k + 1) = 1$$

$$k = \frac{n - 3}{2}$$

$$T(n) = T\left(n - 2\left(\frac{n - 3}{2} + 1\right)\right) + \left(\frac{n - 3}{2} + 1\right)$$

$$1 + \frac{n - 3}{2} + 1$$

$$T(n) = \Theta(n)$$

## 1E

Recurrence relation:

$$T(n) = 3T\left(\frac{n}{2}\right) + 1$$

Running time:

$$T(n) = \Theta\left(n^{\frac{\log_3(n)}{\log_2(n)}}\right)$$

Generic formula after k substitutions:

$$T(n) = 3^{k+1}T\left(\frac{n}{2^{k+1}}\right) + \sum_{i=0}^k 3^i \quad (8)$$

$$\text{Stop when } \frac{n}{2^{k+1}} = 1$$

$$k = \log_2(n) - 1$$

$$T(n) = 3^{(\log_2(n)-1+1)}T\left(\frac{n}{2^{(\log_2(n)-1+1)}}\right) + \sum_{i=0}^{(\log_2(n)-1)} 3^i \quad (9)$$

$$3^{\log_2(n)} \cdot 1 + \frac{1 - 3^{\log_2(n)}}{-2}$$

$$T(n) = \Theta\left(n^{\frac{\log_3(n)}{\log_2(n)}}\right)$$

## Problem 3

```

int BinarySearch(A, i, j, k)
    if i > j
        index = -1
    else
        midpt1 = (i+j)/3
        midpt2 = midpt1
        if k = A[midpt1]
            index = midpt1
        else if k < A[midpt1]
            index = BinarySearch(A, i, midpt1 - 1, k)
        else if k = A[midpt2]
            index = midpt2
        else if k > A[midpt2]
            index = BinarySearch(A, midpt2 + 1, j, k)
        else
            index = BinarySearch(A, midpt1 + 1, midpt2 - 1, k)
    return index

```

TrinarySearch:

Recurrence relation:

$$T(n) = T\left(\frac{n}{3}\right) + C$$

Running time:

$$T(n) = \Theta(\log(n))$$

Generic formula after k substitutions:

$$T(n) = T\left(\frac{n}{3^{k+1}}\right) + (k + 1) \quad (10)$$

Stop when  $\frac{n}{3^{k+1}} = 1$

$$k = \log_3(n) - 1$$

$$T(n) = T\left(\frac{n}{3^{(\log_3(n)-1+1)}}\right) + (\log_3(n) - 1 + 1) \quad (11)$$

$$1 + \log_3(n)$$

$$T(n) = \Theta(\log(n))$$

BinarySearch (from class notes):

Recurrence relation:

$$T(n) = T\left(\frac{n}{2}\right) + C$$

Running time:

$$T(n) = \Theta(\log(n))$$

TrinarySearch has same running time as BinarySearch.