CSE 2321 Homework 9 Template

Problem 1

```
DFSModified (G)
    boolean hamiltonian = false
    node s = u \# s is the start node
    for each u in V
        u.color = white
        u.p = nil # p is for parent
    for each u in V
        if u.color == white
            hamiltonian = Visit (G, u, s)
        if u.color == black
            u.color == white # revisit past nodes
    return hamiltonian
Visit (G, u, s)
    boolean hamiltonian = false
    for each v in G. Adj[u]
        u.color = gray
        if v.color = white
        v.p = u
        Visit (G, v)
        u.color = black
        if v == s \# If the next node is the starting node
            hamiltonian = true
```

Problem 2

```
\begin{aligned} & \operatorname{FindSink}\left(G\right) \\ & \operatorname{int} \ i \ , \ j = 0 \\ & \operatorname{boolean} \ \operatorname{hasSink} = \operatorname{false} \\ & \operatorname{for} \ i \ \operatorname{to} \ |V| \ \&\& \ ! \operatorname{hasSink} \\ & \operatorname{if} \ ! ( \ i == \ j ) \ \&\& \ G[ \ i \ ][ \ j \ ] == 0 \end{aligned}
```

```
# j is not a sink, next column
            j++;
            hasSink = isSink(i)
        else if !(i = j) \&\& G[i][j] = 1
         # i is not a sink, next row and column
            i++
            j++
            hasSink = isSink(j)
        else if i > |V| \mid |j > |V|
            hasSink = false
    return hasSink
isSink(G, x)
    boolean isSink = true
    for int i = 0 to |V|
        # The xth row should be all 0
        if G[k][i] != 0
            isSink = false
        i++
    for int i = 0 to |V|
        # The xth column should be all 1s
        if i != k && G[i][k] != 1
            isSink = false
        else if !(i = j) \&\& G[i][j] = 1
            isSink == false
        i++
    return isSink
```

Problem 3

```
int length = BFSModified(G, i)
            # Note: if two nodes are not connected, infinity
            if length != infinity && length > max
                max = length
            j++
        i++
    return max
BFSModified (G, s, e)
    for each v in V
        v.dist = infinity
        s.dist = 0
        q.enqueue(s)
        while q != 0
            v. dequeue
                for each w s.t. (v,w) in E
                     if w.dist = infinity
                         q.enqueue(w)
    return w.dist # Modified to return length of path
```

The running time of this algorithm is $O(|V|^2 + |V||E|)$. BFSModified has the same running time as BFS, which is O(|V| + |E|). The for loops have the running time of $O(|V|^2)$.