

CSE 2321   Foundations I   Spring, 2024   Dr. Estill  
Homework 7   Due: Friday, March 22

**Theorem 1** (Master Theorem).

Let  $a \geq 1$  and  $b > 1$  be constants, let  $f(n)$  be a function  $f : \mathbb{N} \rightarrow \mathbb{R}^+$ , and let  $T(n)$  be defined on the nonnegative integers by the recurrence

$$T(n) = aT(n/b) + f(n).$$

Then  $T(n)$  has the following asymptotic bounds:

1. if  $f(n) \in O(n^{\log_b a - \varepsilon})$  for some constant  $\varepsilon > 0$ , then  $T(n) \in \Theta(n^{\log_b a})$ ,
2. if  $f(n) \in \Theta(n^{\log_b a})$  then  $T(n) \in \Theta(n^{\log_b a} \log n)$ , and
3. if  $f(n) \in \Omega(n^{\log_b a + \varepsilon})$  for some  $\varepsilon > 0$  and if  $af(n/b) \leq df(n)$  for some constant  $d < 1$  and all sufficiently large  $n$ , then  $T(n) \in \Theta(f(n))$ .

Use the Master Theorem above to solve the following recurrences when possible. If you need to confirm the regularity condition ( $af(n/b) \leq df(n)$  for  $d < 1$ ), work should be shown, but otherwise answers are all that is needed. (Note that not every blank needs to be filled in in every problem.):

(20 points each)

1.)  $T(n) = T(n/3) + c$

$f(n) = \underline{\hspace{2cm}}$  versus  $n^{\log_b a} = \underline{\hspace{2cm}}$

Which is growing faster:  $f(n)$  or  $n^{\log_b a}$ ?  $\underline{\hspace{2cm}}$

Which case of the Master Theorem does that potentially put us in?  $\underline{\hspace{2cm}}$

If you're potentially in case one or three, is it possible to find an epsilon which makes either  $f(n) \in O(n^{\log_b a - \varepsilon})$  (if you're in case one) or  $f(n) \in \Omega(n^{\log_b a + \varepsilon})$  (if you're in case three) true? Choose one or show an inequality.  $\underline{\hspace{2cm}}$

If you're potentially in case three and there is an  $\varepsilon$ , try to find a constant  $d < 1$  such that  $af(n/b) \leq df(n)$  for large enough  $n$ 's.  $\underline{\hspace{2cm}}$

What can you conclude?  $\underline{\hspace{2cm}}$

2.)  $T(n) = T(n/3) + c \log_2 n$

$f(n) = \underline{\hspace{2cm}}$  versus  $n^{\log_b a} = \underline{\hspace{2cm}}$

Which is growing faster:  $f(n)$  or  $n^{\log_b a}$ ?  $\underline{\hspace{2cm}}$

Which case of the Master Theorem does that potentially put us in?  $\underline{\hspace{2cm}}$

If you're potentially in case one or three, is it possible to find an epsilon which makes either  $f(n) \in O(n^{\log_b a - \varepsilon})$  (if you're in case one) or  $f(n) \in \Omega(n^{\log_b a + \varepsilon})$  (if you're in case three) true? Choose one or show an inequality.  $\underline{\hspace{2cm}}$

If you're potentially in case three and there is an  $\varepsilon$ , try to find a constant  $d < 1$  such

that  $af(n/b) \leq df(n)$  for large enough  $n$ 's. \_\_\_\_\_

What can you conclude? \_\_\_\_\_

3.)  $T(n) = 4T(n/2) + cn$

$f(n) =$  \_\_\_\_\_ versus  $n^{\log_b a} =$  \_\_\_\_\_

Which is growing faster:  $f(n)$  or  $n^{\log_b a}$ ? \_\_\_\_\_

Which case of the Master Theorem does that potentially put us in? \_\_\_\_\_

If you're potentially in case one or three, is it possible to find an epsilon which makes either  $f(n) \in O(n^{\log_b a - \varepsilon})$  (if you're in case one) or  $f(n) \in \Omega(n^{\log_b a + \varepsilon})$  (if you're in case three) true? Choose one or show an inequality. \_\_\_\_\_

If you're potentially in case three and there is an  $\varepsilon$ , try to find a constant  $d < 1$  such that  $af(n/b) \leq df(n)$  for large enough  $n$ 's. \_\_\_\_\_

What can you conclude? \_\_\_\_\_

4.)  $T(n) = 4T(n/2) + cn^3$

$f(n) =$  \_\_\_\_\_ versus  $n^{\log_b a} =$  \_\_\_\_\_

Which is growing faster:  $f(n)$  or  $n^{\log_b a}$ ? \_\_\_\_\_

Which case of the Master Theorem does that potentially put us in? \_\_\_\_\_

If you're potentially in case one or three, is it possible to find an epsilon which makes either  $f(n) \in O(n^{\log_b a - \varepsilon})$  (if you're in case one) or  $f(n) \in \Omega(n^{\log_b a + \varepsilon})$  (if you're in case three) true? Choose one or show an inequality. \_\_\_\_\_

If you're potentially in case three and there is an  $\varepsilon$ , try to find a constant  $d < 1$  such that  $af(n/b) \leq df(n)$  for large enough  $n$ 's. \_\_\_\_\_

What can you conclude? \_\_\_\_\_

5.)  $T(n) = 2T(n/6) + \sqrt{n}$

$f(n) =$  \_\_\_\_\_ versus  $n^{\log_b a} =$  \_\_\_\_\_

Which is growing faster:  $f(n)$  or  $n^{\log_b a}$ ? \_\_\_\_\_

Which case of the Master Theorem does that potentially put us in? \_\_\_\_\_

If you're potentially in case one or three, is it possible to find an epsilon which makes either  $f(n) \in O(n^{\log_b a - \varepsilon})$  (if you're in case one) or  $f(n) \in \Omega(n^{\log_b a + \varepsilon})$  (if you're in case three) true? Choose one or show an inequality. \_\_\_\_\_

If you're potentially in case three and there is an  $\varepsilon$ , try to find a constant  $d < 1$  such that  $af(n/b) \leq df(n)$  for large enough  $n$ 's. \_\_\_\_\_

What can you conclude? \_\_\_\_\_