SP25 CSE2431 HOMEWORK 1 SOLUTION

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Electronic Submission: 11:59 pm, Monday, Feb 03, 2025

Point: 25 points (5% of total grade)

Instructions:

This homework will be submitted **online**. Student will do their homework on paper/word processor and scan/take photos/convert to **pdf file** and submit on Carmen on or before the due date. **Only pdf** will be accepted.

Question 1 (3 points): What is the output of the following program? List ALL the possible output(s).

```
int num = 99;

int main(int argc, char**argv){
    printf("start num is %d\n", num);
    num++;
    int rc = fork();
    if(rc>0){
        num+=10;
        printf("Parent's num is %d\n", num);
    }
    else if(rc==0){
        num+=30;
        printf("Child's num is %d\n", num);
    }
    num--;
    printf("end num is %d\n", num);
}
```

ANSWER:

1st possible output: start num is 99 Parent's num is 110 end num is 109 Child's num is 130 end num is 129 2nd possible output: start num is 99 Child's num is 130 end num is 129 Parent's num is 110 end num is 109

Question 2 (10 points): Which of the following scheduling algorithms could result in starvation? For those algorithms that might result in starvation, describe a situation in which starvation is likely to occur? (2 points)

- 1. First-come, first-served (FCFS)
- 2. Shortest Job First (SJF)
- 3. Round Robin
- 4. Basic MLFQ with first 4 rules.

Given the following mix of job, job lengths, and arrival times, assume a time slice of 10 seconds and compute the <u>completion time</u> for each job, <u>average turnaround time</u>, and <u>average response time</u> for the FIFO, RR, and STCF algorithms. (For FIFO, the order of job is the order of scheduling) (8 points)

| Job | Length (s) | Arrival Time |
|-----|------------|--------------|
| J1 | 85 | 0 |
| J2 | 30 | 10 |
| J3 | 35 | 10 |
| J4 | 20 | 80 |
| J5 | 50 | 85 |

A note: There might be different scheduling order, you only need to give one solution scheduling for each of the algorithms.

ANSWER:

- +++ Among the listed scheduling algorithms, starvation could happen for:
 - 1. Shortest Job First (SJF): Although Shortest Job First is a non-preemptive scheduling, there is still a possibility that starvation can happen. Even though a long process may be ready to get scheduled, if short processes are continually added, this long process may never been scheduled. A simple example: although its arrival time is at 0, J2 will not have a chance to run, since shorter-time job such as J3, J4, J5, ... keep coming

and flushing the processor. Another example is in slide 20 of the lecture note.

| Job | Length (s) | Arrival time |
|-----|------------|--------------|
| J1 | 10 | 0 |
| J2 | 100 | 0 |
| J3 | 10 | 10 |
| J4 | 10 | 10 |
| J5 | 10 | 20 |
| ••• | 10 | 30 |
| | | |

2. <u>Basic MLFQ with first 4 rules:</u> In multi-level feedback queue with first 4 rules, if we do not have a priority reset after some period of time, short jobs keep coming and long jobs (with low priority) has to wait and could not be executed.

+++ Here I calculate the Response time = $T_{first-run} - T_{arrival}$ (The time that the process arrived till the time that the process got first scheduled)

+++ One schedule result in:

| | | | Scheduling Algorithms | | | | | |
|-----------------|-------------------|---------|-----------------------|----------|------------|-----------|------------|----------|
| Langeth Austral | | FIFO | | RR | | SRTF/STCF | | |
| Job | Length | Arrival | Completion | Response | Completion | Response | Completion | Response |
| | (s) | time | time | time | Time | Time | Time | Time |
| J1 | 85 | 0 | 85 | 0 | 220 | 0 | 220 | 0 |
| J2 | 30 | 10 | 115 | (85-10) | 80 | (10-10) | 40 | (10-10) |
| J3 | 35 | 10 | 150 | (115-10) | 125 | (20-10) | 75 | (40-10) |
| J4 | 20 | 80 | 170 | (150-80) | 145 | (100-80) | 100 | (80-80) |
| J5 | 50 | 85 | 220 | (170-85) | 215 | (110-85) | 150 | (100-85) |
| Aver | age turn- time | -around | 111s | | 120s | | 80s | |
| Av | erage res time | ponse | | 67s | | 11s | | 9s |

$$Avg T_{turn-around}^{FIFO} = \frac{(85-0) + (115-10) + (150-10) + (170-80) + (220-85)}{5}$$
= 111s

$$Avg T_{response}^{FIFO} = \frac{(0-0) + (85-10) + (115-10) + (150-80) + (170-85)}{5} = 67s$$

$$Avg T_{turn-around}^{RR} = \frac{(220-0) + (80-10) + (125-10) + (145-80) + (215-85)}{5}$$
= 120s

$$Avg T_{response}^{RR} = \frac{(0-0) + (10-10) + (20-10) + (100-80) + (110-85)}{5} = 11s$$

$$Avg T_{turn-around}^{SRTF/STCF} = \frac{(220-0) + (40-10) + (75-10) + (100-80) + (150-85)}{5}$$
= 80s

$$Avg T_{response}^{SRTF/STCF} = \frac{(0-0) + (10-10) + (40-10) + (80-80) + (100-85)}{5} = 9s$$

Schedule for RR (Time slice 10s) (1 possible scheduling)

| Time | Job |
|------|-----|
| 0 | 1 |
| 10 | 2 |
| 20 | 3 |
| 30 | 1 |
| 40 | 2 |
| 50 | 3 |
| 60 | 1 |
| 70 | 2 |

| 80 | 3 → J2 completes |
|-----|---------------------|
| 90 | 1 |
| 100 | 4 |
| 110 | 5 |
| 120 | 3 |
| 125 | 1 → J3 completes |
| 135 | 4 |
| 145 | 5 → J4 completes |
| 155 | 1 |
| 165 | 5 |
| 175 | 1 |
| 185 | 5 |
| 195 | 1 |
| 205 | 5 |
| 215 | 1 → Job 5 completes |
| 220 | J1 completes |

Queue right before 80 (70-80): J3 J1 (J2 being executed then complete at t=80)

Queue at 80: J1 J4 (J4 arrives at 80, J3 is being execute)

Queue at 85: J1 J4 J5

Queue at 90: J4 J5 J3 (J1 is executing)

*** Note: At t=10, there is another possiblity that: J1 finished its timeslice, and then placed ahead of the queue compared to J2 and J3 which arrived at t=10. So queue would be: J1 - J2 - J3 and thus J1 would be executed again!

Schedule for STCF/SRTF

| Time | Job |
|------|------------------|
| 0 | 1 |
| 10 | 2 |
| 20 | 2 |
| 30 | 2 |
| 40 | 3 – J2 completes |
| 50 | 3 |
| 60 | 3 |
| 75 | 1 → J3 completes |
| 80 | 4 |
| 90 | 4 |

| 100 | 5 → J4 completes |
|-----|------------------|
| 110 | 5 |
| 120 | 5 |
| 130 | 5 |
| 140 | 5 |
| 150 | 1 |
| 160 | 1 |
| 170 | 1 |
| 180 | 1 |
| 190 | 1 |
| 200 | 1 |
| 210 | 1 |
| 220 | J1 completes |

Question 3 (8 points): Analyze scheduling algorithms for the following five processes given each process' length and arrival time.

| Processes | Length (s) | Arrival Time |
|-----------|------------|--------------|
| P1 | 16 | 0 |
| P2 | 7 | 2 |
| Р3 | 2 | 4 |
| P4 | 4 | 6 |
| P5 | 22 | 8 |

For each of the following scheduling algorithms draw a diagram about how processes are scheduled, <u>calculate</u> the turnaround time for <u>each process</u> and calculate the <u>average</u> <u>turnaround time</u>.

- 1) First In First Out (FIFO) [2 points]
- 2) Shortest Job First (SJF) (Assume OS first schedules at time 0 and then schedules whenever a process terminates.) [2 points]
- 3) Shortest Time to Completion First (STCF) [2 points]
- 4) Round Robin Scheduling (Assume that the time slice is 5 seconds. Assume that if a new process arrives as the time slice of the executing process expires the executing process is put at the end of the *current* waiting queue. Assume that if a process terminates before its current time slice expires, the next process will immediately get a new and full slice. E.g. if a process starts at 10 and has only 2 remaining jobs, the next process will start at 12 and continue till 12+5. [2 points]

ANSWER:

1) First In First Out (FIFO)

P1 [0 16] P2 [16 23] P3 [23 25] P4 [25 29] P5 [29 51]

Turnaround time for each process:

$$T_{P_1} = 16 - 0 = 16(s)$$

$$T_{P_2} = 23 - 2 = 21 (s)$$

$$T_{P_2} = 25 - 4 = 21 (s)$$

$$T_{P_A} = 29 - 6 = 23 (s)$$

$$T_{P_5} = 51 - 8 = 43 (s)$$

Average turnaround time = [(16-0) + (23-2) + (25-4) + (29-6) + (51-8)]/5 = 24.8

 Shortest Job First (SJF) (Assume OS first schedules at time 0 and then schedules whenever a process terminates.)

Turnaround time for each process:

$$T_{P_1} = 16 - 0 = 16(s)$$

$$T_{P_2} = 29 - 2 = 27 (s)$$

$$T_{P_2} = 18 - 4 = 14 (s)$$

$$T_{P_4} = 22 - 6 = 16$$
 (s)

$$T_{P_5} = 51 - 8 = 43 (s)$$

Average turnaround time = [(16-0) + (18-4) + (22-6) + (29-2) + (51-8)]/5 = 23.2

3) Shortest Time to Completion First (STCF)

Turnaround time for each process:

$$T_{P_1} = 29 - 0 = 29(s)$$

$$T_{P_2} = 15 - 2 = 13$$
 (s)

$$T_{P_3} = 6 - 4 = 2 (s)$$

$$T_{P_4} = 10 - 6 = 4 (s)$$

$$T_{P_5} = 51 - 8 = 43 (s)$$

Average turnaround time = [(29-0) + (15-2) + (6-4) + (10-6) + (51-8)]/5 = 18.2

4) Round Robin Scheduling (Assume that the time quantum is 5 time units. Assume that if a new process arrives as the time slice of the executing process expires the executing process is put at the end of the *current* waiting queue. Assume that if a process terminates before its current time slice expires, the next process will immediately get a new and full slice. E.g. if a process starts at 10 and has only 2 remaining work, the next process will start at 12 and continue till 12+5.)

P1 [0 5] P2 [5 10] P3 [10 12] P1 [12 17] P4 [17 21] P5 [21 26] P2 [26 28] P1 [28 33] P5 [33 38] P1 [38 39] P5 [39 44] P5 [44 49] P5 [49 51]

Turnaround time for each process:

$$T_{P_1} = 39 - 0 = 39(s)$$

 $T_{P_2} = 28 - 2 = 26(s)$
 $T_{P_3} = 12 - 4 = 8(s)$
 $T_{P_4} = 21 - 6 = 15(s)$
 $T_{P_5} = 51 - 8 = 43(s)$

Average turnaround time = [(39-0) + (28-2) + (12-4) + (21-6) + (51-8)]/5 = 26.2

Question 4 (4 points): Fix memory-related bugs in the following programs (one program may have more than one bugs). You can try to run it with valgrind, but remember that valgrind is not a panacea.

```
int *add(int *a, int *b){
    int *ret = malloc(sizeof(int));
    if(a==NULL || b==NULL)
        return NULL;
    *ret = *a + *b;
    return ret;
}

int main(int argc, char**argv){
    int a = 3;
    int b = 4;
    int * ret = add(&a, &b);
    if(ret==NULL)
        printf("Error\n");
    else
        printf("3+4=%d\n", *ret);
}
```

ANSWER:

```
int *add(int *a, int *b){
    int *ret = malloc(sizeof(int));
 if(a==NULL || b==NULL) {
     free(ret);
      //or move the if check before malloc
        return NULL; }
    *ret = *a + *b;
    return ret;
}
int main(int argc, char**argv){
    int a = 3;
    int b = 4;
    int * ret = add(&a, &b);
    if(ret==NULL)
        printf("Error\n");
    else{
        printf("3+4=%d\n", *ret);
        free(ret);
     }
}
```