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The given CFG is,

$$S \rightarrow aTb \mid b$$
$$T \rightarrow Ta \mid \varepsilon$$

The rules for this CFG, are applied as,

$\delta(q, \varepsilon, S) = (q, B)$ where S and $T \rightarrow B$ is production grammer.

$\delta(q, \varepsilon, T) = (q, B)$

These are applied for each S and T

For each terminal a and b,

$\delta(q, a, a) = (q, \varepsilon)$ which have to be poped

$\delta(q, b, b) = (q, \varepsilon)$

Let's explain, $\delta(q, \varepsilon, A) = (q, B)$

So,
Variables = S and T
Terminals = a and b

q ← State, ε ← Input, A ← Variable, q ← State, B ← Production rule of grammer

First the production and pop rules as follows :

(i) $\delta(q_0, \varepsilon, \varepsilon) = (q_0, \varepsilon)$

(ii) $\delta(q_0, \varepsilon, S) = (q_0, aTb)$

(iii) $\delta(q_0, \varepsilon, S) = (q_0, b)$

(iv) $\delta(q_0, \varepsilon, T) = (q_0, Ta)$

(v) $\delta(q_0, \varepsilon, T) = (q_0, \varepsilon)$

(vi) $\delta(q_0, a, a) = (q_0, \varepsilon)$

(vii) $\delta(q_0, b, b) = (q_0, \varepsilon)$

So the input would be

$aTb = aTab = aaab = aab$

$S \rightarrow aTb \mid b$
$T \rightarrow Ta \mid \varepsilon$

Transition Table:

Serial	State	Unread input	Stack	Transition
1	q_0	aab	$	1
2	q_0	aab	S	1
3	q_0	aab	aTb	2
4	q_0	aab (poped)	aTb	3
5	q_0	ab	Tb	
6	q_0	ab	Tab	4
7.	q_0	ab	ab	5
8.	q_0	ab (poped)	ab	6
9	q_0	b	b	
10	q_0	b (poped)	b	7
11	q_0	ε	ε	Accept

So the pdA,