Administrivia

- Student hours (or by appointment):
 - Listed on Canvas in Syllabus
 - -MTWTh 10:30 11:30 a.m. via **Zoom**
- Assignment five posted
 - —Due Friday 14 July by 10pm

CptS 355- Programming Language Design

Object Oriented Programming and Object Oriented Languages

Instructor: Jeremy E. Thompson



5 Insights About 00 Programming

- 1. Data and operations belong together
- 2. Abstraction vs. interface: implementation *details* are hidden
- 3. Subtyping
- 4. Dynamic Dispatch
- 5. Inheritance

- 1. Data and operations belong together
 - By packaging data and operations/methods
 - clearer and easier to write and understand programs
 - legal operations for data are more easily enforced
 - This is closely related to "data abstraction" or "abstract data types"

- 2. Abstraction and interface: implementation details are hidden
 - access to the data should be only via the defined (abstract) operations
- Interface: a set of operations expressed in application-level terms rather than implementation-level terms
 - Includes: operations, their arguments, their results, and their meaning
- In Java, the use of the term "interface" is little bit different
 - a reference type in Java
 - collection of abstract methods—a class implements an interface
 - along with <u>abstract methods</u>, an <u>interface</u> may also contain <u>constants</u>, <u>default methods</u>, <u>static methods</u>, <u>and nested types</u>
 - unless the class that implements the interface is abstract, all methods of the interface must be defined in the class

- 3. Subtyping
 - A is a subtype of type B when:
 - methods (A) \supseteq methods (B), and
 - fields(A) \supseteq fields (B)
 - Notation: A<:B</p>
 - If A<:B then:</p>
 - A can be used whenever B can
- If A<:B, then "A provides all operations B provides, in terms of <u>signatures</u>."
 - However, the operations may not have the same meaning
 - Consequently, languages require programmer to <u>explicitly declare</u> subtype relationships

- 4. Dynamic and Static Dispatch (polymorphism)
 - Dynamic Dispatch:
 - When we have subtyping, the method to call depends on the actual value contained in variable, not its type
 - Binding of the method is determined at <u>run time</u> depending on the type of the object pointed to

Dynamic vs Static Dispatch

C++ Example

```
class Person
{
  public:
    void setSSN(std::string myssn) {
       ssn=myssn;
    }
    void print() {
       std::cout << ssn << std::endl;
    }
  private:
    std::string ssn;
};</pre>
```

```
class Student : public Person
{
   public:
     void setGPA(float mygpa) {
        gpa=mygpa;
     }
     void print() {
        std::cout << gpa << std::endl;
     }
   private:
     double gpa;
};</pre>
```

```
int main() {

1    Student s ;
    s.setSSN("999-99-9999");

3    s.setGPA(3.41);
    Person *p = &s;
    s.print();
    (*p).print();
}
```

The <u>compile-time</u> (static) type of (*p) is Person The <u>run-time</u> (dynamic) type of (*p) is Student

Dynamic Dispatch

```
int main() {
    Student s ;
    s.setSSN("999-99-9999");
    s.setGPA(3.41);
    Person *p = &s;
    s.print();
    (*p).print();
}
```

- In C++, by default, the decision on which member function to invoke (base or overridden) is made on the basis of the <u>compile-time</u> type
 - This is called *static* dispatch
- In C++, the decision is made based on the run-time type when the member function is defined as a <u>virtual</u> function
 - This is called *dynamic* dispatch

Dynamic Dispatch

```
class Person
{
   public:
     string& setSSN(std::string myssn) {
        ssn=myssn;
   }
   virtual void print() {
        std::cout << ssn << std::endl;
    }
   private:
     std::string ssn;
};</pre>
```

C++ Example

```
class Student : public Person
{
  public:
    string& setGPA(float mygpa) {
        gpa=mygpa;
        virtual void print() {
        std::cout << gpa << std::endl;
        }
  private:
        double gpa;
};</pre>
```

```
int main() {
    Student s;
    s.setSSN("999-99-9999");
    s.setGPA(3.41);
    Person *p = &s;
    s.print();
    (*p).print();
}
```

- A *non-virtual* function uses static dispatch
- A *virtual* function uses dynamic dispatch
- here, (*p).print() uses gpa not ssn

Dynamic Dispatch

```
int main() {
    Student s ;
    s.setSSN("999-99-9999");
    s.setGPA(3.41);
    Person *p = &s;
    s.print();
    (*p).print();
}
```

- In C++, by default, the decision on which member function to invoke (base or overridden) is made on the basis of the <u>compile-time</u> type
 - This is called *static* dispatch
- The decision is made based on the run-time type, when the member function is defined as a <u>virtual</u> function
 - This is called *dynamic* dispatch
- How about Java?
- Why is static dispatch the default in C++?

5. Inheritance

 The implementations of methods of a supertype are available (and used) in a subtype <u>unless</u> they are overridden

- Note: Subtyping vs inheritance
 - subtyping is about compatibility of <u>interfaces</u> (in the general sense)
 - inheritance is about re-use of implementations

Questions?



RECALL: Dynamic Dispatch

```
int main() {
    Student s ;
    s.setSSN("999-99-9999");
    s.setGPA(3.41);
    Person *p = &s;
    s.print();
    (*p).print();
}
```

- In C++, by default, the decision on which member function to invoke (base or overridden) is made on the basis of the <u>compile-time</u> type
 - This is called *static* dispatch
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 - This is called *dynamic* dispatch
- How about Java?
- Why is static dispatch the default in C++?

How are *virtual* methods implemented in C++?

- On a per-class basis (not per-instance)
 - run-time data structure called a v-table that contains pointers to the code for virtual methods

How are objects and virtual methods implemented in C++?

• Example:

```
class B
{
  public:
    int f1;
    int f2;
    int f3;
    virtual void m1{ some code }
    virtual void m2{ some code }
    virtual void m3{ some code }
    virtual void m3{ some code }
    void m4{ some code }
};
```

```
class A : public B
{
  public:
    int f4;
    virtual void m2{ some code }
    virtual void m3{ some code }
    virtual void m5{ some code }
    void m4{ some code }
};
```

```
int main() {

1     B p ;

2     p.m1();

3     B *q = new A();

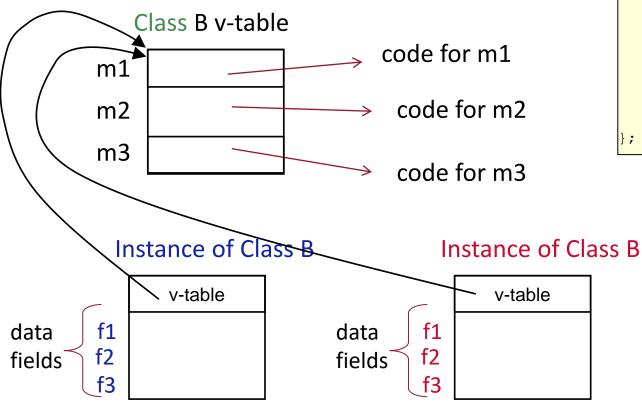
4     q->m1();

5     q->m2();

6     q->m4();
}
```

How are objects and virtual methods implemented in C++?

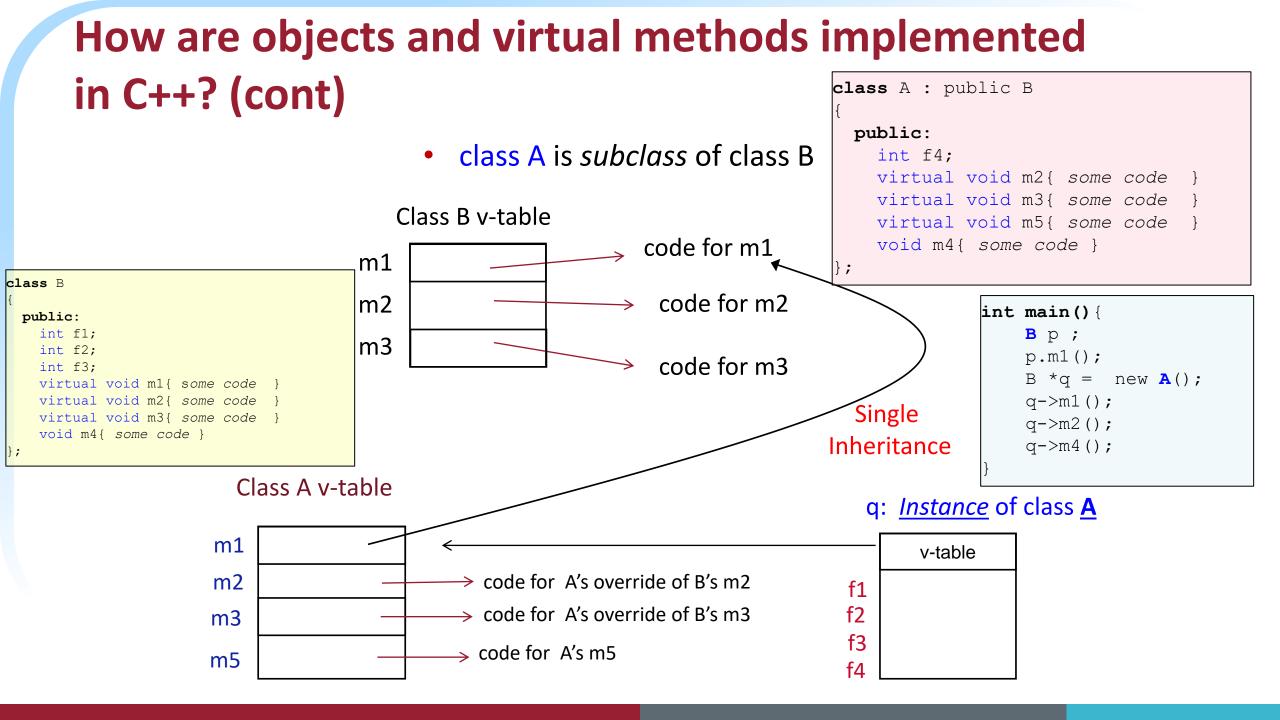
 v-table: run-time data structure which contains pointers to the code for the virtual methods



```
class B
{
  public:
    int f1;
    int f2;
    int f3;
    virtual void m1{ some code }
    virtual void m2{ some code }
    virtual void m3{ some code }
    void m4{ some code }
};
```

```
int main() {
    B x ;
    x.m1();
    B *y = new B();
    y->m2();
}
```

Since m4 is a non-virtual method, what happens?



How are objects and methods implemented in Java?

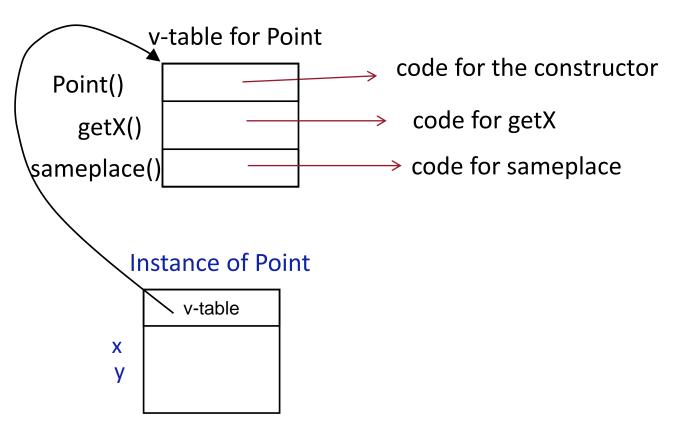
• Example:

```
class Point
 public:
    double x;
    double y;
    Point (double x, double y) {
       this.x = ; this.y=y;
    double getX() {
       return x;
    boolean sameplace (Point p) {
       return (x==p.x) \&\& (y==p.y)
```

```
class PtSubClass extends Point
public:
   int aNewField;
   PtSubClass(double x, double y) {
      super(x,y);
   boolean sameplace (Point p) {
      return false;
   void sayHi () {
      System.out.println("hello!");
```

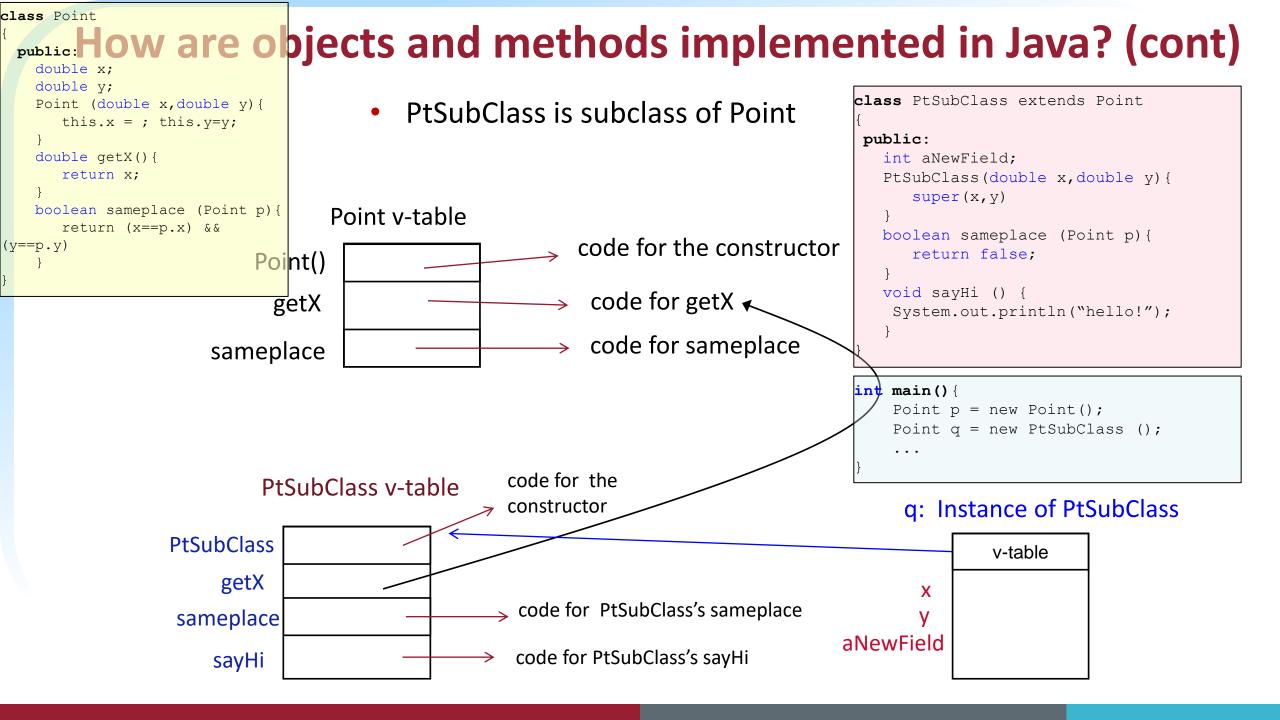
```
int main() {
    Point p = new Point();
    Point q = new PtSubClass ();
    ...
}
```

How are objects and methods implemented in Java?



```
class Point
 public:
    double x;
    double y;
    Point (double x, double y) {
       this.x = ; this.y=y;
    double getX() {
       return x;
    boolean sameplace (Point p) {
       return (x==p.x) \&\& (y==p.y)
```

- V-table is shared across all objects in class!
- If the object instance of Point is no longer needed, what will happen to its v-table?



Questions?

