

Administrivia

- Student hours (or by appointment):
 - Listed on Canvas in Syllabus
 - MTWTh 10:30 – 11:30 a.m. via Zoom
- **Note**: this coming Monday student hour is **canceled**
- **Assignment four posted**
 - Due Tuesday 4 July by **10pm**

CptS 355- Programming Language Design

Types and Type Checking

Instructor: Jeremy E. Thompson

Types

- How would you define a type? What is it and what does it do?
 - Type defines a collection of **values** that share a common property - usually a *common set of operations*
 - Type tells you what is *legal* to do with some value in the language

Types

What is a **type error**?

- An attempt to use a value in an operation *inconsistent* with the value's type
- Examples:
 - The following *will* produce a type error in many languages
x = 17
x ();
 - The following *may* produce a type error in *some* programming languages
3 + 4.5

Compile-time (**static**) type checking *versus* Run-time (**dynamic**) type checking

- Compile-time (**static**) type checking:
 - Examples (Haskell):
 - Checks that *all* return values of a function have the same type
 - Checks that *all* patterns are exhaustive
 - **Compile-time** type checking is necessarily *conservative*:
 - it may flag as an **error** for something that would not actually cause a run-time error
 - **Advantages**:
 - **Less** runtime *overhead*
 - You only do the type-checking *once* (compile-time)
 - The *entire* program is checked

Compile-time (**static**) type checking *versus* Run-time (**dynamic**) type checking (cont.)

- Run-time (**dynamic**) type checking:
 - Run-time type checking is *expensive*
 - must be done **each time** the program is executed
 - Advantages:
 - Allows certain **programming styles** not possible with compile-time type checking
 - More *flexible* data structures
 - For example, lists in Python may contain values of any type whereas lists in Haskell must have elements of the same type

Type Safety: strong typing

- Strong typing ensures that *every* use of a value is compatible with its type
 - Requires *explicit* conversion
 - Static strong typing → compiler error
 - Example : Haskell
 - Dynamic strong typing → error at the point of misuse
 - Example: Python
- In *type-safe* languages, values are managed "*from the cradle to the grave*":
 - Objects are created and initialized in a type-safe way
 - An object cannot be corrupted during its lifetime
 - its *representation* is in accordance with its *type*
 - Objects are *destroyed*, and their memory reclaimed, in a *type-safe* way

Type Safety: strong typing

- C doesn't have type safety
 - C heap values are created in a **type-unsafe** way
 - C *casts* and unchecked array accesses can **corrupt memory** during its lifetime

```
void f(char* char_ptr) {  
    double* d_ptr = (double*)char_ptr;  
    (*d_ptr) = 3.5;  
}
```

← Breakout: What happens **here**?

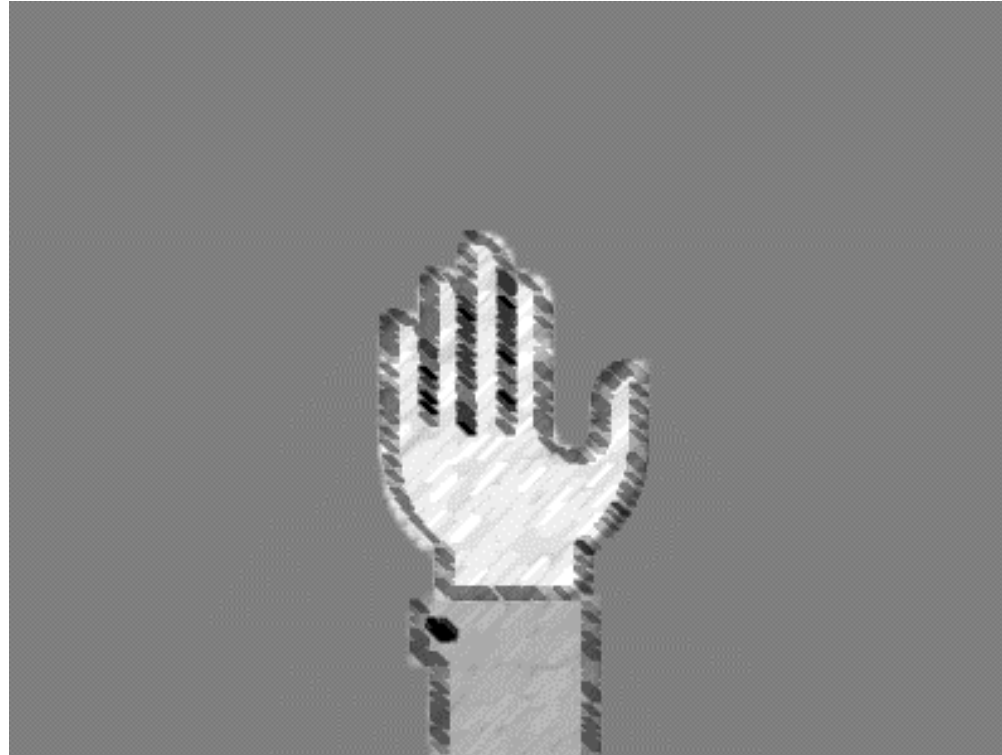
- C deallocation is unsafe, and can lead to **dangling** pointers

```
int *i = (int*)malloc(sizeof(int));  
int j = 0;  
*i = 4;  
free(i);  
/* ... some code that might allocate memory, then: */  
*i = j;
```

← And **here**?

Questions?

Quiz!



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