Reverse Engineering

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Part I

What is reverse engineering?

Computer Underground

- Cracking software copy protection
 - PC software and games
 - modding the Xbox and Playstation
- Exploit development
 - advanced exploitation
- Reversing undocumented operating system APIs
 - virus writers
 - o spyware, keyloggers, malware, rootkits

Security Industry

- Virus and malware analysis
 - AV and Anti-Spyware companies
 - forensics
- Patch analysis, vulnerability analysis
 - IDS, IPS companies
- Binary code auditing
 - discovering new vulnerabilities
- Exploit development
 - penetration testing

Roadmap for today

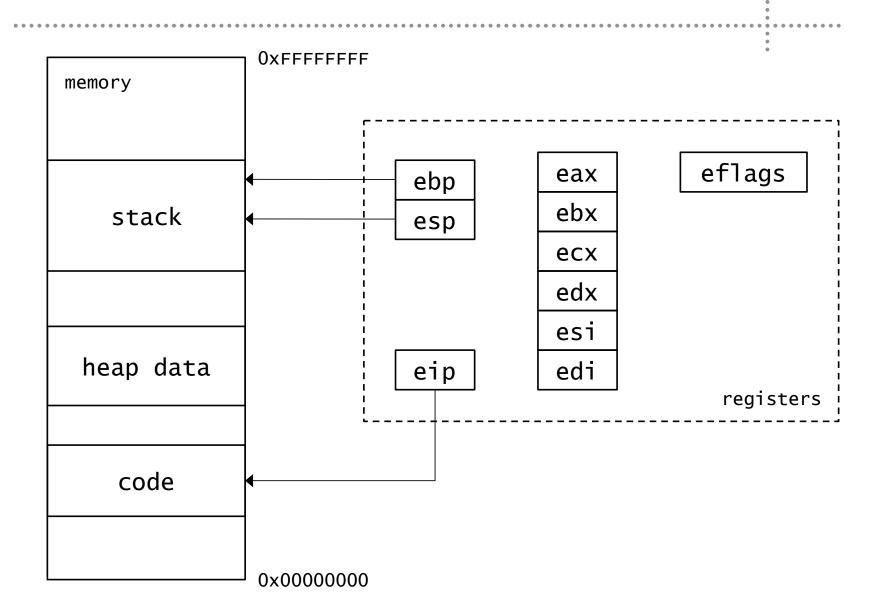
The goal for today: turn a x86 binary executable back into C source code.

- x86 assembly in 10 minutes
- compiler design and code optimizations
- binary reverse engineering

We'll use the IDA Pro disassembler and a sample executable from a reversing contest.

Part II x86 assembly in 10 minutes

CPU architecture



Arithmetic instructions

; ebx = ebx - 2

```
mov eax, 2 ; eax = 2

mov ebx, 3 ; ebx = 3

add eax, ebx ; eax = eax + ebx
```

sub ebx, 2

Accessing memory

```
mov eax, [1234] ; eax = *(int*)1234

mov ebx, 1234 ; ebx = 1234

mov eax, [ebx] ; eax = *ebx

mov [ebx], eax ; *ebx = eax
```

Conditional branches

```
cmp eax, 2; compare eax with 2
je label1
              ; if (eax == 2) goto label1
ja label2
              ; if (eax > 2) goto label2
              ; if (eax < 2) goto label3
jb label3
              ; if (eax <= 2) goto label4
jbe label4
              ; if (eax != 2) goto label5
ine label5
jmp label6; unconditional goto label6
```

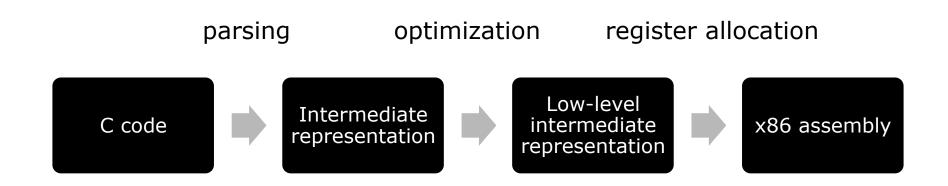
Function calls

```
call func
               ; store return address on
               ; the stack and jump to func
func:
     push esi ; save esi
               ; restore esi
     pop esi
               ; read return address from
     ret
               ; the stack and jump to it
```

Part III

Compiler design and code optimizations

Modern compiler architecture



High-level Optimizations

- Inlining
- Loop unrolling
- Loop-invariant code motion
- Common subexpression elimination
- Constant folding and propagation
- Dead code elimintation

Inlining

```
int foo(int a, int b) {
     return a + b;
c = foo(a, b+1);
c = a + b + 1;
```

Loop unrolling

```
for (i = 0; i < 2; i++) {
     a[i] = 0;
a[0] = 0;
a[1] = 0;
```

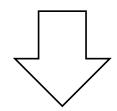
Loop-invariant code motion

```
for (i = 0; i < 2; i++) {
     a[i] = p + q;
temp = p + q;
for (i = 0; i < 2; i++) {
     a[i] = temp;
```

Common subexpression elimination

$$a = b + (z + 1)$$

$$p = q + (z + 1)$$



$$temp = z + 1$$

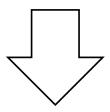
$$a = b + z$$

$$p = q + z$$

Constant folding and propagation

$$a = 3 + 5$$

 $b = a + 1$
func(b)



func(9)

Dead code elimination

```
a = 1
if (a < 0) {
    printf("ERROR!")
}</pre>
```

Low-level optimizations

- Strength reduction
- Code block reordering
- Register allocation
- Instruction scheduling

Strength reduction

$$y = x * 2$$

 $y = x * 15$



$$y = x + x$$

 $y = (x << 4) - x$

Code block reordering

Register allocation

- Memory access is slower than registers
- Try to fit as many local variables as possible in registers
- The mapping of local variables to stack locations and registers is not constant

Instruction scheduling

```
mov eax, [esi]
add eax, 1
mov ebx, [edi]
add ebx, 1
mov eax, [esi]
mov ebx, [edi]
add eax, 1
add ebx, 1
```

Part IV

Binary reverse engineering

Overview

We need to understand:

- How the compiler turns C code into assembly code
- Low-level OS structures and executable file format

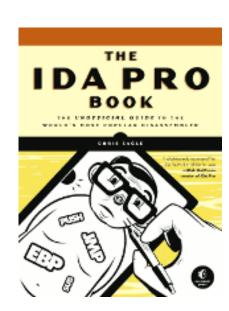
We need to think like the compiler, but in reverse!

Tools

IDA Pro disassembler from hex-rays.com

Limited demo and freeware versions are available.

Almost no documentation!



Diving into IDA Pro