

“Hollywood Style Decryption” on Block Cipher-CBC

Rizki Wicaksono / ilmuHacking.com

Rizki Wicaksono

- Penetration tester
- Programming, application security, cryptography
- S1 Teknik Informatika ITB, ECSP, OSWP, ITIL-F
- ilmuHacking.com , facebook.com/ilmuHacking

Hollywood Style Password Cracking

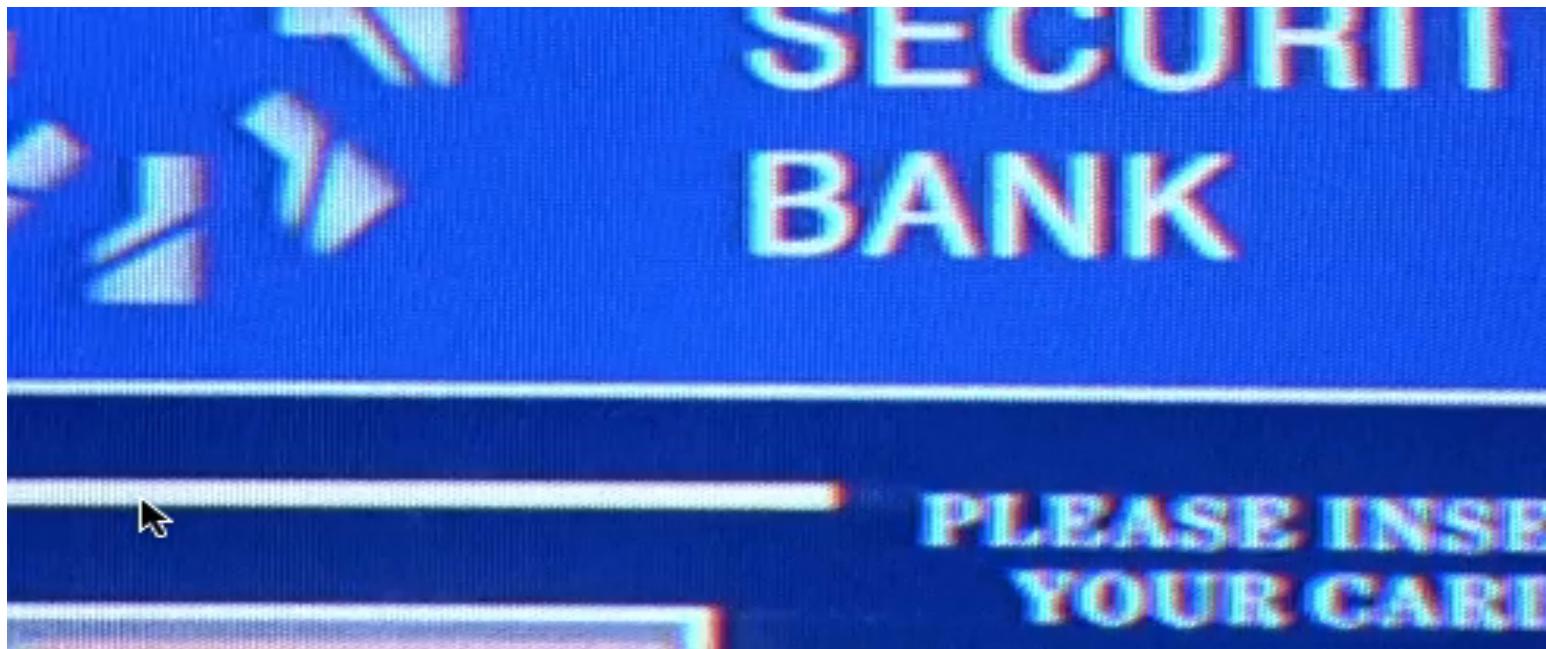


Lets Watch Some Movies

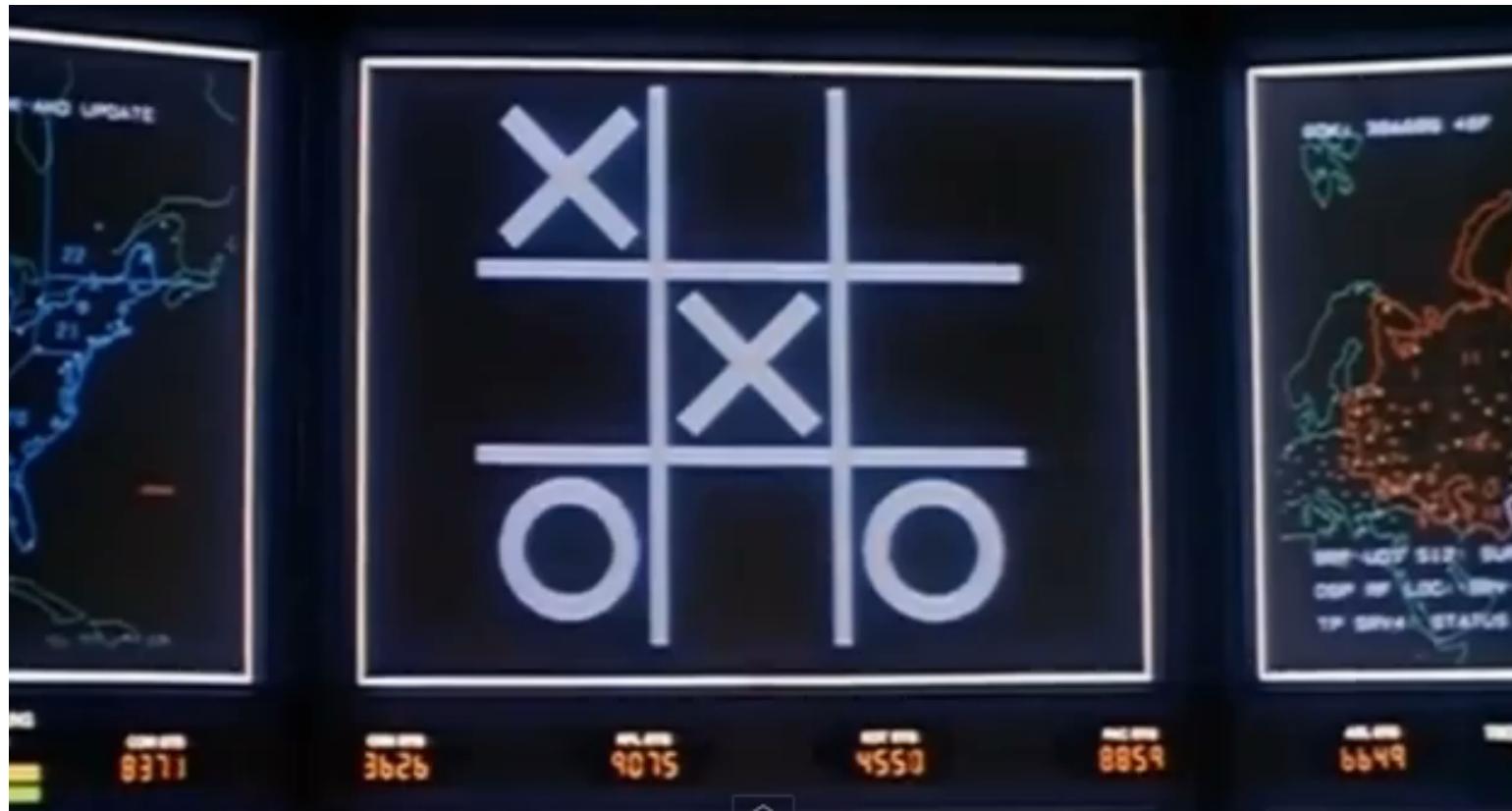
Resident Evil Breaking Door Key Scene



Terminator 2 ATM PIN Cracking Scene



Wargame Launch Code Hacking Scene



The Matrix Beginning Scene

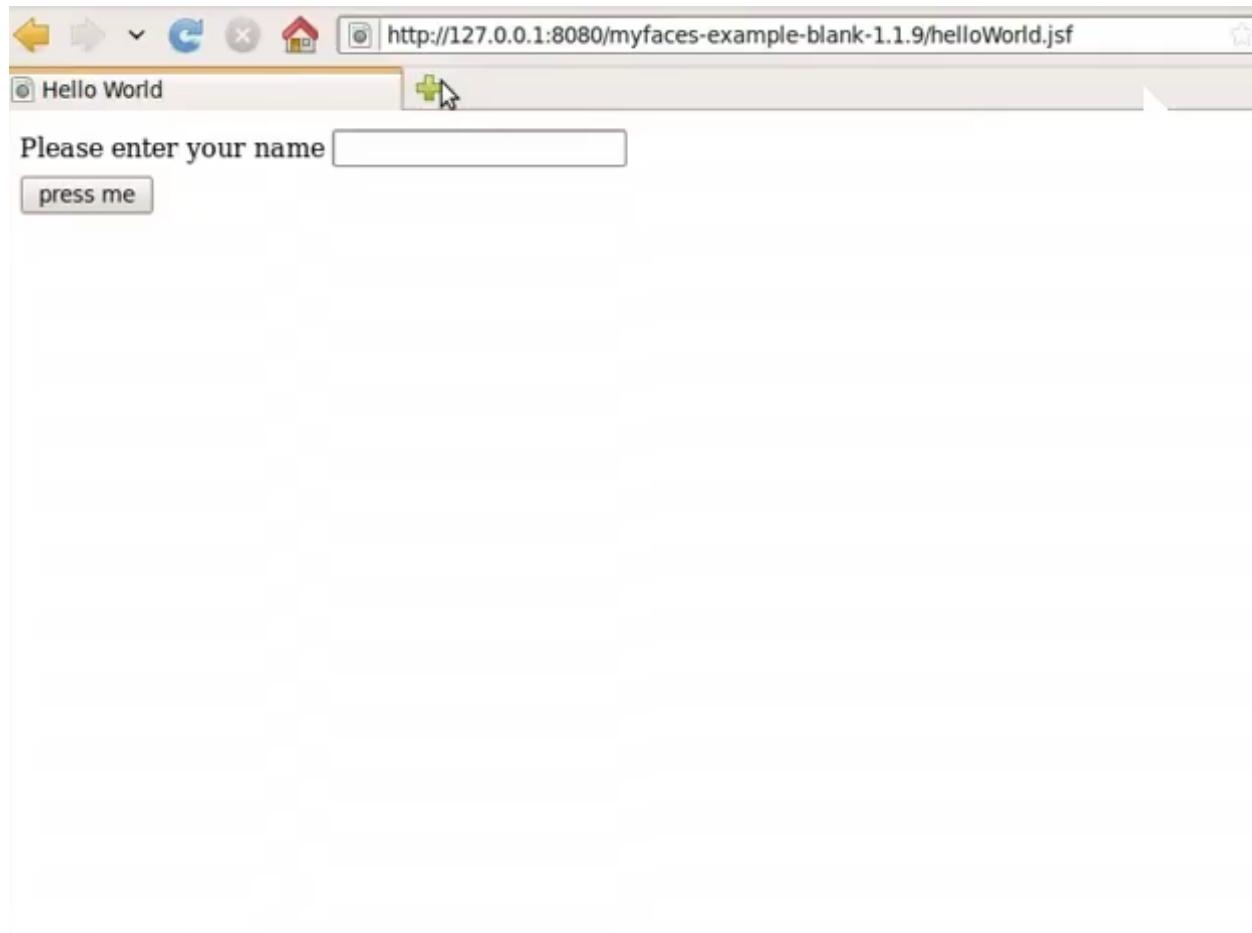


Bloodfist IV Passcode Breaking Scene

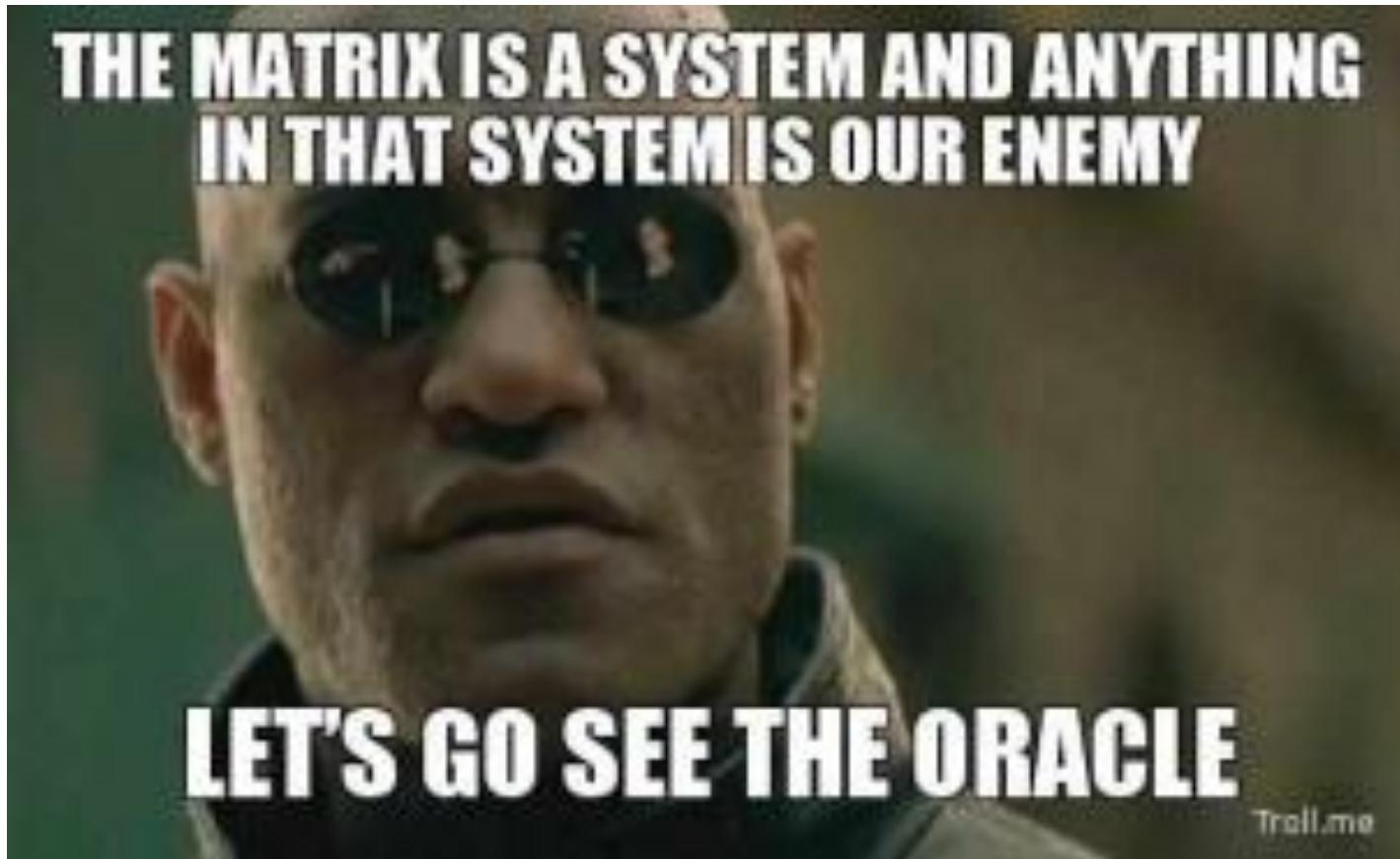


Hollywood Style Decryption with Padding Oracle Attack

Sample Real Attack



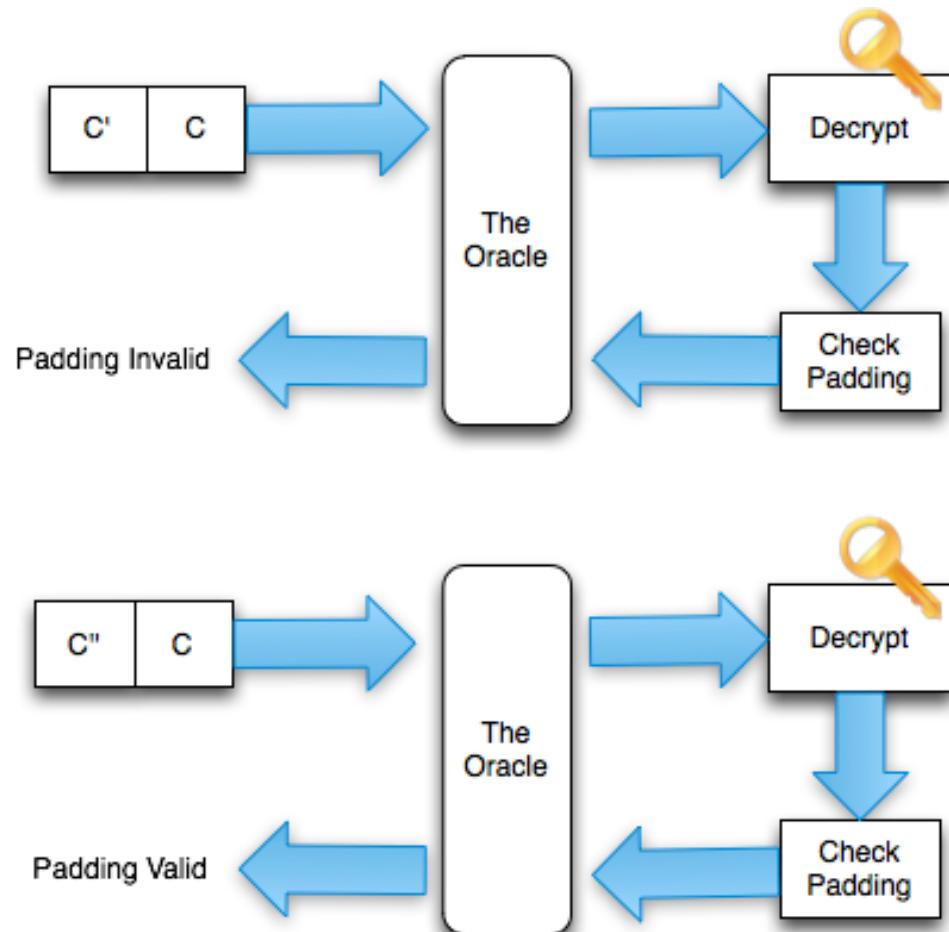
Morpheus: Let's Go See the Oracle



The Oracle



Padding Oracle: Valid/Invalid Pad



1 Bit Information Leakage



PKCS#7 Valid Padding

'A'	'B'	'C'
41	42	43

'A'	'B'	'C'	'D'				
41	42	43	44	04	04	04	04

'A'	'B'	'C'	'D'	'E'
41	42	43	44	45

'A'	'B'	'C'	'D'	'E'	'F'
41	42	43	44	45	46

'A'	'B'	'C'	'D'	'E'	'F'	'G'
41	42	43	44	45	46	47

'A'	'B'	'C'	'D'	'E'	'F'	'G'	'H'
41	42	43	44	45	46	47	48

08 08 08 08 08 08 08 08

PKCS#7 Invalid Padding

01	01	01	01	01	01	01	3A
----	----	----	----	----	----	----	----

3A di luar range valid padding

41	42	43	02	FF	5C	1F	02
----	----	----	----	----	----	----	----

2 byte terakhir bukan 02

41	42	43	03	03	03	03	00
----	----	----	----	----	----	----	----

00 di luar range valid padding

41	42	43	02	08	08	08	08
----	----	----	----	----	----	----	----

8 byte terakhir bukan 08

Cipher Block Chaining

$$C_1 = E_k(P_1 \oplus IV)$$

$$C_2 = E_k(P_2 \oplus C_1)$$

$$C_3 = E_k(P_3 \oplus C_2)$$

$$C_4 = E_k(P_4 \oplus C_3)$$



$$P_1 = D_k(C_1) \oplus IV$$

$$P_2 = D_k(C_2) \oplus C_1$$

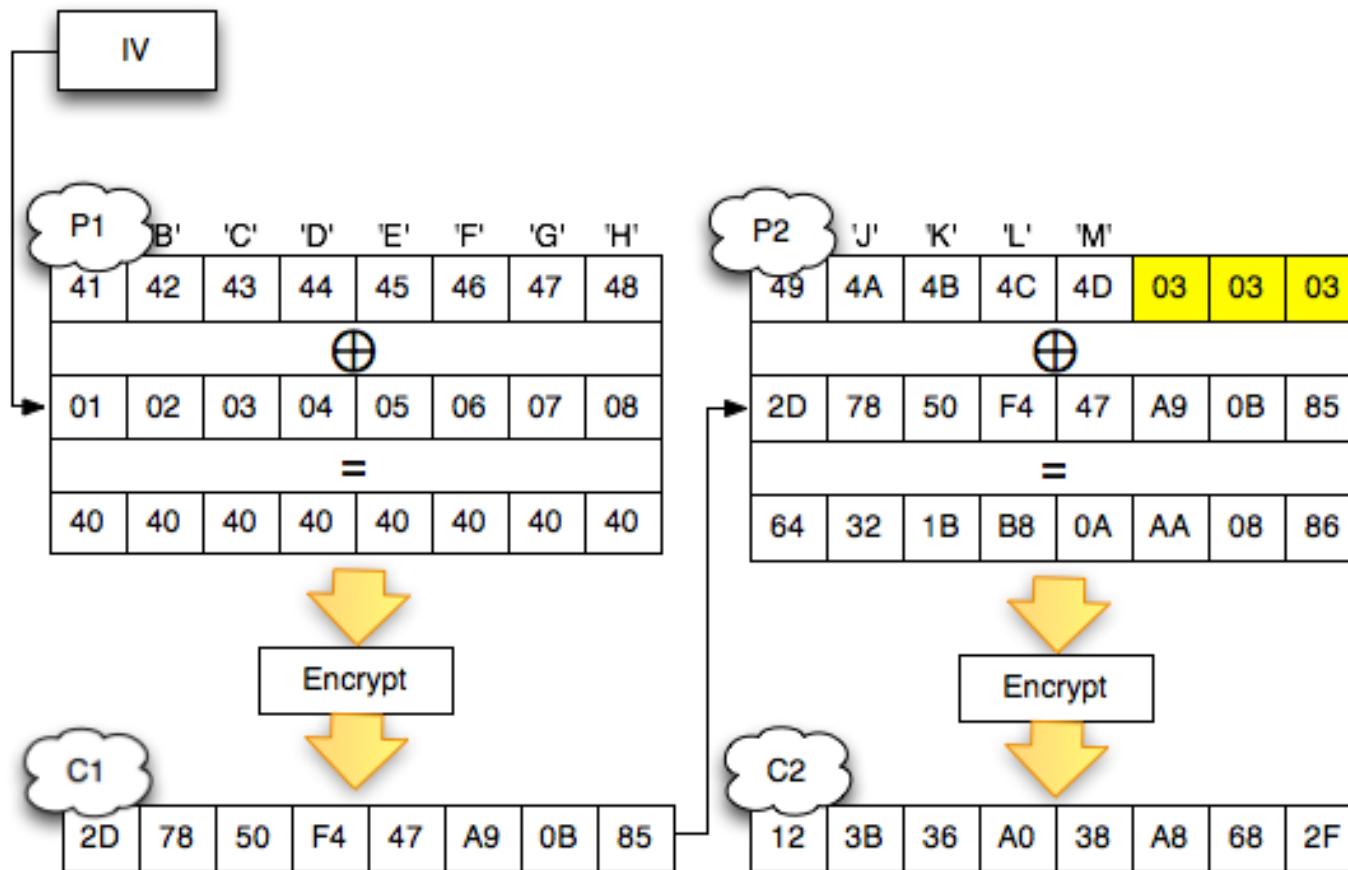
$$P_3 = D_k(C_3) \oplus C_2$$

$$P_4 = D_k(C_4) \oplus C_3$$

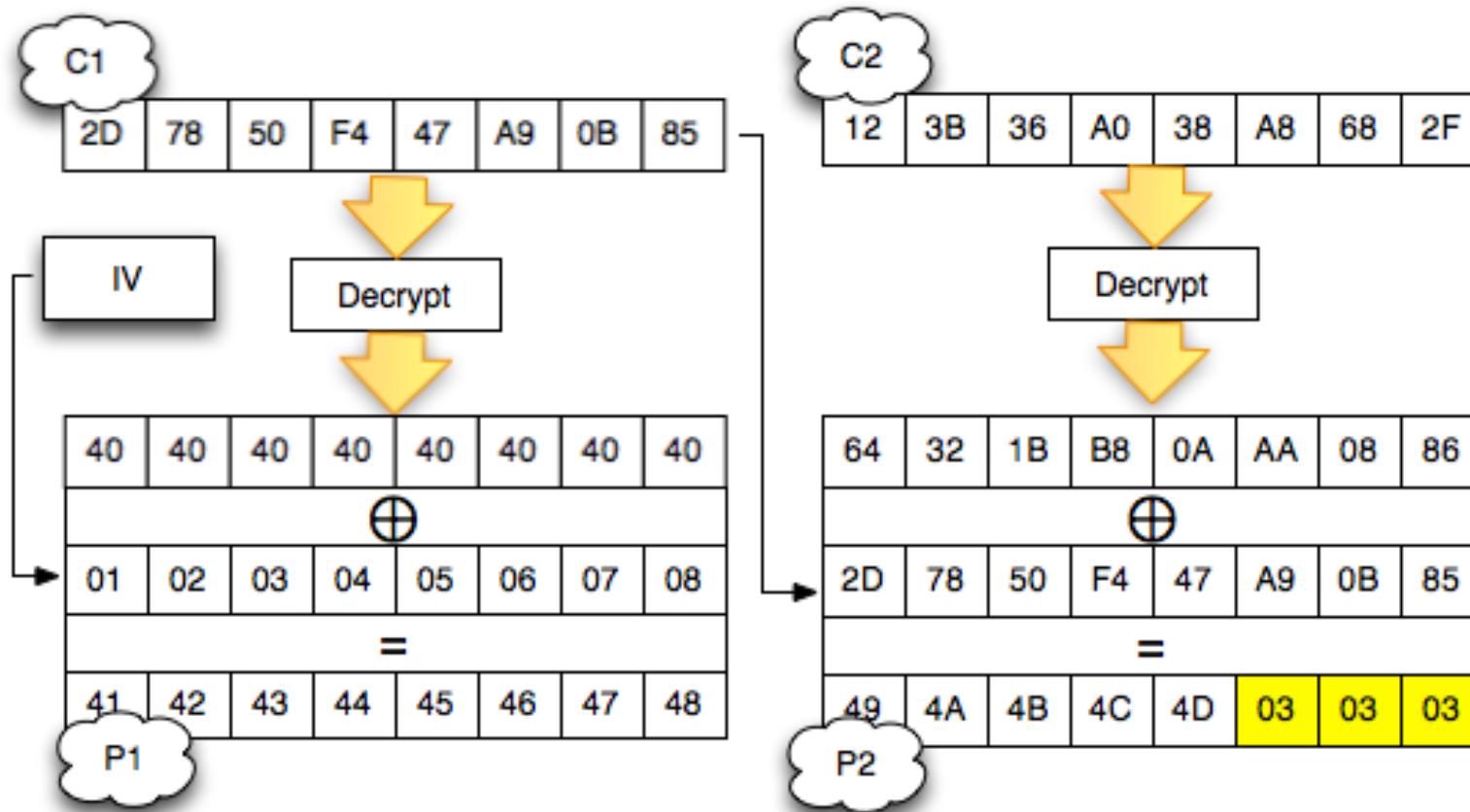
$$C_n = E_k(P_n \oplus C_{n-1})$$

$$P_n = D_k(C_n) \oplus C_{n-1}$$

CBC Mode Encryption



CBC Mode Decryption



Malleability

before (original ciphertext)

Ciphertext	2D	78	50	F4	47	A9	0B	85	12	3B	36	A0	38	A8	68	2F
Plaintext	41	42	43	44	45	46	47	48	49	4A	4B	4C	4D	03	03	03

No padding	41	42	43	44	45	46	47	48	49	4A	4B	4C	4D			
------------	----	----	----	----	----	----	----	----	----	----	----	----	----	--	--	--

after (tampered ciphertext)

Ciphertext	2D	78	50	F4	47	A9	0B	87	12	3B	36	A0	38	A8	68	
Plaintext	14	E2	4D	AC	A9	10	F6	0B	49	4A	4B	4C	4D	03	03	01

85 --> 87

VALID

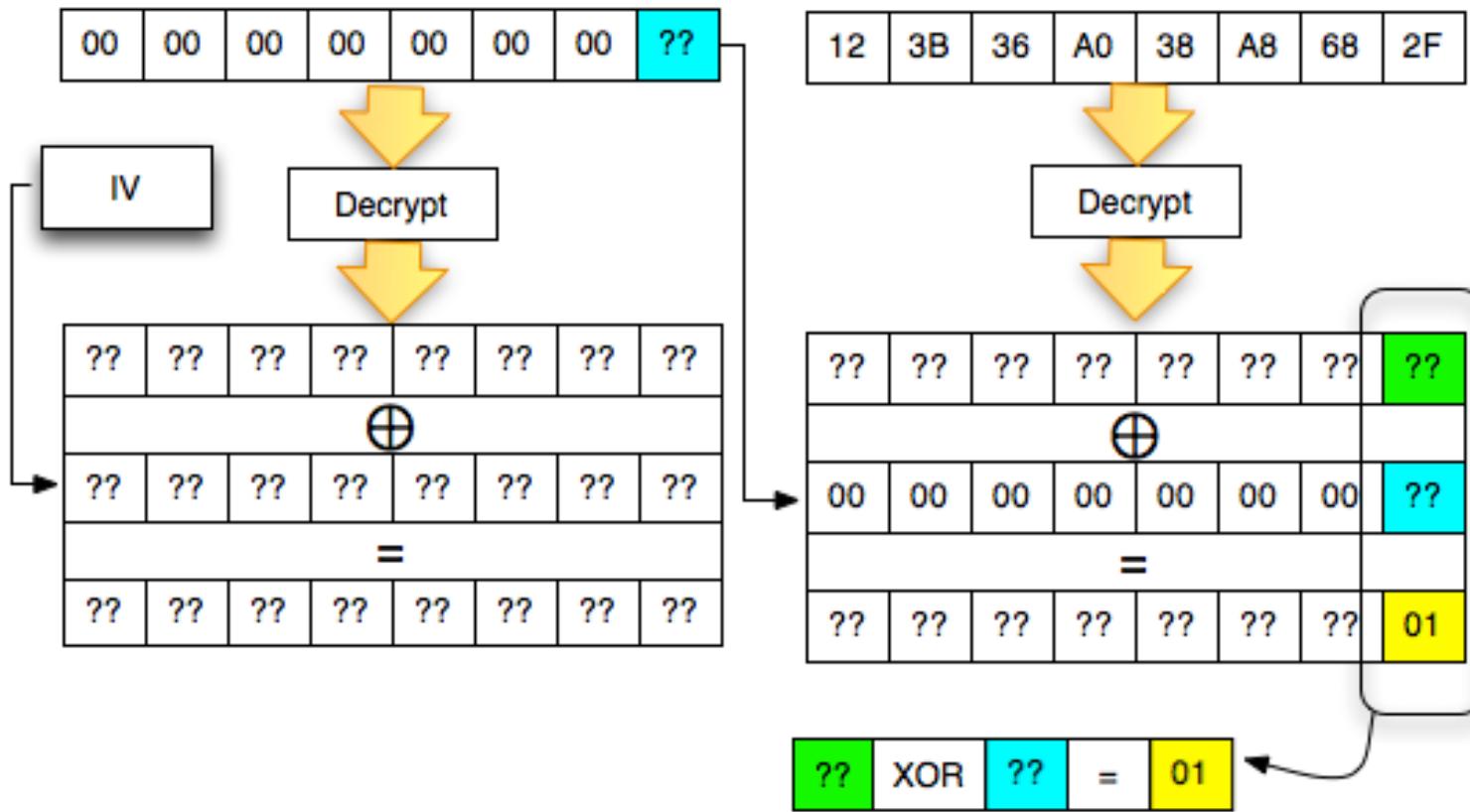
No padding	14	E2	4D	AC	A9	10	F6	0B	49	4A	4B	4C	4D	03	03	
------------	----	----	----	----	----	----	----	----	----	----	----	----	----	----	----	--

Enough Talking, Start
Cracking!

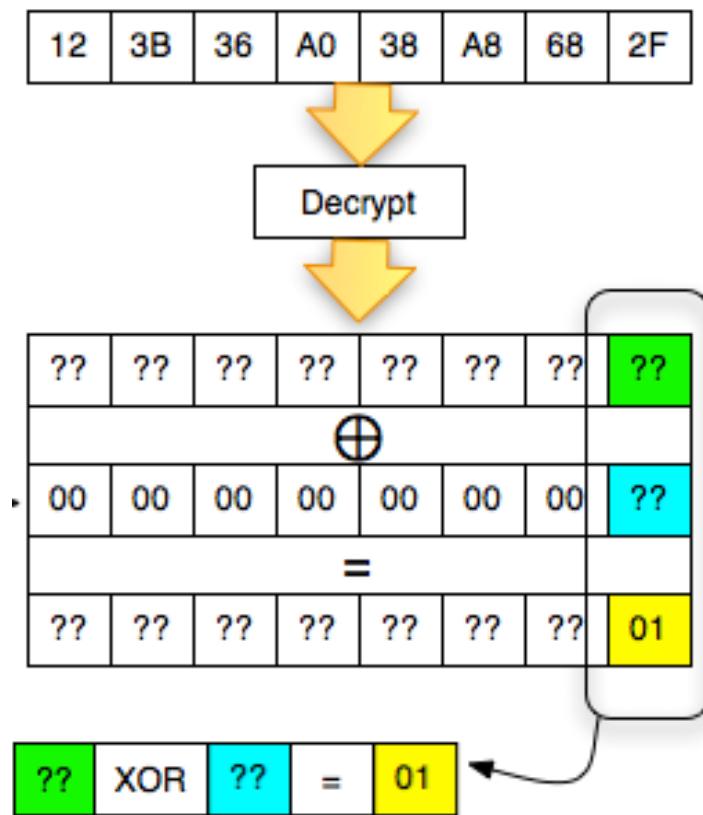
Sample Case

- Decrypt this: 2D7850F447A90B87123B36A038A8682F
- Split into two 8 byte blocks:
 - $C_1 = 2D7850F447A90B87$
 - $C_2 = 123B36A038A8682F$
- Decrypt C_2 first, send two block to oracle:
 - One block + 123B36A038A8682F
- Decrypt one byte at a time (“hollywood style”) starting from the last byte

Decrypt Last Byte

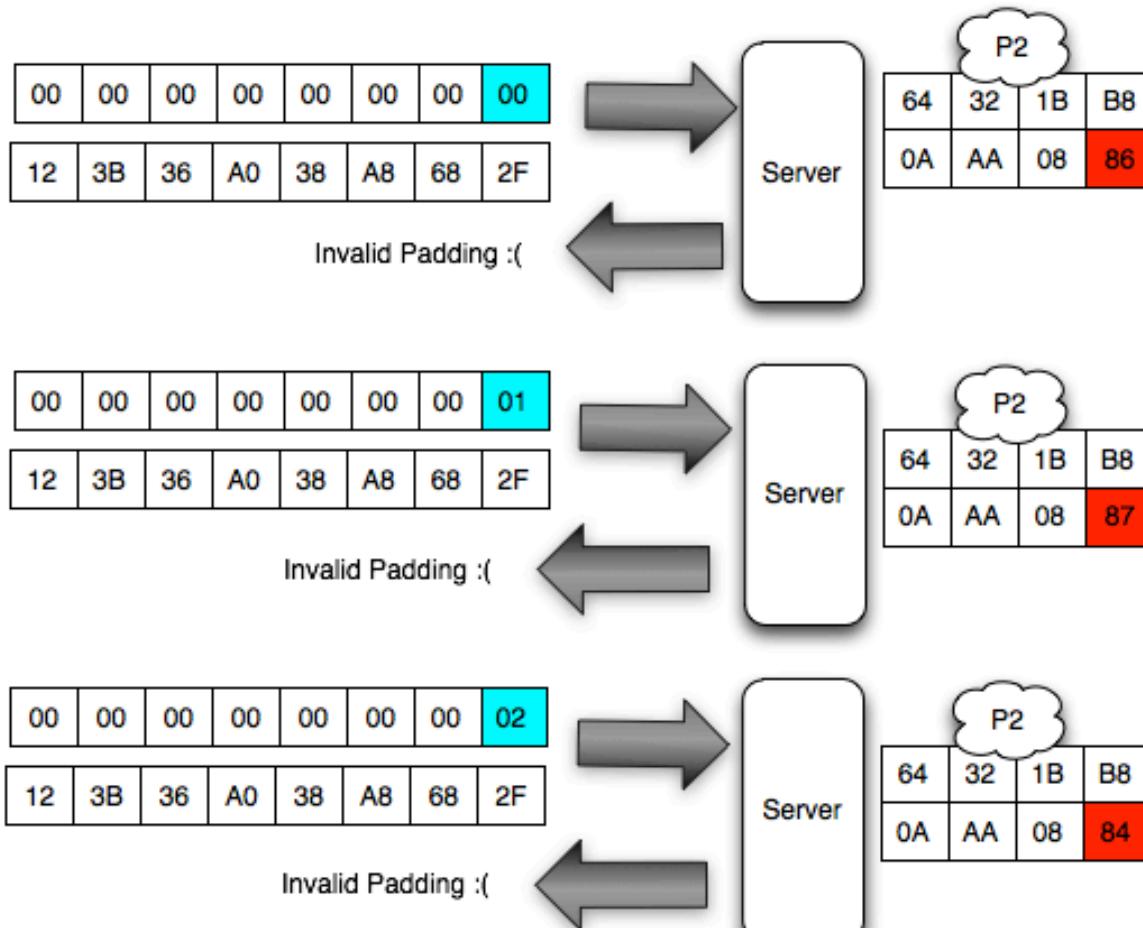


Ask the Oracle

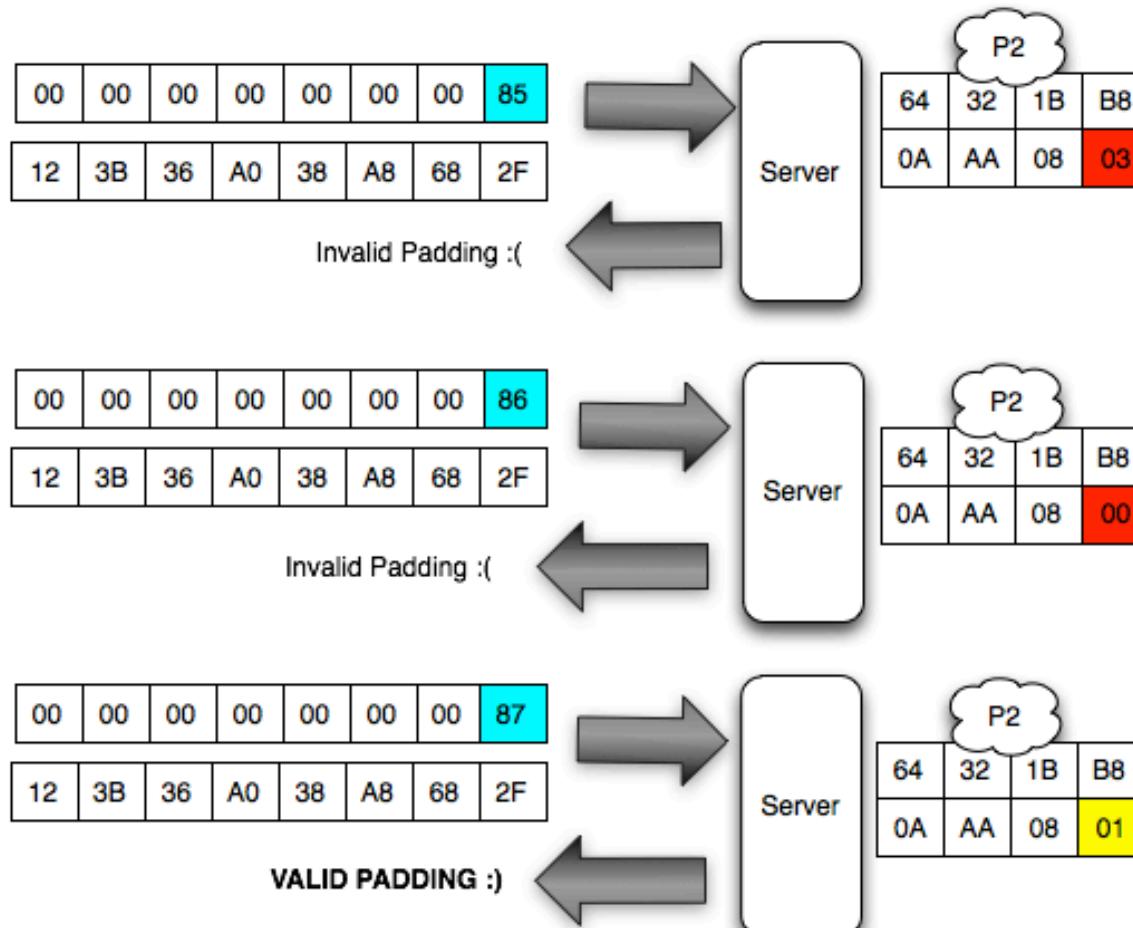


- $A \text{ xor } B = 01$. Find A and B!
- Ask the Oracle:
 - $A \text{ xor } 0 = 01$?
 - $A \text{ xor } 1 = 01$?
 -
 - $A \text{ xor } 255 = 01$?
- Oracle answer:
 - Valid pad = Yes
 - Invalid pad = No

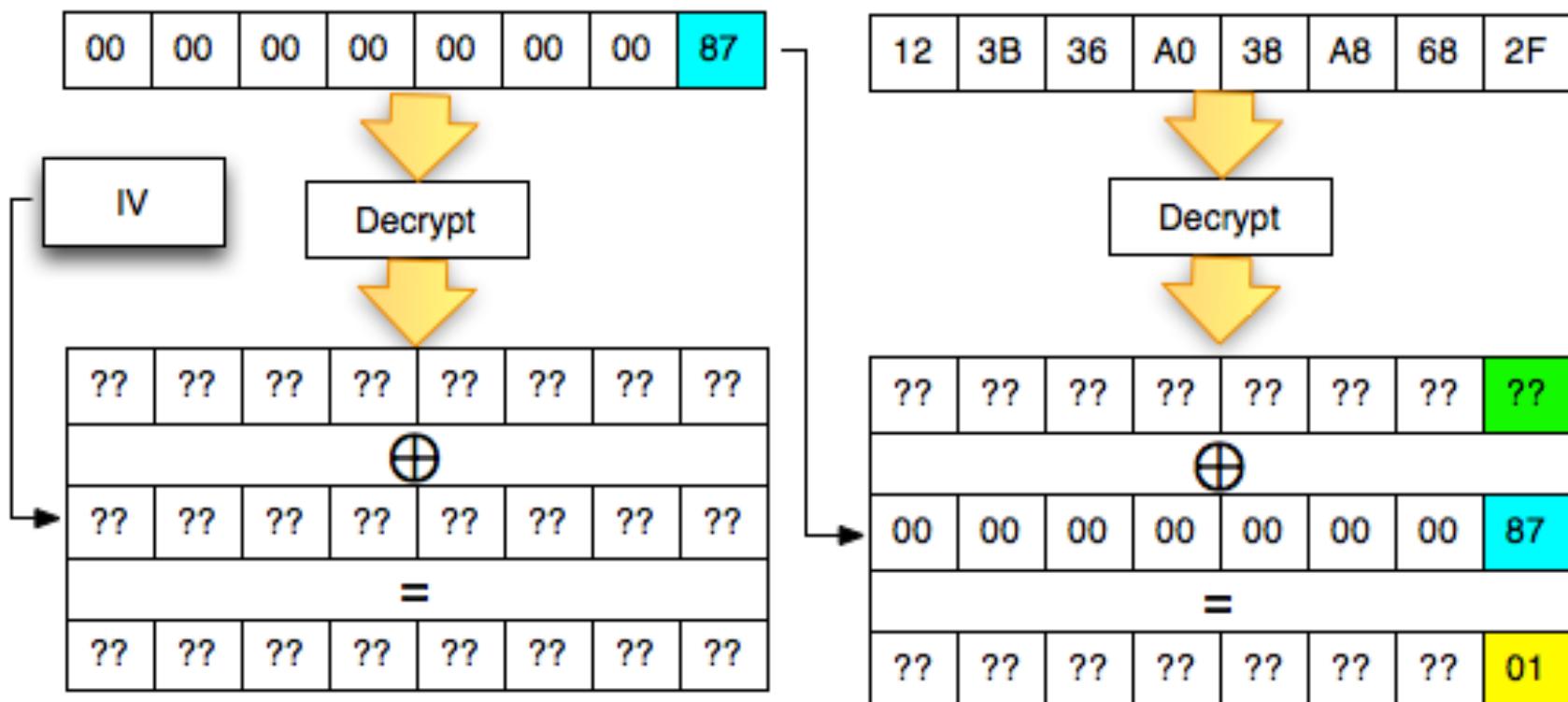
Look for Valid Single Byte Pad



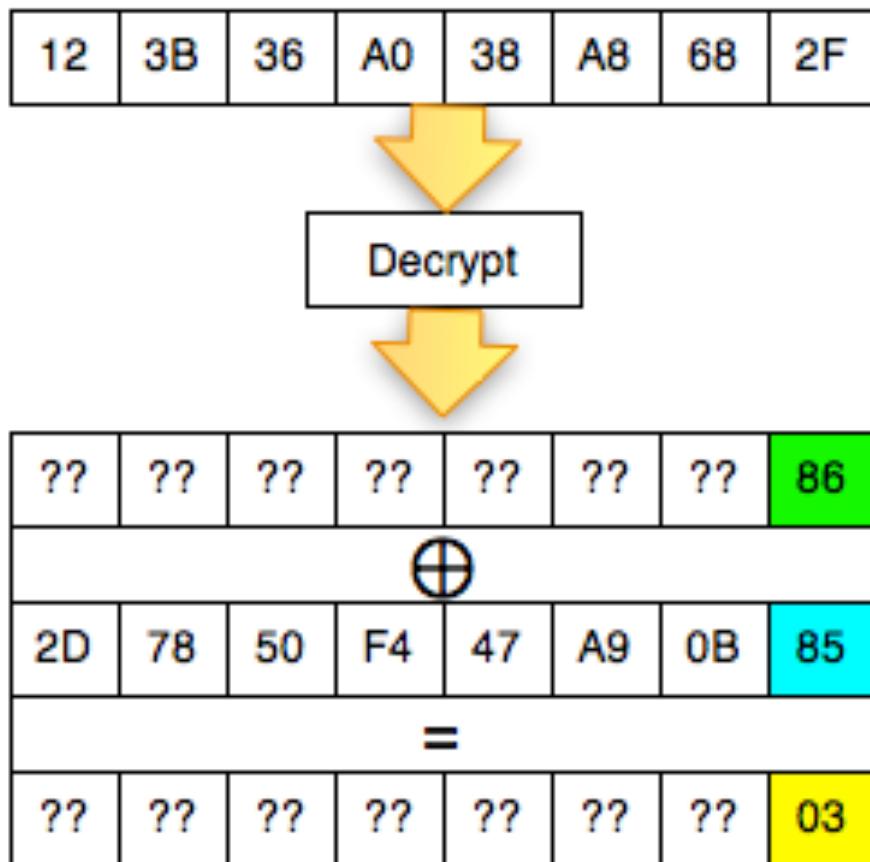
Valid Single Byte Pad Found!



Last Byte Decrypted



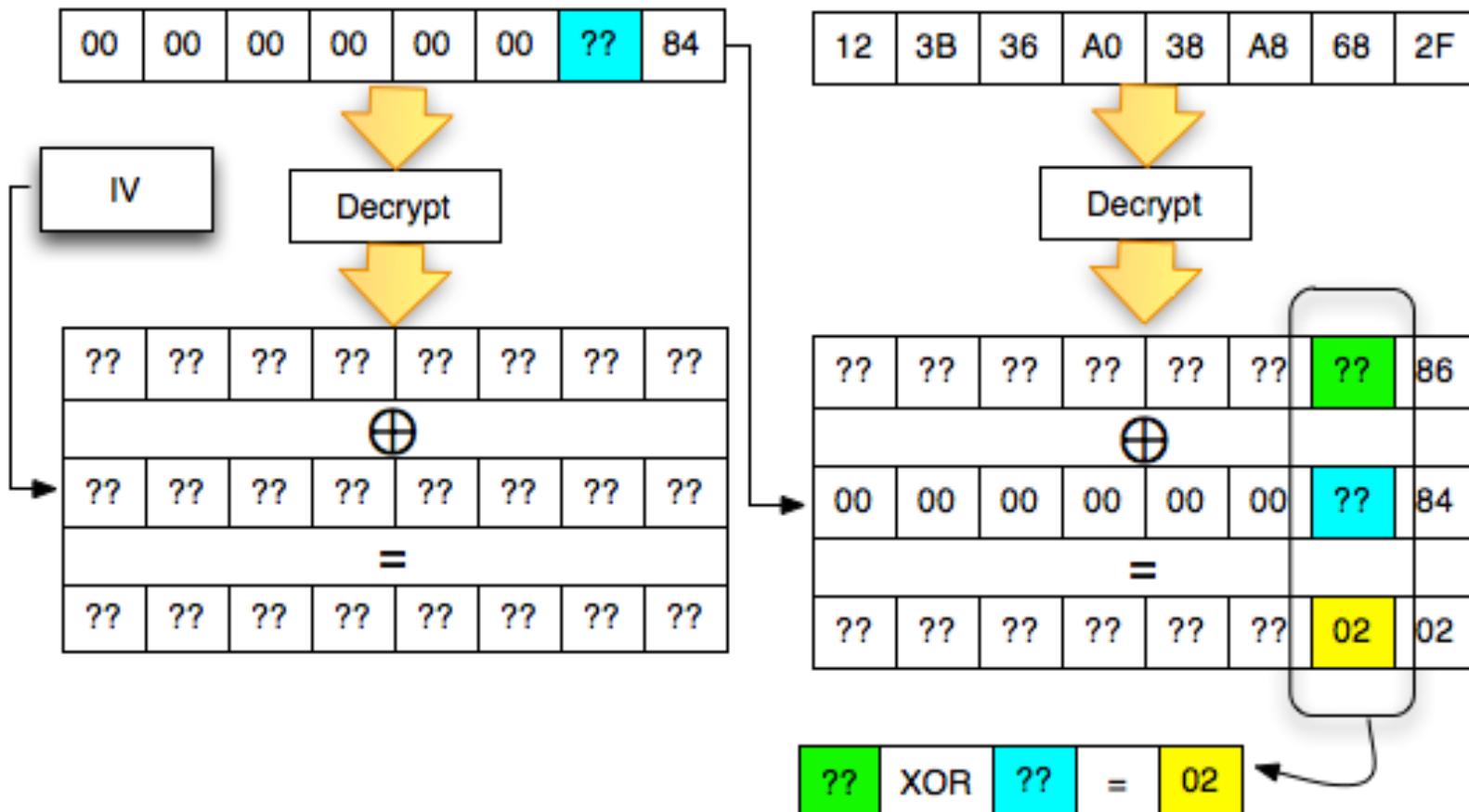
Last Byte = 0x86



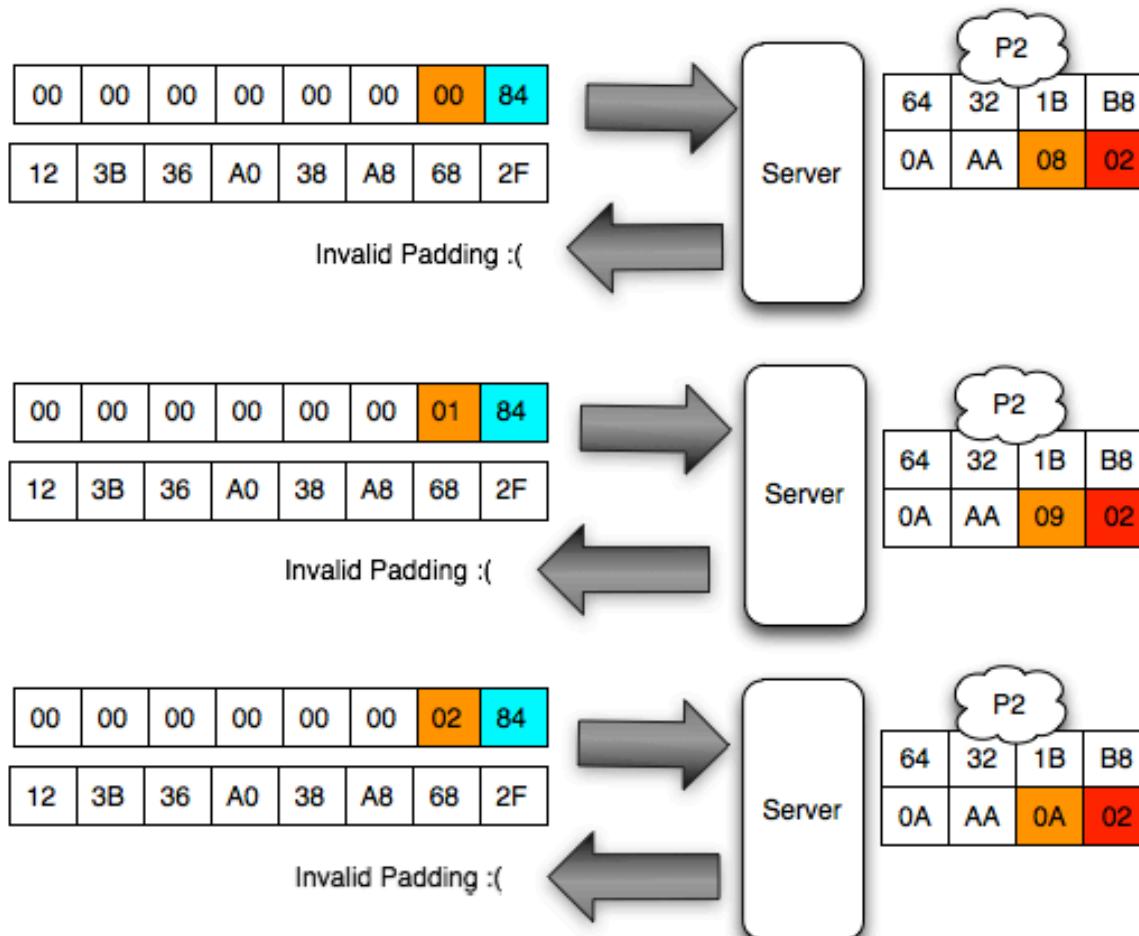
- $A \text{ xor } B = 01$. Find A and B!
- Ask the Oracle:
 - $A \text{ xor } 0x85 = 01$?
- Oracle answer:
 - Valid pad = Yes
- A must be 0x86

Decrypt 7th Byte

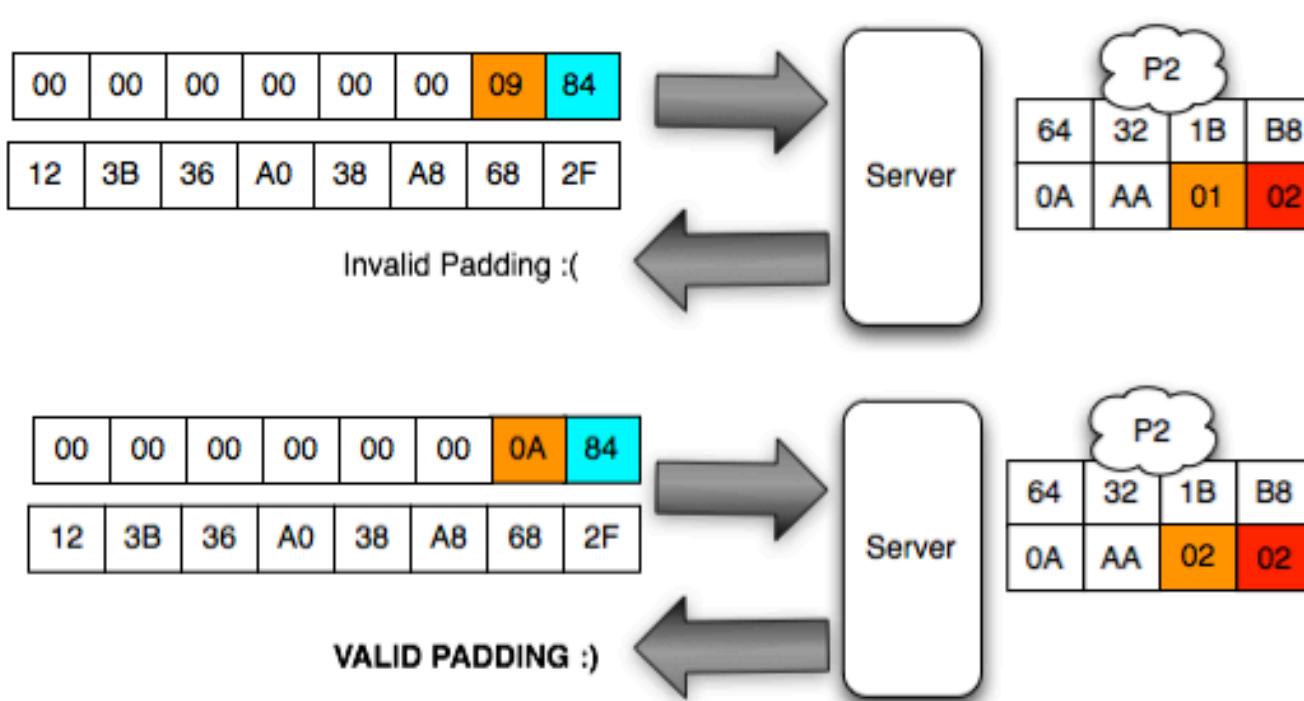
Decrypt 7th Byte



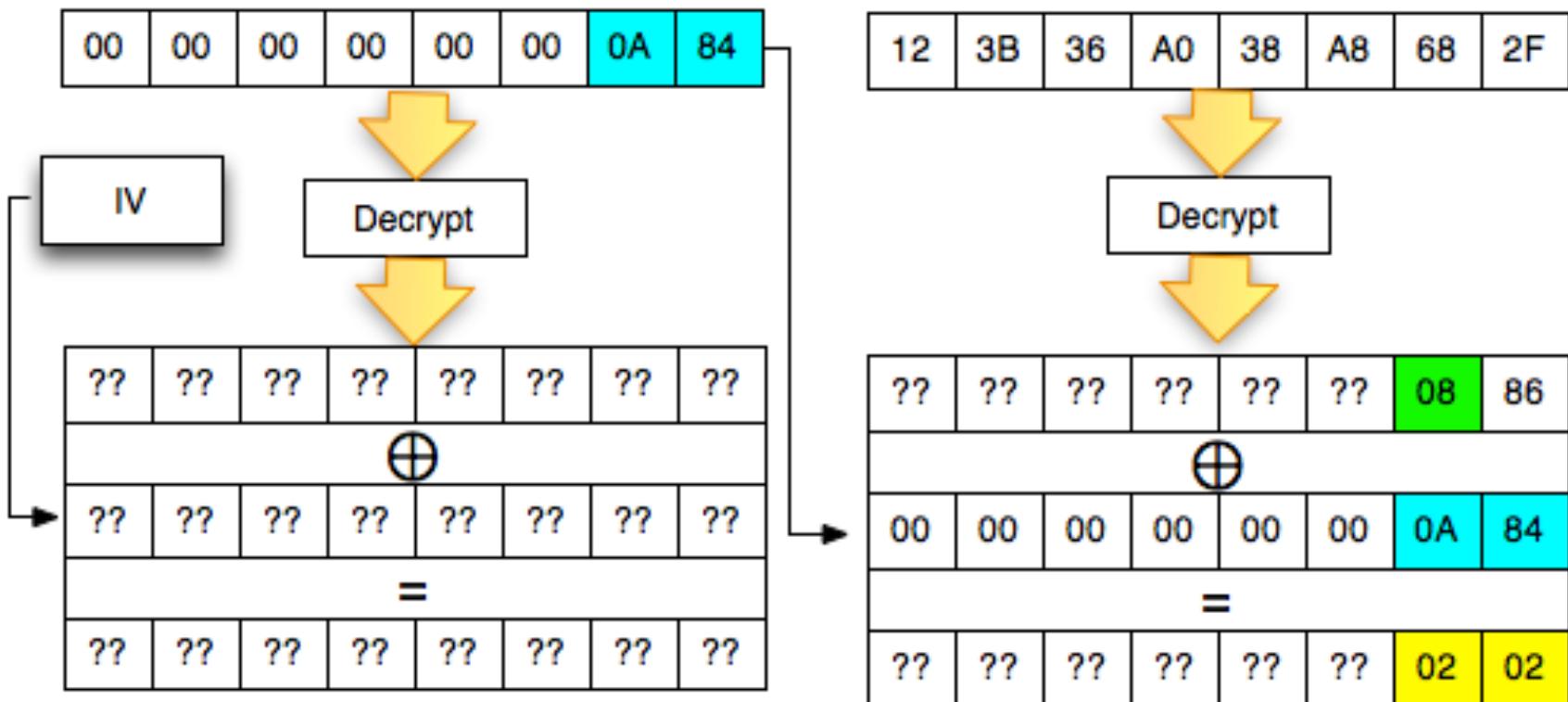
Look for Valid 2 Byte Pad



Valid 2 Byte Pad Found!

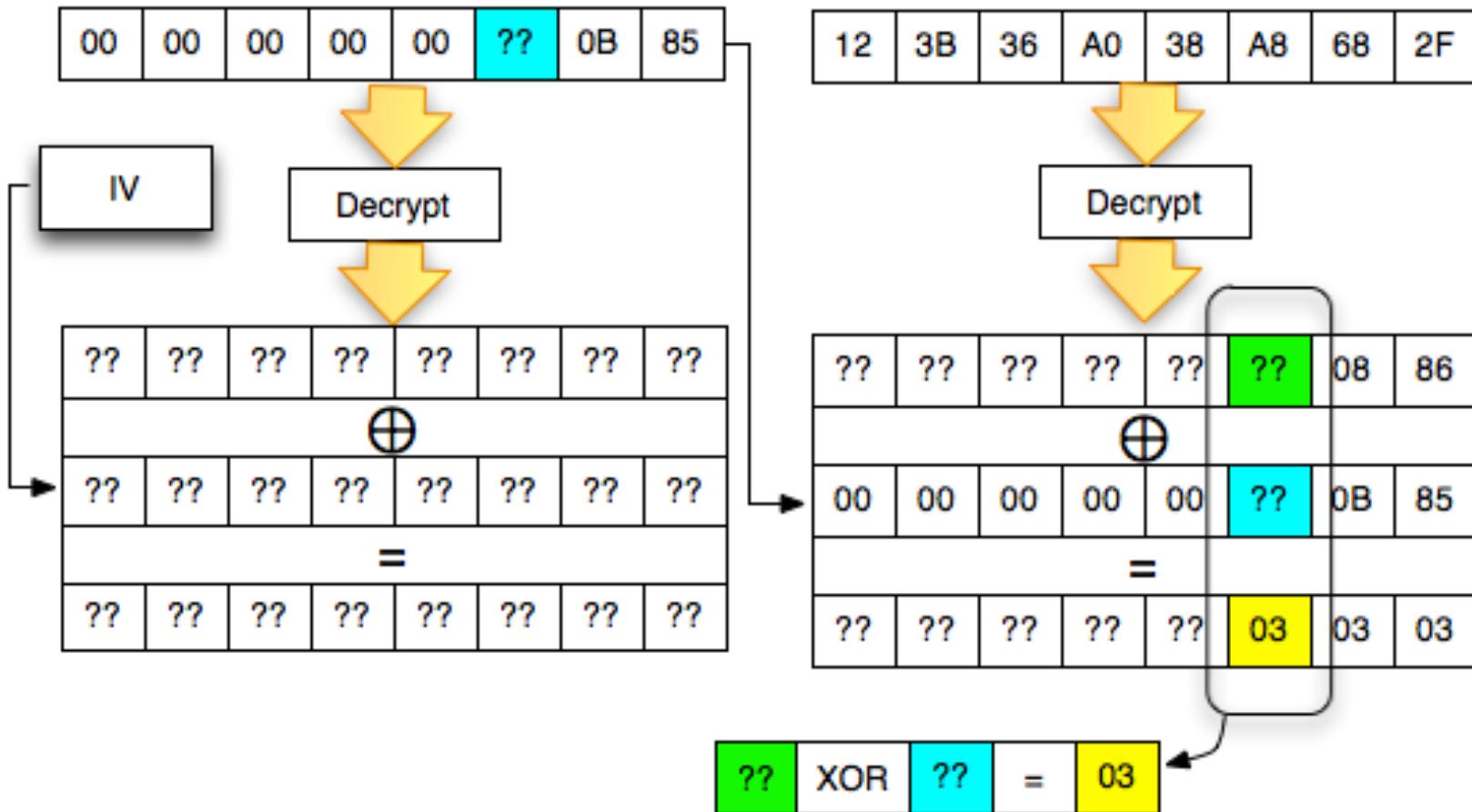


7th Byte Decrypted

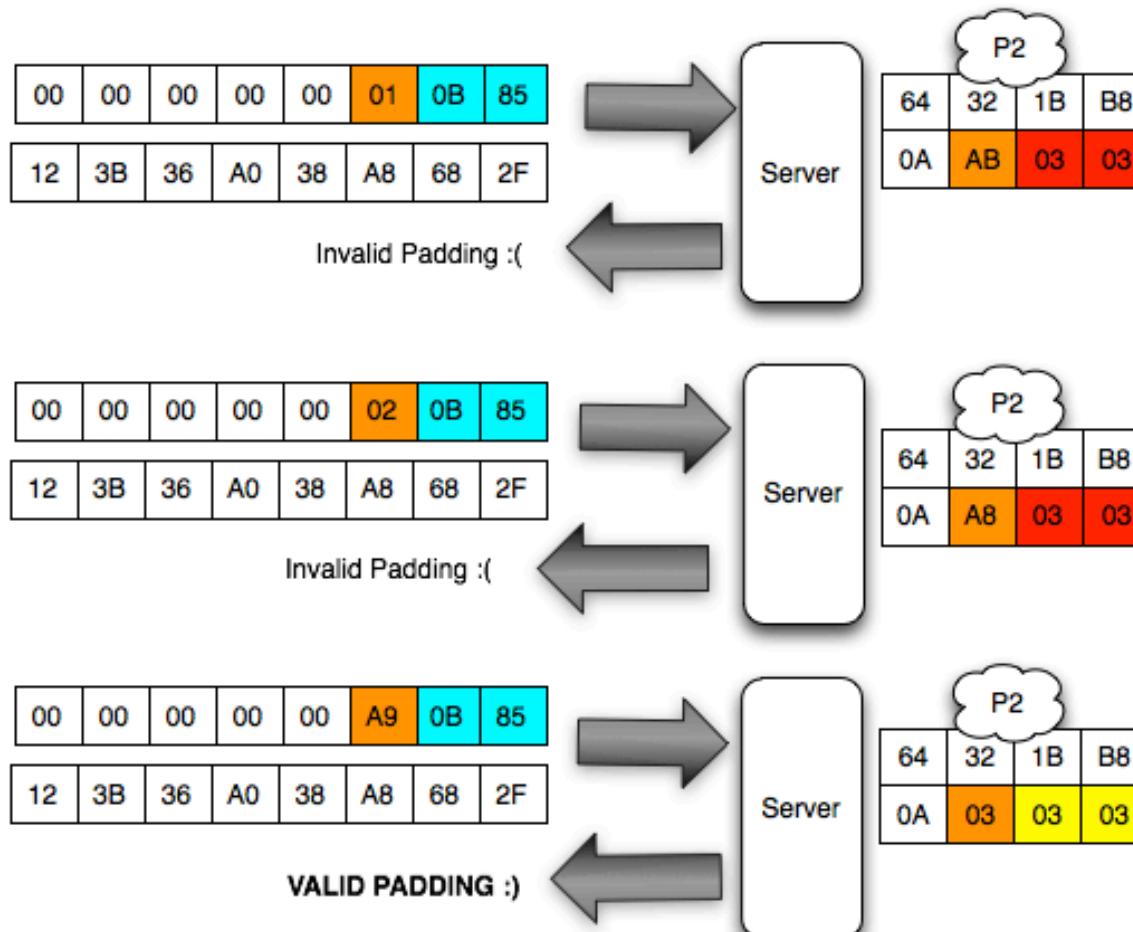


Decrypt 6th Byte

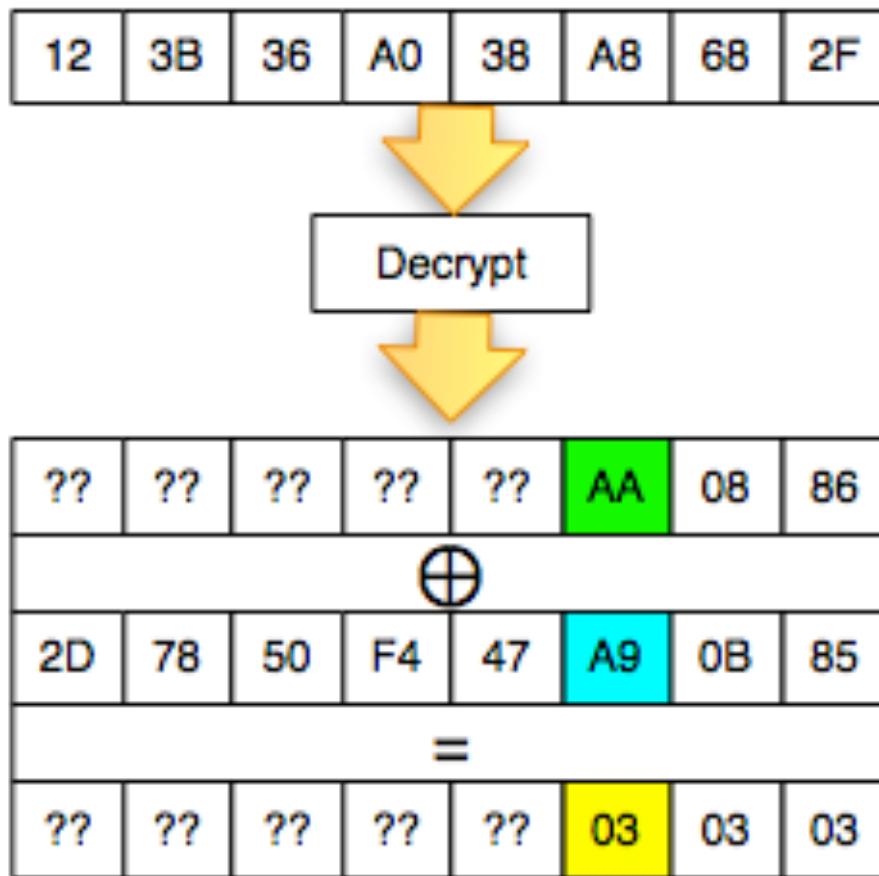
Decrypt 6th Byte



Valid 3 Byte Pad Found

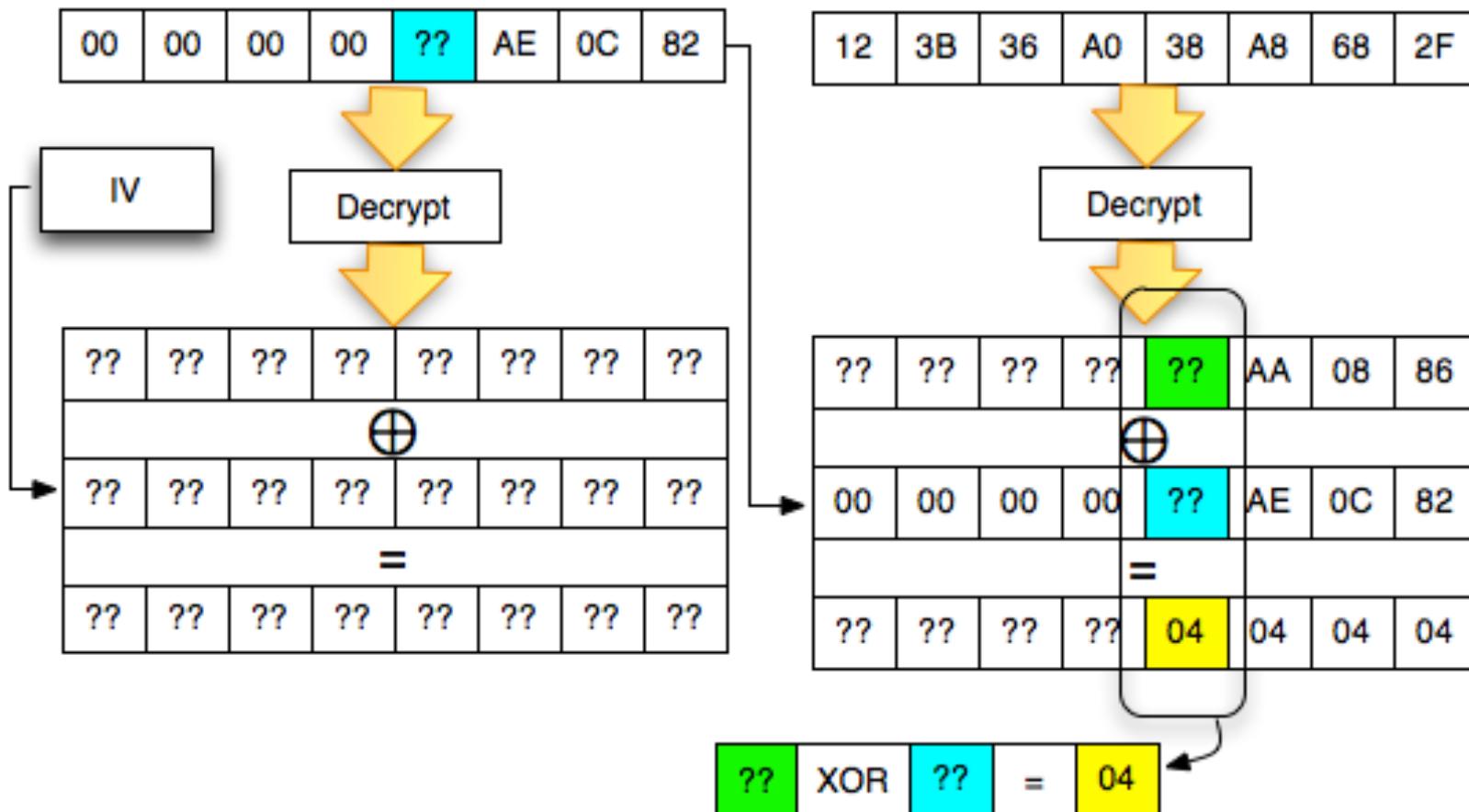


6th Byte Decrypted

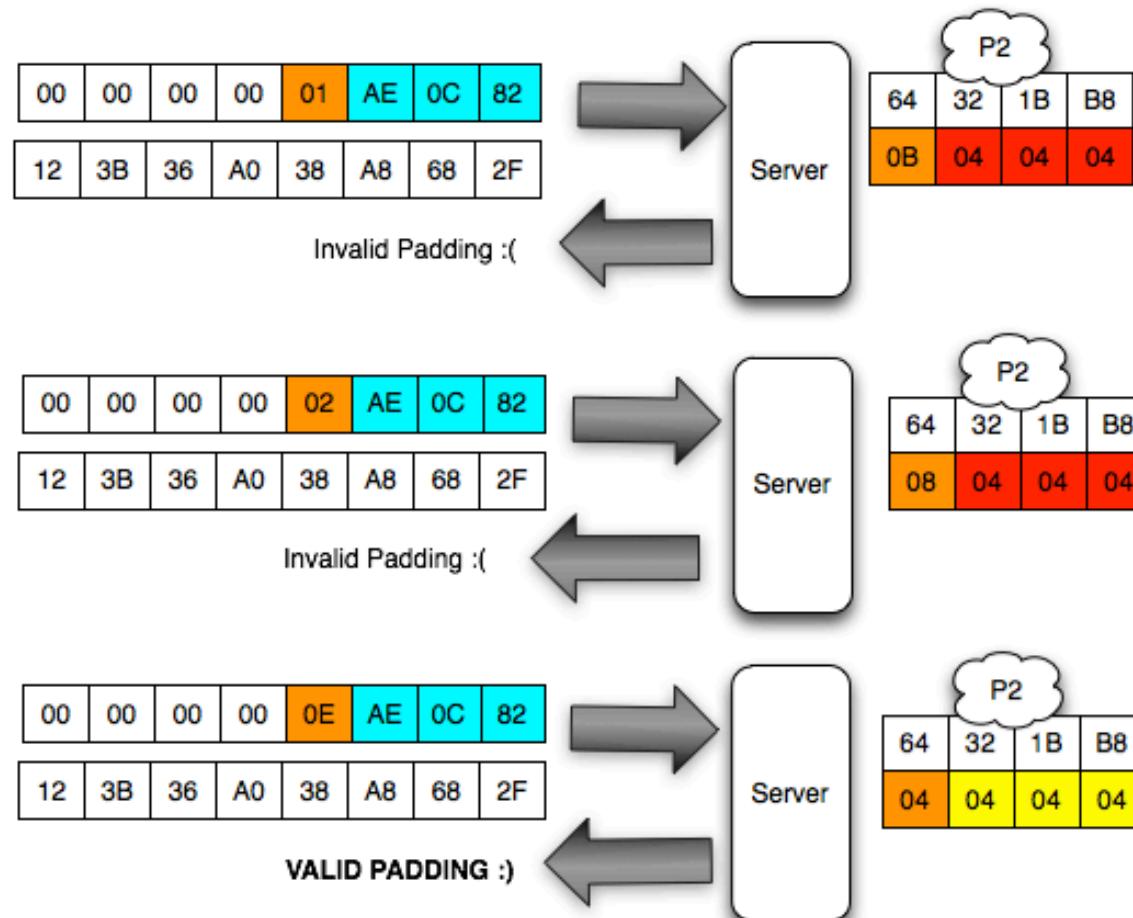


Decrypt 5th Byte

Decrypt 5th Byte

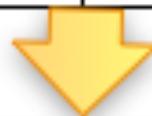


Valid 4 Byte Pad Found



5th Byte Decrypted

12	3B	36	A0	38	A8	68	2F
----	----	----	----	----	----	----	----



Decrypt



??	??	??	??	0A	AA	08	86
----	----	----	----	----	----	----	----



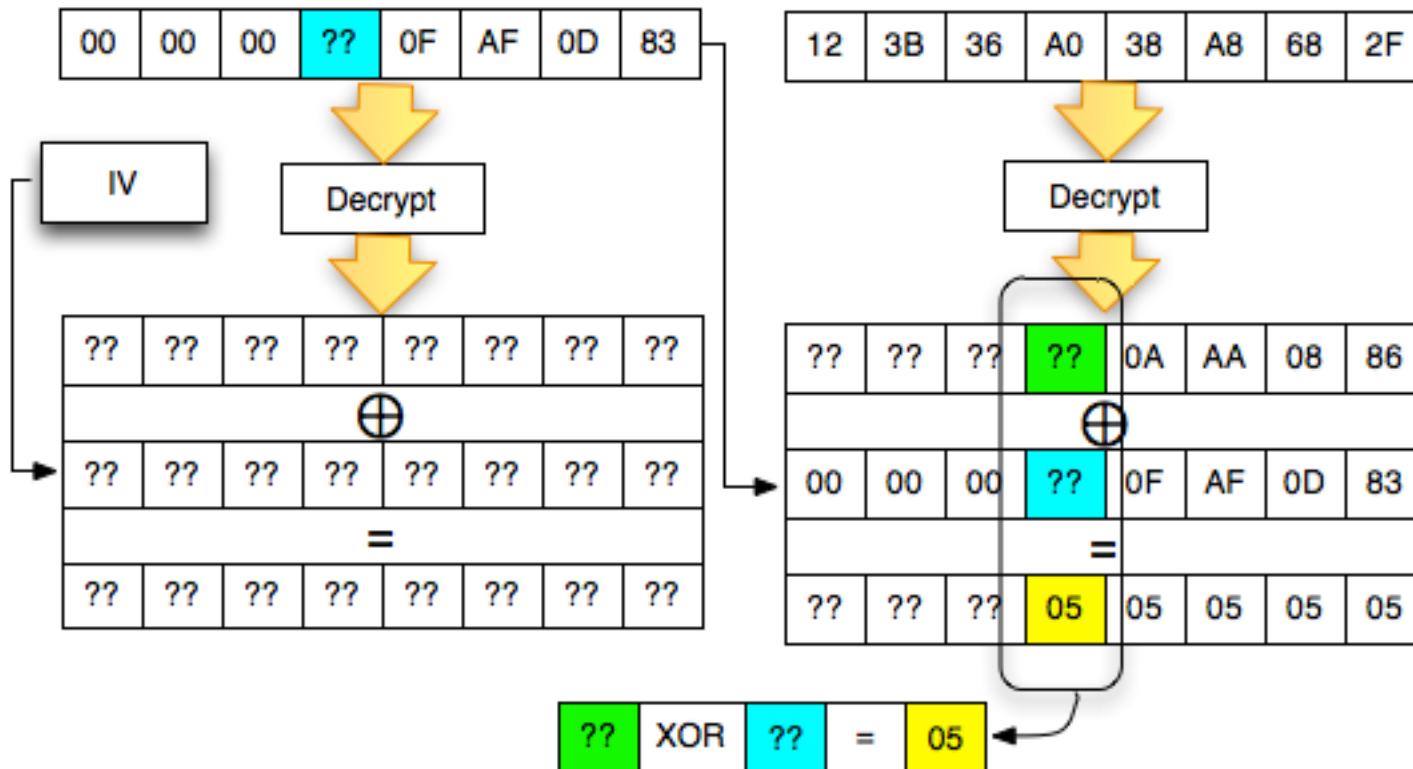
2D	78	50	F4	47	A9	0B	85
----	----	----	----	----	----	----	----



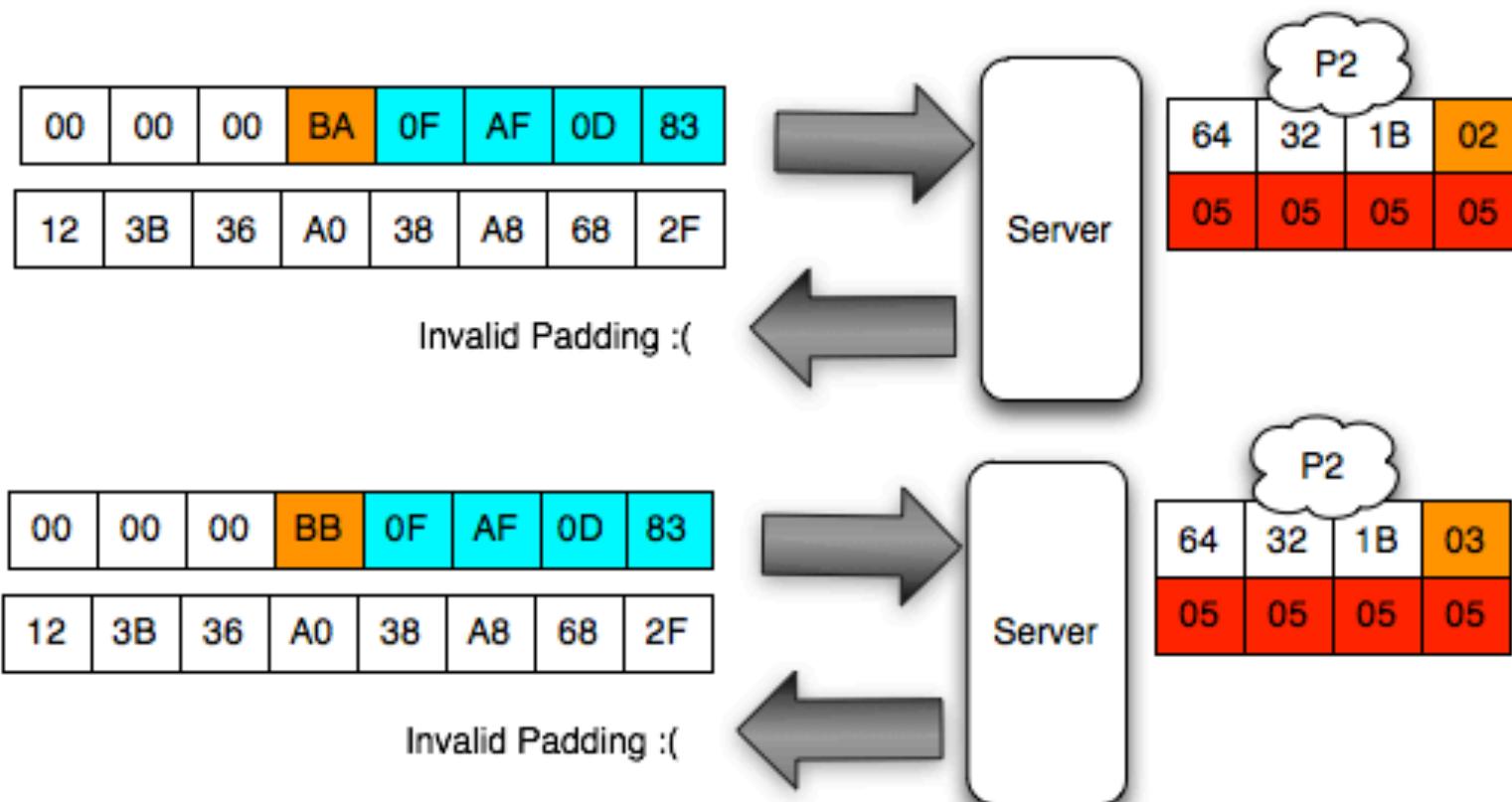
??	??	??	??	4D	03	03	03
----	----	----	----	----	----	----	----

Decrypt 4th Byte

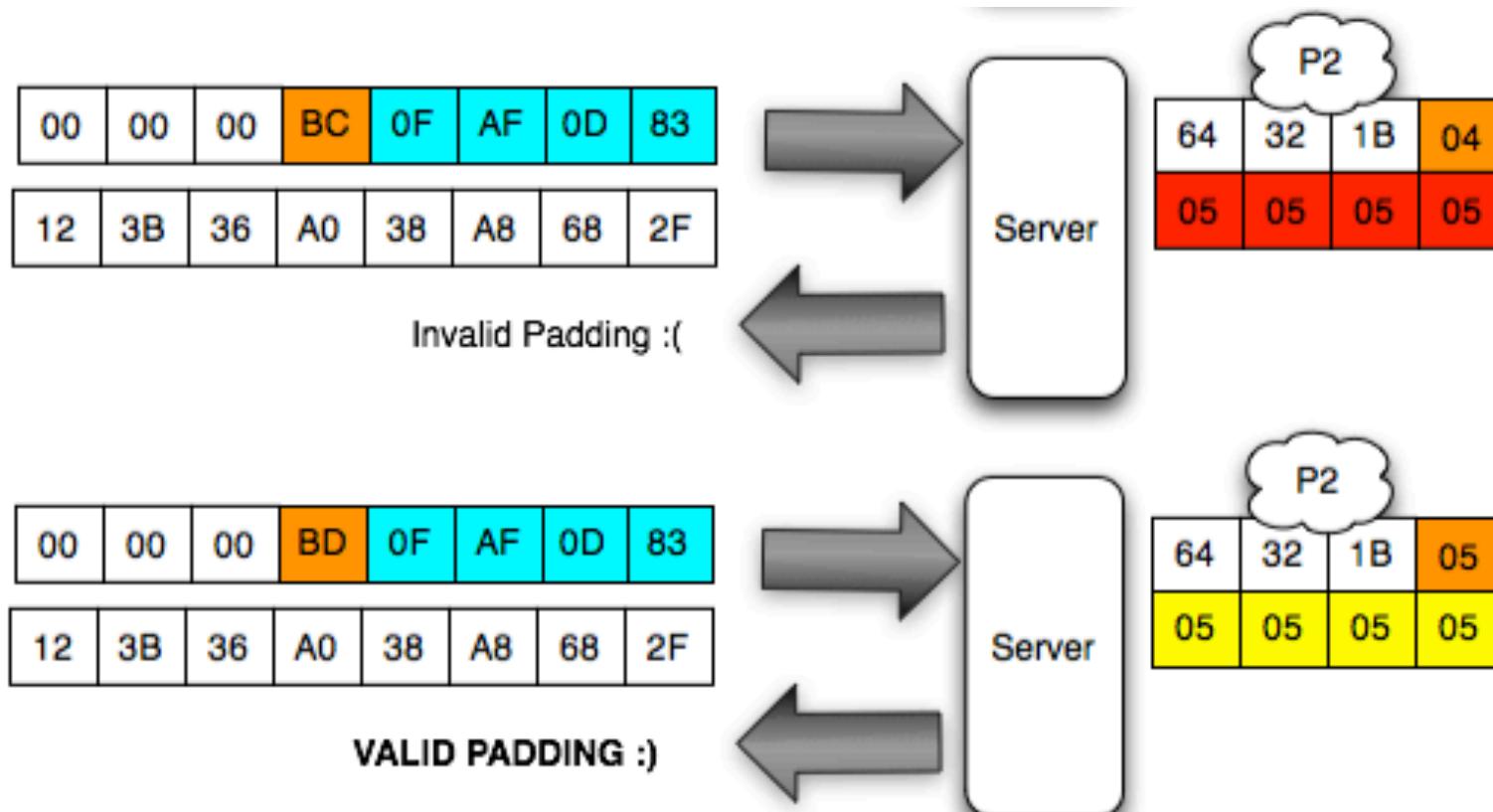
Decrypt 4th Byte



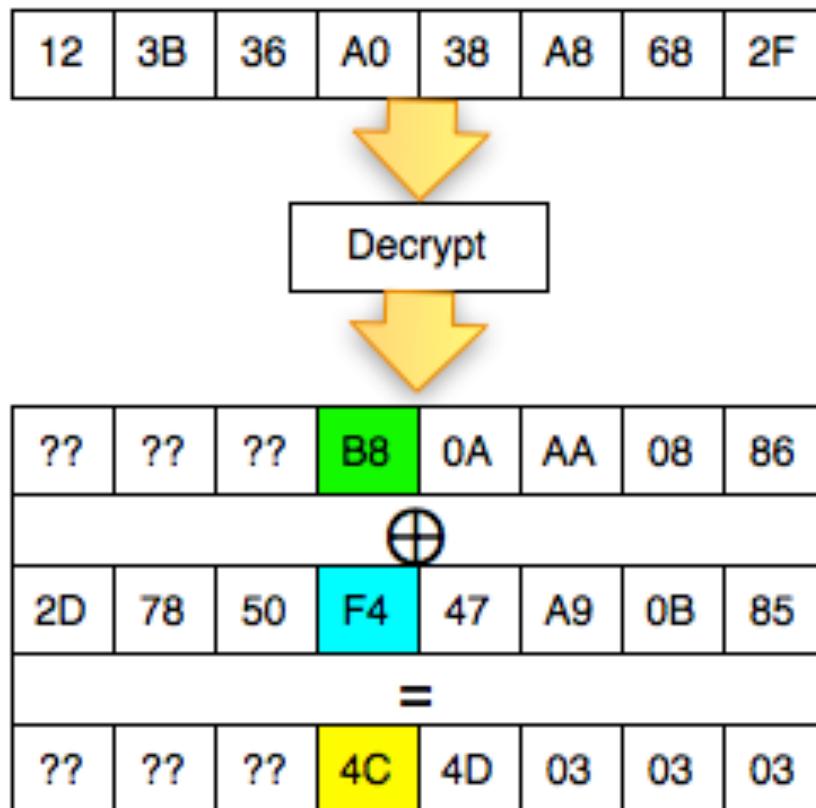
Look for Valid 5 Byte Pad



Valid 5 Byte Pad Found

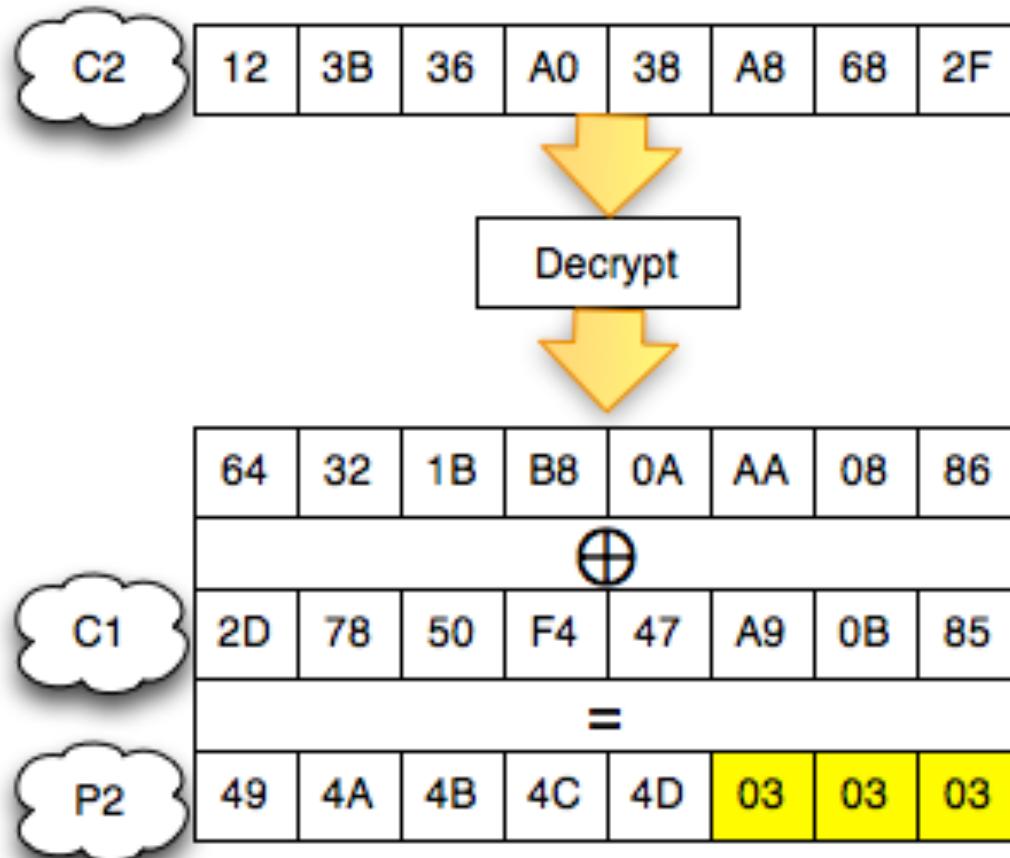


4th Byte Decrypted

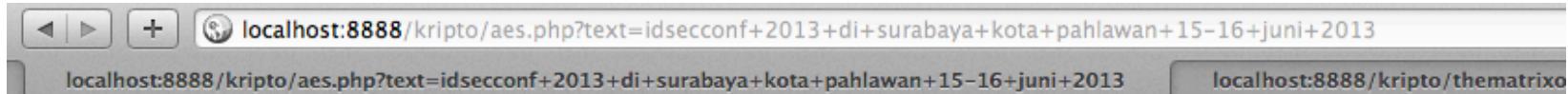


Full Block Decrypted

C₂ Block Decrypted



Case



IV

4A | 32 | B2 | 9B | 21 | 32 | 17 | B9 | 3A | F2 | 7A | FB | 3B | BD | 72 | 7E

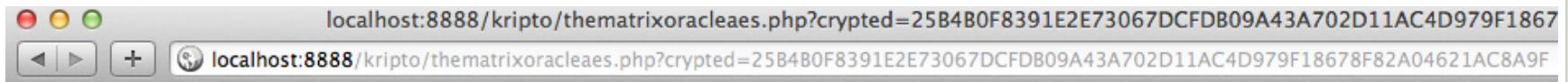
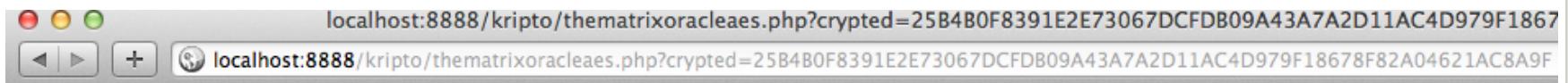
Plaintext dan Ciphertext

69	64	73	65	63	63	6F	6E	66	f	20	32	30	31	33	20	64	69	20	73	75	72	61	62	61	79	61	20	6B	6F	74	61	20	7
i	d	s	e	c	c	o	n		2	0	1	3			d	i		s	u	r	a	b	a	y	a		k	o	t	a		p	
CA	26	07	76	F0	9F	52	5A	19	3C	EE	4F	48	A0	06	23	47	1C	33	5F	00	18	FD	CD	D2	75	64	DF	94	65	12	74	2	

IV + Ciphertext

4A32B29B213217B93AF27AFB3BBD727ECA260776F09F525A193CEE4F48A00623471C335F0018FDCDD27564DF94651:

The Oracle



Decryptor

```
Pucuk:idsecconf rizki$ ./padfinder-aes.py 4A32B29B213217B93AF27AFB3BBD727ECA2607  
76F09F525A193CEE4F48A00623471C335F0018FDCDD27564DF9465127425B4B0F8391E2E73067DCF  
DB09A43A7A2D11AC4D979F18678F82A04621AC8A9F
```

64203331303220666e6f636365736469	Plaintext Blok ke-1:idsecconf 2013 d
2061746f6b2061796162617275732069	Plaintext Blok ke-2:i surabaya kota
6a2036312d3531206e6177616c686170	Plaintext Blok ke-3:pahlawan 15-16 j
080808080808083331303220696e75	Plaintext Blok ke-4:uni 2013

```
Plaintext: 'idsecconf 2013 di surabaya kota pahlawan 15-16 juni 2013'  
Pucuk:idsecconf rizki$ █
```

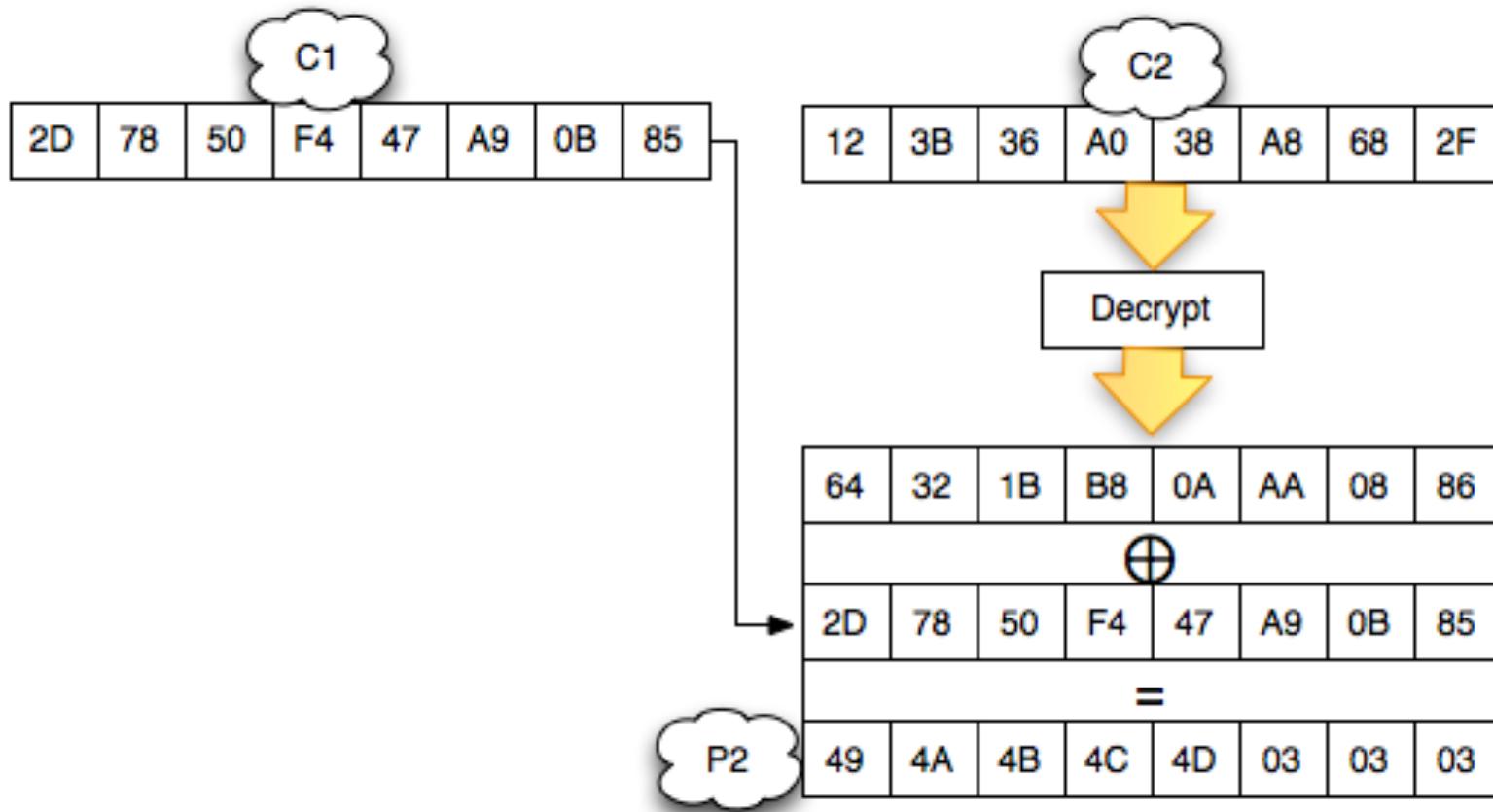
Decryption Demo

Encrypt Fake Message

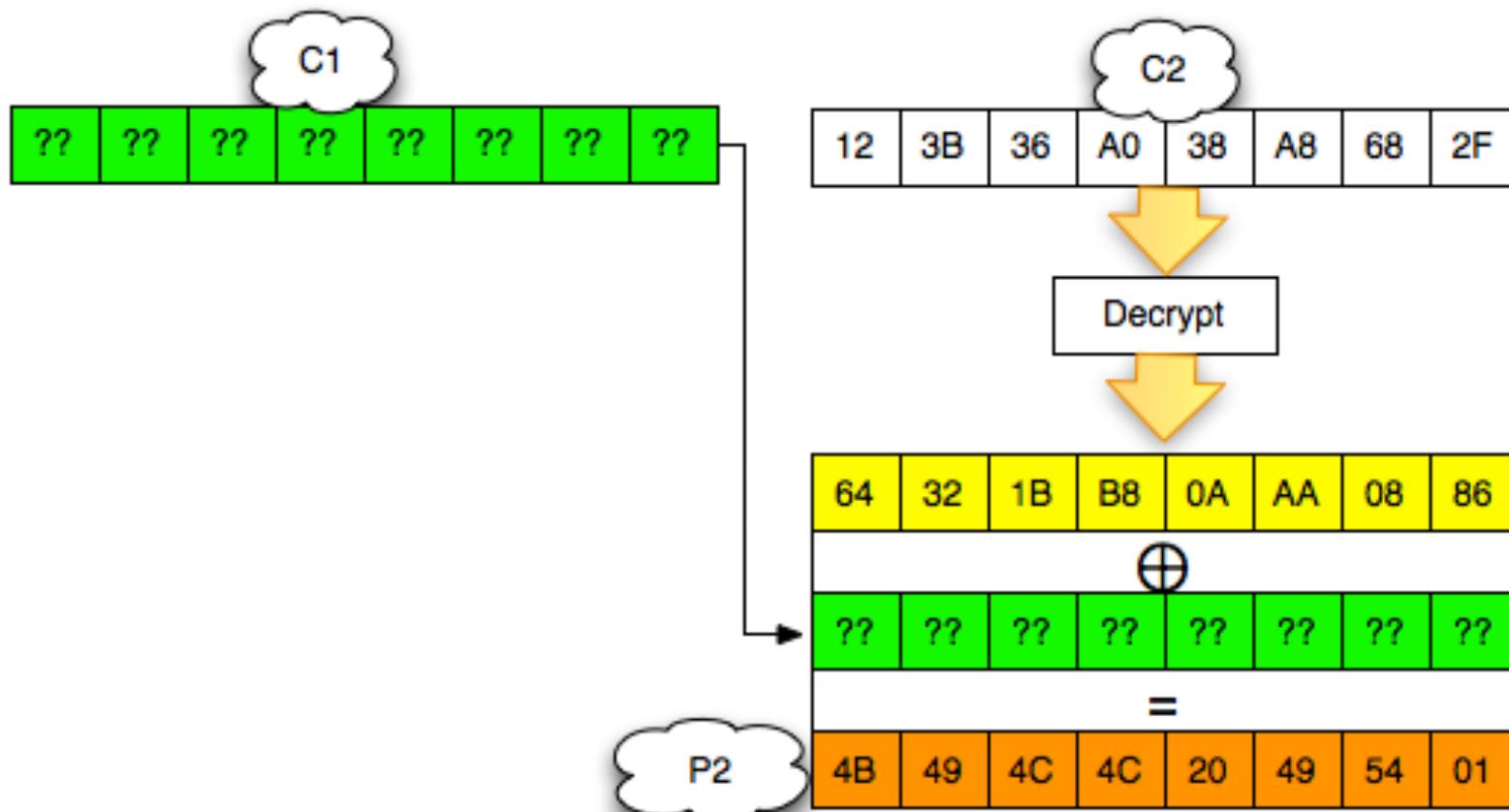
Encrypt without Knowing the Key

- You can make cipher text say whatever you want when decrypted
- Property of CBC mode

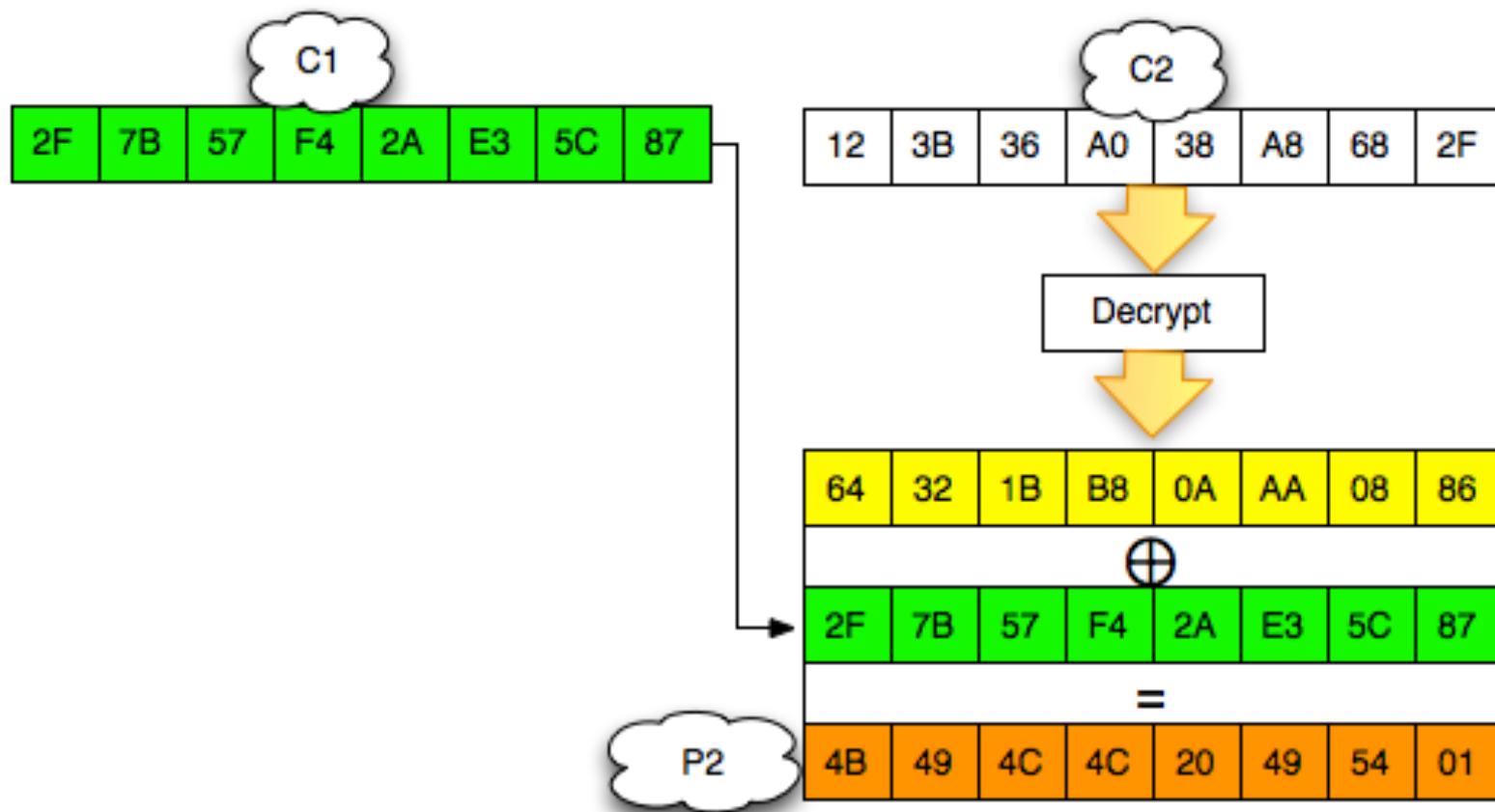
P_2 depends on C_1



“KILL IT”



“KILL IT”



Encryption Procedure

- Encrypt: “BESOK PAGI SERANGAN UMUM IWO JIMA”
- Split plaintext into blocks:
 - $P_1 = \text{'BESOK PA'}$
 - $P_2 = \text{'GI SERAN'}$
 - $P_3 = \text{'GAN UMUM'}$
 - $P_4 = \text{' IWO JIM'}$
 - $P_5 = \text{'A'}+07+07+07+07+07+07$

Encryption Procedure

- Choose C_5 all-zeros
- Use padding oracle attack to find $\text{Decrypt}(C_i)$
- $C_4 = \text{Decrypt}(C_5) \text{ XOR } P_5$
- $C_3 = \text{Decrypt}(C_4) \text{ XOR } P_4$
- $C_2 = \text{Decrypt}(C_3) \text{ XOR } P_3$
- $C_1 = \text{Decrypt}(C_2) \text{ XOR } P_2$
- $\text{IV} = \text{Decrypt}(C_1) \text{ XOR } P_1$

Encryption Demo

```
$ □
```



Authenticated Encryption

Authenticate before Decrypt

- Why we need to authenticate/verify encrypted message before decrypting it ? It's already encrypted with shared secret key, after all.
- Imagine that only Alice and Bob know the key. If Bob could decrypt a cipher text with the secret key and get a clean and understandable plain text, then Bob know it only could be encrypted by Alice
- Many people have thought that, but they were wrong
- Without message authentication, active attacker could use padding oracle attack to decrypt and also encrypt without knowing the key

Encryption and MAC

- Encryption provides confidentiality, it doesn't provide integrity and authenticity
- Don't use encryption without message authentication
- Encrypt your message then calculate MAC
- Never decrypt message without checking MAC
- Decrypt only when ciphertext is MAC-authenticated