

Q 2.1)

g1 part: The type of g1 is:

$$('a \rightarrow \text{bool}) \rightarrow 'a \text{ list} \rightarrow 'a \text{ list}$$

A brief argument (that is not required):

From the first clause we see $p: 'a \rightarrow \text{bool}$
and $x: 'a$ (from the 'when p x' part).

Then $x::xs : 'a \text{ list}$ and $x::g1\ p\ xs: 'a \text{ list}$.

$g1\ p\ xs$ gives the longest prefix of xs where every element x_i satisfies p , i.e. $p\ x_i = \text{true}$.

g2 part: The type of g2 is:

$$\underbrace{('a \rightarrow 'a)}_f \rightarrow \underbrace{('a \rightarrow 'a)}_h \rightarrow \underbrace{\text{int}}_n \rightarrow \underbrace{'a}_x \rightarrow 'a$$

Brief (not required) argument:

The match clause shows directly that $n: \text{int}$.

Since x and $f\ x$ both occur as last argument of $g2$, we have $f: 'a \rightarrow 'a$.

f and h exchange position in calls of $g2$
so $h: 'a \rightarrow 'a$.

$$g2\ f\ h\ n\ x =$$

$\left\{ \begin{array}{l} \text{fails with an exception if } n < 0 \\ f(h(f(\dots f\ x)\dots)) \\ \text{if } n > 0, n \text{ is odd and there are } \frac{n+1}{2} \text{ applications of } f \text{ and } \frac{n-1}{2} \text{ of } h \\ h(f(\dots h(f(x))\dots)) \\ \text{if } n > 0, n \text{ is even and there are } \frac{n}{2} \text{ applications of } f \text{ and of } h \end{array} \right.$

Q 2.2)

2

let rec glA p xs acc =
 match xs with
 | x::rest when p x → glA p rest (x::acc)
 | _ → List.rev acc

let gl p xs = glA p xs []

let rec glC p xs k =
 match xs with
 | x::rest when p x
 → glC p rest (fun v → k(x::v))
 | _ → k []

let gl p xs = glC p xs id

Q 2.3)

g2 is tail recursive because the only recursive call of g2 in the last clause

$!n \rightarrow \underline{g2} \ n \ f \ (n-1) \ (fx)$

is in a tail position. When that call returns a value, there is no work left to be done.