Data Management in Large-Scale Distributed Systems

MapReduce and Hadoop

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2021

References

- Coursera Big Data, University of California San Diego
- The lecture notes of V. Leroy
- Designing Data-Intensive Applications by Martin Kleppmann
- Mining of Massive Datasets by Leskovec et al.

In this course

- History of MapReduce
- Overview of the Hadoop Eco-system
- Description of HDFS and Hadoop MapReduce
- Our first MapReduce programs

Agenda

Introduction to MapReduce

The Hadoop Eco-System

HDFS

Hadoop MapReduce

MapReduce at Google

Publication

- The Google file system, S. Ghemawat et al. SOSP 2003.
- MapReduce: simplified data processing on large clusters, J. Dean and S. Ghemawat. OSDI 2004.

Main ideas

- Data represented as key-value pairs
- Two main operations on data: Map and Reduce
- A distributed file system
 - Compute where the data are located

Use of MapReduce at Google

- Used to implement several tasks:
 - ▶ Building the indexing system for Google Search
 - Extracting properties of web pages
 - Graph processing
 - etc.
- Google does not use MapReduce anymore¹
 - Moved on to more efficient technologies
 - We will study BigTable (data storage) in this course
 - The main principles are still valid

https://www.datacenterknowledge.com/archives/2014/06/25/
google-dumps-mapreduce-favor-new-hyper-scale-analytics-system

MapReduce

The Map operation

- Transformation operation
 - ► A function is applied to each element of the input set
- $map(f)[x_0,...,x_n] = [f(x_0),...,f(x_n)]$
- map(*2)[2,3,6] = [4,6,12]

The Reduce operation

- Aggregation operation (fold)
- reduce $(f)[x_0,...,x_n] = [f((x_0),f((x_1),...,f(x_{n-1},x_n)))]$
- reduce(+)[2,3,6] = (2+(3+6)) = 11
- In MapReduce, Reduce is applied to all the elements with the same key

Why MapReduce became very popular?

Main advantages

- Simple to program
- Scales to large number of nodes
 - Targets scale out (share-nothing) infrastructures
- Handles failures automatically

Simple to program

Provides a distributed computing execution framework

- Simplifies parallelization
 - Defines a programming model
 - Handles distribution of the data and the computation
- Fault tolerant
 - Detects failures
 - Automatically takes corrective actions
- Code once (expert), benefit to all

Limits the operations that a user can run on data

- Inspired from functional programming (MapReduce)
- Allows expressing several algorithms
 - ▶ But not all algorithms can be implemented in this way

Scales to large number of nodes

Data parallelism

- Running the same task on different (distributed) data pieces in parallel.
- As opposed to Task parallelism that runs different tasks in parallel (e.g., in a pipeline)

Move the computation instead of the data

- The distributed file system is central to the framework
 - ► GFS in the case of Google
 - ► Heavy use of partitioning
- The tasks are executed where the data are stored
 - Moving data is costly

Fault tolerance

Motivations

- Failures are the norm rather than the exception¹.
 - In Google datacenters, jobs can be preempted at any time
 - MapReduce jobs have low priority and have high chances of being preempted
 - A 1-hour task has 5% chances of being preempted
- Dealing with stragglers (slow machines)

¹The Google file system, S. Ghemawat et al, 2003

Fault tolerance

Mechanisms

- Data are replicated in the distributed file system
- Results of computation are written to disk
- Failed tasks are re-executed on other nodes
- Tasks can be executed multiple times in parallel to deal with stragglers
 - Towards the end of a computation phase

Word Count

Description

- Input: A set of lines including words
 - ► Pairs < line number, line content >
 - ► The initial keys are ignored in this example
- Output: A set of pairs < word, nb of occurrences >

Input

- < 1, "aaa bb ccc" >
- < 2, "aaa bb" >

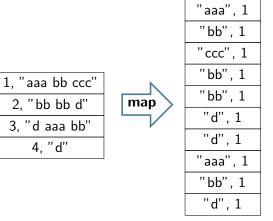
Output

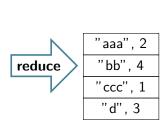
- "aaa", 2 >
- < "bb", 2 >
- < "ccc", 1 >

Word Count

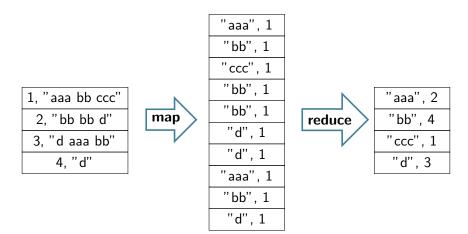
```
map(key, value): /* pairs of {line num, content} */
  foreach word in value.split():
    emit(word, 1)
reduce(key, values): /* {word, list nb occurences} */
  result = 0
  for value in values:
    result += value
  emit(key, result) /* -> {word, nb occurences} */
```

Word Count





Word Count



Question:

How is it implemented in a distributed environment? (stay tuned)

Example: Web index

Description

Construct an index of the pages in which a word appears.

- Input: A set of web pages
 - ► Pairs < URL, content of the page >
- Output: A set of pairs < word, set of URLs >

Example: Web index

```
map(key, value): /* pairs of {URL, page_content} */
  foreach word in value.parse():
    emit(word, key)
reduce(key, values): /* {word, URLs} */
 list=[]
 for value in values:
    list.add(value)
  emit(key, list) /* {word, list of URLs} */
```

About batch and stream processing

Batch processing

- A batch processing system takes a large amount of input data, runs a job to process it, and produces some output data.
- Offline system
 - All inputs are already available when the computation starts
- In this lecture, we are discussing batch processing.

Stream processing

- A stream processing system processes data shortly after they have been received
- Near real-time system
- The amount of data to process is unbounded
 - Data arrives gradually over time

Agenda

Introduction to MapReduce

The Hadoop Eco-System

HDFS

Hadoop MapReduce

Apache Hadoop



History

Open source implementation of a MapReduce framework

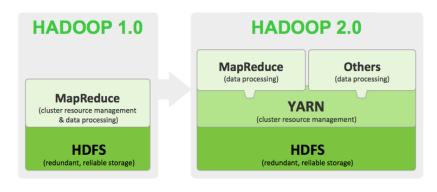
- Implemented by people working at Yahoo!
- Inspired from the publications of Google
- Released in 2006

Evolution

- A full ecosystem
- Used by many companies
 - Facebook Big Data stack is still inspired by (and even making use of) Hadoop¹

¹https://www.datanami.com/2020/08/31/
how-facebook-accelerates-sql-at-extreme-scale/

Hadoop evolution



The Hadoop ecosystem

The main blocks

- HDFS: The distributed file system
- Yarn: The cluster resource manager
- MapReduce: The processing engine

The Hadoop ecosystem

The main blocks

- HDFS: The distributed file system
- Yarn: The cluster resource manager
- MapReduce: The processing engine

Other blocks

- Hive: Provide SQL-like query language
- Pig: High-level language to create MapReduce applications
 - Notion of Pipeline
- Giraph: Graph processing
- etc.

A few words about Yarn

A resource management framework

- Dynamically allocates the resources of a cluster to jobs
- Allows multiple engines to run in parallel on the cluster
 - Not all jobs have to be MapReduce jobs
 - Increases resource usage
- Main components of the system
 - ResourceManager: Allocates resources to applications and monitors the available nodes
 - ApplicationMaster: Negotiates resources access for one application with the RM; Coordinates the application's tasks execution
 - ► The NodeManager: Launches tasks on nodes and monitors resource usage
- Has been replaced by other frameworks (Mesos, Kubernetes)

Agenda

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Hadoop MapReduce

Hadoop Distributed File System

Purpose

Store and provide access to large datasets in a share-nothing infrastructure

Challenges

- Scalability
- Fault tolerance

¹http:

^{//}yahoohadoop.tumblr.com/post/138739227316/hadoop-turns-10

Hadoop Distributed File System

Purpose

Store and provide access to large datasets in a share-nothing infrastructure

Challenges

- Scalability
- Fault tolerance

Example of large scale deployment

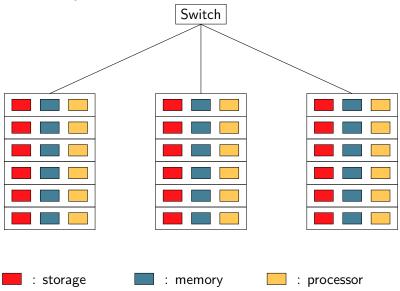
At Yahoo!: 600PB of data on 35K servers¹

¹http:

^{//}yahoohadoop.tumblr.com/post/138739227316/hadoop-turns-10

Target infrastructure (recall)

Cluster of commodity machines



Main principles

Achieving scalability and fault tolerance

Main assumptions

- Storing large datasets
 - Provide large aggregated bandwidth
 - ► Allow storing large amount of files (millions)
- Batch processing (i.e., simple access patterns)
 - ► The file system is not POSIX-compliant
 - Assumes sequential read and writes (no random accesses)
 - Write-once-read-many file accesses
 - Supported write operations: Append and Truncate
 - Stream reading
- Optimized for throughput (not latency)

Random vs Sequential disk access

- Example
 - DB 100M users
 - 100B/user
 - Alter 1% records
- Random access
 - Seek, read, write: 30mS
 - 1M users → 8h20
- Sequential access
 - Read ALL Write ALL
 - 2x 10GB @ 100MB/S → 3 minutes
- → It is often faster to read all and write all sequentially

Main principles

Achieving scalability and fault tolerance

Main principles

Achieving scalability and fault tolerance

Partitioning

- Files are partitioned into blocks
- Blocks are distributed over the nodes of the system
- Default block size in recent versions: 128MB

Replication

- Multiple replicas of each block are created
- Replication is topology aware (rack awareness)
- Default replication degree is 3

A Master-Slave architecture

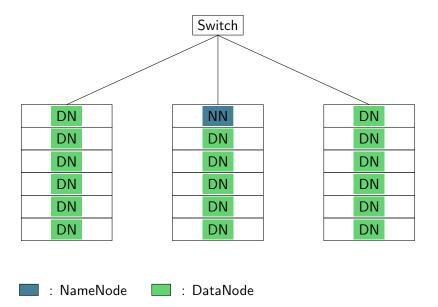
A set of DataNodes

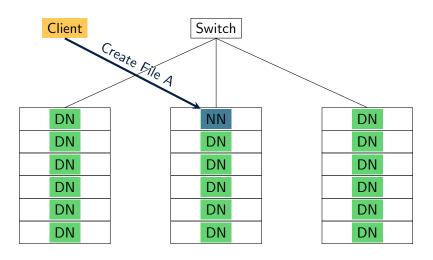
- One *daemon* per node in the system
- A network service allowing to access the file blocks stored on that node
 - ▶ It is responsible for serving read and write requests

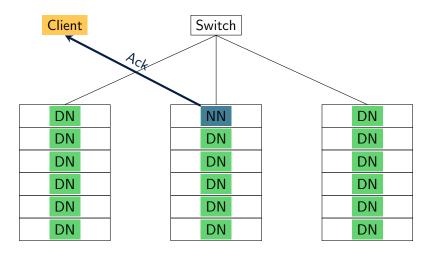
One NameNode

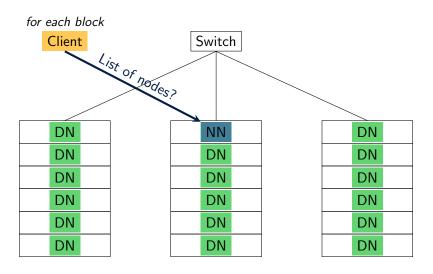
- Keeps track of where blocks are stored
- Monitors the DataNodes
- Entry point for clients

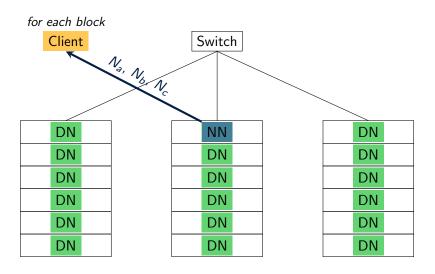
HDFS architecture

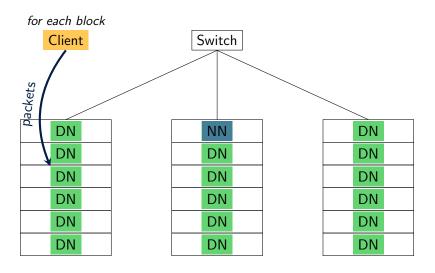


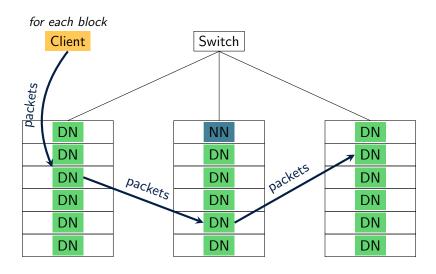


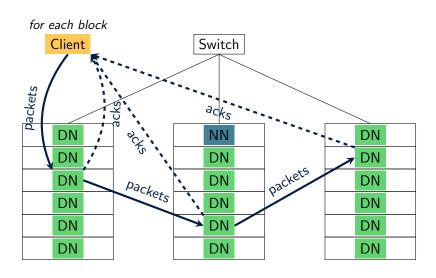


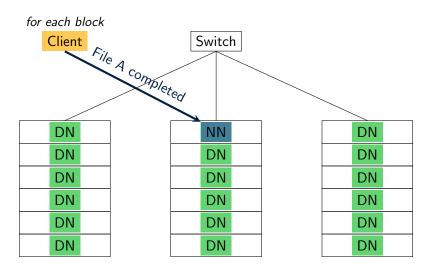






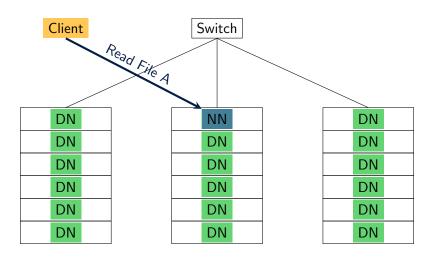


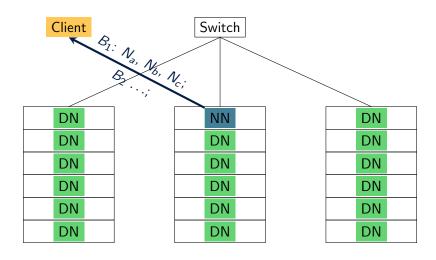


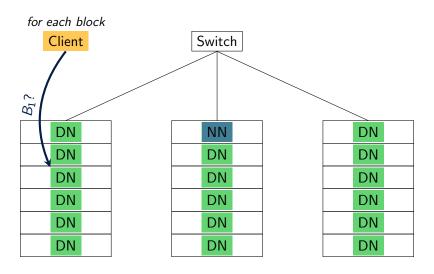


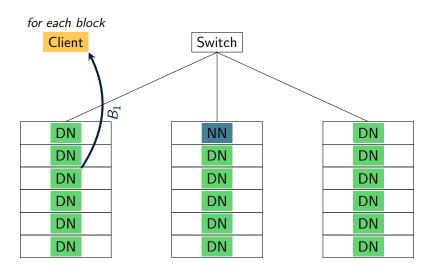
Writing a file: Summary

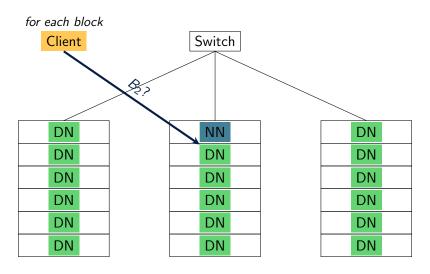
- The client contacts the NameNode to request new file creation
 - The NameNode makes all required checks (Permissions, file does not exists, etc.)
- 2. The NameNode allows the client to write the file
- 3. The client splits the data to be written into blocks
 - For each block, it asks the NameNode for a list of destination nodes
 - The returned list is sorted in increasing distance from the client
- 4. Each block is written in a pipeline
 - The client picks the closest node to write the block
 - ► The DataNode receives the packets (*portions*) and forwards them to the next DataNode in the list
- Once all blocks have been created with a sufficient replication degree, the client acknowledges file creation completion to the name node.
- The NameNode flushes information about the file to disk

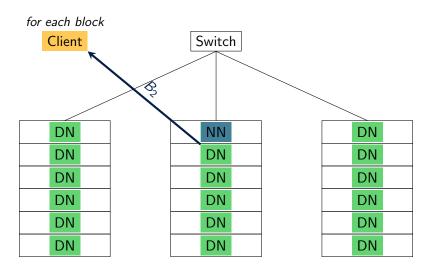












Reading a file: Summary

- 1. The client contacts the NameNode to have info about a file
- 2. The NameNode returns the list of all blocks
 - For each block, it provides a list of nodes hosting the block
 - ► The list is sorted according to the distance from the client
- 3. The client can start reading the blocks sequentially in order
 - By default, contacts the closest DataNode
 - If the node is down, contacts the next one in the list

Supported file formats

- Text/CSV files
- JSON records
- Sequence files (binary key-value pairs)
 - Can be used to store photos, videos, etc
- Defining custom formats
 - Avro
 - Parquet
 - ▶ ORC

Agenda

Introduction to MapReduce

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Hadoop MapReduce

In a Nutshell

A distributed MapReduce framework

- Map and Reduce tasks are distributed over the nodes of the system
- Runs on top of HDFS
 - Move the computation instead of the data
- Fault tolerant

2 main primitives

- Map (transformation)
- Reduce (aggregation)

In a nutshell

Key/Value pairs

- MapReduce manipulate sets of Key/Value pairs
- Keys and values can be of any types

Functions to apply

- The user defines the functions to apply
- In Map, the function is applied independently to each pair
- In Reduce, the function is applied to all values with the same key

MapReduce operations

About the Map operation

- A given input pair may map to zero, one, or many output pairs
- Output pairs need not be of the same type as input pairs

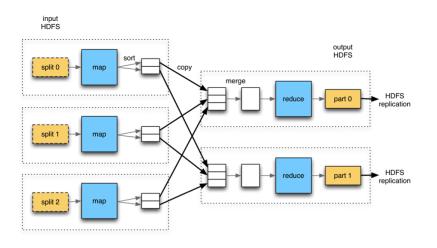
About the Reduce operation

- Applies operation to all pairs with the same key
- 3 steps:
 - Shuffle and Sort: Groups and merges the output of mappers by key
 - Reduce: Applies the reduce operation to the new key/value pairs

Distributed execution

Figure from

https://www.supinfo.com/articles/single/2807-introduction-to-the-mapreduce-life-cycle



Distributed execution: the details

Map tasks

- As many as the number of blocks to process
- Executed on a node hosting a block (when possible)
- Data read from HDFS

Reduce tasks

- Number selected by the programmer
- Key-value pairs are distributed over the reducers using a hash of the key
- The output is stored in HDFS

Data management

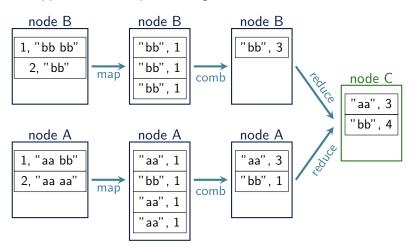
Moving data from the Map to the Reduce tasks

- 1. Output of map tasks are partitioned. The result is stored locally
 - As many partitions are created as the number of reducers
 - By default, a partitioning function based on the hash of the key is used
 - ► The user can specify its own partitioning function
- 2. The reducers fetch the data from the map tasks
 - They connect to the map nodes to fetch data (shuffle)
 - ► This can start as soon as some map tasks finish (customizable)
- 3. The reducers sort the data by key (sort)
 - Can start only when all map tasks are finished

Reducing the amount of data transferred

Combiner

- User-defined function for local aggregation on the map tasks
- Applied after the partitioning function



About more complex programs Workflows

Sequence of Map and Reduce operations

- The output of one job is the input of the next job
- Example: Getting the word that occurs to most often in a text
 - ▶ Job 1: counting the number of occurrence of each word
 - ▶ Job 2: Find the word with the highest count

Implementation

- No specific support in Hadoop
- Data simply go through HDFS

Additional references

Mandatory reading

 MapReduce: Simplified Data Processing on Large Clusters, by J. Dean and S. Ghemawat.

Suggested reading

- Chapter 10 of Designing Data-Intensive Applications by Martin Kleppmann
- HDFS Carton: https://wiki.scc.kit.edu/ gridkaschool/upload/1/18/Hdfs-cartoon.pdf
- MapReduce illustration: https: //words.sdsc.edu/words-data-science/mapreduce