

Weakly-affine monads

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Abstract

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1 Introduction

For context:

► **Proposition 1.1.** *A monoid (M, m, e) is a group if and only if the associativity square*

$$\begin{array}{ccc} M \times M \times M & \xrightarrow{m \times \text{id}} & M \times M \\ \downarrow \text{id} \times m & & \downarrow m \\ M \times M & \xrightarrow{m} & M \end{array} \quad (1)$$

is a pullback.

The same can be said more generally for a monoid object in a cartesian monoidal category.

Proof. The square (1) is a pullback, both of sets and of groups, if and only if given $a, g, h, c \in M$ such that $ag = hc$, there exists a unique $b \in M$ such that $g = bc$ and $h = ab$. First, suppose that g is a group. The only possible choice of b is

$$b = a^{-1}h = gc^{-1},$$

which is unique by uniqueness of inverses.

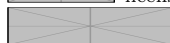
Conversely, suppose that (1) is a pullback. We can set $g, h = e$ and $c = a$ so that $ae = ea = a$. Instantiating the pullback property, there is a unique b such that $ab = e$ and $ba = e$, that is, $b = a^{-1}$. ◀

► **Definition 1.1.** *Let X be a set. Denote by MX the set of finitely supported measures on X , i.e. functions $m : X \rightarrow [0, \infty)$ which are zero for all but a finite number of $x \in X$. Given a function $f : X \rightarrow Y$, denote by $Mf : MX \rightarrow MY$ the function sending $m \in MX$ to the assignment*

$$f_*m : y \mapsto \sum_{x \in f^{-1}(y)} p(x).$$



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34 This way M is a functor, and even a monad with the following unit and multiplication maps.

$$35 \quad \begin{array}{ccc} X & \xrightarrow{\delta} & MX \\ x & \longmapsto & \delta_x, \end{array} \quad \begin{array}{ccc} MMX & \xrightarrow{E} & MX \\ \xi & \longmapsto & E\xi, \end{array}$$

36 where

$$37 \quad \delta_x(x') = \begin{cases} 1 & x = x'; \\ 0 & x \neq x', \end{cases} \quad E\xi(x) = \sum_{m \in MX} \xi(m) m(x).$$

38 Call M the measure monad on **Set**.

39 Denote also by $DX \subseteq MX$ the subset of probability measures, i.e. those finitely supported
40 $p : X \rightarrow [0, \infty)$ such that

$$41 \quad \sum_{x \in X} p(x) = 1.$$

42 D forms a submonad of M called the distribution monad.

43 1.1 GS and Markov categories

44 The notion of *gs-monoidal category* has been originally introduced in the context of algebraic
45 approaches to term graph rewriting [?], and then developed in a series of papers [?, ?, ?].
46 We recall here the basic definitions adopting the graphical formalism of string diagrams,
47 referring to [?] for an overview of various notions of monoidal categories and their associated
48 diagrammatic calculus.

49 ► **Definition 1.2.** A *gs-monoidal category* is a symmetric monoidal category $(\mathcal{C}, \otimes, I)$
50 with a commutative comonoid structure on each object X , consisting of a comultiplication
51 and counit,

$$52 \quad \text{copy}_X = \begin{array}{c} \text{---} \cup \text{---} \\ \bullet \\ | \\ X \end{array} \quad \text{del}_X = \begin{array}{c} \bullet \\ | \\ X \end{array}$$

53 which satisfy the commutative comonoid equations:

$$54 \quad \begin{array}{c} \text{---} \cup \text{---} \\ \bullet \\ | \\ X \end{array} = \begin{array}{c} \text{---} \cup \text{---} \\ \bullet \\ | \\ X \end{array} \quad \begin{array}{c} \bullet \\ \text{---} \cup \text{---} \\ \bullet \\ | \\ X \end{array} = \begin{array}{c} | \\ X \end{array} \quad \begin{array}{c} \text{---} \cup \text{---} \\ \bullet \\ | \\ X \end{array} = \begin{array}{c} \text{---} \cup \text{---} \\ \bullet \\ | \\ X \end{array}$$

55 These comonoid structures must be multiplicative with respect to the monoidal structure:

$$56 \quad \begin{array}{c} \text{---} \cup \text{---} \\ \bullet \\ | \\ X \otimes Y \end{array} = \begin{array}{c} \text{---} \cup \text{---} \\ \bullet \quad \bullet \\ | \quad | \\ X \quad Y \end{array} \quad \begin{array}{c} \bullet \\ | \\ X \otimes Y \end{array} = \begin{array}{c} \bullet \\ | \\ X \end{array} \quad \begin{array}{c} \bullet \\ | \\ Y \end{array}$$

$$\begin{array}{c} \text{---} \cup \text{---} \\ \bullet \\ | \\ I \end{array} = \begin{array}{c} \text{---} \cup \text{---} \\ \bullet \\ | \\ I \end{array} \quad \begin{array}{c} \bullet \\ | \\ I \end{array} = \begin{array}{c} \text{---} \cup \text{---} \\ \bullet \\ | \\ I \end{array}$$

57 ► **Definition 1.3.** A morphism $f : X \rightarrow Y$ in a gs-monoidal category is called *copyable* or
 58 *functional* if and only if

$$59 \quad \begin{array}{c} Y \quad Y \\ \diagdown \quad \diagup \\ \bullet \\ | \\ \boxed{f} \\ | \\ X \end{array} = \begin{array}{c} Y \quad Y \\ \diagup \quad \diagdown \\ \bullet \\ \diagup \quad \diagdown \\ \boxed{f} \quad \boxed{f} \\ | \\ X \end{array}$$

60 It is called *discardable* or *full* if

$$61 \quad \begin{array}{c} \bullet \\ | \\ \boxed{f} \\ | \\ X \end{array} = \begin{array}{c} \bullet \\ | \\ X \end{array}$$

62 ► **Example 1.4.** The category **Rel** of sets and relations with the monoidal operation
 63 $\otimes : \mathbf{Rel} \times \mathbf{Rel} \rightarrow \mathbf{Rel}$ given by the direct product of sets is a gs-monoidal
 64 category [?]. In this gs-monoidal category, copyable arrows are precisely partial functions, and
 65 discardable arrows are total relations.

66 ► **Remark 1.5.** It is well-known that if duplicators and dischargers of a given gs-monoidal
 67 category \mathcal{C} satisfy naturality, then the monoidal product is the categorical product, and
 68 thus the category is cartesian monoidal [?], i.e. the following conditions are equivalent for a
 69 gs-monoidal category \mathcal{C} :

- 70 ■ \mathcal{C} is cartesian monoidal;
- 71 ■ every morphism is copyable and discardable;
- 72 ■ the copy and discard maps are natural.

73 In recent works [?] it has been proved that the gs-monoidal structure naturally arises in
 74 several situations, such as when Kleisli categories of commutative monads or span categories.
 75 In the following proposition, we recall the result regarding Kleisli categories:

76 ► **Proposition 1.6.** Let T be a monoidal (equivalently, commutative) monad on a cartesian
 77 monoidal category \mathcal{D} . Then $\mathbf{Kl}(T)$ is canonically a gs-monoidal category with copy and
 78 discard structure induced by those of \mathcal{D} .

79 ► **Definition 1.7.** A gs-monoidal category is called *Markov* if any (hence all) of the following
 80 equivalent conditions are satisfied:

- 81 ■ The monoidal unit is terminal;
- 82 ■ The discard map is natural;
- 83 ■ Every morphism is discardable.

84 ► **Definition 1.8.** A monad T on a cartesian monoidal category is called *affine* if $T1 \cong 1$.

85 ► **Proposition 1.9.** In the hypotheses of Proposition 1.6, we have that $\mathbf{Kl}(T)$ is Markov if
 86 and only if T is affine.

87 In this work we want to introduce an intermediate level between GS and Markov, which
 88 we call *weakly Markov*, and its corresponding notion for monads, which we call *weakly affine*.

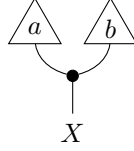
2 Weakly Markov categories and weakly affine monads

2.1 The monoid of effects

In a gs-monoidal category \mathcal{C} we call a *state* a morphism from the monoidal unit $p : I \rightarrow X$, and *effect* or *co-state* a morphism to the monoidal unit $a : X \rightarrow I$. We represent them as triangles as follows.



Effects, i.e. elements of the set $\mathcal{C}(X, I)$, form canonically a commutative monoid as follows: the monoidal unit is the discard map $X \rightarrow I$, and given $a, b : X \rightarrow I$, their product ab is given by copying, as follows.



If a morphism $f : X \rightarrow Y$ is copyable and discardable, precomposition with f induces a morphism of monoids $\mathcal{C}(Y, I) \rightarrow \mathcal{C}(X, I)$.

Let's now consider the case where the GS structure comes from a commutative monad on a cartesian monoidal category \mathcal{D} . In this case, the monoid structure of Kleisli morphisms $X \rightarrow 1$ comes from the following canonical internal monoid structure of $T1$ in \mathcal{D} , given by

$$1 \xrightarrow{\eta} T1, \quad T1 \times T1 \xrightarrow{c} T(1 \times 1) \xrightarrow{\cong} T1.$$

For example, for the monad of measures M , we have that $M1 = [0, \infty)$ with its usual multiplication.

The monoid structure of Kleisli morphisms $X \rightarrow 1$ is now given as follows. The unit is given by

$$X \xrightarrow{!} 1 \xrightarrow{\eta} T1,$$

and the multiplication of the morphisms $f^\sharp, g^\sharp : X \rightarrow T1$ is

$$X \xrightarrow{\text{copy}} X \times X \xrightarrow{f^\sharp \times g^\sharp} T1 \times T1 \xrightarrow{c} T(1 \times 1) \xrightarrow{\cong} T1.$$

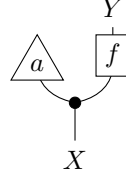
For the monad of measures M , Kleisli morphisms $X \rightarrow 1$ are functions $X \rightarrow [0, \infty)$, and their monoidal structure is their pointwise product.

2.2 Main definitions

► **Definition 2.1.** A GS-category \mathcal{C} is called *weakly Markov* if for every object X , the monoid $\mathcal{C}(X, I)$ is a group.

Every Markov category is weakly Markov: for each X , the monoid $\mathcal{C}(X, I)$ is the trivial group.

Note that the group $\mathcal{C}(X, I)$ (or monoid, in the general case) acts on the set $\mathcal{C}(X, Y)$: given $a : X \rightarrow I$ and $f : X \rightarrow Y$, $a \cdot f$ is given as follows,



and the product $(f, g) \mapsto f \cdot g := (f \otimes g) \circ \text{copy}_X$ is equivariant for this action in both variables (separately).

► **Definition 2.2.** Given two parallel morphisms $f, g : X \rightarrow Y$ in a weakly Markov GS-category \mathcal{C} , we say that f and g are *equivalent*, and write $f \sim g$, if they lie in the same orbit for the action of $\mathcal{C}(X, I)$, i.e. if there is $a \in \mathcal{C}(X, I)$ such that $a \cdot f = g$. We say they are *uniquely equivalent* if there is a unique $a \in \mathcal{C}(X, I)$ such that $a \cdot f = g$.

Let's now look at the Kleisli case.

► **Definition 2.3.** A commutative monad T on a cartesian monoidal category is called *weakly affine* if $T1$ with its canonical internal monoid structure is a group.

This choice of terminology is motivated by the following proposition, which can be seen as a weak version of Proposition 1.9.

► **Proposition 2.4.** Let \mathcal{D} be a cartesian monoidal category, and let T be a commutative monad on \mathcal{D} . The Kleisli category of T is weakly Markov if and only if T is weakly affine.

Proof. First, suppose that $T1$ is an internal group, and denote by $\iota : T1 \rightarrow T1$ its inversion map. The inverse of the morphism $f^\sharp : X \rightarrow T1$ in $\text{Kl}_T(X, 1)$ is given by $\iota \circ f$: indeed, the following diagram commutes,

$$\begin{array}{ccccc}
 X & \xrightarrow{\text{copy}} & X \times X & & \\
 f^\sharp \downarrow & & f^\sharp \times f^\sharp \downarrow & \searrow f \times (\iota \circ f) & \\
 T1 & \xrightarrow{\text{copy}} & T1 \times T1 & \xrightarrow{\text{id} \times \iota} & T1 \times T1 \xrightarrow{c} T(1 \times 1) \\
 \downarrow ! & & & & \downarrow \cong \\
 1 & \xrightarrow{\eta} & & & T1
 \end{array}$$

where the bottom rectangle commutes since ι is the inversion map for $T1$. The analogous diagram with $\iota \times \text{id}$ in place of $\text{id} \times \iota$ commutes analogously.

Conversely, suppose that for every X , the monoid structure on $\text{Kl}_T(X, 1)$ has inverses. Then in particular we can take $X = T1$, and the inverse of the Kleisli morphism $\text{id} : T1 \rightarrow T1$ is an inversion map for $T1$. ◀

This result can be seen in terms of the Yoneda embedding, see the details in Appendix A.

2.3 Examples of weakly affine monads

Every affine monad is a weakly affine monad. Here are less trivial examples.

► **Example 2.5.** Let $M^* : \mathbf{Set} \rightarrow \mathbf{Set}$ be the monad assigning to every set the set of finitely supported discrete *nonzero* measures on M^* , or equivalently let $M^*(X)$ for any set X be the set of nonzero finitely supported functions $X \rightarrow [0, \infty)$. The monad structure is defined in terms of the same formulas as for the monad of measures M (Definition 1.1) and the components $c_{X,Y}$ are also given by the formation of product measures, or equivalently point-wise products of functions $X \rightarrow [0, \infty)$.

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Since $M^*1 \cong (0, \infty) \not\cong 1$, this monad is not affine. However the monoid structure of $(0, \infty)$ induced by M^* is the usual multiplication of positive real numbers, which form a group. Therefore M^* is weakly affine, and its Kleisli category is weakly Markov.

On the other hand, if the zero measure is included, we have $M1 \cong [0, \infty)$ which is not a group under multiplication, so M is not weakly affine.

► **Example 2.6.** Let A be a commutative monoid. Then the functor $T_A := A \times -$ on **Set** has a canonical structure of commutative monad, where the lax structure components $c_{X,Y}$ are given by multiplying elements in A while carrying the elements of X and Y along.

Since $T_A(1) \cong A$, the monad T_A is weakly affine if and only if A is a group, and affine if and only if $A \cong 1$.

► **Example 2.7.** Here is a negative example. Consider the free abelian group monad F on **Set**. Its functor takes a set X and forms the set FX of finite multisets (with repetition, where order does not matter) of elements of X and their formal inverses. We have that $F1 \cong \mathbb{Z}$, which is an abelian group under addition. However, the monoid structure on $F1$ induced by the monoidal structure of the monad corresponds to the *multiplication* in \mathbb{Z} , which does not have inverses. Therefore F is not weakly affine.

3 Conditional independence in weakly Markov categories

Markov categories have a rich theory of conditional dependence and independence [?]. Some of those ideas can be translated and generalized to the setting of weakly Markov categories.

► **Definition 3.1.** A morphism $f : A \rightarrow X_1 \otimes \cdots \otimes X_n$ in a GS-category \mathcal{C} is said to exhibit *conditional independence of the X_i given A* if and only if it can be expressed as a product of the following form.



Note that this is slightly different from [?, Definition 6.6], although it is equivalent for the case of Markov categories.

Here is what conditional independence looks like in the Kleisli case.

► **Proposition 3.2.** Let \mathcal{D} be a cartesian monoidal category, and let T be a commutative monad on \mathcal{D} . A Kleisli morphism $f^\sharp : A \rightarrow T(X_1 \times \cdots \times X_n)$ exhibits conditional independence of the X_i given A if and only if it factors as follows

$$\begin{array}{ccc}
 A & & \\
 \downarrow (g_1^\sharp, \dots, g_n^\sharp) & \searrow f^\sharp & \\
 TX_1 \times \cdots \times TX_n & \xrightarrow{c} & T(X_1 \times \cdots \times X_n),
 \end{array}$$

for some Kleisli maps $g_i^\sharp : A \rightarrow TX_i$, where the map c above is the one obtained by iterating the multiplication of the monoidal structure (such a map is unique by associativity).

Proof. In terms of the base category \mathcal{D} , a Kleisli morphism in the form of Definition 3.1 reads as follows.

$$A \xrightarrow{\text{copy}} A \times \cdots \times A \xrightarrow{g_1^\sharp \times \cdots \times g_n^\sharp} TX_1 \times \cdots \times TX_n \xrightarrow{c} T(X_1 \times \cdots \times X_n).$$

188 Therefore $f^\# : A \rightarrow T(X_1 \times \cdots \times X_n)$ is exhibiting conditional independence if and only if it
 189 is in the form above. ◀

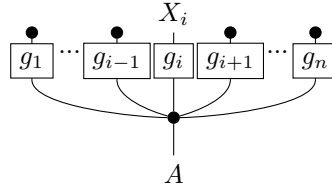
190 ▶ **Example 3.3.** In the Kleisli category of the distribution monad D , which is Markov, a
 191 morphism $f : A \rightarrow X \otimes Y$ exhibits conditional independence if and only if it is the product
 192 of its marginals [?, Section 12].

193 ▶ **Example 3.4.** In the Kleisli category of the measure monad M , the zero measure always
 194 displays conditional independence of its outputs given its inputs: for example, for $A = 1$, the
 195 zero measure on $X \times Y$ is the product of the zero measure on X and the zero (or any other)
 196 measure on Y . Notice that both marginals of the zero measure are zero measures—therefore,
 197 the factors appearing in the product are not necessarily related to the marginals.

198 In a weakly Markov category, the situation is similar to the Markov case, but up to
 199 equivalence.

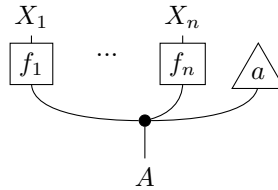
200 ▶ **Proposition 3.5.** Let $f : A \rightarrow X_1 \otimes \cdots \otimes X_n$ be a morphism in a weakly Markov GS-category
 201 \mathcal{C} . Then f exhibits conditional independence of the X_i given A if and only if it is equivalent
 202 to the product of all its marginals. Moreover, in that case f is uniquely equivalent to the
 203 product of its marginals.

204 **Proof.** Denote the marginals of f by f_1, \dots, f_n . Suppose that f is a product as in Defini-
 205 tion 3.1. For each $i = 1, \dots, n$, by marginalizing, we get that f_i is equal to the following.

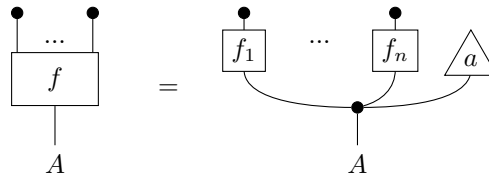


207 Therefore for each i we have that $f_i \sim g_i$.

208 Conversely, suppose that f is equivalent to the product of its marginals, i.e. that there
 209 exists $a : X \rightarrow I$ such that f is equal to the following.

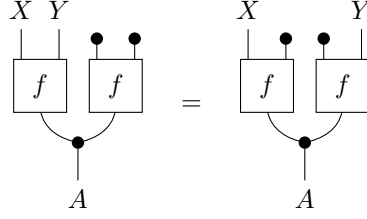


211 One can then choose $g_i = f_i$ for all $i < n$, and $g_n = a \cdot f_n$, so that f is in the form of
 212 Definition 3.1. Moreover, by marginalizing over all the X_i at once, we see that



214 so that a is uniquely determined. ◀

215 ▶ **Remark 3.6.** For $n = 2$, a morphism $f : A \rightarrow X \otimes Y$ in a weakly Markov GS-category \mathcal{C}
 216 exhibits conditional independence of X and Y given A if and only if the following equation
 217 holds.



218

219 3.1 Main result

220 The concept of conditional independence for general weakly Markov categories allow us to
 221 give an equivalent characterization of affine monads. The condition is in terms of a pullback
 222 condition on the associativity diagram, and can be seen as a generalization of Proposition 1.1.

223 ► **Theorem 3.7.** *Let \mathcal{D} be a cartesian monoidal category, and let T be a commutative monad*
 224 *on \mathcal{D} . The following conditions are equivalent.*

- 225 1. T is weakly affine;
- 226 2. $\text{Kl}(T)$ is weakly Markov;
- 227 3. For all objects X, Y , and Z , the following associativity diagram is a pullback.

$$\begin{array}{ccc}
 T(X) \times T(Y) \times T(Z) & \xrightarrow{\text{id} \times c_{Y,Z}} & T(X) \times T(Y \times Z) \\
 c_{X,Y} \times \text{id} \downarrow & & \downarrow c_{X,Y \times Z} \\
 T(X \times Y) \times T(Z) & \xrightarrow{c_{X \times Y, Z}} & T(X \times Y \times Z)
 \end{array} \quad (2)$$

229 We prove the theorem by means of the following property of weakly Markov categories.

230 ► **Lemma 3.8** (localized independence property). *Let \mathcal{C} be a weakly Markov GS-category.*
 231 *Whenever a morphism $f : A \rightarrow X \otimes Y \otimes Z$ exhibits conditional independence of $X \otimes Y$*
 232 *(jointly) and Z given A , as well as conditional independence of X and $Y \otimes Z$ given A , then*
 233 *it exhibits conditional independence of X, Y and Z given A .*

234 **Proof of Lemma 3.8.** Suppose $f : A \rightarrow X \otimes Y \otimes Z$ exhibits conditional independence of
 235 $X \otimes Y$ (jointly) and Z given A , as well as conditional independence of X and $Y \otimes Z$ given
 236 A . By marginalizing out X , we have that f_{YZ} exhibits conditional independence of Y and Z
 237 given A . Since by hypothesis f exhibits conditional independence of X and $Y \otimes Z$ given A ,
 238 by Proposition 3.5 we have that f is equivalent to the product of f_X and f_{YZ} . But, again
 239 by Proposition 3.5, f_{YZ} is equivalent to the product of f_Y and f_Z , so we have that f is
 240 equivalent to the product of all its marginals. Using Proposition 3.5 in the other direction,
 241 this means that f exhibits conditional independence of X, Y and Z given A . ◀

242 We are now ready to prove the theorem.

243 **Proof of Theorem 3.7.** $1 \Leftrightarrow 2$: see Proposition 2.4.

244 $1 \Rightarrow 3$: By the universal property of products, a cone over the cospan in (2) consists of
 245 maps $g_1^\# : A \rightarrow TX$, $g_{23}^\# : A \rightarrow T(Y \times Z)$, $g_{12}^\# : A \rightarrow T(X \times Y)$ and $g_3^\# : A \rightarrow TZ$ such that
 246 the following diagram commutes.

$$\begin{array}{ccc}
 A & \xrightarrow{(g_1^\#, g_{23}^\#)} & T(X) \times T(Y \times Z) \\
 (g_{12}^\#, g_3^\#) \searrow & & \downarrow c_{X,Y} \times \text{id} \quad \downarrow c_{X,Y \times Z} \\
 T(X \times Y) \times T(Z) & \xrightarrow{c_{X \times Y, Z}} & T(X \times Y \times Z)
 \end{array}$$

247

By Proposition 3.2, this amounts to a Kleisli map $f^\# : A \rightarrow T(X \times Y \times Z)$ exhibiting conditional independence of X and $Y \otimes Z$ given A , as well as of $X \otimes Y$ and Z given A . By the localized independence property (Lemma 3.8), we then have that f exhibits conditional independence of all X , Y and Z given A , and so, again by Proposition 3.2, $f^\#$ factors through the product $TX \times TY \times TZ$. More specifically, by marginalizing over Z , we have that $g_{12}^\#$ factors through $TX \times TY$, i.e. the following diagram on the left commutes for some $h_1^\# : A \rightarrow TX$ and $h_2^\# : A \rightarrow TY$, and similarly, by marginalizing over X , the diagram on the right commutes for some $\ell_2^\# : A \rightarrow TY$ and $\ell_3^\# : A \rightarrow TZ$.

$$\begin{array}{ccc} A & \xrightarrow{g_{12}^\#} & TX \times TY \\ (h_1^\#, h_2^\#) \downarrow & \searrow & \downarrow c \\ TX \times TY & \xrightarrow{c} & T(X \times Y) \end{array} \quad \begin{array}{ccc} A & \xrightarrow{g_{23}^\#} & TY \times TZ \\ (\ell_2^\#, \ell_3^\#) \downarrow & \searrow & \downarrow c \\ TY \times TZ & \xrightarrow{c} & T(Y \times Z) \end{array}$$

In other words, the upper and the left curved triangles in the following diagram commute.

$$\begin{array}{ccc} A & \xrightarrow{(g_1^\#, g_{23}^\#)} & T(X) \times T(Y \times Z) \\ \downarrow (g_{12}^\#, g_3^\#) & \searrow (g_1^\#, \ell_2^\#, \ell_3^\#) & \downarrow c_{X,Y \times Z} \\ T(X) \times T(Y) \times T(Z) & \xrightarrow{\text{id} \times c_{Y,Z}} & T(X) \times T(Y \times Z) \\ \downarrow c_{X,Y} \times \text{id} & & \downarrow c_{X,Y \times Z} \\ T(X \times Y) \times T(Z) & \xrightarrow{c_{X \times Y, Z}} & T(X \times Y \times Z) \end{array}$$

By marginalizing over Y and Z , and by weak affinity of T , there exists a unique $a^\# : A \rightarrow T1$ such that $h_1 = a \cdot g_1$. Therefore

$$g_{12} = h_1 \cdot h_2 = (a \cdot g_1) \cdot h_2 = g_1 \cdot (a \cdot h_2),$$

and so in the diagram above we can equivalently replace h_1 and h_2 with g_1 and $a \cdot h_2$. Similarly by marginalizing over X and Y , there exists a unique $c^\# : A \rightarrow T1$ such that $\ell_3 = c \cdot g_3$, so that

$$g_{23} = \ell_2 \cdot \ell_3 = \ell_2 \cdot (c \cdot g_3) = (c \cdot \ell_2) \cdot g_3$$

and in the diagram above we can replace ℓ_2 and ℓ_3 with $c \cdot \ell_2$ and g_3 , as follows.

$$\begin{array}{ccc} A & \xrightarrow{(g_1^\#, g_{23}^\#)} & T(X) \times T(Y \times Z) \\ \downarrow (g_{12}^\#, g_3^\#) & \searrow (g_1^\#, (c \cdot \ell_2)^\#, g_3^\#) & \downarrow c_{X,Y \times Z} \\ T(X) \times T(Y) \times T(Z) & \xrightarrow{\text{id} \times c_{Y,Z}} & T(X) \times T(Y \times Z) \\ \downarrow c_{X,Y} \times \text{id} & & \downarrow c_{X,Y \times Z} \\ T(X \times Y) \times T(Z) & \xrightarrow{c_{X \times Y, Z}} & T(X \times Y \times Z) \end{array}$$

Now, marginalizing over Z and Z , we see that necessarily $a \cdot h_2 = c \cdot \ell_2$. Therefore there is a unique map $A \rightarrow TX \times TY \times TZ$ making the whole diagram commute, which means that (2) is a pullback.

$3 \Rightarrow 1$: If T is weakly affine, then taking $X = Y = Z = 1$ in (2) shows that this monoid must be an abelian group: we obtain a unique arrow $\iota : T(1) \rightarrow T(1)$ making the following

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273 diagram commute,

$$\begin{array}{ccccc}
 T1 & & \xrightarrow{(id, \eta_1!)} & & T1 \times T1 \\
 & \searrow (id, \iota, id) & & \searrow id \times c_{1,1} & \\
 & & T1 \times T1 \times T1 & \xrightarrow{id \times c_{1,1}} & T1 \times T(1 \times 1) \xrightarrow{\cong} T1 \times T1 \\
 & & \downarrow c_{1,1} \times id & & \downarrow c_{1,1} \times 1 \\
 & & T(1 \times 1) \times T1 & \xrightarrow{c_{1 \times 1, 1}} & T(1 \times 1 \times 1) \xrightarrow{\cong} T(1 \times 1) \\
 & & \downarrow \cong & & \downarrow \cong \\
 & & T1 \times T1 & \xrightarrow{c_{1,1}} & T(1 \times 1) \xrightarrow{\cong} T1
 \end{array}$$

274

275 and the commutativity shows that ι satisfies the equations making it the inversion map for a
 276 group structure. \blacktriangleleft

277 **► Example 3.9.** In the Kleisli category of the measure monad $Kl(M)$ (which is not weakly
 278 affine) consider the following diagram.

$$\begin{array}{ccc}
 MX \times MY \times MZ & \xrightarrow{id \times c_{Y,Z}} & MX \times M(Y \times Z) \\
 c_{X,Y} \times id \downarrow & & \downarrow c_{X,Y \times Z} \\
 M(X \times Y) \times MZ & \xrightarrow{c_{X \times Y, Z}} & M(X \times Y \times Z)
 \end{array}$$

279

280 In the top-right corner $MX \times M(Y \times Z)$ take the pair $(0, p)$ where p is a nonzero measure
 281 on $Y \times Z$, and similarly, in the bottom-left corner take the pair $(q, 0)$ where q is a nonzero
 282 measure on $X \times Y$. Following the diagram, both pairs are mapped to the zero measure in
 283 the bottom-right corner. If the diagram were a pullback, we would be able to express the
 284 top-right and bottom-left corners as coming from the same triplet in $MX \times MY \times MZ$, that
 285 is, there would exist a measure m on Y such that $m \cdot 0 = p$ and $0 \cdot m = q$. Since p and q are
 286 nonzero, this is not possible.

287 4 Further results

288 **► Proposition 4.1.** *Let T be a weakly affine monad. If the diagram*

$$\begin{array}{ccc}
 T(1) & \xrightarrow{id} & T(1) \\
 \iota \downarrow & & \downarrow \eta_{T1} \\
 T(1) & \xrightarrow{T(\eta_1)} & T^2(1)
 \end{array}$$

289

290 *commutes, then $T^2(1) \cong T(1)$ in \mathcal{A} .*

291 **Proof.** To prove the result it is enough to show that $T(1) \cong 1$ in the Kleisli category \mathcal{A}_T .
 292 We know from Lemma that $T(1)$ is a group in \mathcal{A} , where the arrow $\eta_1: 1 \rightarrow T(1)$ is the
 293 unit of the group, and $\iota: T(1) \rightarrow T(1)$ is the inversion map. Therefore, we have that the
 294 composition $\iota\eta_1: 1 \rightarrow T(1)$ has to be equal to η_1 . Therefore, we can consider the arrows
 295 $1 \rightarrow T(1)$ and $T(1) \rightarrow 1$ in the Kleisli category \mathcal{A}_T given by $T(\eta_1)\eta_1$ and ι respectively. The
 296 composition $T(\eta_1)\eta_1$ with ι in \mathcal{A}_T is given by $\mu_{1,1}T(\iota)T(\eta_1)\eta_1$. Employing the naturality
 297 of η_1 and the fact that $\iota\eta_1 = \eta_1$, it is direct to check that $\mu T(\iota)T(\eta_1)\eta_1 = \eta_1$, that is the
 298 identity $1 \rightarrow 1$ in \mathcal{A}_T . Now to show that the other composition gives the identity on $T(1)$ in
 299 \mathcal{A}_T , it is enough to show that $T(\eta_1)\iota = \eta_{T(1)}$, but this follows by hypothesis. \blacktriangleleft

300 ► **Corollary 4.2.** *Let T be a weakly affine monad. If the diagram*

$$\begin{array}{ccc}
 T(1) & \xrightarrow{\text{id}} & T(1) \\
 \downarrow \iota & & \downarrow \eta_{T1} \\
 T(1) & \xrightarrow{T(\eta_1)} & T^2(1)
 \end{array}$$

302 *commutes, then $T(1)$ is an idempotent group, namely $\iota = \text{id}_{T1}$.*

303 **Proof.** By weak affineness, $T(1)$ is a group. If $\eta_{T1} = T(\eta_1)\iota$, then we can apply the
 304 multiplication of the monad to both sides, obtaining $\iota = \text{id}_{T1}$. ◀

305 ► **Remark 4.3.** Bart Jacobs calls a strong monad T on a cartesian monoidal category *strongly*
 306 *affine* [?] if for every pair of objects X and Y , the following diagram is a pullback,

$$\begin{array}{ccc}
 X \times TY & \xrightarrow{s} & T(X \times Y) \\
 \downarrow \pi_1 & & \downarrow T\pi_1 \\
 X & \xrightarrow{\eta} & TX
 \end{array}$$

308 where s denotes the strength and η denotes the unit of the monad. Every strongly affine
 309 monad is affine. The corresponding condition on the (Markov) category $\text{Kl}(T)$ is called
 310 *positivity* [?, Section 2].

311 Note that for a generic commutative monad, the diagram above may even fail to commute
 312 (take for example the measure monad M , and start with $(x, 0)$ in the top left corner). One can
 313 however consider the following diagram, which reduces to the one above (up to isomorphism)
 314 in the affine case,

$$\begin{array}{ccc}
 X \times TY & \xrightarrow{s} & T(X \times Y) \\
 \downarrow \text{id} \times T! & & \downarrow T(\text{id} \times !) \\
 X \times T1 & \xrightarrow{s} & T(X \times 1) \cong TX
 \end{array}$$

316 and which always commutes by naturality of the strength. One can then call the monad
 317 T *positive* if this second diagram is a pullback (and possibly define *positive GS categories*
 318 analogously to positive Markov categories). All the examples of weakly affine monads that
 319 we have are positive in this sense, so one may wonder if every weakly affine monad is positive.
 320 For now, this remains an open question.

321 **A Yoneda embedding interpretation of Proposition 2.4**

322 We can interpret Proposition 2.4 more abstractly in terms of presheaves. Let \mathcal{D} be a
 323 cartesian monoidal category. Consider the presheaf category $[\mathcal{D}^{\text{op}}, \mathbf{Set}]$, equipped with the
 324 Day convolution product,

$$325 \quad F \boxtimes G \cong \int^{A, B \in \mathcal{D}} \mathcal{D}(-, A \times B) \times F(A) \times G(B).$$

326 The Yoneda embedding $\mathcal{D} \rightarrow [\mathcal{D}^{\text{op}}, \mathbf{Set}]$ is strong monoidal: indeed, for each X ,

$$327 \quad 1 \cong \mathcal{D}(X, 1),$$

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since 1 is terminal, and for each X and Y , by Yoneda reduction,

$$\begin{aligned} \mathcal{D}(-, X) \boxtimes \mathcal{D}(-, Y) &\cong \int^{A, B \in \mathcal{D}} \mathcal{D}(-, A \times B) \times \mathcal{D}(-, X) \times \mathcal{D}(-, Y) \\ &\cong \mathcal{D}(-, X \times Y). \end{aligned}$$

Therefore, and by the universal property of products, at the level of individual hom-sets the Day convolution product of representable presheaves just takes the cartesian products of sets:

$$(\mathcal{D}(-, X) \boxtimes \mathcal{D}(-, Y))(A) \cong \mathcal{D}(A, X \times Y) \cong \mathcal{D}(A, X) \times \mathcal{D}(A, Y).$$

Take now an object M of \mathcal{D} . Since the Yoneda embedding is fully faithful and strong monoidal, a monoid structure (M, m, e) on M is equivalently a monoid structure on the representable presheaf $\mathcal{D}(-, M)$. This makes the individual hom-sets monoids, with unit and multiplication as follows for each object X :

$$\begin{aligned} 1 &\xrightarrow{\cong} \mathcal{D}(X, 1) \xrightarrow{e_*} \mathcal{D}(X, M) \\ \mathcal{D}(X, M) \times \mathcal{D}(X, M) &\xrightarrow{\cong} \mathcal{D}(X, M \times M) \xrightarrow{m_*} \mathcal{D}(X, M) \end{aligned}$$

This is precisely the monoid structure that we have defined in Section 2.1 for $M = T1$.

► **Proposition A.1.** *M is an internal group if and only if all the monoids $\mathcal{D}(X, M)$ are groups.*

Proof. By Proposition 1.1, M is a group object if and only if its associativity square (1) is a pullback. Since the hom-functor preserves and reflects all limits in its second argument, we have that (1) is a pullback if and only if for each object X , the following diagram (or equivalently, its bottom right square) is a pullback,

$$\begin{array}{ccccc} \mathcal{D}(X, M) \times \mathcal{D}(X, M) \times \mathcal{D}(X, M) & \xrightarrow{\quad} & \mathcal{D}(X, M) \times \mathcal{D}(X, M) & & \\ \downarrow & \searrow \cong & \downarrow \cong & & \\ & \mathcal{D}(X, M \times M \times M) & \xrightarrow{(m \times \text{id})_*} & \mathcal{D}(X, M \times M) & \\ & \downarrow (\text{id} \times m)_* & & \downarrow m_* & \\ \mathcal{D}(X, M) \times \mathcal{D}(X, M) & \xrightarrow{\cong} & \mathcal{D}(X, M \times M) & \xrightarrow{m_*} & \mathcal{D}(X, M) \end{array}$$

where the unlabelled arrows are the unique ones that make the diagram commute. Again by Proposition 1.1, the diagram above is a pullback if and only if $\mathcal{D}(X, M)$ is a group. ◀