### Planar Graphs Have Bounded Queue-Number

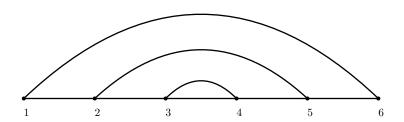
Shengzhe Wang

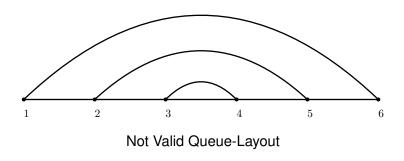
ETH Zürich

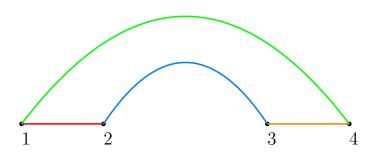
April 6, 2023

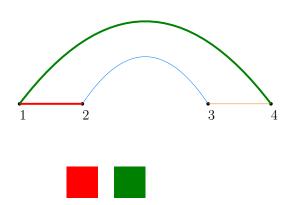


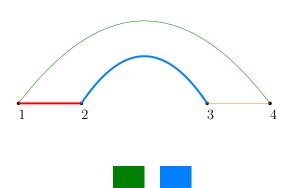


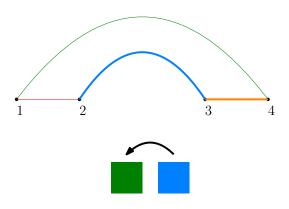






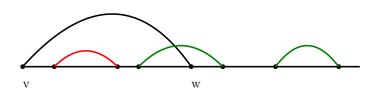






#### **Definition: Queue**

Let G = (V, E), consider a linear odering  $\leq$  of V, a queue of G is a set of edges  $E' \subseteq E$  such that any disjoint edges  $vw, xy \in E'$ , w.l.o.g,  $v \prec w, x \prec y$  and  $v \prec x$ , we have  $w \prec y$ .



### Queue-Number

#### Definition: K-Queue Layout

Let G = (V, E), consider a linear odering  $\leq$  of V, for an integer  $k \geq 0$  a k-queue layout of G is a partition of E into  $E_1, E_2, ..., E_k$  such that each  $E_i$  is a queue of G.

### Queue-Number

#### Definition: K-Queue Layout

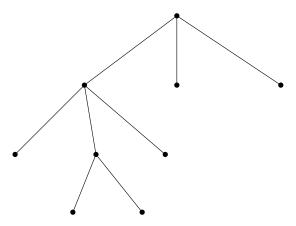
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#### **Definition: Queue-Number**

The queue-number of G, denoted by qn(G), is the minimum integer k such that G has a k-queue layout.

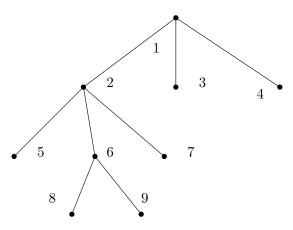
### Queue-Number: Tree

What is the queue-number of tree?

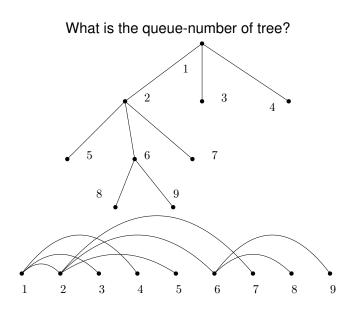


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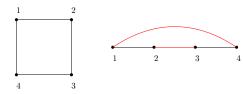
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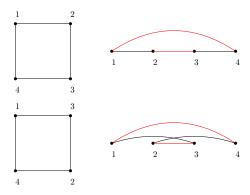
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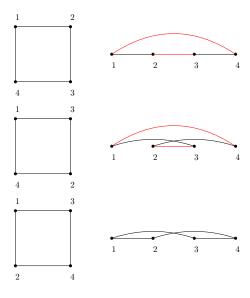
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### Treewidth

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### **Treewidth**

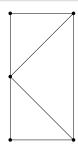
Do we have some tools to help bound the queue-number?

#### Theorem (Veit, 2017)

Every graph with treewidth k has queue-number at most  $2^k - 1$ .

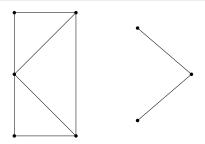
#### Definition: Tree-decomposition

- $\forall \{v, w\} \in E(G)$ , there exists  $x \in V(T)$  with  $v, w \in B_x$
- $\forall v \in V(G)$ , the set  $\{x | x \in V(T) \land v \in B_x\}$  induces a non -empty connected subtree of T.



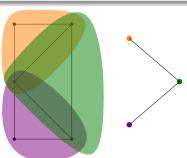
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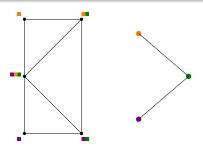
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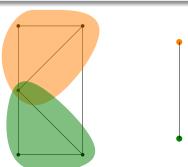
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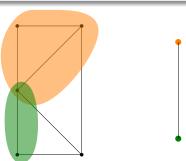
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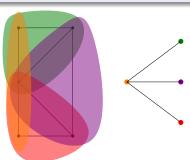
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#### **Treewidth**

#### Definition: Tree-decomposition

A tree-decomposition of a graph G is a pair (B, T). T is a tree and  $B = \{B_x | x \in V(T)\}$  where each  $B_x$  is a subset of V(G) for every vertex x in V(T) such that

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The width of a tree-decomposition of G is  $\max_{x \in V(T)} |B_x| - 1$ 

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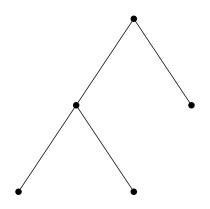
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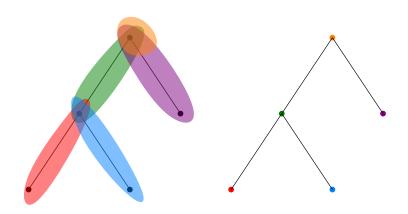
#### **Definition: Treewidth**

The treewidth of a graph *G* is the minimum width of all tree-decomposition of *G*.

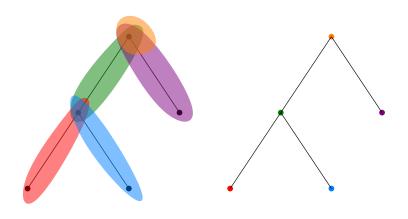
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Tree has treewidth 1

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#### Theorem

Planar graph on n vertices has treewidth  $O(\sqrt{n})$  and the bound is tight.

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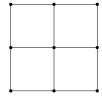
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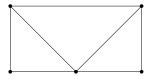
### Theorem (Vida et al.,2020)

For a graph G, if G has an H-partition of layered width I and H has treewidth k, then

$$qn(G) \leq 3l(2^k-1) + \lfloor \frac{3}{2}l \rfloor$$

#### **Definition: Partition and Quotient**

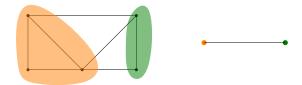
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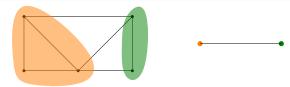


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A partition of G is a set  $\mathcal{P} = \{P_1, ..., P_n\}$  of non-empty subsets of V(G) and each vertex of G is in exactly one element (part) of  $\mathcal{P}$ . The quotient of  $\mathcal{P}$  is a graph, denoted by  $G/\mathcal{P}$ , with each vertex  $v_i$  corresponds  $P_i$ . For any two vertices  $v_i, v_j$  in  $G/\mathcal{P}$ , they are connected if and only if some vertex in  $P_i$  is connected to some vertex in  $P_j$  in graph G.

### Definition: *H*-partition

A *H*-partition of a graph *G* is a pair (A, H).  $A = \{A_x | x \in V(H)\}$  is a partition of V(G) and *H* is a graph isomorphic to the quotient G/A.

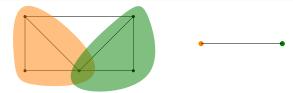


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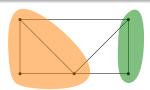


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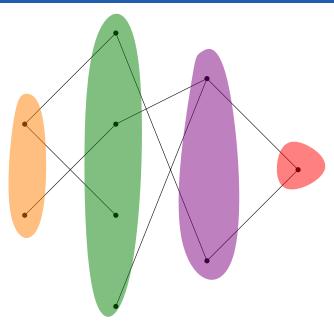
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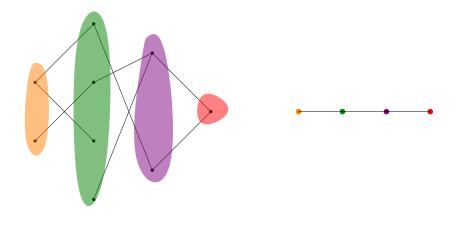
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# Layering

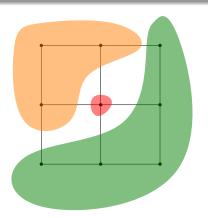


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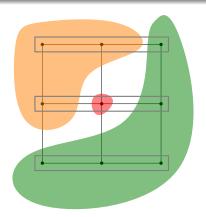


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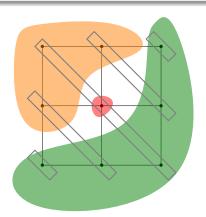
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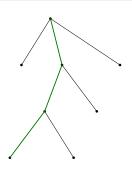


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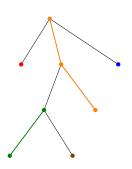
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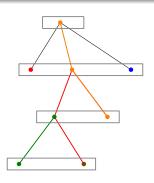
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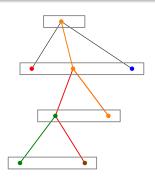


#### **Definition: Vertical Path**



#### **Definition: Vertical Path**

Let T be a tree rooted at a vertex r, a non-empty path  $(x_1, ..., x_p)$  in T is vertical if for some  $d \ge 0$  and for all  $1 \le i \le p$  we have  $\operatorname{dist}_T(x_i, r) = d + i$ .



A partition of the tree where each part is a vertical path has layered width 1.

### Theorem (Vida et al.,2020)

For a graph G, if G has an H-partition of layered width I and H has treewidth k, then

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Every planar graph G has queue-number at most

$$3(2^8-1)+\lfloor\frac{3}{2}\rfloor=766$$

#### Lemma (Vida et al.,2020)

- $|B_x| \le 9$  for any  $x \in V(T)$
- $\exists x \in V(T)$ , all vertices correspond to  $P_1, ..., P_k$  in H is in  $B_x$

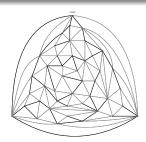
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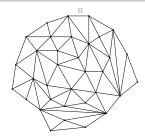
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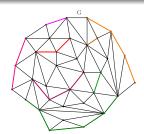
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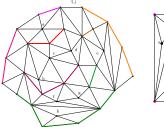
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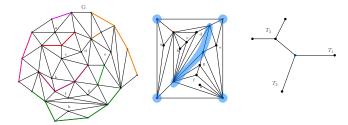
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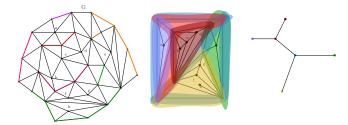
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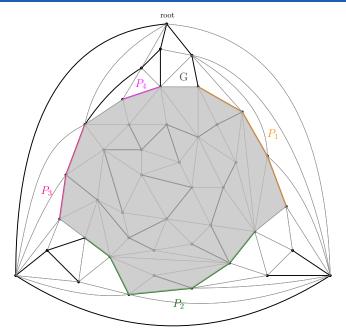
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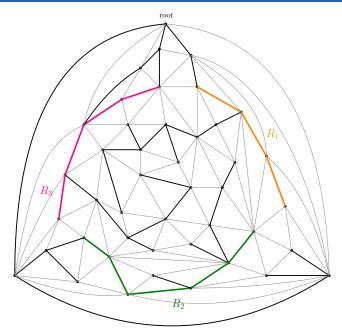


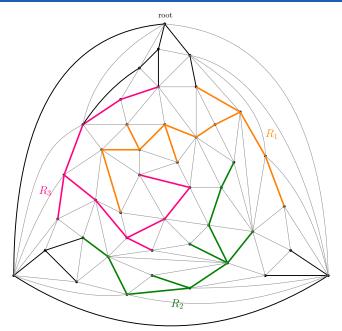
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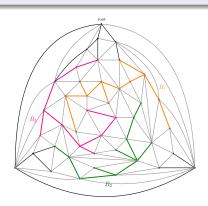


#### Lemma (Sperner's Lemma)

Let G be an internally triangulated graph whose vertices are colored 1,2,3 with the outer-face  $F = [P_1, P_2, P_3]$  where each vertex in  $P_i$  is colored i. Then G contains an internal face whose vertices are colored 1,2,3.

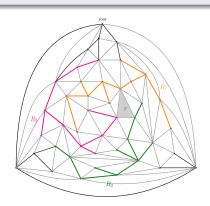
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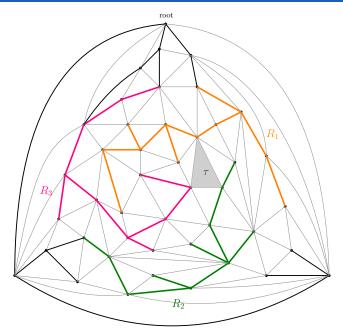
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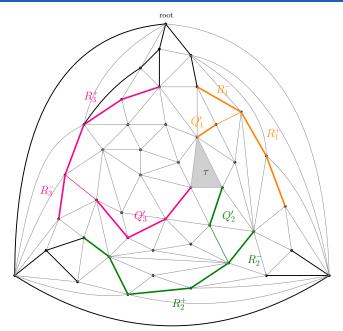


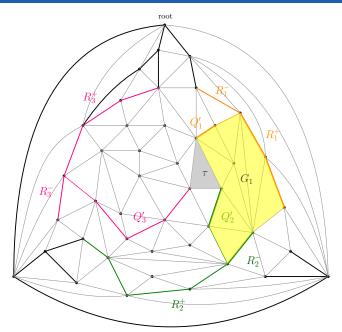
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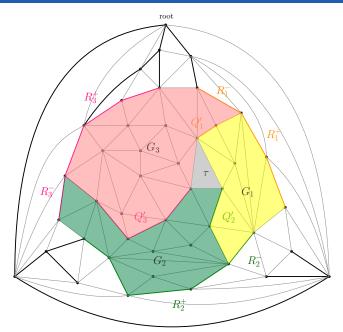
Let G be an internally triangulated graph whose vertices are colored 1,2,3 with the outer-face  $F = [P_1, P_2, P_3]$  where each vertex in  $P_i$  is colored i. Then G contains an internal face whose vertices are colored 1,2,3.

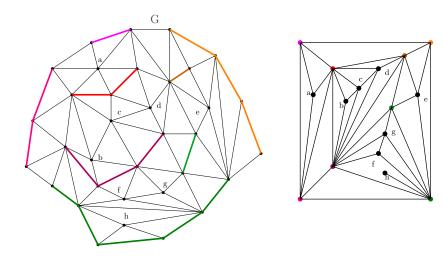


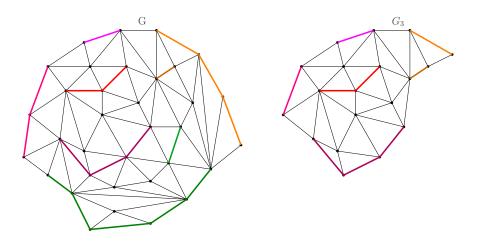


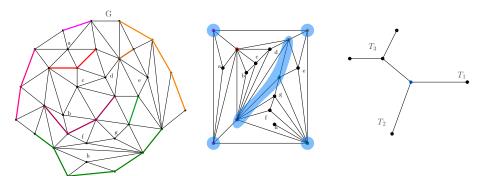


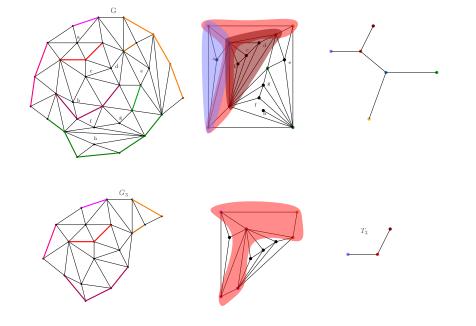


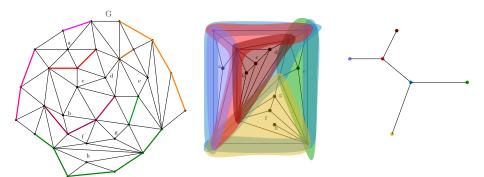












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Every planar graph G has a connected partition  $\mathcal{P}$  with layered width 1 such that  $H = G/\mathcal{P}$  has treewidth at most 8.

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$$qn(G) \leq 3l(2^k - 1) + \lfloor \frac{3}{2}l \rfloor$$

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#### Theorem (Vida et al.,2020)

Every planar graph G has queue-number at most

$$3(2^8-1)+\lfloor\frac{3}{2}\rfloor=766$$

#### Portal

- Introduction to Queue-Number
  - Queue Layout
  - Queue-Number
- Introduction to Treewidth
  - Treewidth
- Introduction to Partitions
  - Partitions
  - Layering
  - Layered Width
  - Vertical Path
- Planar Graph Decomposition
  - The Decomposition Lemma
  - Induction
- Summary