Planar Graphs Have Bounded Queue-Number

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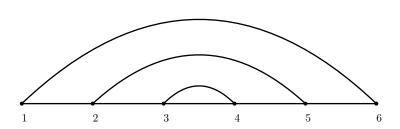
Geometry: Combinatorics and Algorithms Seminar Shengzhe Wang April 21, 2023

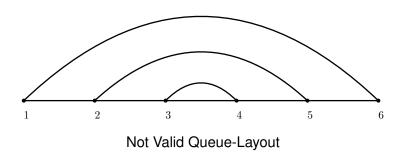


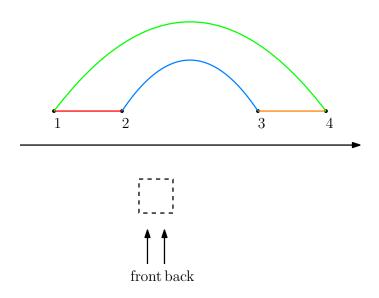
A path graph with 6 vertices

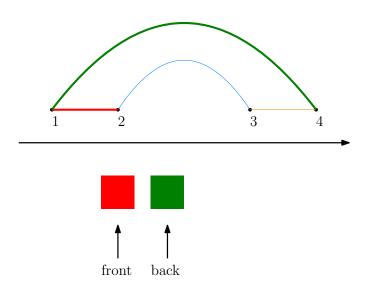


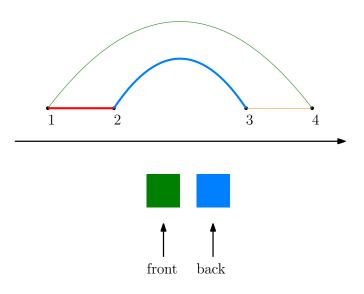
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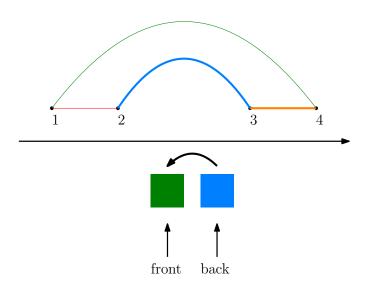






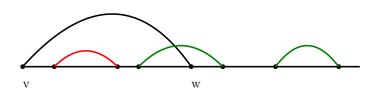






Definition: Queue

Let G = (V, E), consider a linear odering \leq of V, a queue of G is a set of edges $E' \subseteq E$ such that any disjoint edges $vw, xy \in E'$, w.l.o.g, $v \prec w, x \prec y$ and $v \prec x$, we have $w \prec y$.



Queue-Number

Definition: K-Queue Layout

Let G = (V, E), consider a linear odering \leq of V, for an integer $k \geq 0$ a k-queue layout of G is a partition of E into $E_1, E_2, ..., E_k$ such that each E_i is a queue of G with respect to the ordering \leq .

Queue-Number

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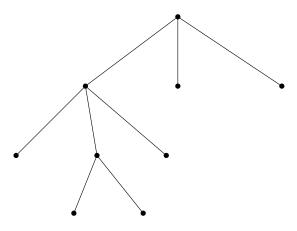
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Definition: Queue-Number

The queue-number of G, denoted by qn(G), is the minimum integer k such that G has a k-queue layout for some ordering \leq .

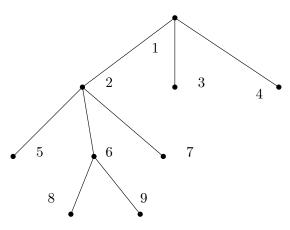
Queue-Number: Tree

What is the queue-number of a tree?

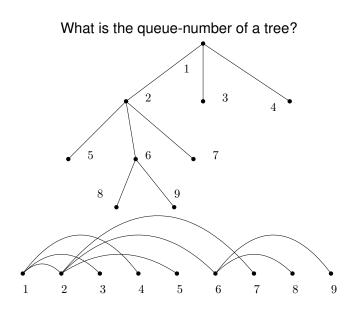


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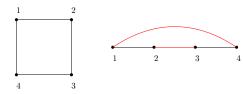
What is the queue-number of a tree?



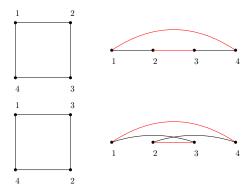
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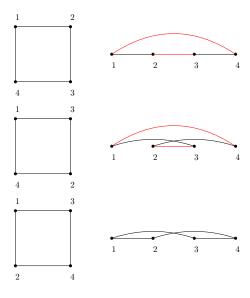
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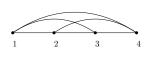


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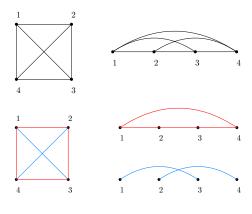


Queue Number: K₄

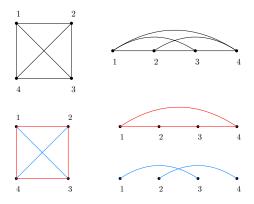




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Theorem (Heath, Rosenberg, 1992)

The complete graph K_n has queue number $\lfloor \frac{n}{2} \rfloor$.

Treewidth

Do we have some tools to help bound the queue-number?

Treewidth

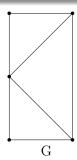
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Theorem (Wiechert, 2017)

Every graph with treewidth k has queue-number at most $2^k - 1$.

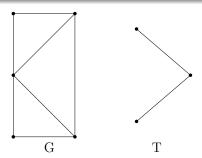
Definition: Tree-decomposition

- $\forall \{v, w\} \in E(G)$, there exists $x \in V(T)$ with $v, w \in B_x$
- $\forall v \in V(G)$, the set $\{x | x \in V(T) \land v \in B_x\}$ induces a non -empty connected subtree of T.



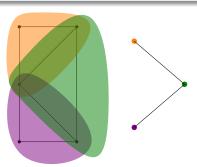
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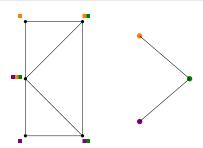
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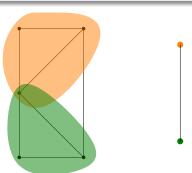
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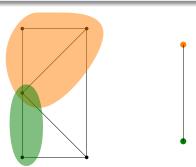
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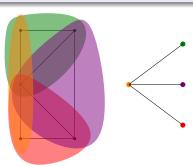
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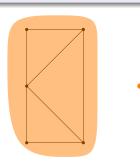
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Treewidth

Definition: Tree-decomposition

A tree-decomposition of a graph G is a pair (B, T). T is a tree and $B = \{B_x | x \in V(T)\}$ where each B_x is a subset of V(G) for every vertex x in V(T) such that

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The width of a tree-decomposition of G is $\max_{x \in V(T)} |B_x| - 1$

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The width of a tree-decomposition of *G* is $\max_{x \in V(T)} |B_x| - 1$

Definition: Treewidth

The treewidth of a graph *G* is the minimum width of all tree-decomposition of *G*.

Treewidth: Fixed-Parameter Tractability

• Computing a maximum independent set in a graph *G* is NP-hard.

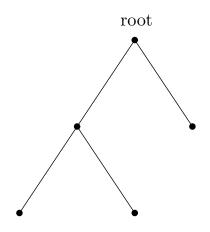
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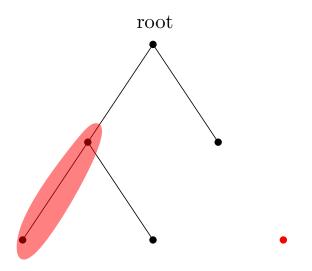
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- If G has treewidth $\leq k$, then a maximum independent set in G can be computed in time $O(k^24^k|V|)$.

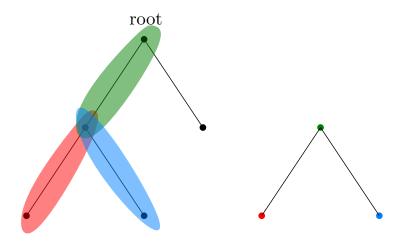
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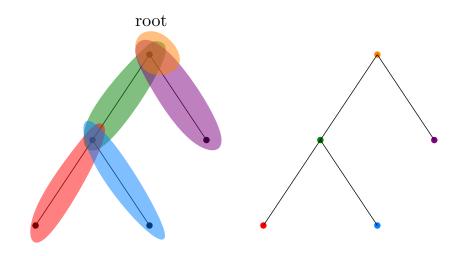
- Computing a maximum independent set in a graph G is NP-hard.
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- Dynamic Programming on trees is relatively fast.

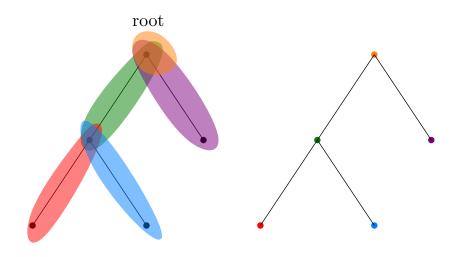
Treewidth: Tree











Tree has treewidth 1

Treewidth: Planar graph

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Theorem (Robertson, Seymour)

A grid graph with size $n \times n$ has treewidth n.

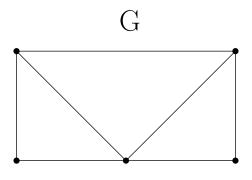
Maybe we need more structure

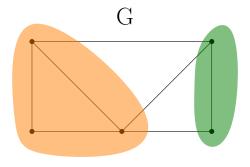
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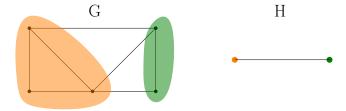
Theorem (Dujmović, Joret, Micek, Morin, Ueckerdt, Wood, 2020)

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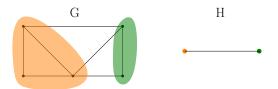






Definition: Partition and Quotient

A partition of G is a set $\mathcal{P} = \{P_1, ..., P_n\}$ of non-empty subsets of V(G) and each vertex of G is in exactly one element (part) of \mathcal{P} . The quotient of \mathcal{P} is a graph, denoted by G/\mathcal{P} , where each vertex v_i corresponds to P_i . For any two vertices v_i, v_j in G/\mathcal{P} , they are connected if and only if some vertex in P_i is connected to some vertex in P_i in graph G.

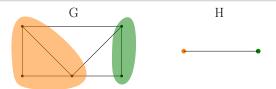


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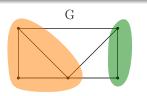


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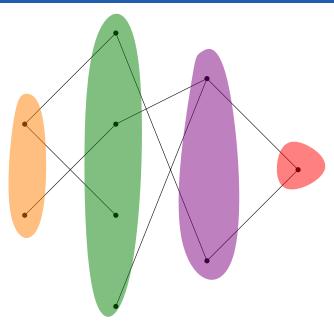
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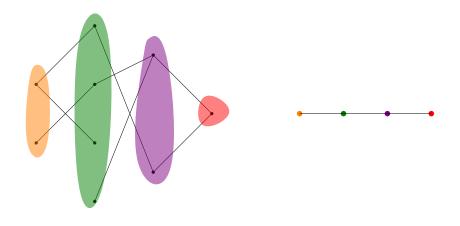


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Layering



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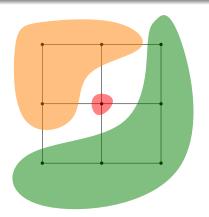


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The layered width of a partition \mathcal{P} of a graph G is the minimum integer ℓ such that there exists a path-partition (layering) of G, s.t. each element in \mathcal{P} has at most ℓ vertices in each element of the path-partition.

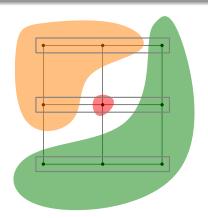
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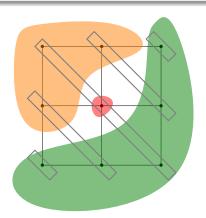
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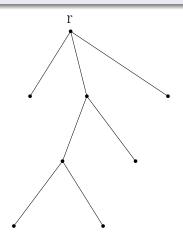
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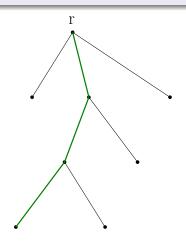


Definition: Vertical Path

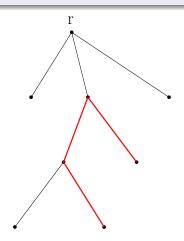
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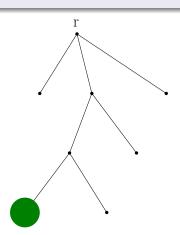
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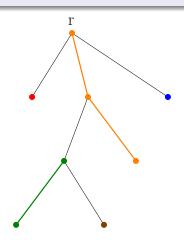
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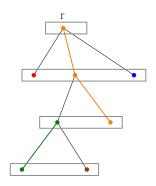
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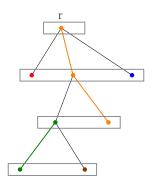


Definition: Vertical Path



Definition: Vertical Path

Let T be a tree rooted at a vertex r, a non-empty path $(x_1, ..., x_p)$ in T is vertical if for some $d \ge 0$ and for all $1 \le i \le p$ we have $\operatorname{dist}_T(x_i, r) = d + i$.



A partition of layered width 1 where each part is a vertical path.

Theorem (Dujmović, Joret, Micek, Morin, Ueckerdt, Wood, 2020)

For a graph G, if G has an H-partition of layered width ℓ and H has treewidth k, then

$$qn(G) \leq 3\ell(2^k-1) + \left\lfloor \frac{3}{2}\ell \right\rfloor$$

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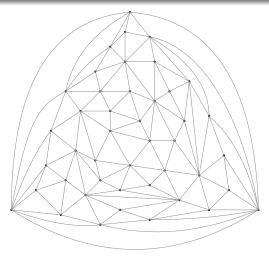
$$3(2^8-1)+\left|\frac{3}{2}\right|=766$$

Lemma (Dujmović, Joret, Micek, Morin, Ueckerdt, Wood, 2020)

(Simplified) Let G^+ be a maximal planar graph, we can find a subgraph G of G^+ , then G has a desired partition.

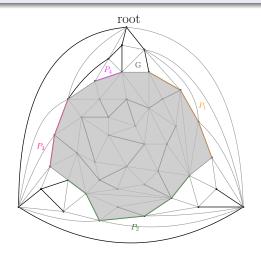
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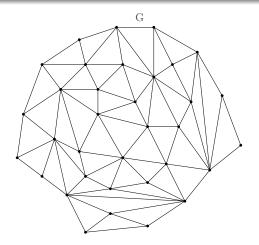
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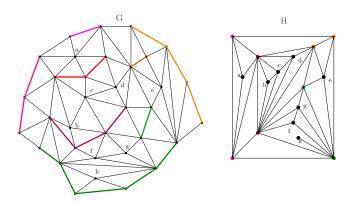
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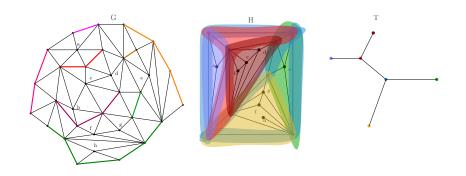
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- $|B_x| \leq 9$ for any $x \in V(T)$
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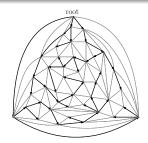
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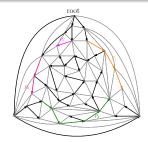
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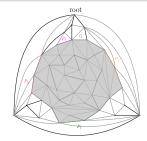
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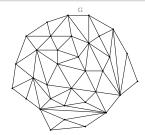
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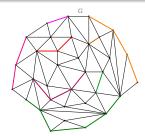
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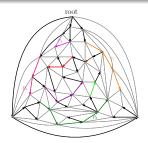
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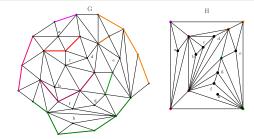
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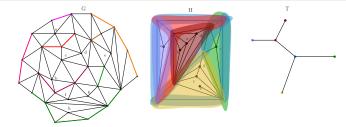
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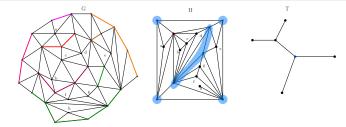
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Lemma (Dujmović, Joret, Micek, Morin, Ueckerdt, Wood, 2020)

Let G^+ be a maximal planar graph, let T be a spanning tree of G^+ rooted at vertex *root* on outerface of G^+ . For any cycle F in G^+ , which can be partitioned into at most 6 pairwise disjoint vertical paths of T, with $F = [P_1, ...P_k]$ and $1 \le k \le 6$. Let G be the internally triangulated subgraph of G^+ which consists of all edges and vertices of G^+ contained in F and the interior of F, then G has a partition $\mathcal P$ into vertical paths of T, and $P_1, ..., P_k \in \mathcal P$, and the quotient graph $H = G/\mathcal P$ has a tree-decomposition (B, T) that

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Note that three vertices on the outer cycle can always be three vertical paths of T.

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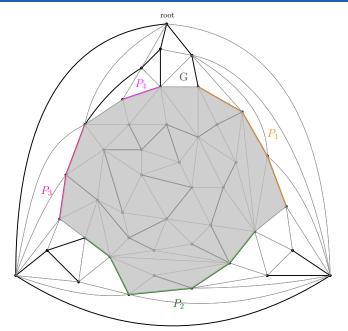
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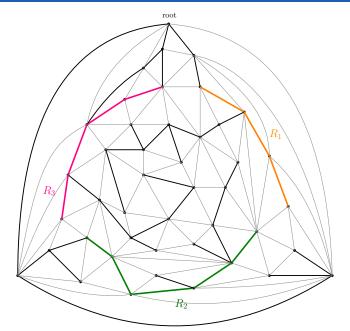
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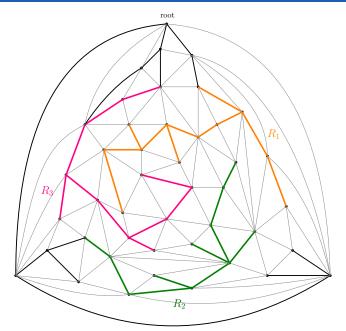
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Every planar graph G has a connected partition \mathcal{P} with layered width 1 such that $H = G/\mathcal{P}$ has treewidth at most 8.





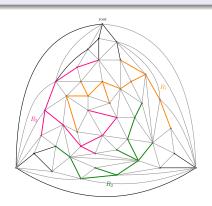


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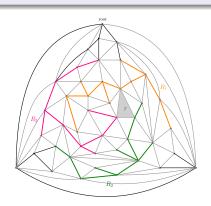
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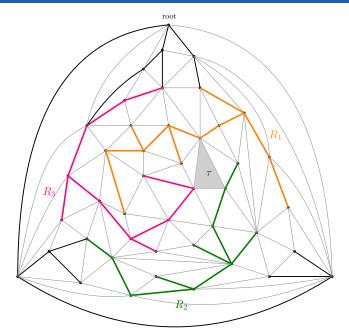
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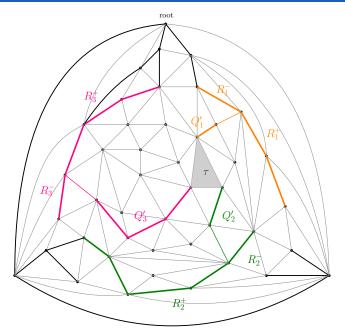


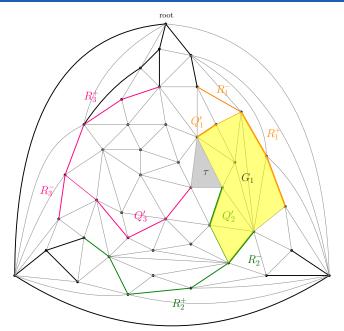
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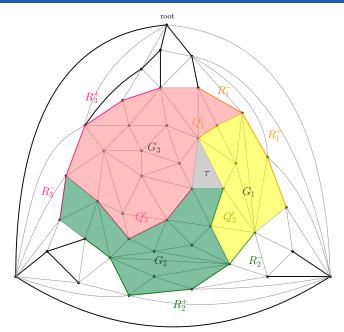
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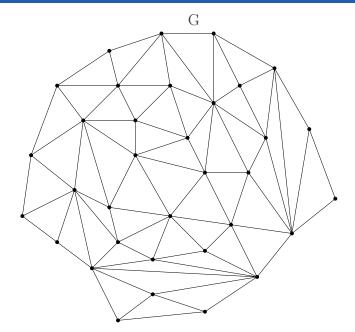


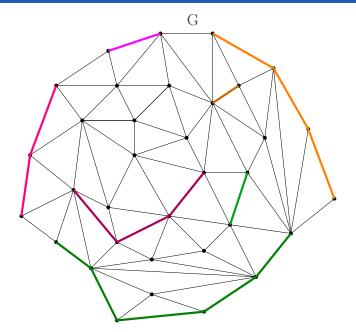


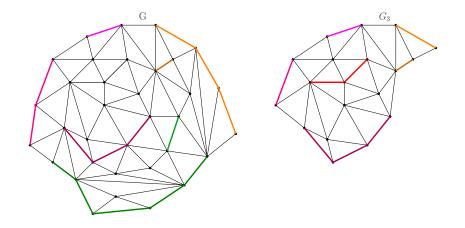


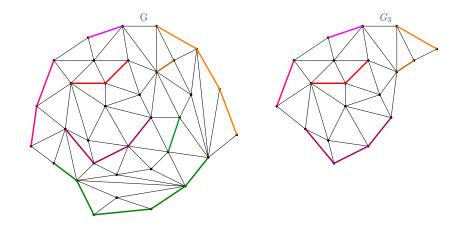


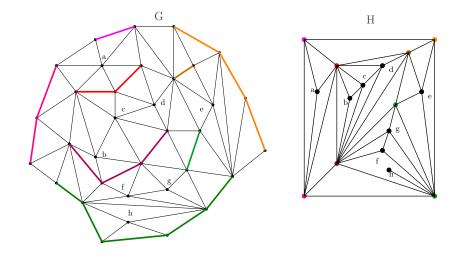


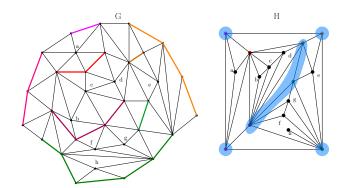


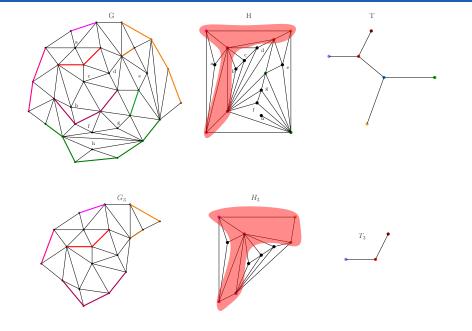


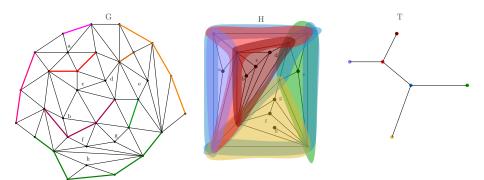




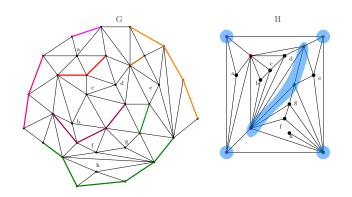




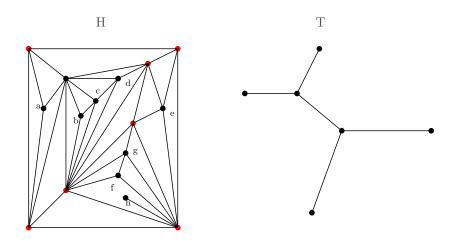


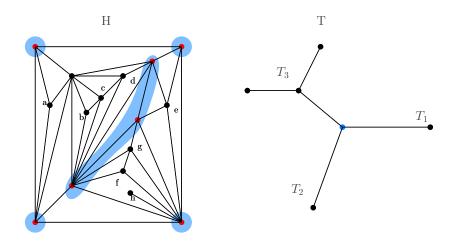


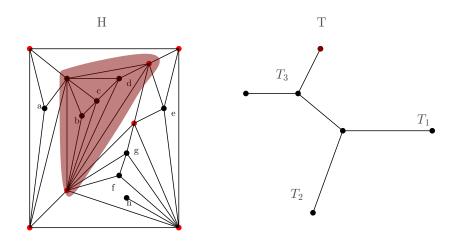
Planar Graph Decomposition: Correctness

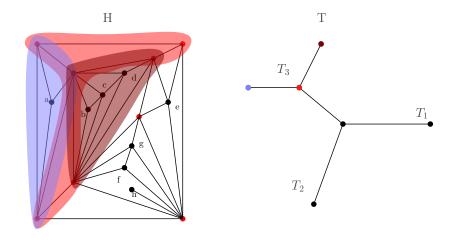


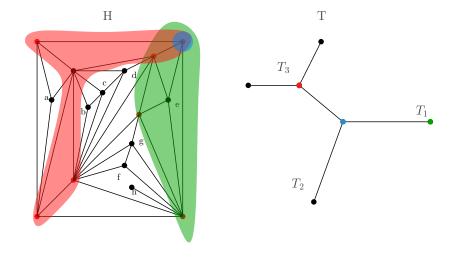
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Every planar graph G has queue-number at most

$$3(2^8-1)+\left|\frac{3}{2}\right|=766$$

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- Relation between treewidth and queue-number
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 - Best known lower bound for queue-number is k + 1.
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...

Portal

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 - Queue-Number
- Introduction to Treewidth
 - Treewidth
 - Introduction to Partitions
 - Partitions
 - Layering
 - Layered Width
 - Vertical Path
 - Planar Graph Decomposition
 - The Decomposition Lemma
 - Induction

 - Construction
 - Correctness
- Summary
 - Wrap-Up
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