

# Recap

• OS

• Interrupts

• Organization

## Computer system architecture

### 1. Single Processor System

→ A single CPU executes general purpose instructions

- Other processors help the CPU

### 2. Multiprocessor System

→ Throughput

- Two or more CPUs

sharing bus, memory

Why?

#### 1. Increased Throughput

• Speedup would be less than  $(n)$

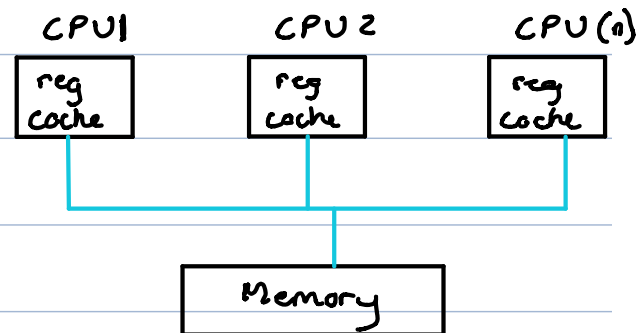
due to there being an overhead in keeping the correct work.

#### 2. More reliable

• System still functions even if some processors fail

#### 3. Less cost

• Share resources such as power supply.



storage, etc.

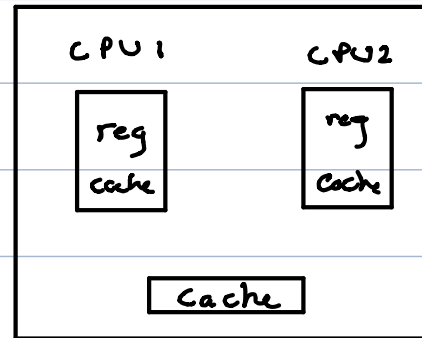
• Asymmetric  
↓  
not peers

vs

Symmetric  
↓  
peers

## Multi-core

Processors are on  
the same chip



Why?

- Faster Communication
- less power
- less space

## Clustered System

• Individual System Connected through a fast local area network (LAN)

- High performance computing clusters

- par for (MATLAB)

↳ each thing will have different interaction  
on a system (different CPU)

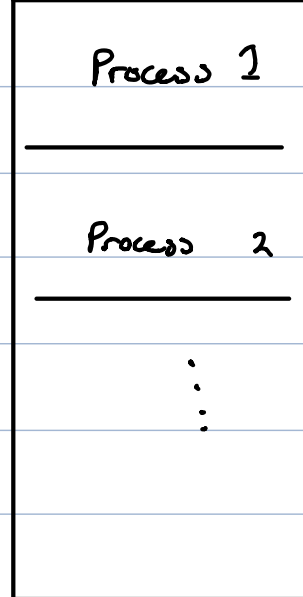
• Each loop on different CPU

## OS Structure

## 1. Multi programming

- keep multiple processes in memory

→ One a process blocks, another process gets the CPU.



## 2. Time sharing

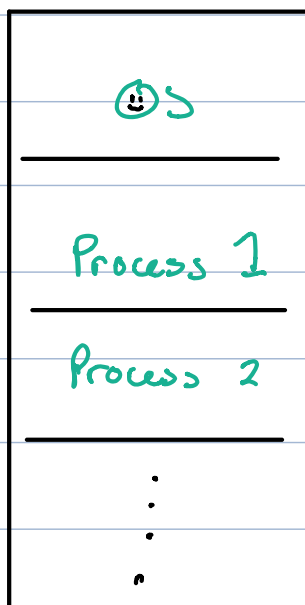
Switching between processes is done very fast

\* Response time | Enter to time you get something from server



The time taken between submitting a task and getting the first response

Turn around time is the whole result

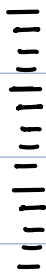


if process 1  
went into process 2  
it would seg fault

OS operations

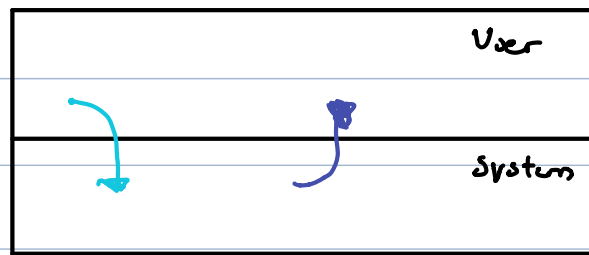
OS interrupt drives an error in a process should not affect the process or the OS / other processes

Process



User mode  
bit = 1

User mode  
bit = 0



Kernel  
supervisor  
System  
Privileged

once OS receives the write bit becomes zero

write() : →



User vs. Kernel bit

Timer

Goal : Make sure the  
OS gets the  
CPU back

Clock + counter.

when counter hits 0 an interrupt occurs (time slices)

OS functionalities to be provided?

• Process management

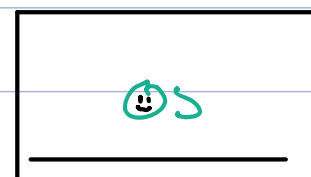
→ Create, terminate, suspend. (Ability to :)

→ Schedule.

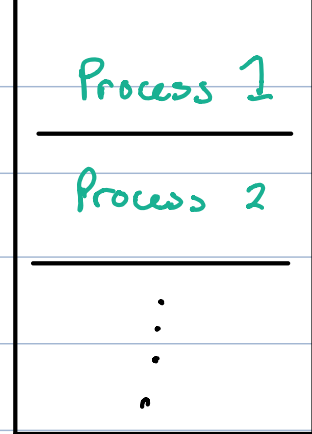
→ Enable communication

(interprocess communication)

→ Process Synchronization



Shared memory  
Message Passing



- Memory management

- Track memory (Free + in use)
- List storage / memory locations belonging to each process -
- Allocating and deallocating
- Read / Write
- Suspend process to / from Disk and unsuspend from Disk



- Storage management

1. File system [Managing files / Directories]
2. I/O devices

Manage devices through drivers

Network, display, mice, cards, printers

- Power Management

- Protection + Security

User groups



Around 200 or 300

system call

- goal is to find the right system call via API

↳ broken into separate ones.

Security / Permissions