

A	B	C
10	5	7

	Allocation		
	A	B	C
P <sub>0</sub>	0	1	0
P <sub>1</sub>	2	0	0
P <sub>2</sub>	3	0	2
P <sub>3</sub>	2	1	1
P <sub>4</sub>	0	0	2

n x m

Max		
A	B	C
2	5	3
3	2	2
1	0	2
2	2	2
4	3	3

Need		
A	B	C
7	4	3
1	2	2
6	0	0
0	1	1
4	3	1

Available

A	B	C
3	3	2
(10-7)	(5-2)	(7-5)

Assume P<sub>1</sub> can finish

P <sub>1</sub> *	available		
	A	B	C
	3	3	2
+	2	0	0
P <sub>3</sub> *	5	3	2
+	2	1	1

→ P <sub>1</sub>	2	0	0
→ P <sub>3</sub>	2	1	1

7 4 3 ⇒ All can finish system is safe

P<sub>1</sub> Request 1 A unit ?

	Allocation			Max	Need
	A	B	C		
P <sub>0</sub>	0	1	0	7	7
P <sub>1</sub>	3	0	0	3	0
P <sub>2</sub>	3	0	2	9	6
P <sub>3</sub>	2	1	1	2	0
P <sub>4</sub>	0	0	2	4	4

n x m

Available

A B C      => safe grant  
2 3 2

P<sub>1</sub> Request 2 additional units of type C

Allocation	Max	Need
3 0 2	3 2 2	0 2 0

Available      still safe

2 3 0  
A ↑ B ↑ C ↑

P<sub>4</sub> request (3, 3, 0)

Allocation      Max      need

$p_4$     3   3   2        4   33        1   0   1

Available

2   3   0         $\Rightarrow$  Deny request

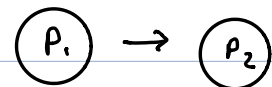
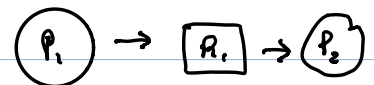
### 3 Detection and Recovery

a Single instant of resource

• wait for graph

If a cycle exists

$\Rightarrow$  System is deadlocked



b Multiple instants

1 Work = Available, Finish[i] = False

2 Find process i s.t.

$\left. \begin{array}{l} \text{finish}[i] == \text{False}; \\ \text{request}_i \leq \text{Work} \end{array} \right\} \rightarrow$

If no such i, go to 4

3 Work = work + allocation<sub>i</sub>

Finish[i] = True;

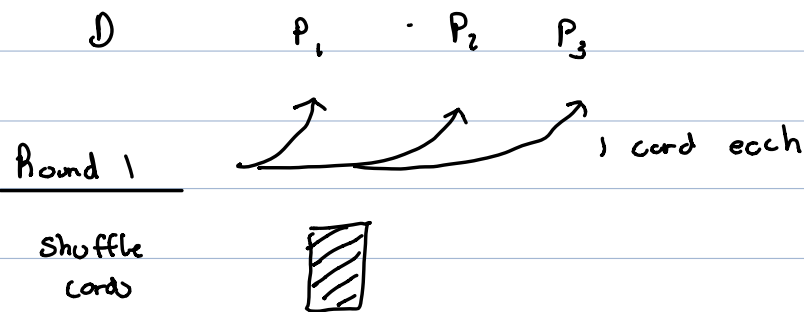
go to 2

4 If finish == false for some (i)  
these are deadlocked

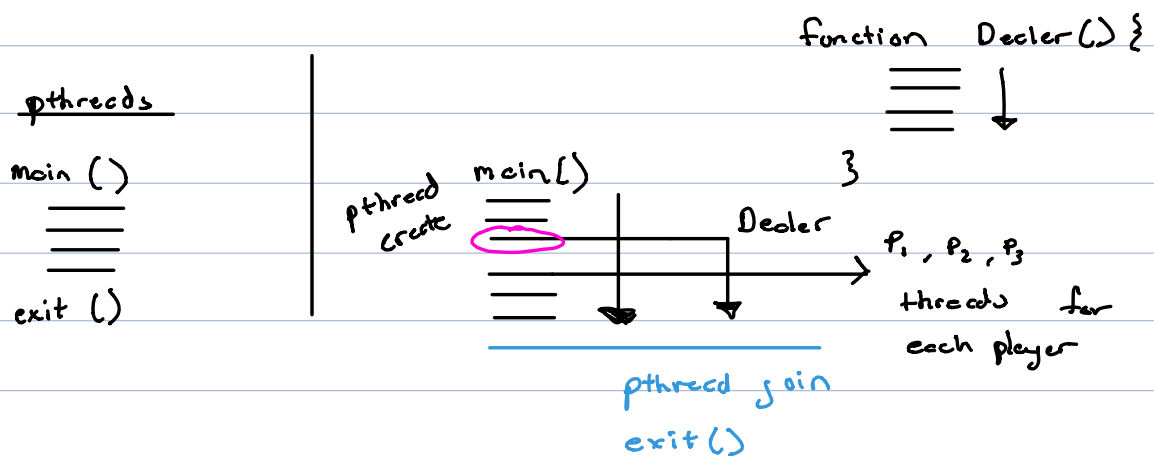
## Project 2

1 Game

3 rounds



• Checking to see if highest card is present



• Dealers run first  
→ wait for  $P_i$

• Must lock & unlock  
resources.

→ wait for  $P_2$

• requires pthread join

....

## Mutex

pthread\_mutex\_t handler;

pthread\_mutex\_init(&);

→ handler.lock();

→ handler.unlock();

## Condition Variable

pthread\_cond\_t y;

pthread\_cond\_wait(&y, &x)

pthread\_cond\_signal(&y);

pthread\_mutex\_lock(&x);

pthread\_cond\_wait(&y, &x)

pthread\_mutex\_unlock(&x)

:

pthread\_mutex\_lock(&x)

pthread\_cond\_signal(&y)

pthread\_mutex\_unlock(&x)

lock i.e.  $P_2$

wait

condition is unlock  
but the signal locks  
it.

2 layers unlocking

Player function (Pass number of players)

2nd argument would tell you if it finished properly

