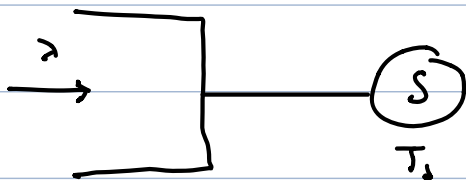


# Queueing Theory



$$\rho = \lambda T_s = \frac{\lambda}{\mu} \quad 0 \leq \rho < 1$$

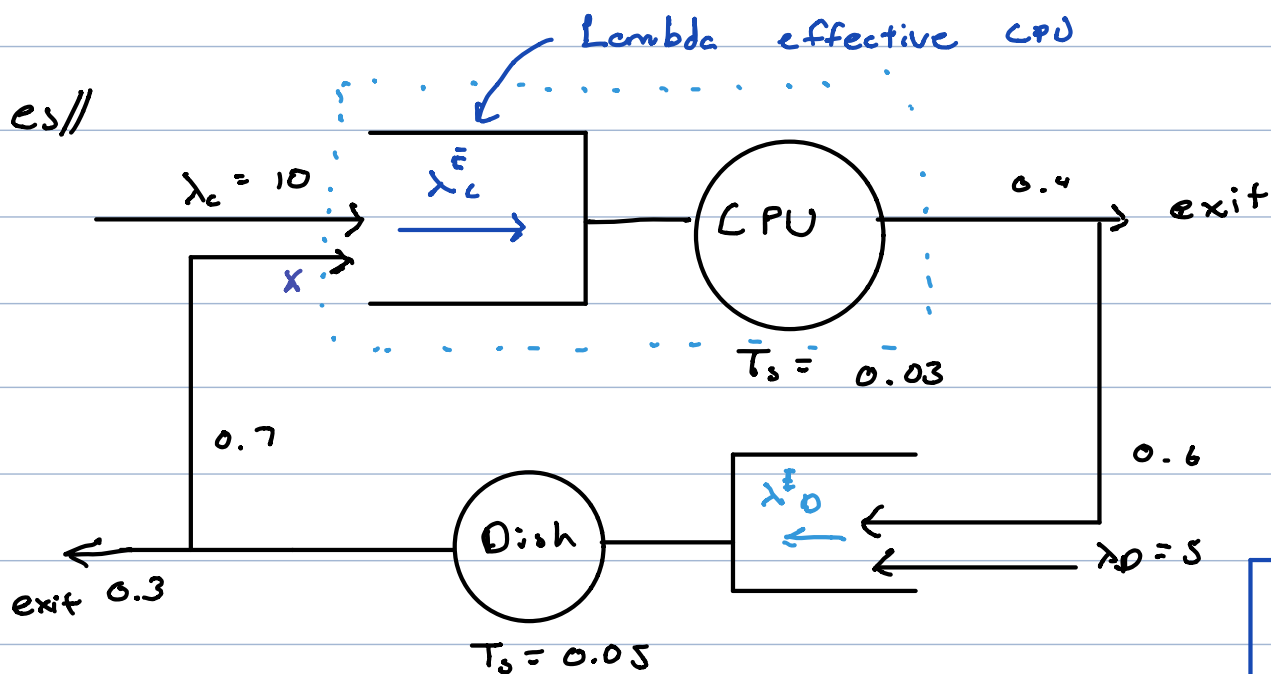
$$q = \lambda T_q$$

$$T_q = T_w + T_s$$

$$w = \lambda T_w$$

$$q = w + s$$

$$q = \frac{\rho}{1 - \rho}$$



Time stamp  
4:07

1  $\rho_{CPU}$   $\rho_{Disk}$

$$\lambda_c^E = 10 + x$$

$$\lambda_D^E = 5 + 0.6 * \lambda_c^E$$

$$x = 0.7 \lambda_D^E$$

process /  
CPU cycle

$$\lambda_c^E = 10 + 0.7 \lambda_D^E$$

$$\lambda_D^E = 5 + 0.6 \lambda_c^E$$

$$\lambda_c^E = 10 + 0.7(5 + 0.6) \lambda_c^E$$

$$\lambda_c^E = 23.2$$

$$\lambda_D^E = 18.9$$

$$\rho_{CPU} = \lambda_L^E * T_{s CPU} = 23.2 * .03 = 0.69$$

$$\rho_{Disk} = \lambda_D^E * T_{s Disk} = 18.9 * 0.05 = 0.945$$

Disk is the bottleneck b/c it's closer to 1 (resource)

$\rho_{Disk}$

$\rho_{CPU}$

$$\rho_{CPU} = \frac{\rho_{CPU}}{1 - \rho_{CPU}} = \frac{0.69}{1 - 0.69} = 2.22 \text{ processes}$$

$$\rho_{Disk} = \frac{0.945}{1 - 0.945} = 17.18$$

$$T_{q_{CPU}} = \frac{\rho_{CPU}}{\lambda_L^E} = \frac{2.22}{23.2} = \text{_____ seconds.}$$

What is the probability that the process arrives and finds the CPU idle.

$$1 - 0.69 = 0.31 \text{ or } 31\% \text{ chance}$$

$T_{2 total}$  = Average turnaround time for a process submitted to the whole system.

$$\lambda_c = 10 \rightarrow$$



$$T_{q, total} = \frac{q_c}{\lambda_c}$$

$$\leftarrow \lambda_D = 5$$

$$\frac{q_{CPU} + q_{Disk}}{15} \quad \leftarrow \quad 5 + 10$$

$$= \frac{2.22 + 17.18}{15} = \text{seconds}$$

What arrival would break the system?

Disk would due to it being closer to 1

$\lambda_D^*$  = break the system

$$\rho_{Disk} = 1$$

$$20 = \lambda_D^* + 24 * 0.6$$

$$\lambda_D^* = 20 - \left( \frac{24 * 6}{10} \right)$$

$$20 - 14.4 = 5.6 \text{ pr/s breaks the system}$$

↓  $\lambda_c^* = \text{breaks the system}$

$\rho_{bwh} = 1$  (bottleneck)

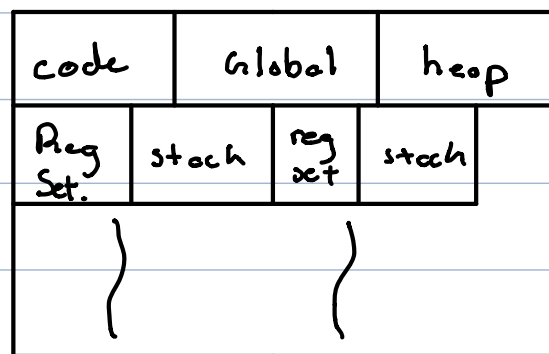
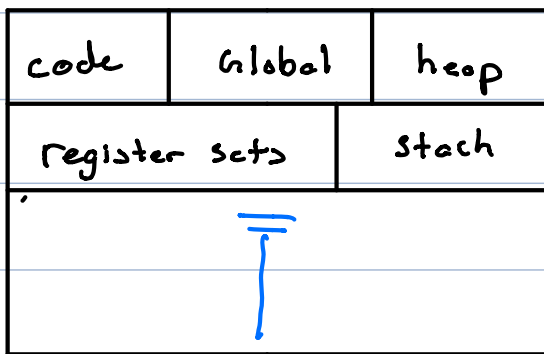
$\lambda_c^* = 11$

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## Threads

**Goal:** Support multiple concurrent paths of execution within a single process.

before



Why?

[1] Scalability: Increased speedup in multi-processor systems.

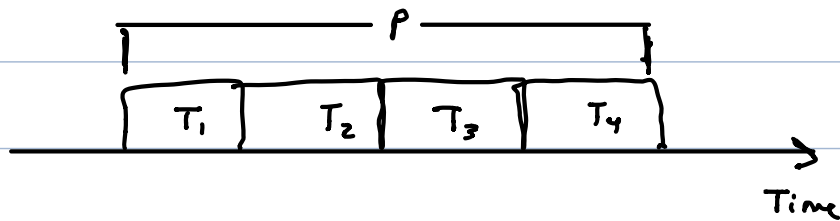
[2] Responsiveness: An application can continue to run / execute, even if part of the program being blocked.

[3] Economy : thread creation is faster + efficient than process creation

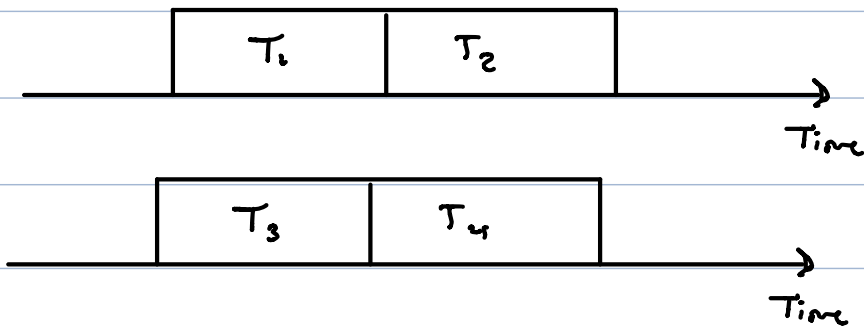
[4] Resource Sharing : By default threads share resources (no need for shared mem / message passing)

$T_1$  —————  $T_4$

Single Processor



Multiple Processor >



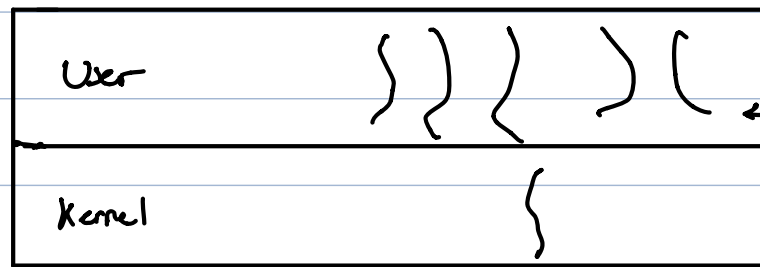
Many to one } are some of the types  
one to one }

Threading Models

⋮

1

Many - to - one model



user threads

Library scheduler

1 : 03 : 00

kernel thread

Thread management done by  
a thread library e.g p threads

Create threads

Kill threads + synchronize

schedule threads

• When one thread  
blocks then the  
entire application  
blocks

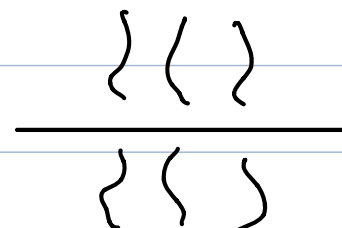
## Drawbacks

- \* process blocks if 1 thread blocks
- \* no speedup on multi processor systems

2 One - to - one Model

run in parallel

Open MP



13 Many-to-Many

Interleaved to work  
with the influx

$\frac{\{ \} )}{\{ | }$