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CHAPTER 4 — FILE SYSTEMS

- >Files
- > Directories
- > File system implementation
- >Example file systems



File Systems

Files File Concept

➤ Long-term Information Storage

- Must store large amounts of data
- Information stored must survive the termination of the process using it
- Multiple processes must be able to access the information concurrently

≻File

- Used to store information on disks and other extenal media in units
- Process can read them and write new ones if need be

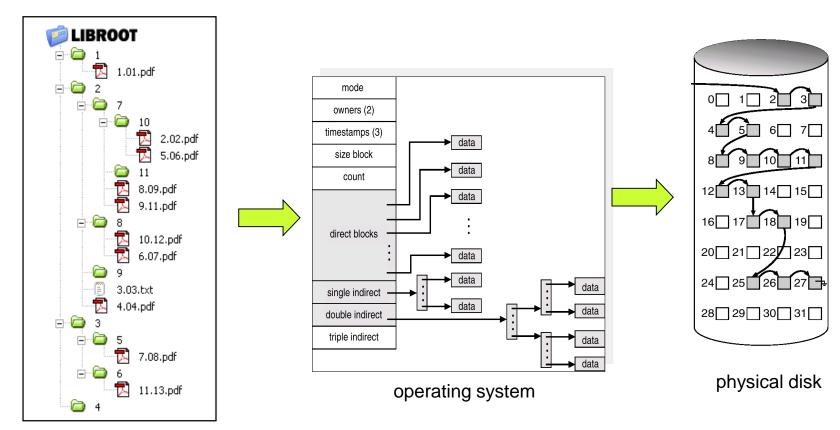
➤ File system

- Part of operating system dealing with files
- Includes two independent parts: set of file and directory structure, organize and provide information about all files in system



User Abstraction		Hardware Resource
Process/Thre ad		CPU
Address Space	⇔ OS ⇒	Memory
Files		Disk

Files File Concept



file system



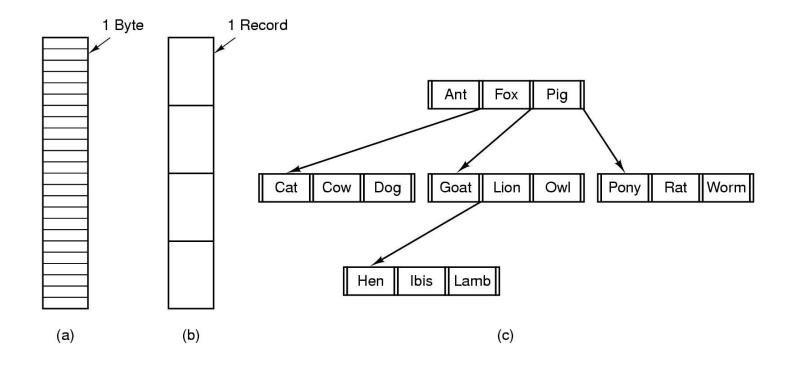
Typical file extensions.

Extension	Meaning	
file.bak	Backup file	
file.c	C source program	
file.gif	Compuserve Graphical Interchange Format image	
file.hlp	Help file	
file.html	World Wide Web HyperText Markup Language document	
file.jpg	Still picture encoded with the JPEG standard	
file.mp3	Music encoded in MPEG layer 3 audio format	
file.mpg	Movie encoded with the MPEG standard	
file.o	Object file (compiler output, not yet linked)	
file.pdf	Portable Document Format file	
file.ps	PostScript file	
file.tex	Input for the TEX formatting program	
file.txt	General text file	
file.zip	Compressed archive	



➤ Three kinds of files

- (a) byte sequence
- (b) record sequence
- **■** (c) tree

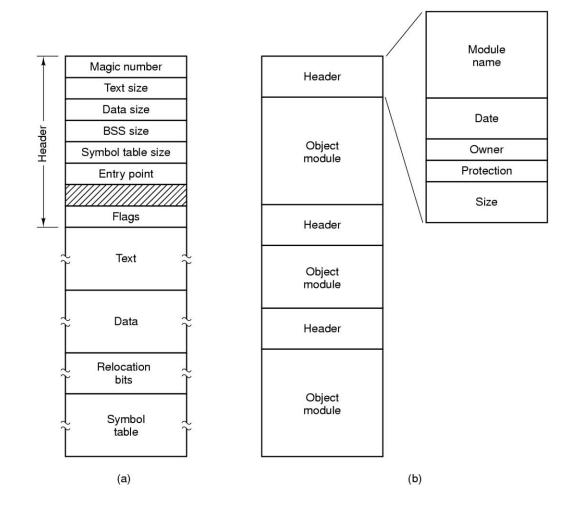


Files File Types

- ➤ Contiguous logical address space
- >Types:
 - Data
 - ✓ numeric
 - √ character
 - √ binary
 - Program
 - ✓ Source
 - ✓ Object
 - ✓ Executable
 - Regular, special (character, block)

Files File Types

(a) An executable file (b) An archive



Files File Access

➤ Sequential access

- read all bytes/records from the beginning
- cannot jump around, could rewind or back up
- convenient when medium was mag tape

> Random access

- bytes/records read in any order
- essential for data base systems
- read can be ...
 - ✓ move file marker (seek), then read or ...
 - ✓ read and then move file marker



Possible file attributes

Attribute	Meaning
Protection	Who can access the file and in what way
Password	Password needed to access the file
Creator	ID of the person who created the file
Owner	Current owner
Read-only flag	0 for read/write; 1 for read only
Hidden flag	0 for normal; 1 for do not display in listings
System flag	0 for normal files; 1 for system file
Archive flag	0 for has been backed up; 1 for needs to be backed up
ASCII/binary flag	0 for ASCII file; 1 for binary file
Random access flag	0 for sequential access only; 1 for random access
Temporary flag	0 for normal; 1 for delete file on process exit
Lock flags	0 for unlocked; nonzero for locked
Record length	Number of bytes in a record
Key position	Offset of the key within each record
Key length	Number of bytes in the key field
Creation time	Date and time the file was created
Time of last access	Date and time the file was last accessed
Time of last change	Date and time the file has last changed
Current size	Number of bytes in the file
Maximum size	Number of bytes the file may grow to

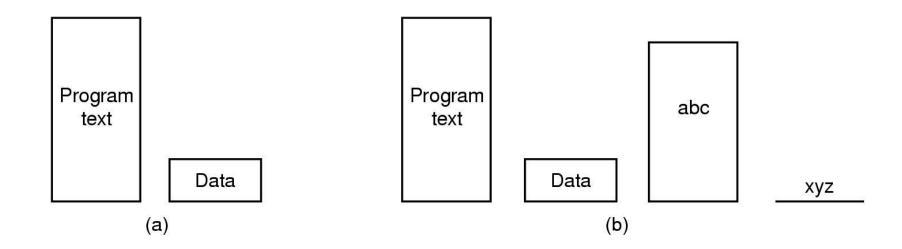
Files File Operations

- 1. Create
- 2. Delete
- 3. Open
- 4. Close
- 5. Read
- 6. Write

- 7. Append
- 8. Seek
- 9. Get attributes
- 10.Set Attributes
- 11.Rename

Files Memory-Mapped Files

- (a) Segmented process before mapping files into its address space
- (b) Process after mapping existing file *abc* into one segment creating new segment for *xyz*

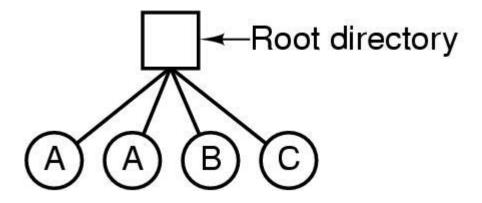


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13

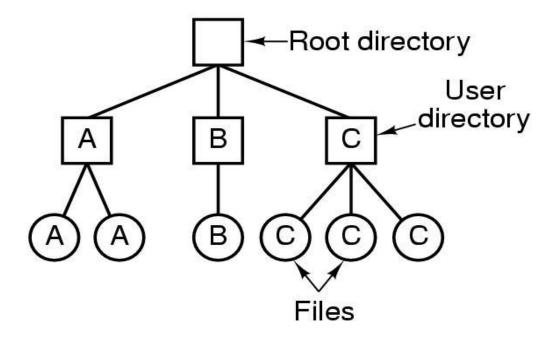
Directories Single-Level Directory Systems

- ➤ A single level directory system
 - contains 4 files
 - owned by 3 different people, A, B, and C



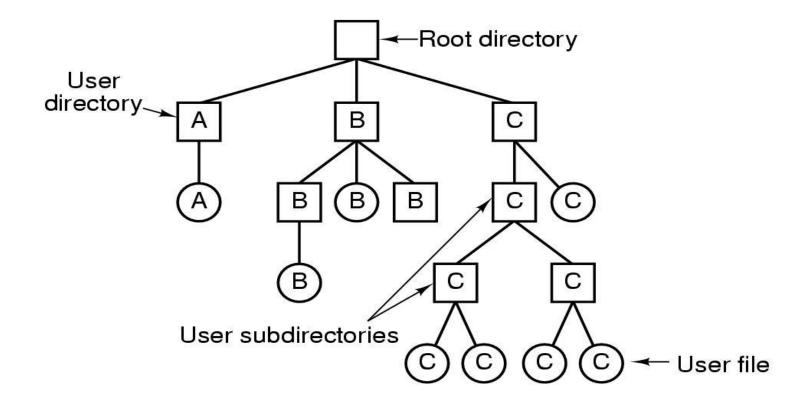
Directories Two-level Directory Systems

Letters indicate *owners* of the directories and files



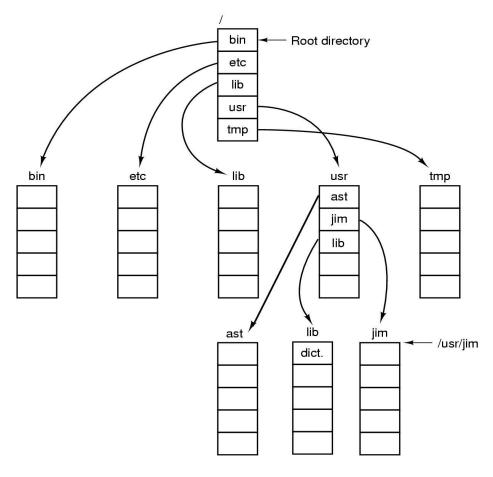


A hierarchical directory system





A UNIX directory tree



Directories Directory Operations

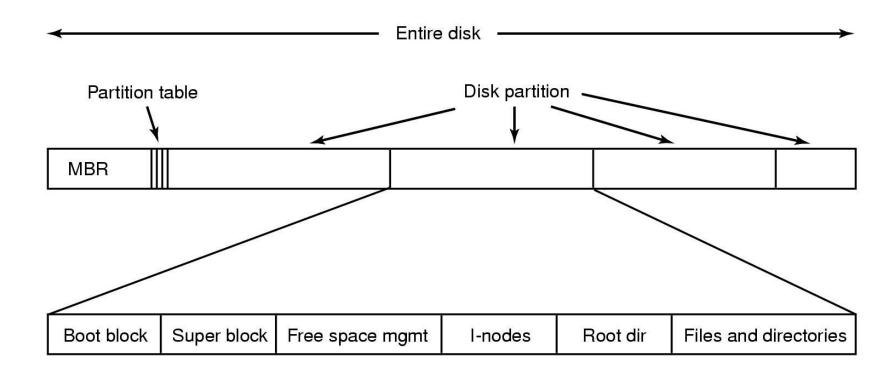
- 1. Create
- 2. Delete
- 3. Opendir
- 4. Closedir

- 5. Readdir
- 6. Rename
- 7. Link
- 8. Unlink



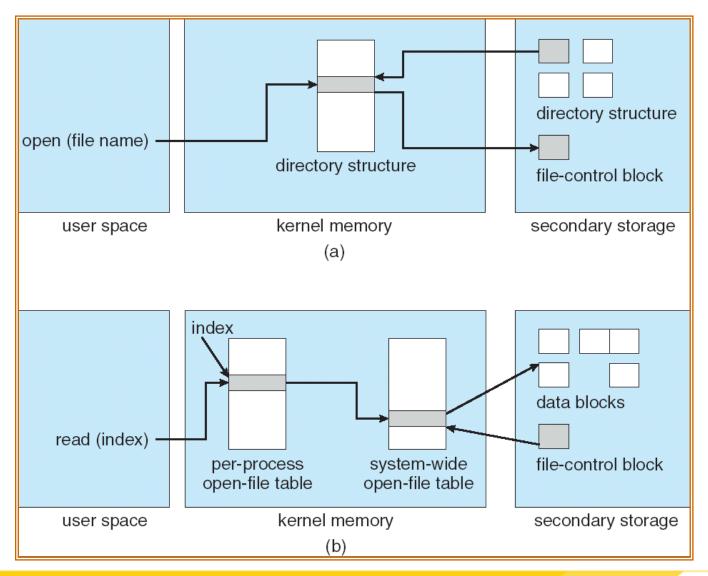


A possible file system layout





In-Memory File System Structures





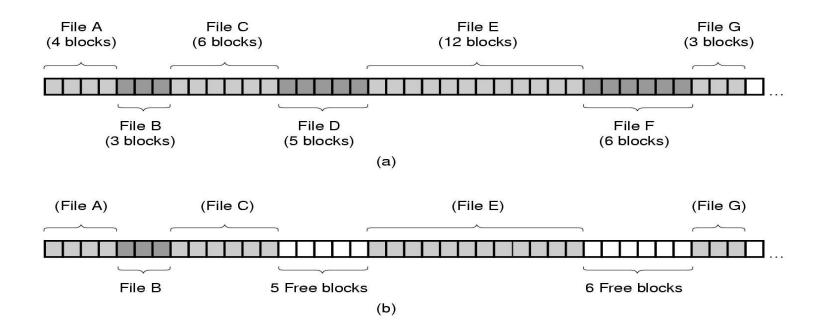
File System Implementation Allocation Methods

- >An allocation method refers to how disk blocks are allocated for files:
- **→** Contiguous allocation
- **►** Linked allocation
- >Indexed allocation



Allocation Methods: Contiguous allocation (1)

- (a) Contiguous allocation of disk space for 7 files
- (b) State of the disk after files D and E have been removed





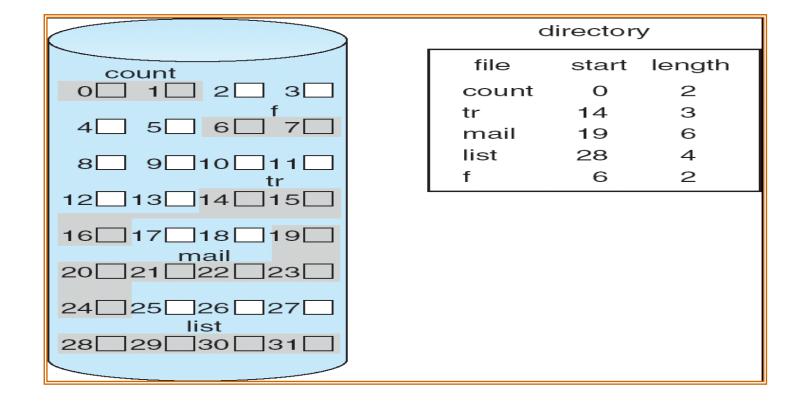
Allocation Methods: Contiguous allocation (2)

- Each file occupies a set of contiguous blocks on the disk
- ➤ Simple only starting location (block #) and length (number of blocks) are required
- ➤ Random access
- ➤ Wasteful of space (dynamic storage-allocation problem)
- ➤ Files cannot grow
- ➤ Use for CD-ROM because the length of file are known in advance and no deletion



Allocation Methods: Contiguous allocation (3)

Contiguous Allocation of Disk Space



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25



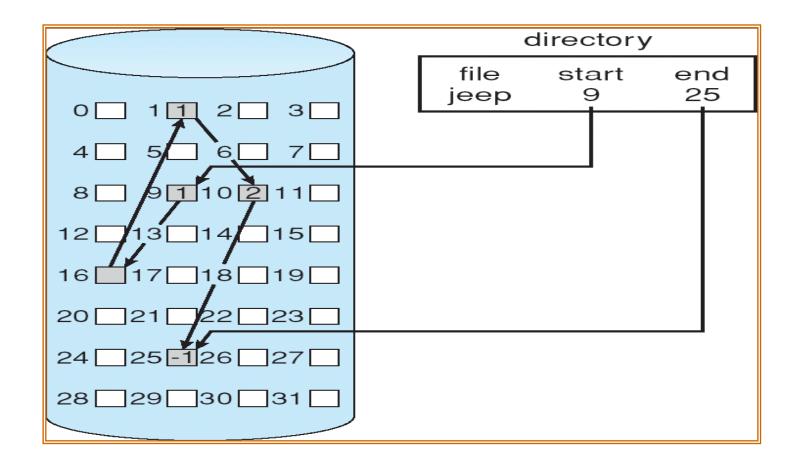
File System Implementation Allocation Methods: Linked allocation (1)

- Each file is a linked list of disk blocks: blocks may be scattered anywhere on the disk.
- ➤ Simple need only starting address
- > Free-space management system no waste of space
- ➤ No random access

block = pointer

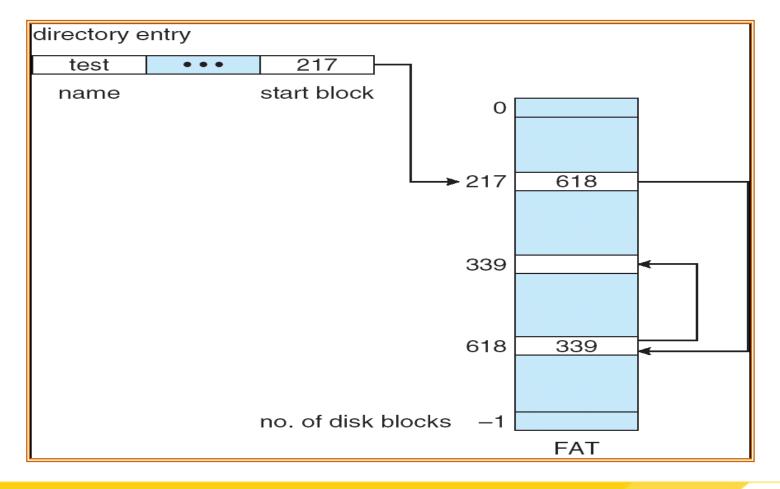


File System Implementation Allocation Methods: Linked allocation (2)



File System Implementation Allocation Methods: Linked allocation (3)

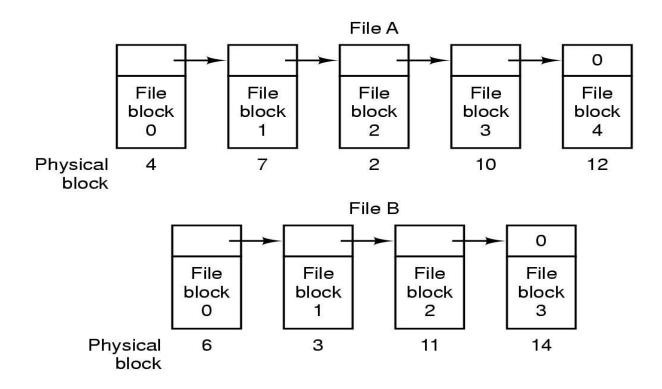
➤ File-Allocation Table





File System Implementation Allocation Methods: Linked allocation (4)

Storing a file as a linked list of disk blocks

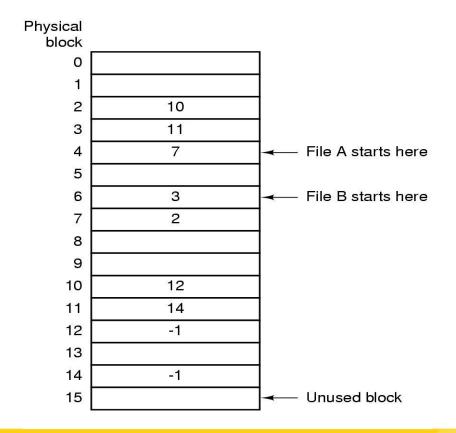




File System Implementation Allocation Methods: Linked allocation (5)

Linked list allocation using a file allocation table in RAM

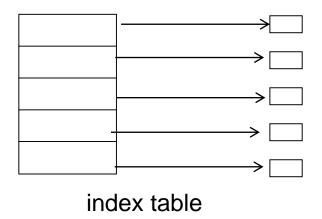
FAT File Allocation Table





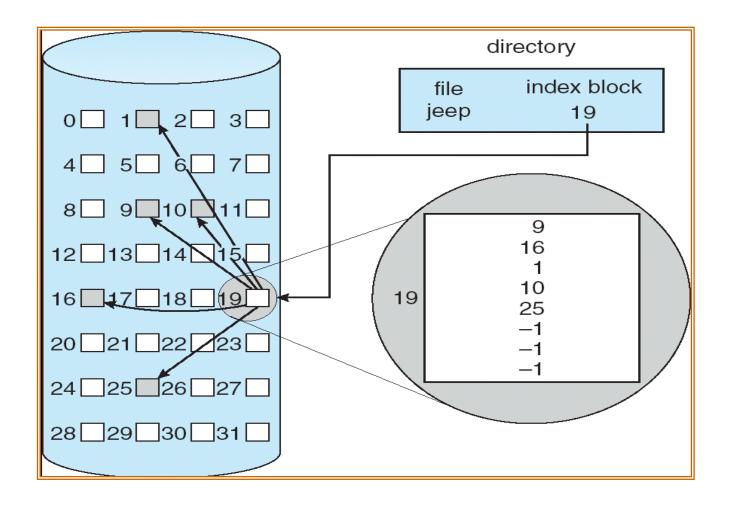
File System Implementation Allocation Methods: Indexed allocation (1)

- ➤ Brings all pointers together into the *index block*.
- ➤ Logical view.
- ➤ Need index table
- ➤ Random access
- >Dynamic access without external fragmentation, but have overhead of index block.





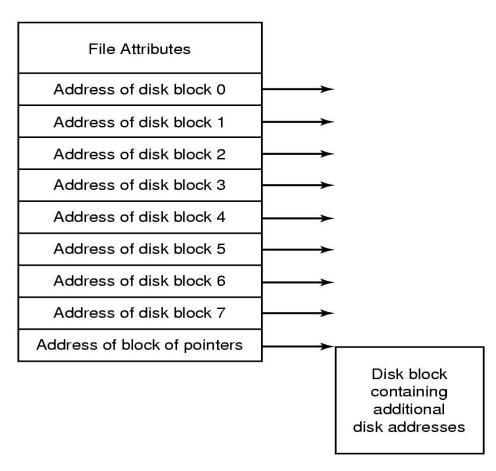
File System Implementation Allocation Methods: Indexed allocation (2)





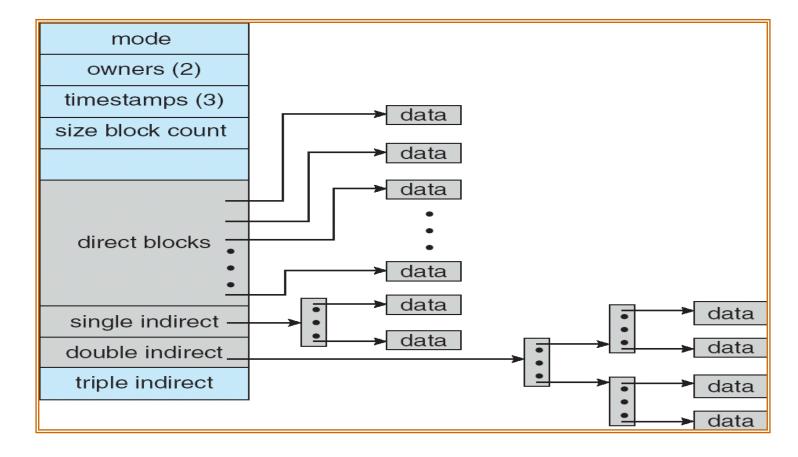
File System Implementation Allocation Methods: Indexed allocation (3)

An example i-node



File System Implementation Allocation Methods: Indexed allocation (4)

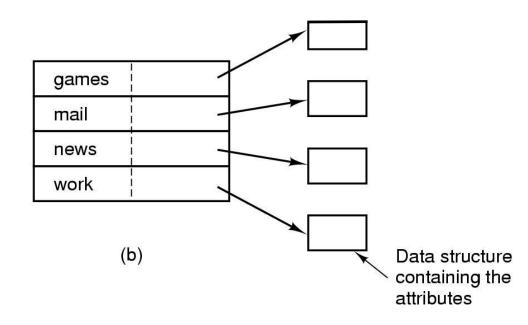
➤ Combined Scheme: UNIX (4K bytes per block)



Implementing Directories (1)

- (a) A simple directory
 fixed size entries
 disk addresses and attributes in directory entry
- (b) Directory in which each entry just refers to an i-node

games	attributes		
mail	attributes		
news	attributes		
work	attributes		
work	¦ attributes		
(a)			
1	(4)		





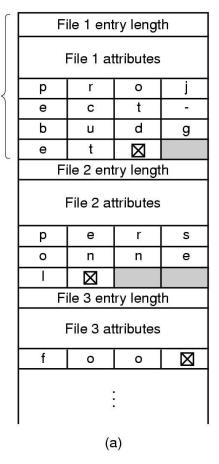
Implementing Directories (2)

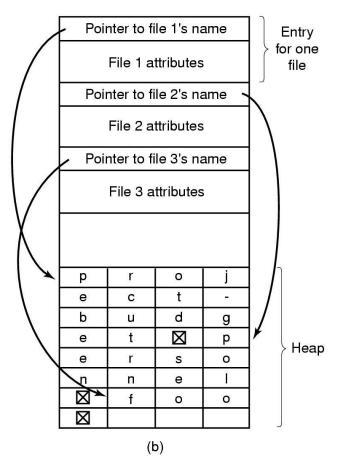
Two ways of handling long file names in directory

Entry

for one

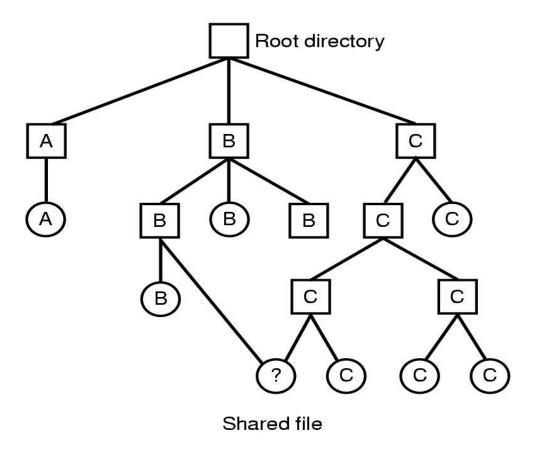
- (a) In-line
- (b) In a heap





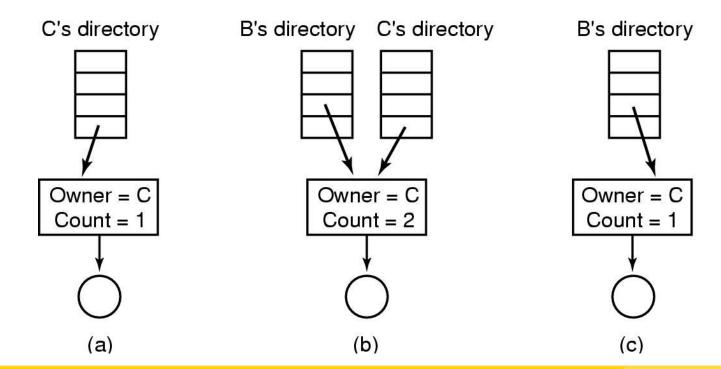


File system containing a shared file



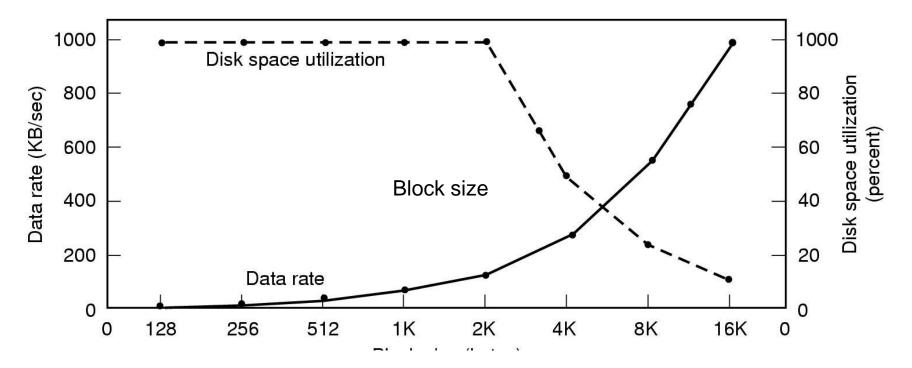
Shared Files (2)

- (a) Situation prior to linking
- (b) After the link is created
- (c)After the original owner removes the file



Disk Space Management (1)

- ➤ Dark line (left hand scale) gives data rate of a disk
- ➤ Dotted line (right hand scale) gives disk space efficiency
- ➤ All files 2KB

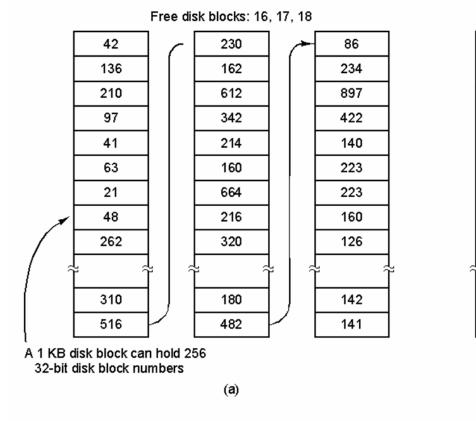


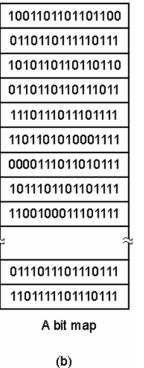


Disk Space Management (2)

(a) Storing the free list on a linked list

(b) A bit map





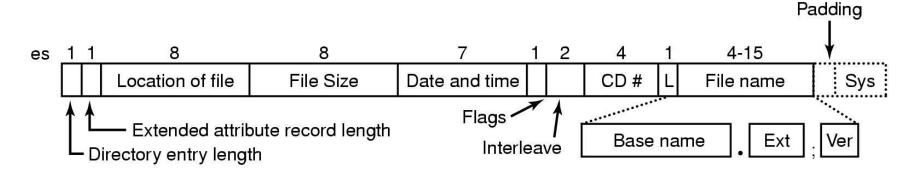


Example File Systems



Example File Systems CD-ROM File Systems

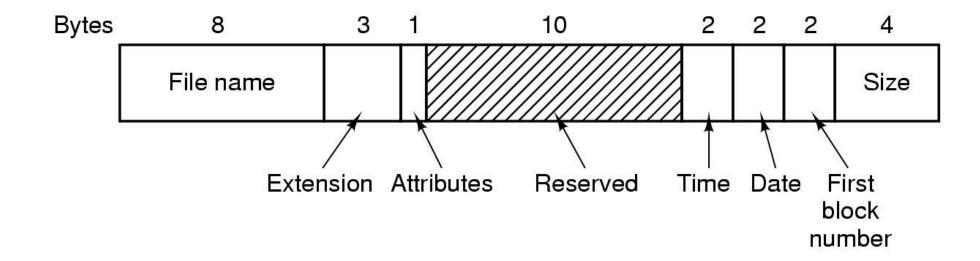
The ISO 9660 directory entry





The MS-DOS File System (1)

The MS-DOS directory entry



The MS-DOS File System (2)

- ➤ Maximum partition for different block sizes
- The empty boxes represent forbidden combinations

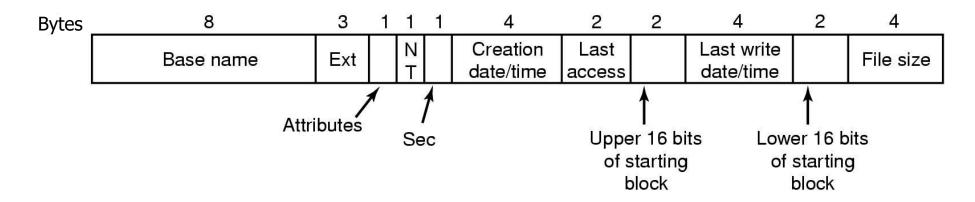
Block size	FAT-12	FAT-16	FAT-32
0.5 KB	2 MB		
1 KB	4 MB		
2 KB	8 MB	128 MB	
4 KB	16 MB	256 MB	1 TB
8 KB		512 MB	2 TB
16 KB		1024 MB	2 TB
32 KB		2048 MB	2 TB



The Windows 98 File System (1)

The extended MOS-DOS

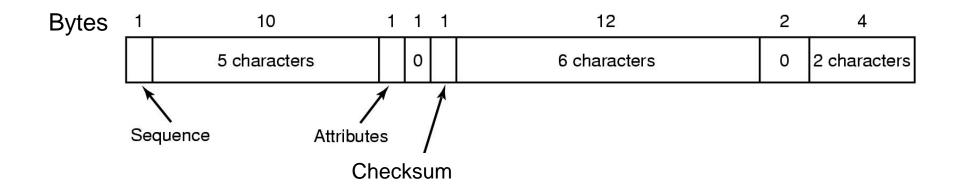
directory entry used in Windows 98





The Windows 98 File System (2)

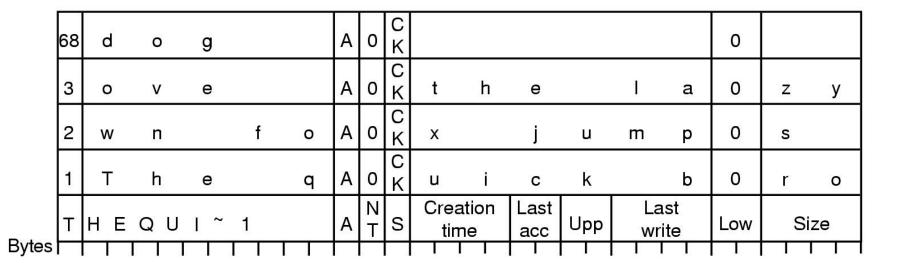
An entry for (part of) a long file name in Windows 98





The Windows 98 File System (3)

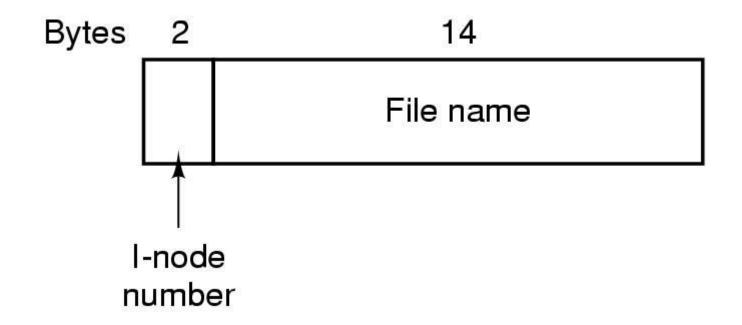
An example of how a long name is stored in Windows 98





The UNIX V7 File System (1)

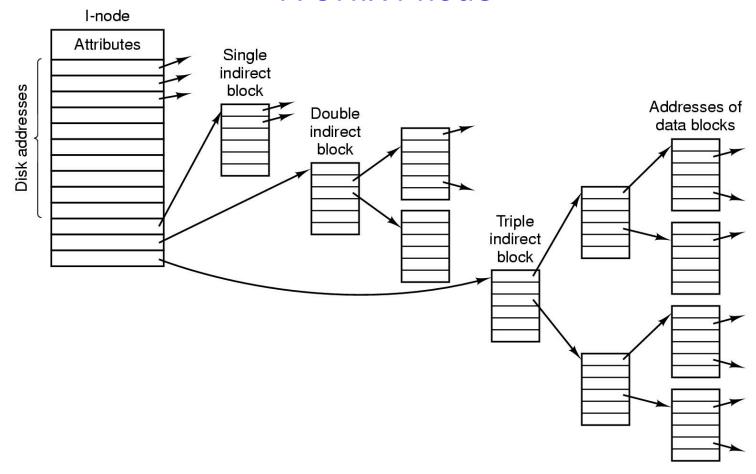
A UNIX V7 directory entry





The UNIX V7 File System (2)

A UNIX i-node





The UNIX V7 File System (3)

The steps in looking up /usr/ast/mbox

Root directory

	. toot an ootory	
1		
1	Total	
4	bin	
7	dev	
14	lib	
9	etc	
6	usr	
8	tmp	

Looking up usr yields i-node 6 I-node 6 is for /usr

15 101 7451
Mode size times
132

I-node 6 says that /usr is in block 132 Block 132 is /usr directory

6	•
1	•
19	dick
30	erik
51	jim
26	ast
45	bal

/usr/ast is i-node 26 I-node 26 is for /usr/ast

_	, 001, 001	
	Mode size times	
	406	

I-node 26 says that /usr/ast is in block 406 Block 406 is /usr/ast directory

26	•
6	••
64	grants
92	books
60	mbox
81	minix
17	src

/usr/ast/mbox is i-node 60



SUMMARY

- > Files
- Directories
- > File system implementation
- > Example file systems