3.Growth of Functions

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3.1 Asymptotic notation

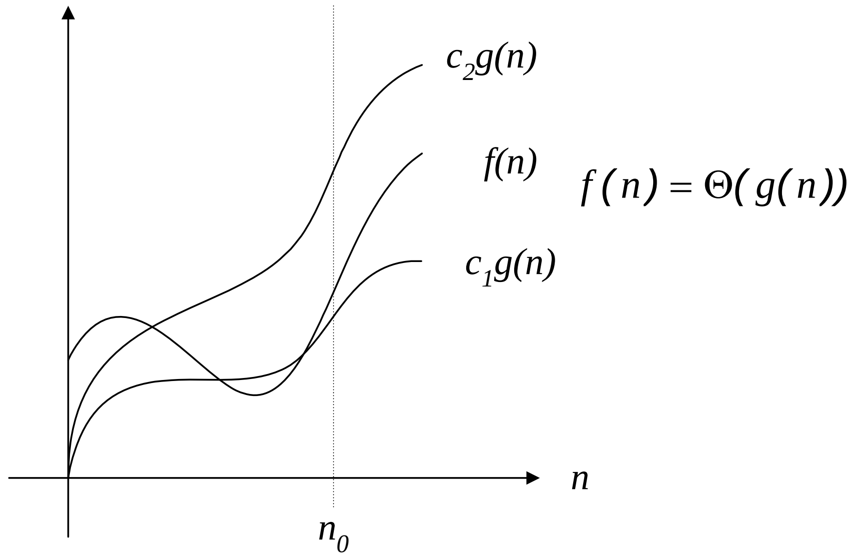
$$\Theta(g(n)) = \{ f(n) \mid \exists c_1, c_2, n_0 \text{ s.t. } 0 \le c_1 g(n) \le f(n) \le c_2 g(n)$$
$$\text{for all } n \ge n_0 \}$$

$$f(n) = \Theta(g(n))$$

$\Rightarrow$ $g(n)$ is an asymptotic tight bound for $f(n)$.

- ``=’’ abuse

- The definition of required every member of be asymptotically nonnegative.



Example:

$$\frac{n^2}{14} \le \frac{n^2}{2} - 3n \le \frac{n^2}{2} \text{ if } n>7.$$

$$6n^3 \ne \Theta(n^2)$$

$$f(n) = an^2 + bn + c,\ a, b, c \text{ constants, } a>0.$$

$$\Rightarrow f(n)=\Theta(n^2).$$

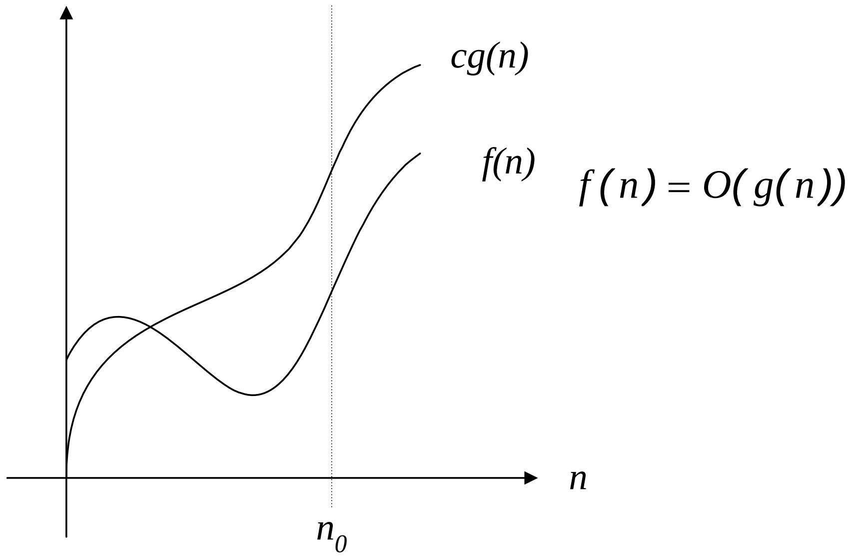
- In general,

$$p(n) = \sum_{i=0}^{d} a_i n^i \text{ where } a_i \text{ are constant with } a_d > 0.$$

$$\text{Then } P(n) = \Theta(n^d).$$

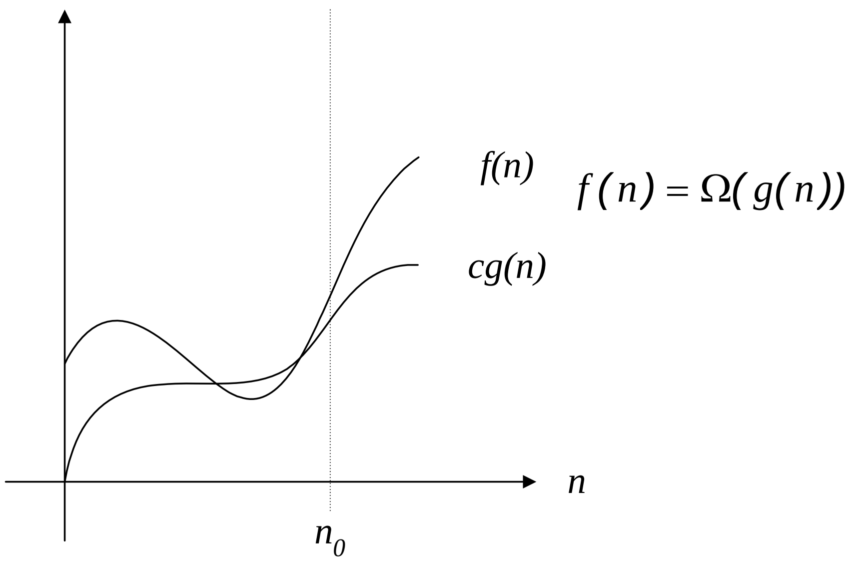
asymptotic upper bound

$$O(g(n)) = \{ f(n) \mid \exists c, n_0 \text{ s.t. } 0 \le f(n) \le cg(n)\ \forall n \ge n_0 \}$$



asymptotic lower bound

$$\Omega(g(n)) = \{ f(n) \mid \exists c, n_0 \text{ s.t. } 0 \le cg(n) \le f(n)\ \forall n \ge n_0 \}$$



Theorem 3.1.

- For any two functions *f*(*n*) and *g*(*n*), $f(n) = \Theta(g(n))$ if and only if $f(n) = O(g(n))$ and $f(n) = \Omega(g(n))$.

- $o(g(n)) = \{ f(n) | \forall c, \exists n_0 \ \forall n > n_0, 0 \leq f(n) \cancel{\leq} cg(n) \}$
- $f(n) = o(g(n)) \Leftrightarrow lim_{n \to \infty} \dfrac{f(n)}{g(n)} = 0$
- $\omega(g(n)) = \{ f(n) | \forall c, \exists n_0 \ \forall n > n_0, 0 \leq cg(n) \cancel{\leq} f(n) \}$
- $f(n) = \omega(g(n)) \Leftrightarrow lim_{n \to \infty} \dfrac{f(n)}{g(n)} = \infty$

- Transitivity

$$f(n)=\Theta(g(n)) \wedge g(n)=\Theta(h(n)) \Rightarrow f(n)=\Theta(h(n))$$
$$f(n)=O(g(n)) \wedge g(n)=O(h(n)) \Rightarrow f(n)=O(h(n))$$
$$f(n)=\Omega(g(n)) \wedge g(n)=\Omega(h(n)) \Rightarrow f(n)=\Omega(h(n))$$
$$f(n)=o(g(n)) \wedge g(n)=o(h(n)) \Rightarrow f(n)=o(h(n))$$
$$f(n)=\omega(g(n)) \wedge g(n)=\omega(h(n)) \Rightarrow f(n)=\omega(h(n))$$

- Reflexivity

$$f(n)=\Theta(f(n))$$
$$f(n)=O(f(n))$$
$$f(n)=\Omega(f(n))$$

- Symmetry

$$f(n)=\Theta(g(n)) \Leftrightarrow g(n)=\Theta(f(n))$$

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- Transpose symmetry

$$f(n)=O(g(n)) \Leftrightarrow g(n)=\Omega(f(n))$$
$$f(n)=o(g(n)) \Leftrightarrow g(n)=\omega(f(n))$$

$$f(n)=O(g(n)) \approx a \le b$$
$$f(n)=\Omega(g(n)) \approx a \ge b$$
$$f(n)=\Theta(g(n)) \approx a = b$$
$$f(n)=o(g(n)) \approx a < b$$
$$f(n)=\omega(g(n)) \approx a > b$$

Trichotomy

- $a < b$, $a = b$, or $a > b$.
- e.g., n, $n^{1+\sin n}$

  n^0 n^2

2.2 Standard notations and common functions

- Monotonicity:
  - A function f is *monotonically increasing* if $m \leq n$ implies $f(m) \leq f(n)$.
  - A function f is *monotonically decreasing* if $m \leq n$ implies $f(m) \geq f(n)$.
  - A function f is *strictly increasing* if $m < n$ implies $f(m) < f(n)$.
  - A function f is *strictly decreasing* if $m > n$ implies $f(m) > f(n)$.

Floor and ceiling

$$x-1<\lfloor x\rfloor\le x\le\lceil x\rceil<x+1$$

$$\lceil n/2\rceil+\lfloor n/2\rfloor=n$$

$$\lceil\lceil n/a\rceil/b\rceil=\lceil n/ab\rceil$$

$$\lfloor\lfloor n/a\rfloor/b\rfloor=\lfloor n/ab\rfloor$$

$$\lceil a/b\rceil\le(a+(b-1))/b$$

$$\lfloor a/b\rfloor\ge(a-(b-1))/b$$

Modular arithmetic

- For any integer a and any positive integer n, the value a mod n is the **remainder** (or **residue**) of the quotient a/n :

  $$a \bmod n = a - \lfloor a/n\rfloor n.$$

- If(a mod n) = (b mod n). We write $a \equiv b$ (mod n) and say that a is **equivalent** to b, modulo n.

- We write $a \not\equiv b$ (mod n) if a is not equivalent to b modulo n.

Polynomials v.s. Exponentials

- Polynomials: $P(n) = \sum_{i=0}^{d} a^i n^i$
  - A function is *polynomial bounded* if $f(n) = n^{O(1)}$.
- Exponentials: $n^b = o(a^n) \quad (a > 1)$
  - Any positive exponential function grows faster than any polynomial.

$$e^x = \sum_{i=0}^{\infty} \frac{x^i}{i!}$$

$$1 + x \le e^x \le 1 + x + \frac{x^2}{2} \text{ if } |x| < 1$$

$$\lim_{n\to\infty} (1 + \frac{x}{n})^n = e^x$$

Logarithms

$$ln(1+x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \frac{x^5}{5} - \ldots \text{ if } |x| < 1$$

$$\frac{x}{1+x} \le ln(1+x) < x$$

- A function $f(n)$ is *polylogarithmically bounded* if $f(n) = \log^{O(1)} n$
- $\log^b n = o(n^a)$ for any constant $a > 0$.
- Any positive polynomial function grows faster than any polylogarithmic function.

Factorials

- Stirling's approximation

$$n! = \sqrt{2\pi n}\left(\frac{n}{e}\right)^n \left(1+\Theta\left(\frac{1}{n}\right)\right)$$

$$n! = o(n^n)$$

$$n! = \omega(2^n)$$

$$\log(n!) = \Theta(n \log n)$$

$$n! = \sqrt{2\pi n}\left(\frac{n}{e}\right)^n e^{\alpha_n}$$

where

$$\frac{1}{12n+1} < \alpha_n < \frac{1}{12n}$$

Function iteration

$$f^{(i)}(n) = \begin{cases} n & \text{if } i = 0, \\ f(f^{(i-1)}(n)) & \text{if } i > 0. \end{cases}$$

For example, if $f(n) = 2n$, then $f^{(i)}(n) = 2^i n$

The iterative logarithm function

$$\lg^{(i)}(n) = \begin{cases} n & \text{if } i = 0 \\ \lg(\lg^{(i-1)} n) & \text{if } i > 0 \text{ and } \lg^{(i-1)} n > 0 \\ \text{undefined} & \text{if } i > 0 \text{ and } \lg^{(i-1)} n \le 0 \\ & \text{or } \lg^{(i-1)} n \text{ is undefined.} \end{cases}$$

$$lg^*(n) = min\{i \ge 0 \mid lg^{(i)} \le 1\}$$

$$lg^* 2 = 1$$
$$lg^* 4 = 2$$
$$lg^* 16 = 3$$
$$lg^* 65536 = 4$$
$$lg^* 2^{65536} = 5$$

- Since the number of atoms in the observable universe is estimated to be about 10^{80}, which is much less than 2^{65536}, we rarely encounter a value of *n* such that $\lg^* n > 5$.

Fibonacci numbers

$$F_0 = 0$$

$$F_1 = 1$$

$$F_i = F_{i-1} + F_{i-2}$$

$$F_i = \frac{\phi^i - \hat{\phi}^i}{\sqrt{5}}$$

$$\phi = \frac{1 + \sqrt{5}}{2} = 1.61803...$$

$$\hat{\phi} = \frac{1 - \sqrt{5}}{2} = -0.61803...$$