

國立台灣海洋大學資訊工程學系博士班



96 學年度第一學期博士班資格考命題卷 (筆試)

科目	演算法	命題教授	吳宗杉老師 翁世光老師	日期	97/01/09
					4.4 .1

- 1. Write an algorithm to find the median of an unsorted array of n elements with time $\Theta(n)$ in average case (8%). Analyze the complexity (average case) of the algorithm (7%).

 O(λ) $T(k) = T(\frac{1}{2}) + \lambda$
- 2. Briefly answer the following questions: (15%)

Cp: N++2+ ~

- (a) Main recurrence theorem
- (b) Supremum
- (c) Optimal substructure property

- M(1-1/h) = 2N
- 3. Compare dynamic programming with divide-and-conquer technique. (10%)
- 4. For the text searching problem, it can be solved by brute force sequential searching in time O(|p| * |t|), where |p| and |t| stand for the length of pattern p and text t, respectively. The Knuth-Morris-Pratt algorithm solves the pattern-matching problem in time O(|p| + |t|). Demonstrate how the Knuth-Morris-Pratt works with pattern p = "pappar" and text t = "pappapapapaparrassanuaragh". (10%)
- 5. Prove that the maximum height of a red-black tree with N nodes is O(2logN) while the minimum height is O(logN). (Hint: consider the maximum and minimum heights of 2-3-4 trees.) (10%)
- 6. When performing hashing by using a hash function based on division method, we have to select the radix base d and the hash table size M carefully. Why should d and M be relatively prime? (Explain the reasons according to the space spanned by the terms: i*d % M and dⁱ % M.) (15%)

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$$i*d\% M$$
 and $d^i\% M$.) (15%)

$$h(x) = (a_i d^i + a_{i-1} d^{i-1} + ... + a_2 d^2 + a_1 d + a_0)\% M$$

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- 7. Let G=(V, E) be a weighted graph, and G is connected. Select vertex <u>V0</u> as the <u>source</u> and construct a shortest path tree of G. Assume the tree is <u>Tshort</u>. Then construct a minimum spanning tree of G and assume the minimum spanning tree is <u>Tmin</u>. Prove that the total edge weights of <u>Tmin</u> are always less than or equal to that of <u>Tshort</u>. Please give an example graph in which both trees have the same total weights. (10%)
- 8. Let A[N][N] be a two dimensional array with N*N integers. Design a sorting



algorithm with $O(N^2 long N)$ steps to sort A[][] such that all elements of A[][] are sorted in vertical, horizontal and diagonal directions. That is A[0][0] is the smallest element and A[N-1][N-1] the largest element. You cannot copy A[][] to an 1-D array, perform the sorting, and copy it back. Instead, your algorithm have to sort A[][] in place. (Hint: use quick-sort. However, sort A[][] along vertical and horizontal directions alternatively. Otherwise, you can enumerate A[][] in a zig-zag manner during the partition stage and split A[][] into 4 blocks.) (15%)

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