

Dynamic Control and Optimization of Wireless Virtual Networks

Quang-Trung Luu

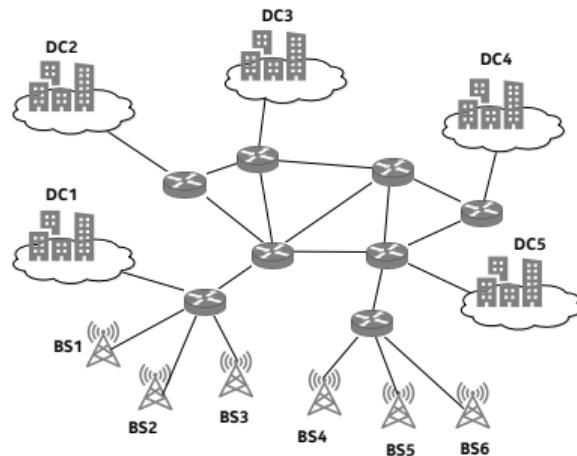
Université Paris-Saclay, CentraleSupélec, CNRS, L2S and Nokia Bell Labs

Advisors

Michel Kieffer, Alexandre Mouradian (Université Paris-Saclay)
Sylvaine Kerboeuf (Nokia Bell Labs)



CentraleSupélec, Gif-sur-Yvette, June 7, 2021

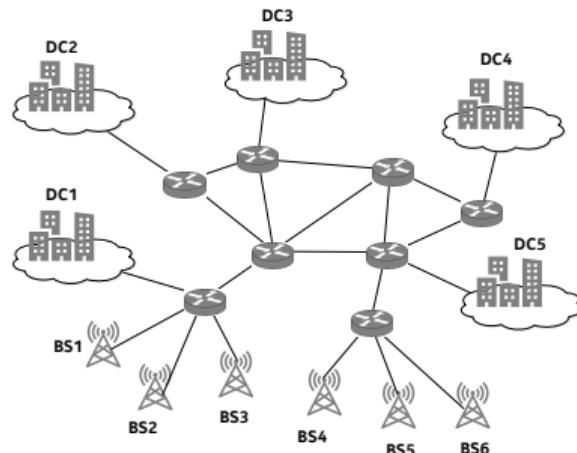


Past: physical dedicated devices: switches, routers, etc.

Network virtualization era:

- SDN (2008): decouples the control and data plane
- NFV (2012): virtualizes network functions → VNF

Network slicing (2015): leverages SDN and NFV [NGMN'15]

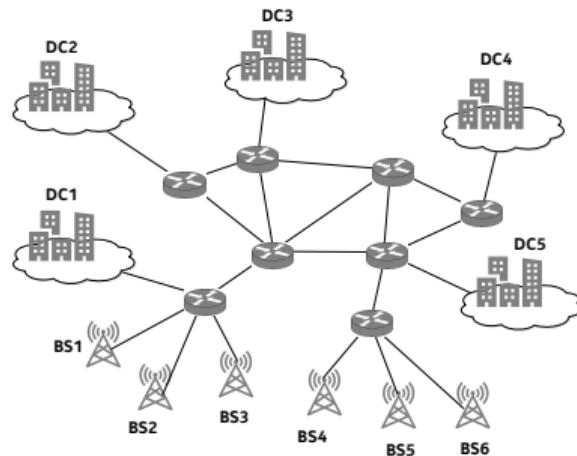


Past: physical dedicated devices: switches, routers, etc.

Network virtualization era:

- SDN (2008): decouples the control and data plane
- NFV (2012): virtualizes network functions → VNF

Network slicing (2015): leverages SDN and NFV [NGMN'15]

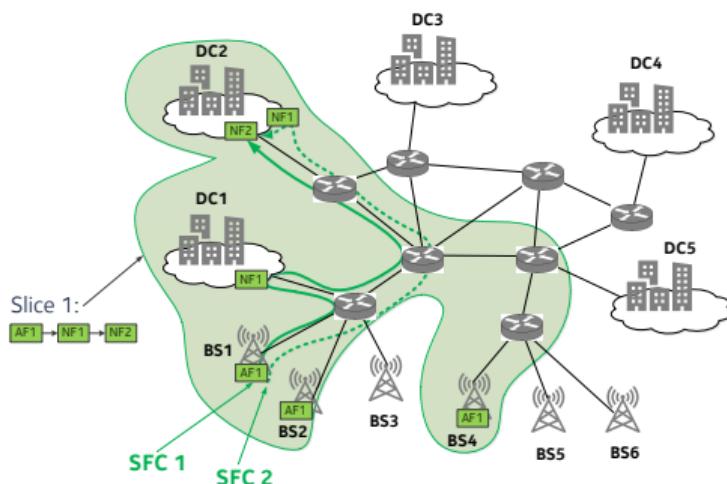


Past: physical dedicated devices: switches, routers, etc.

Network virtualization era:

- SDN (2008): decouples the control and data plane
- NFV (2012): virtualizes network functions → VNF

Network slicing (2015): leverages SDN and NFV [NGMN'15]



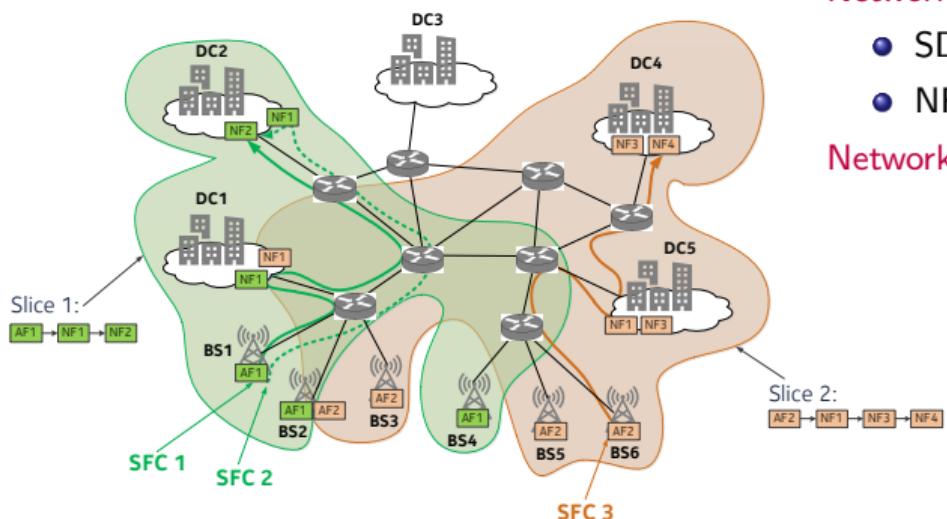
Past: physical dedicated devices: switches, routers, etc.

Network virtualization era:

- SDN (2008): decouples the control and data plane
- NFV (2012): virtualizes network functions → VNF

Network slicing (2015): leverages SDN and NFV [NGMN'15]

SDN: Software Defined Networking NFV: Network Function Virtualization VNF: Virtual Network Function



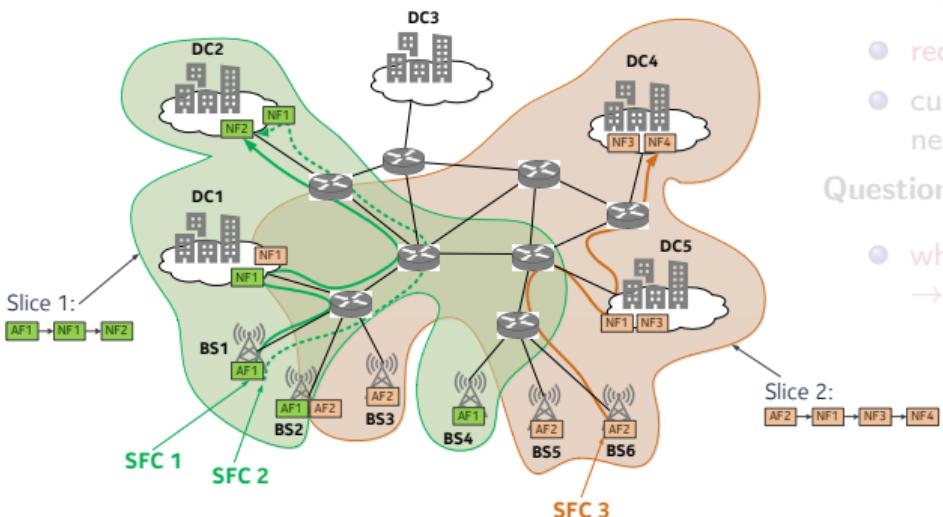
Past: physical dedicated devices: switches, routers, etc.

Network virtualization era:

- SDN (2008): decouples the control and data plane
- NFV (2012): virtualizes network functions → VNF

Network slicing (2015): leverages SDN and NFV [NGMN'15]

SDN: Software Defined Networking NFV: Network Function Virtualization VNF: Virtual Network Function



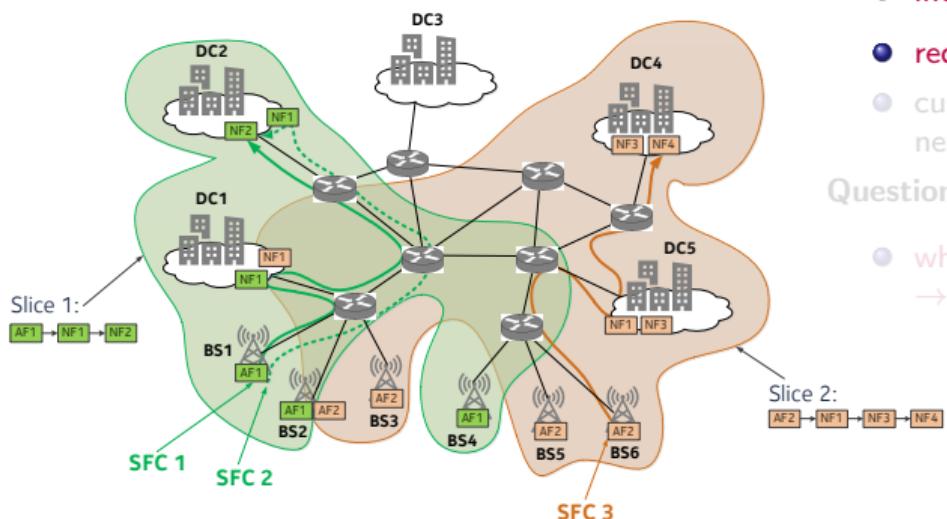
Advantages:

- increase flexibility in network operations [Rost'17]
- reduce overall equipment and management costs [Liang'14]
- customers can manage their own apps by exploiting built-in network slices tailored to their needs [GSM'17]

Question:

- where to place VNFs to deploy SFCs (slice)
→ embedding of slice's SFCs

SDN: Software Defined Networking NFV: Network Function Virtualization VNF: Virtual Network Function



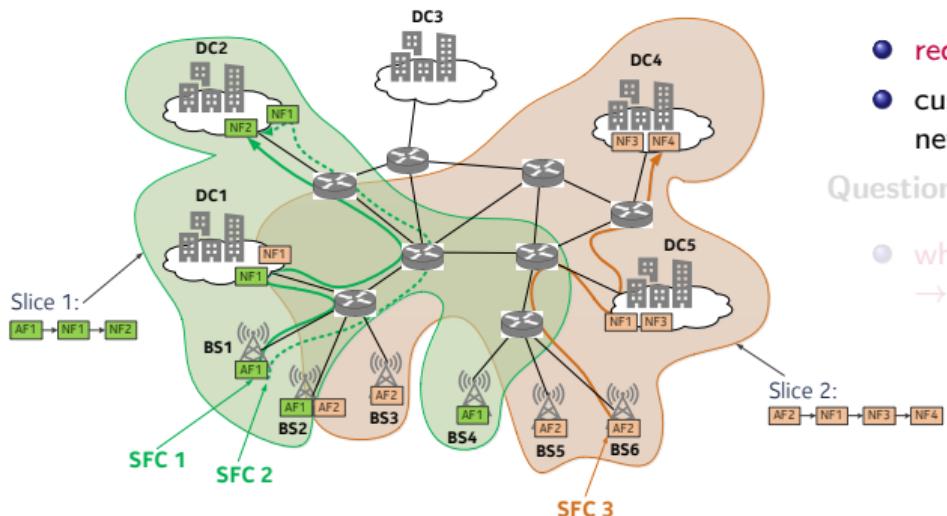
Advantages:

- increase flexibility in network operations [Rost'17]
- reduce overall equipment and management costs [Liang'14]
- customers can manage their own apps by exploiting built-in network slices tailored to their needs [GSM'17]

Question:

- where to place VNFs to deploy SFCs (slice)
→ embedding of slice's SFCs

SDN: Software Defined Networking NFV: Network Function Virtualization VNF: Virtual Network Function



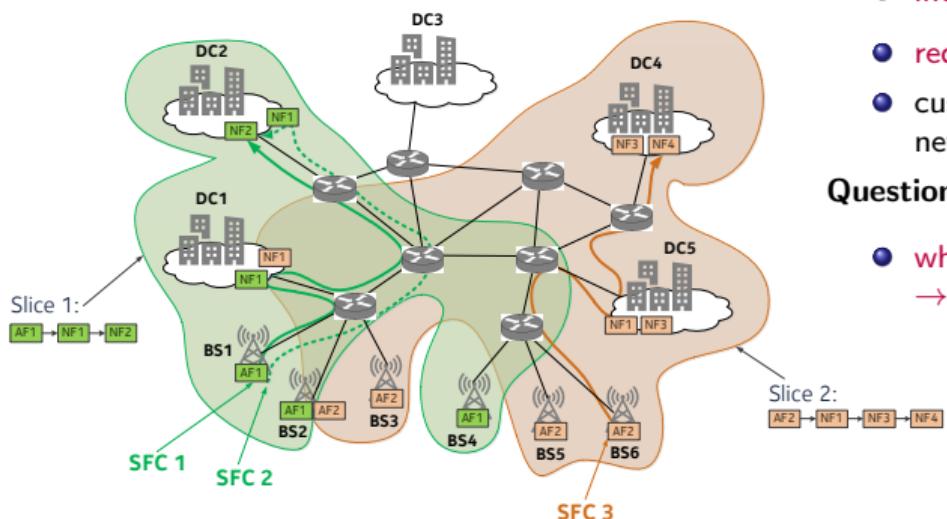
Advantages:

- increase flexibility in network operations [Rost'17]
- reduce overall equipment and management costs [Liang'14]
- customers can manage their own apps by exploiting built-in network slices tailored to their needs [GSM'17]

Question:

- where to place VNFs to deploy SFCs (slice)
→ embedding of slice's SFCs

SDN: Software Defined Networking NFV: Network Function Virtualization VNF: Virtual Network Function



Advantages:

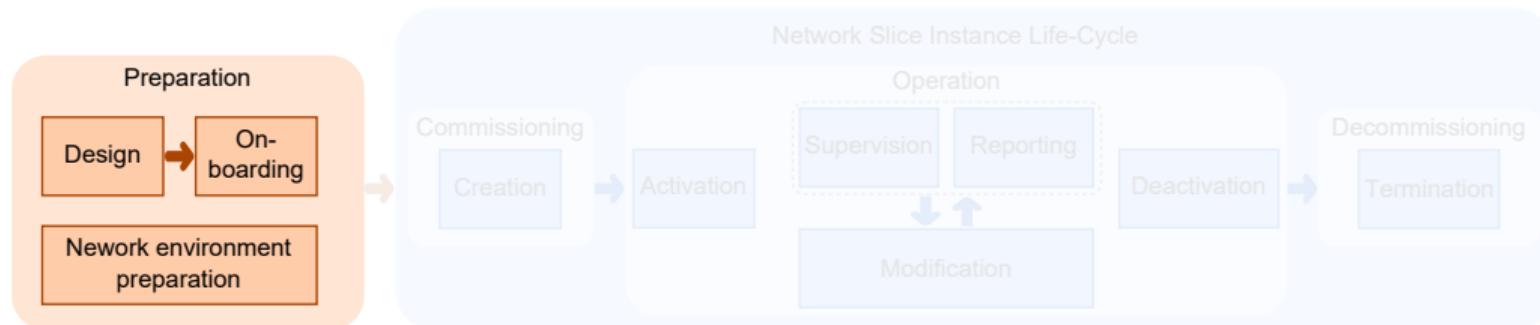
- increase flexibility in network operations [Rost'17]
- reduce overall equipment and management costs [Liang'14]
- customers can manage their own apps by exploiting built-in network slices tailored to their needs [GSM'17]

Question:

- where to place VNFs to deploy SFCs (slice)
→ embedding of slice's SFCs

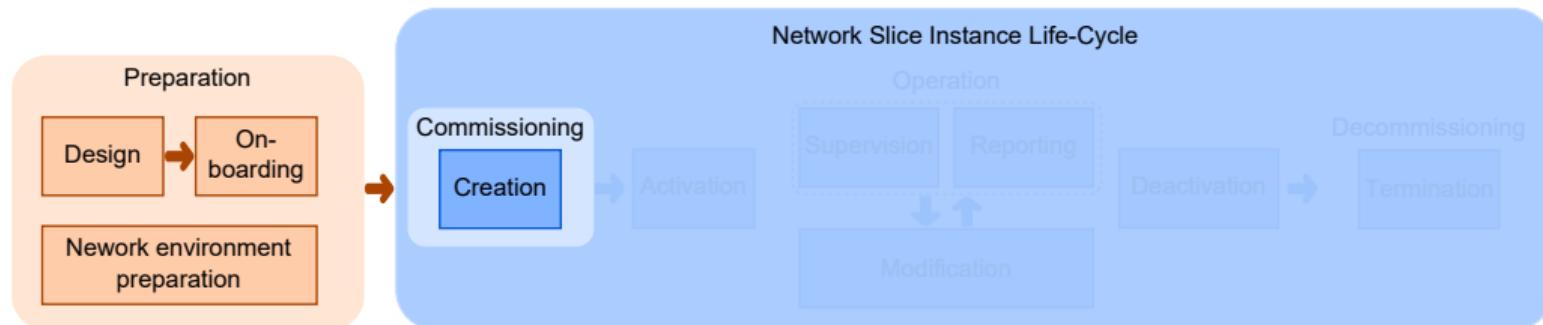
SDN: Software Defined Networking NFV: Network Function Virtualization VNF: Virtual Network Function

3GPP view on network slicing managements aspects



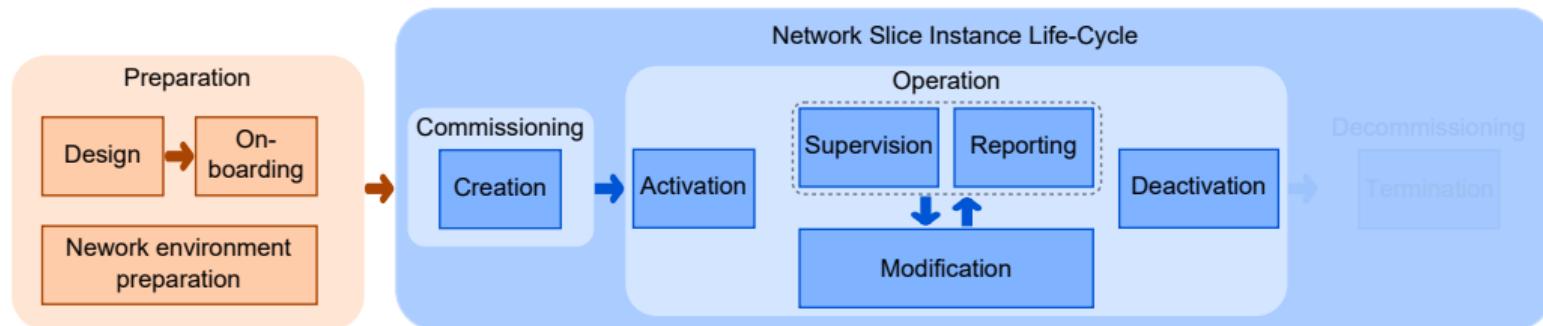
3GPP (2020). Management and Orchestration; Concepts, Use Cases and Requirements. 3GPP TS 28.530 V17.0.0.

3GPP view on network slicing managements aspects



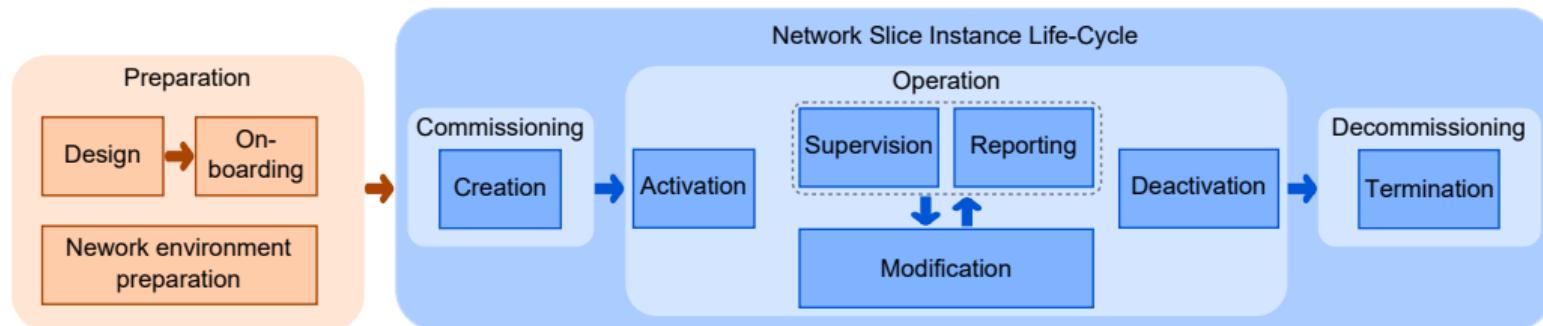
3GPP (2020). Management and Orchestration; Concepts, Use Cases and Requirements. 3GPP TS 28.530 V17.0.0.

3GPP view on network slicing managements aspects



3GPP (2020). Management and Orchestration; Concepts, Use Cases and Requirements. 3GPP TS 28.530 V17.0.0.

3GPP view on network slicing managements aspects



3GPP (2020). Management and Orchestration; Concepts, Use Cases and Requirements. 3GPP TS 28.530 V17.0.0.

- SFC resource allocation (embedding): [Riggio'16], [Vizarreta'17], [Halabian'19]
- Slice resource allocation with coverage constraints: [Lee'16], [Chatterjee'18], [Teague'19]
- Uncertainty-aware slice resource allocation: [Huin'17], [Mireslami'19], [Fendt'19]
- Slice admission control: [Noroozi'19]
 - ▶ SAC with dynamic slice requests: [Bega'17], [Bega'20], [Ebrahimi'20], [Han'20]
- Slice resource provisioning
 - ▶ joint slice resource provisioning and allocation: [Xiong'19], [Sun'19] (radio resource only)

- SFC resource allocation (embedding): [Riggio'16], [Vizarreta'17], [Halabian'19]
- Slice resource allocation with coverage constraints: [Lee'16], [Chatterjee'18], [Teague'19]
- Uncertainty-aware slice resource allocation: [Huin'17], [Mireslami'19], [Fendt'19]
- Slice admission control: [Noroozi'19]
 - ▶ SAC with dynamic slice requests: [Bega'17], [Bega'20], [Ebrahimi'20], [Han'20]
- Slice resource provisioning
 - ▶ joint slice resource provisioning and allocation: [Xiong'19], [Sun'19] (radio resource only)

- SFC resource allocation (embedding): [Riggio'16], [Vizarreta'17], [Halabian'19]
- Slice resource allocation with coverage constraints: [Lee'16], [Chatterjee'18], [Teague'19]
- Uncertainty-aware slice resource allocation: [Huin'17], [Mireslami'19], [Fendt'19]
- Slice admission control: [Noroozi'19]
 - ▶ SAC with dynamic slice requests: [Bega'17], [Bega'20], [Ebrahimi'20], [Han'20]
- Slice resource provisioning
 - ▶ joint slice resource provisioning and allocation: [Xiong'19], [Sun'19] (radio resource only)

- SFC resource allocation (embedding): [Riggio'16], [Vizarreta'17], [Halabian'19]
- Slice resource allocation with coverage constraints: [Lee'16], [Chatterjee'18], [Teague'19]
- Uncertainty-aware slice resource allocation: [Huin'17], [Mireslami'19], [Fendt'19]
- Slice admission control: [Noroozi'19]
 - ▶ SAC with dynamic slice requests: [Bega'17], [Bega'20], [Ebrahimi'20], [Han'20]
- Slice resource provisioning
 - ▶ joint slice resource provisioning and allocation: [Xiong'19], [Sun'19] (radio resource only)

- SFC resource allocation (embedding): [Riggio'16], [Vizarreta'17], [Halabian'19]
- Slice resource allocation with coverage constraints: [Lee'16], [Chatterjee'18], [Teague'19]
- Uncertainty-aware slice resource allocation: [Huin'17], [Mireslami'19], [Fendt'19]
- Slice admission control: [Noroozi'19]
 - ▶ SAC with dynamic slice requests: [Bega'17], [Bega'20], [Ebrahimi'20], [Han'20]
- Slice resource provisioning
 - ▶ joint slice resource provisioning and allocation: [Xiong'19], [Sun'19] (radio resource only)

Challenges addressed:

-  Provision **enough infrastructure resources** to satisfy slice resource demands
-  Account for **coverage constraints**
-  Account for **uncertainties**
-  Account for *in-advance* slice requests (**anticipated** provisioning)

Challenges addressed:

-  Provision enough infrastructure resources to satisfy slice resource demands
-  Account for coverage constraints
-  Account for uncertainties
-  Account for *in-advance* slice requests (*anticipated* provisioning)

Challenges addressed:

-  Provision enough infrastructure resources to satisfy slice resource demands
-  Account for coverage constraints
-  Account for uncertainties
-  Account for *in-advance* slice requests (*anticipated* provisioning)

Challenges addressed:

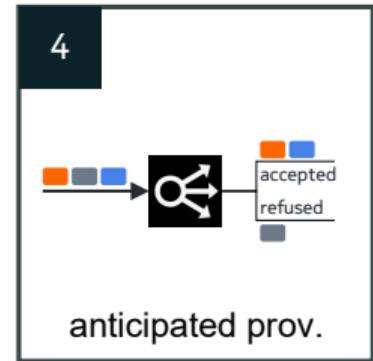
-  Provision enough infrastructure resources to satisfy slice resource demands
-  Account for coverage constraints
-  Account for uncertainties
-  Account for *in-advance* slice requests (**anticipated** provisioning)

Contributions

Part I. Deterministic slice resource demands

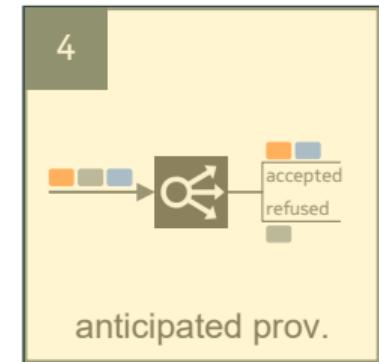
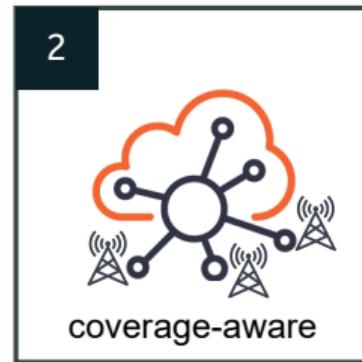
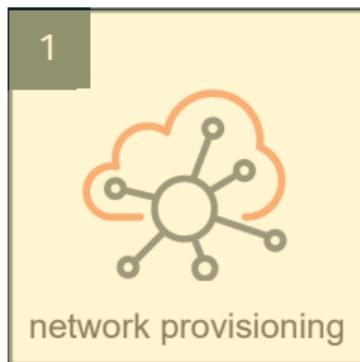


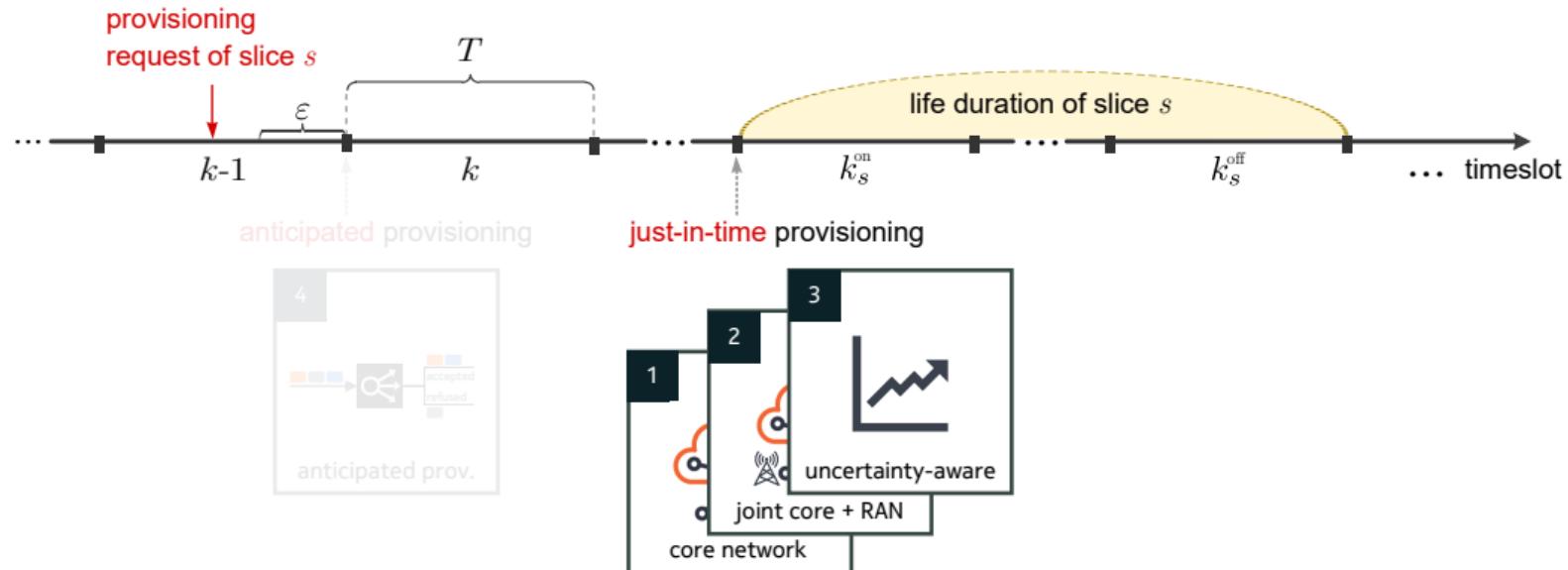
Part II. Random slice resource demands



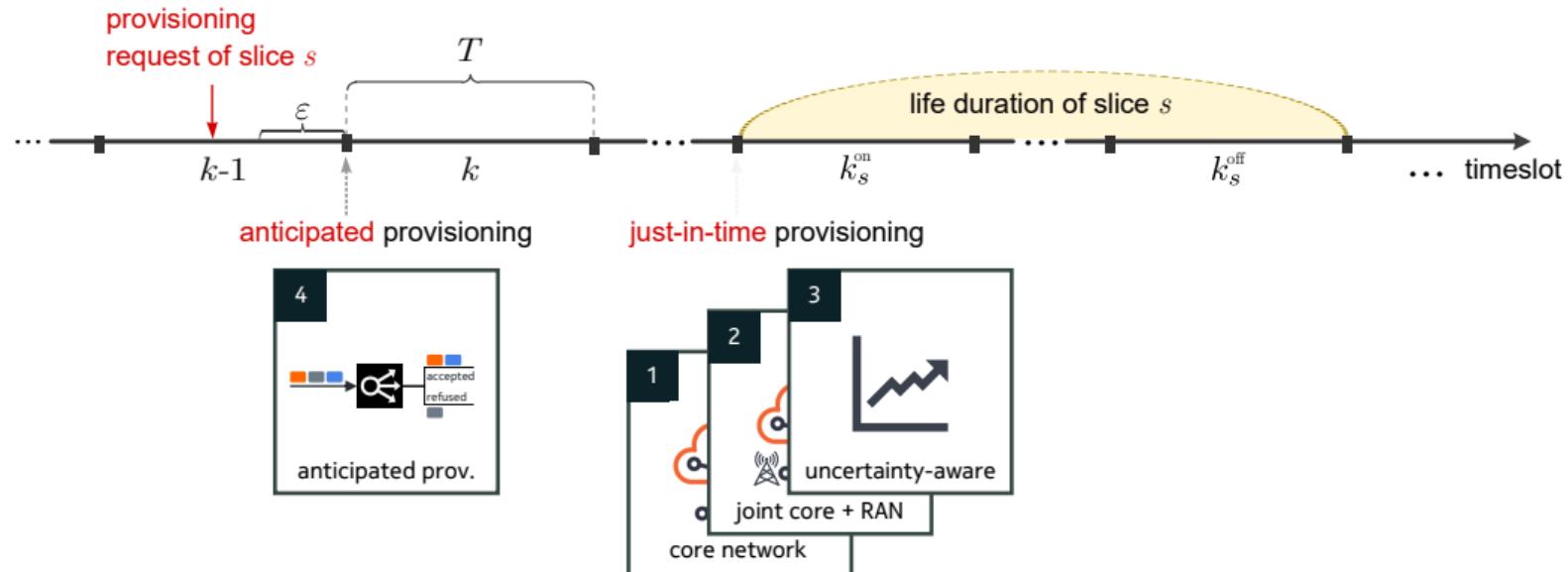
Contributions

Part I. Deterministic slice resource demands





just-in-time: performed only for slices that need to be activated in the next time slot



just-in-time: performed only for slices that need to be activated in the next time slot

anticipated: if not enough resources to satisfy all requests → useful to perform provisioning *in advance* to response to the request

1 System model

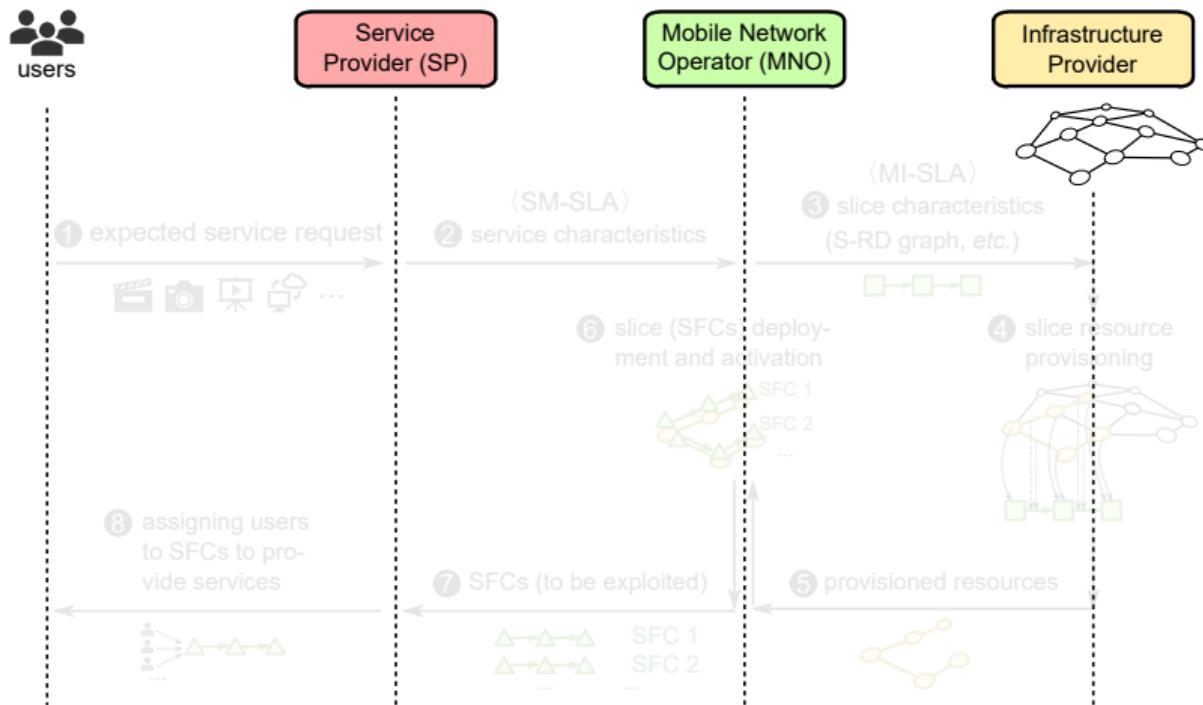
2 Main contributions

- Contribution 1: Network resource provisioning
- Contribution 4: Anticipated provisioning

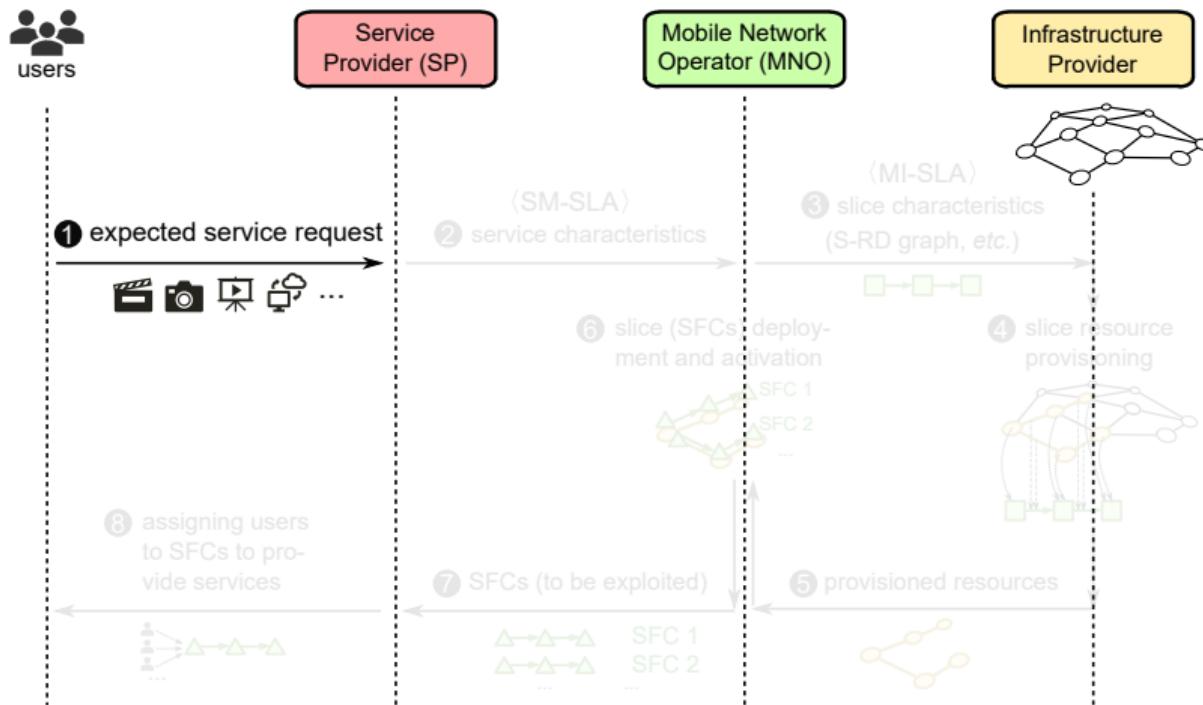
3 Conclusions and Perspective

- Contribution 2: Coverage-aware resource provisioning
- Contribution 3: Uncertainty-aware resource provisioning

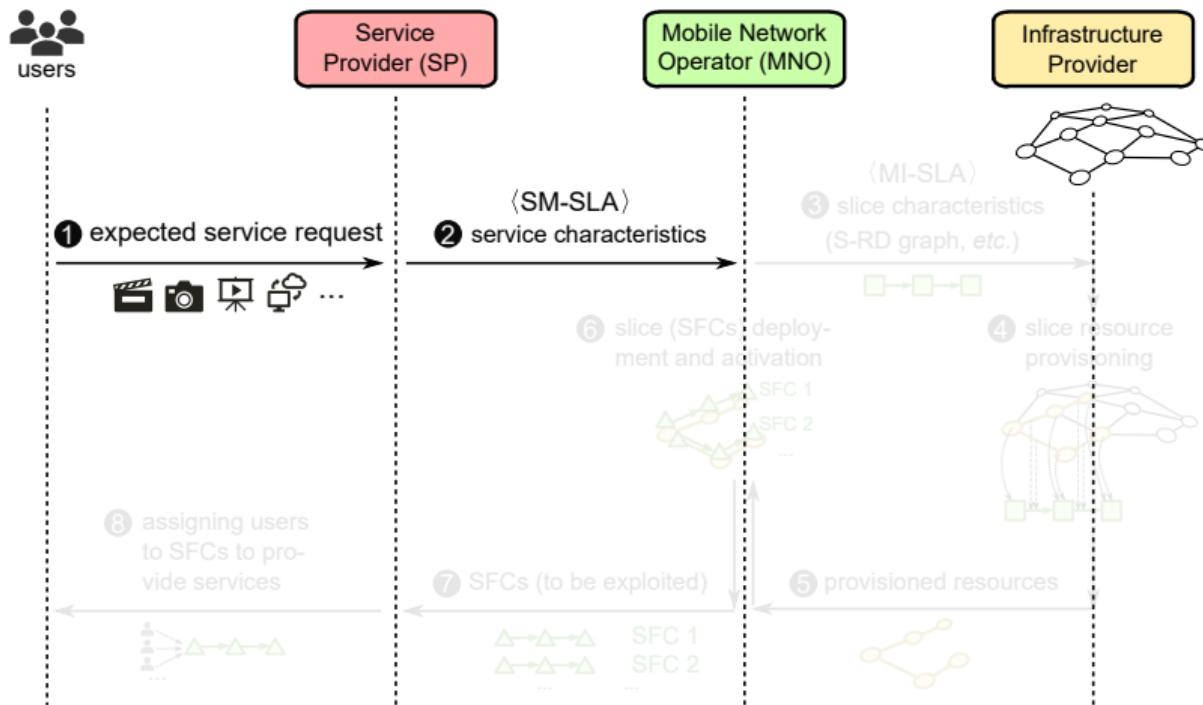
System Model | Slice request processing



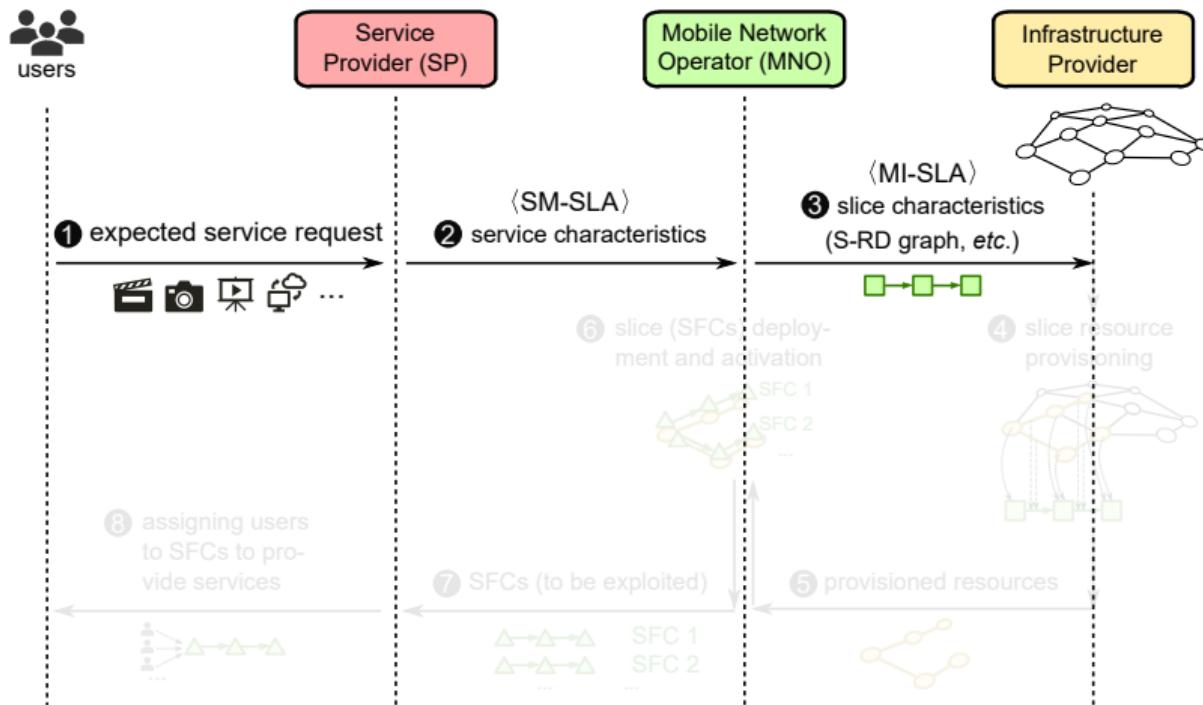
System Model | Slice request processing



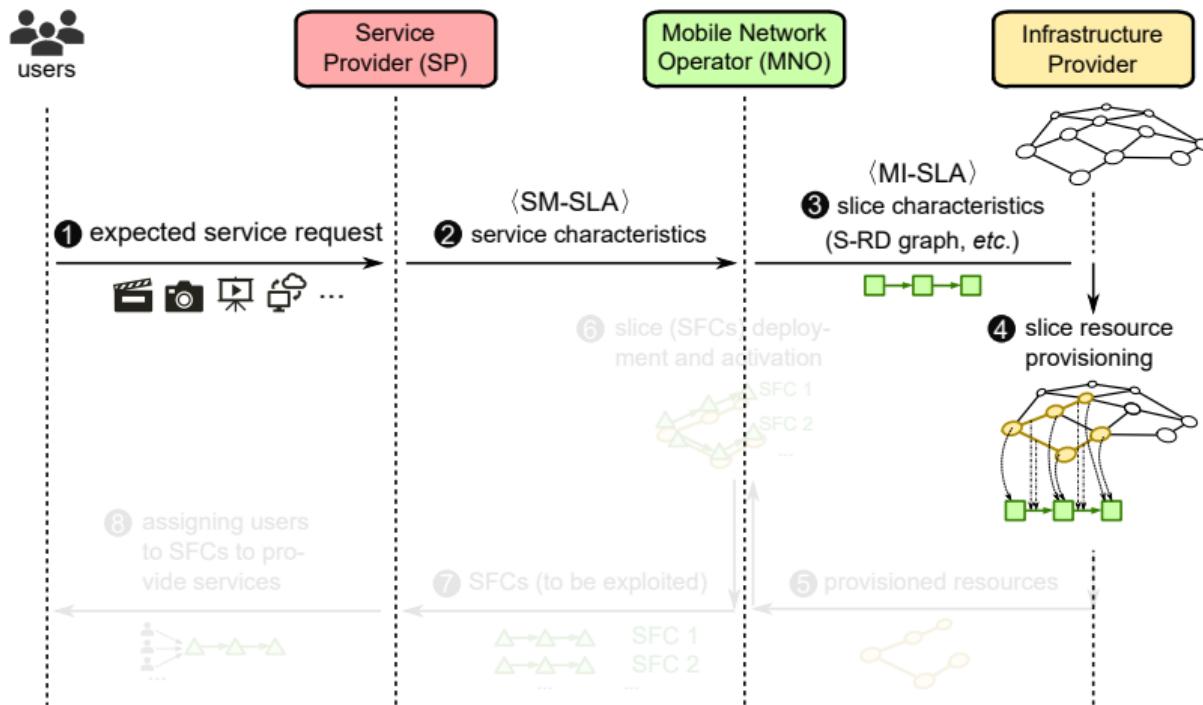
System Model | Slice request processing



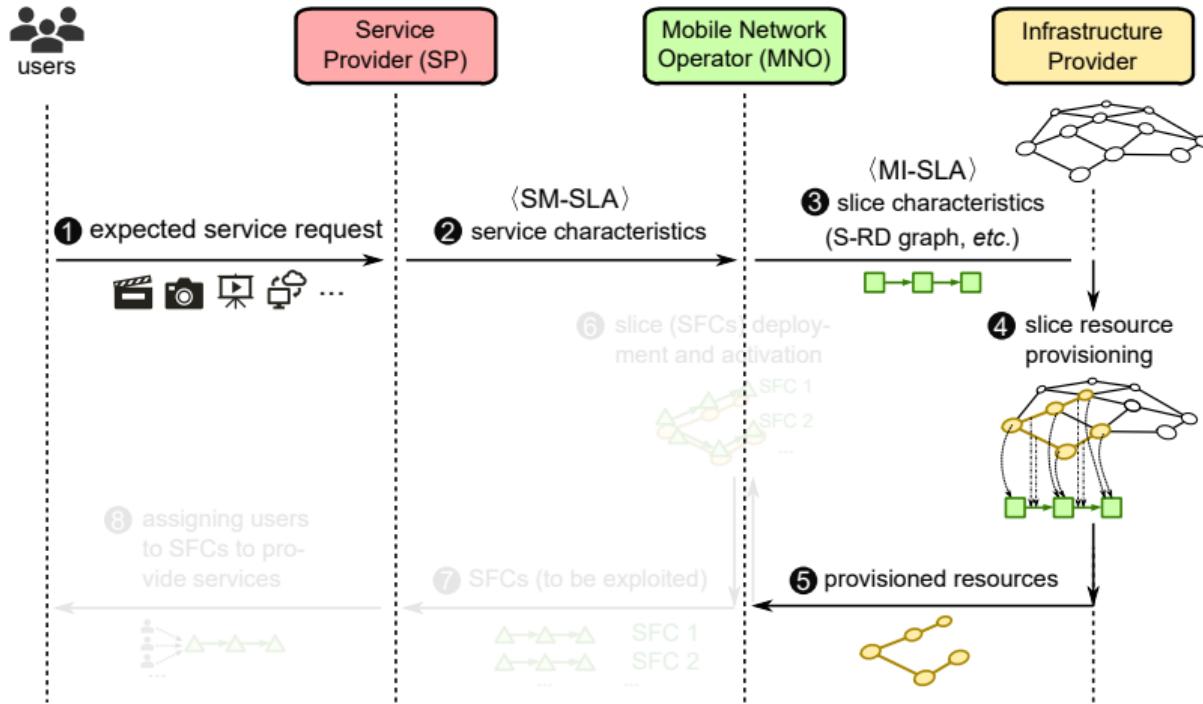
System Model | Slice request processing



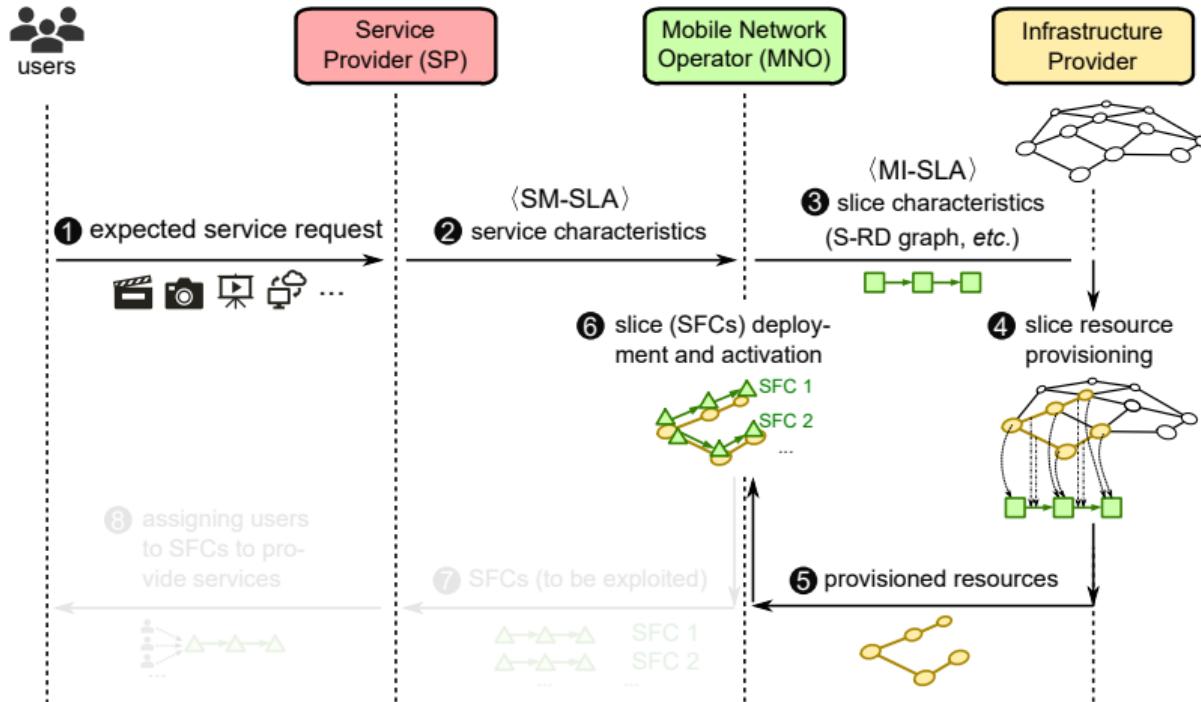
System Model | Slice request processing



System Model | Slice request processing

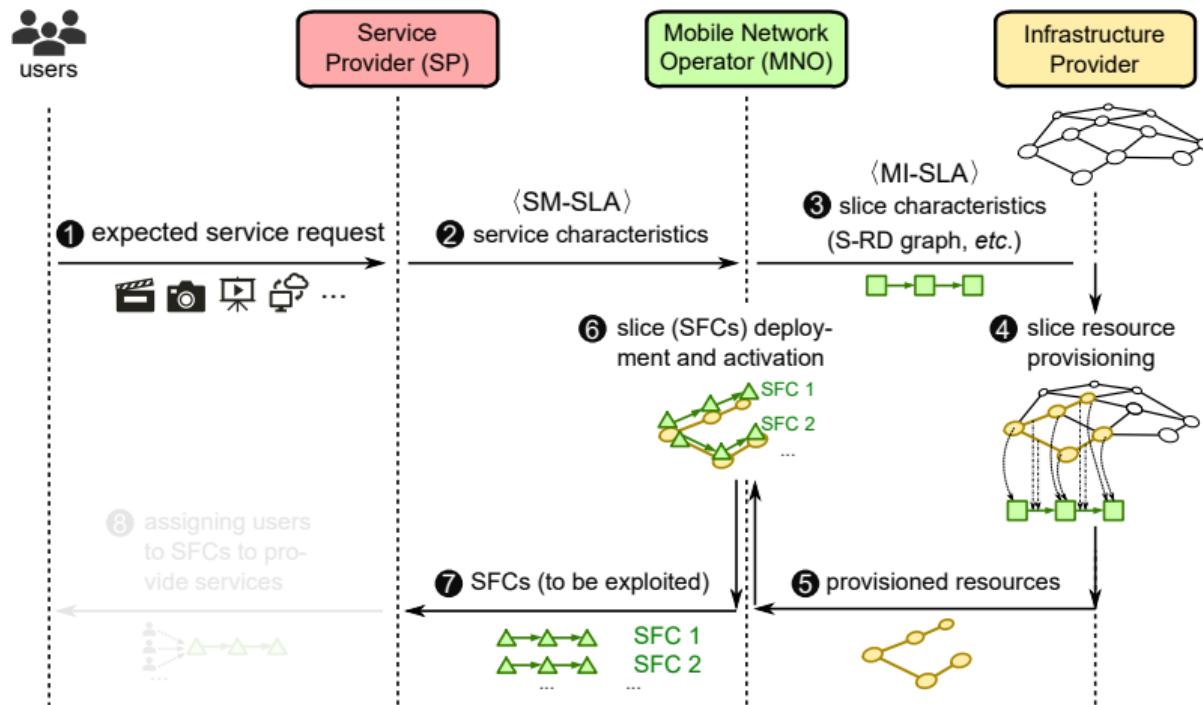


System Model | Slice request processing



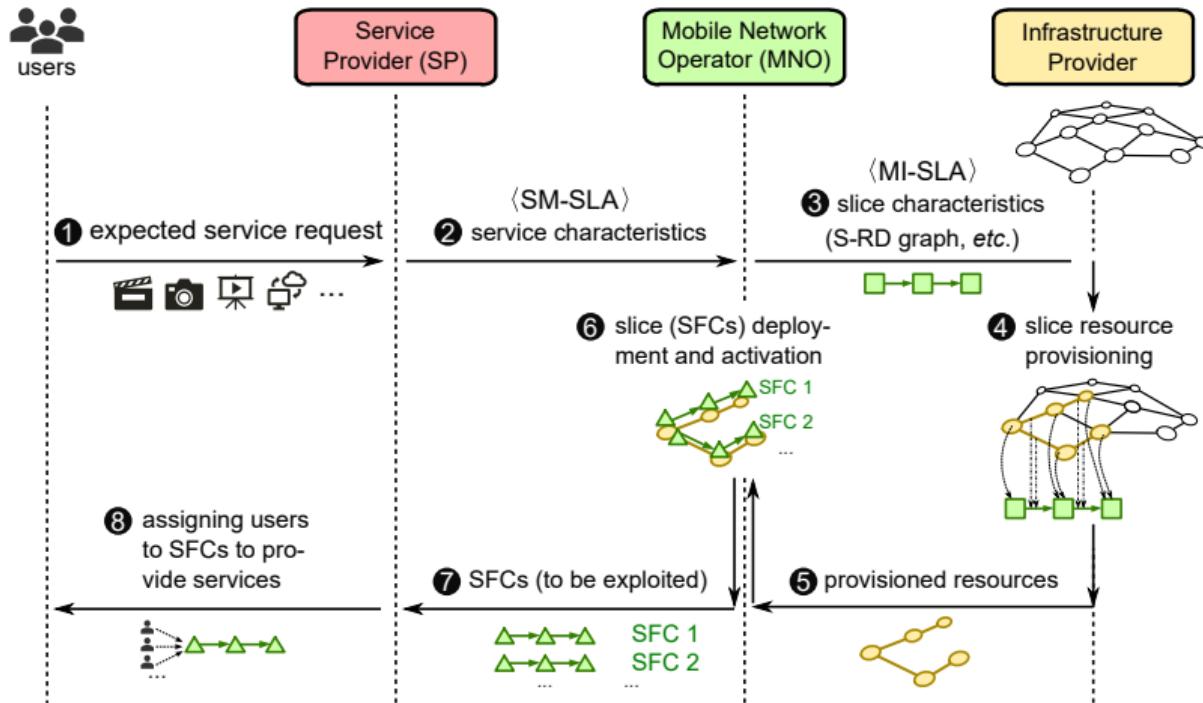
System Model | Slice request processing

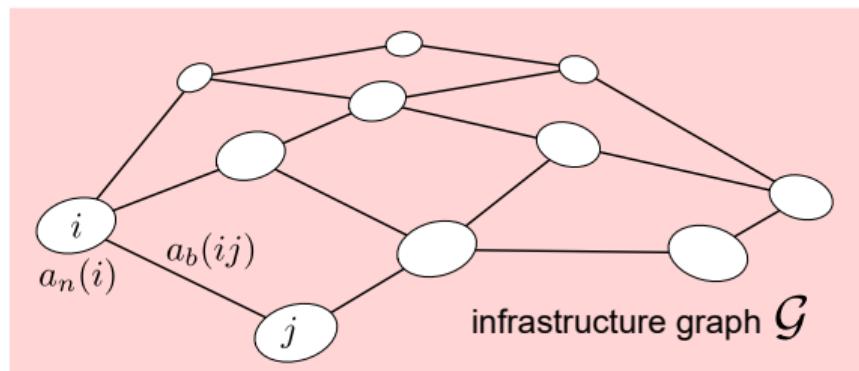
NOKIA Bell Labs



System Model | Slice request processing

NOKIA Bell Labs



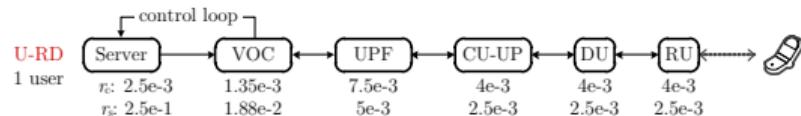


Infrastructure provides resources with associated costs

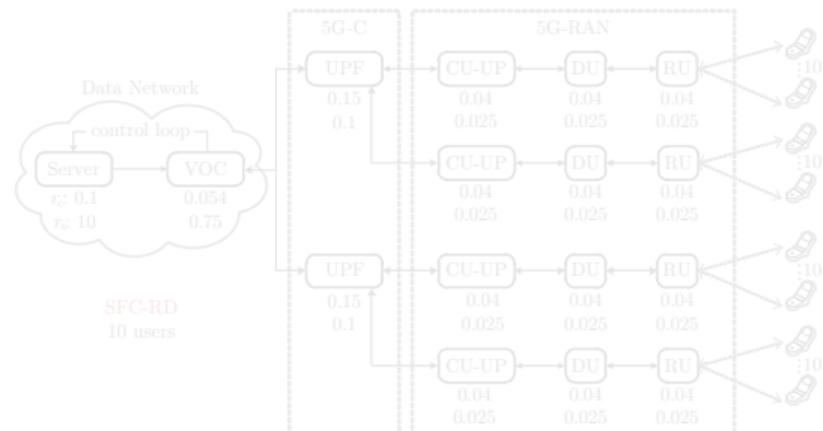
- node $i \in \mathcal{N}$: computing $a_c(i)$, memory $a_m(i)$, wireless $a_w(i)$
- link $ij \in \mathcal{E}$: bandwidth $a_b(ij)$
- costs: $c_c(i)$, $c_m(i)$, $c_w(i)$, $c_b(ij)$

System Model | Virtual graphs of resource demand

Resource demand of a typical user (U-RD): U_s



Resource demand of an SFC (SFC-RD): r_s

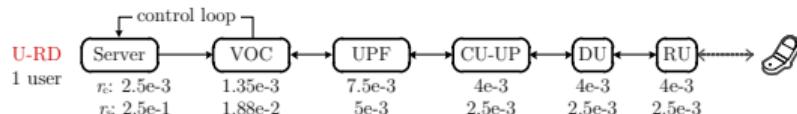


Resource demand of a slice (S-RD): R_s

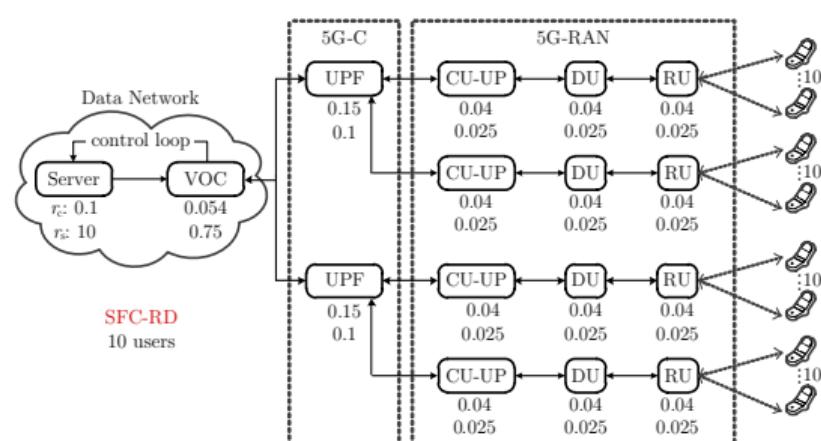


System Model | Virtual graphs of resource demand

Resource demand of a typical user (U-RD): U_s



Resource demand of an SFC (SFC-RD): r_s

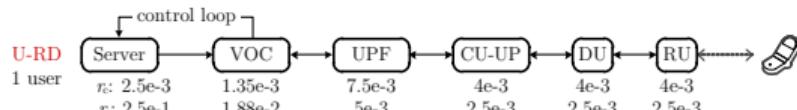


Resource demand of a slice (S-RD): R_s

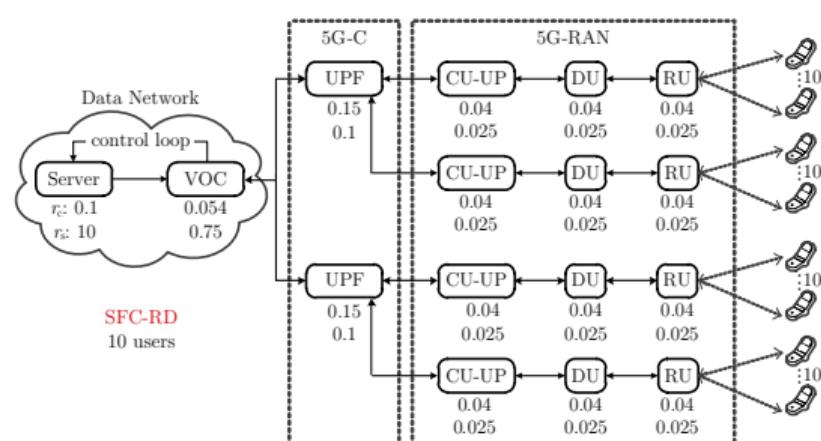


System Model | Virtual graphs of resource demand

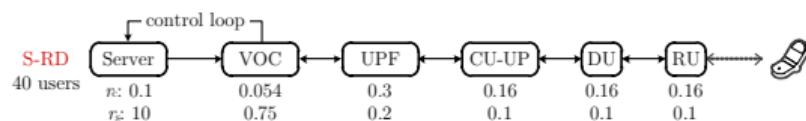
Resource demand of a typical user (U-RD): U_s



Resource demand of an SFC (SFC-RD): r_s



Resource demand of a slice (S-RD): R_s



1 System model

2 Main contributions

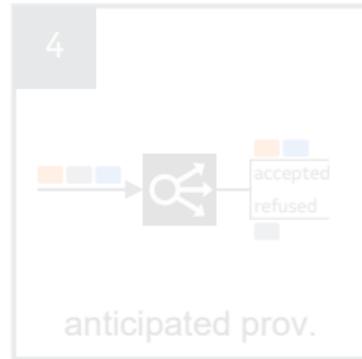
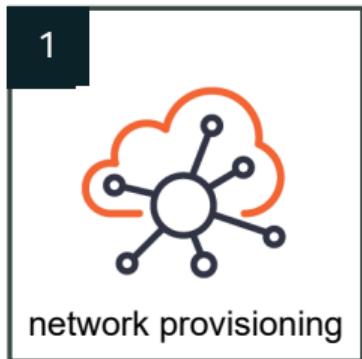
- Contribution 1: Network resource provisioning
- Contribution 4: Anticipated provisioning

3 Conclusions and Perspective

- Contribution 2: Coverage-aware resource provisioning
- Contribution 3: Uncertainty-aware resource provisioning

Contributions

Part I. Deterministic slice resource demands



(1) Network provisioning | Problem statement

Direct embedding vs provisioned embedding

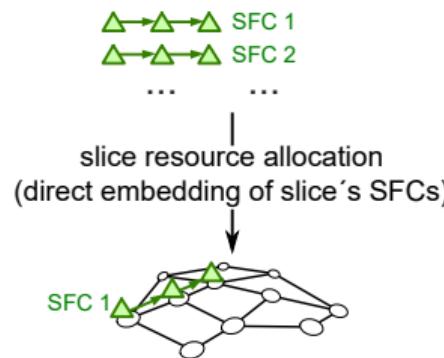
Prior works: Direct embedding (**one-to-one** mapping)



(1) Network provisioning | Problem statement

Direct embedding vs provisioned embedding

Prior works: Direct embedding (**one-to-one** mapping)



(1) Network provisioning | Problem statement

Direct embedding vs provisioned embedding

Prior works: Direct embedding (one-to-one mapping)



Often formulated as (Mixed) Integer Linear Programming and solved either directly or using heuristics

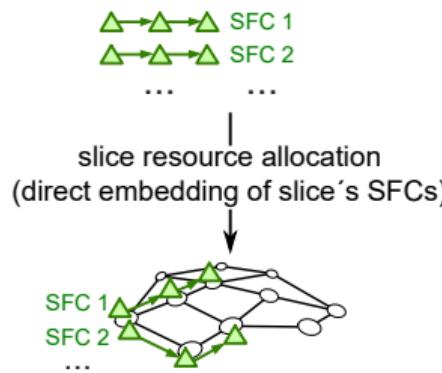
ILP: [Cohen2015], [Riera2016], [Vizarreta2017]

heuristics: column generation [Liu2017] , eigendecomposition [Mechtri2016]

(1) Network provisioning | Problem statement

Direct embedding vs provisioned embedding

Prior works: Direct embedding (one-to-one mapping)



Often formulated as (Mixed) Integer Linear Programming and solved either directly or using heuristics

ILP: [Cohen2015], [Riera2016], [Vizarreta2017]

heuristics: column generation [Liu2017] , eigendecomposition [Mechtri2016]

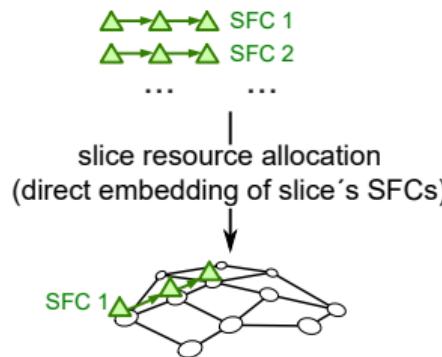
Slice resource provisioning: (many-to-one mapping)



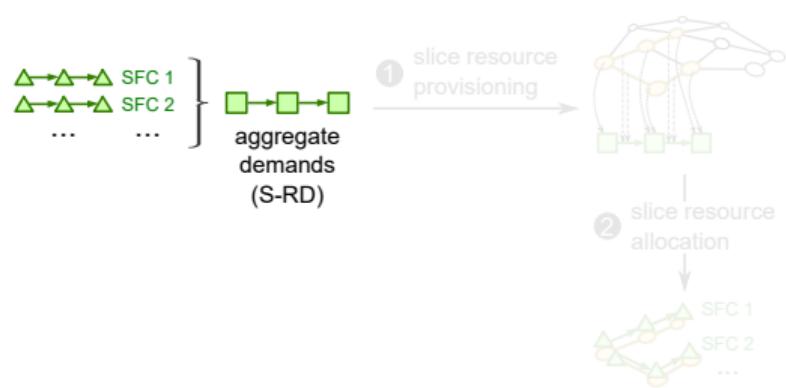
(1) Network provisioning | Problem statement

Direct embedding vs provisioned embedding

Prior works: Direct embedding (one-to-one mapping)



Slice resource provisioning: (many-to-one mapping)



Often formulated as (Mixed) Integer Linear Programming and solved either directly or using heuristics

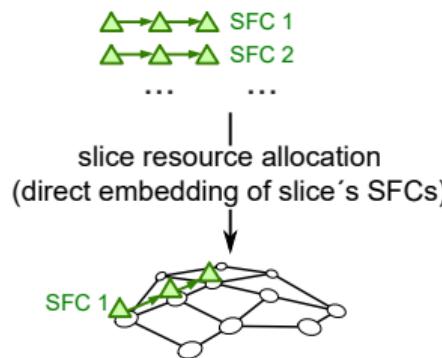
ILP: [Cohen2015], [Riera2016], [Vizarreta2017]

heuristics: column generation [Liu2017] , eigendecomposition [Mechtri2016]

(1) Network provisioning | Problem statement

Direct embedding vs provisioned embedding

Prior works: Direct embedding (one-to-one mapping)

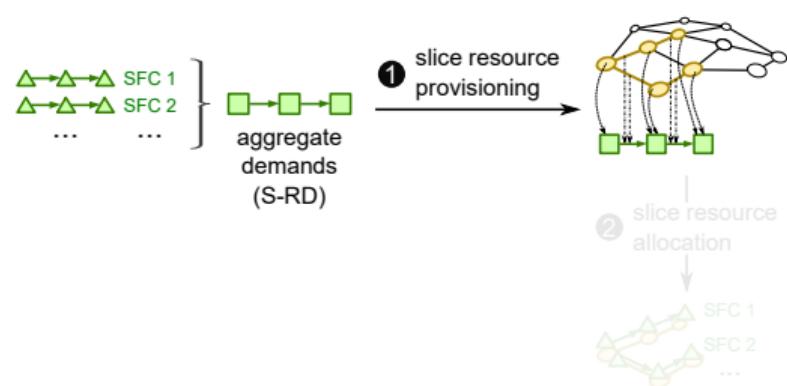


Often formulated as (Mixed) Integer Linear Programming and solved either directly or using heuristics

ILP: [Cohen2015], [Riera2016], [Vizarreta2017]

heuristics: column generation [Liu2017] , eigendecomposition [Mechtri2016]

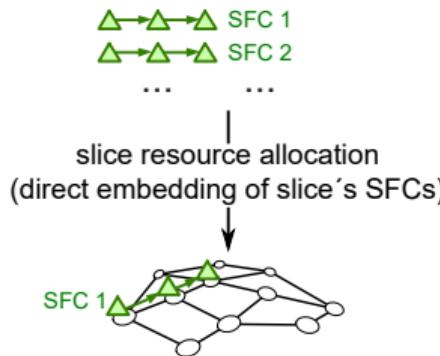
Slice resource provisioning: (many-to-one mapping)



(1) Network provisioning | Problem statement

Direct embedding vs provisioned embedding

Prior works: Direct embedding (one-to-one mapping)

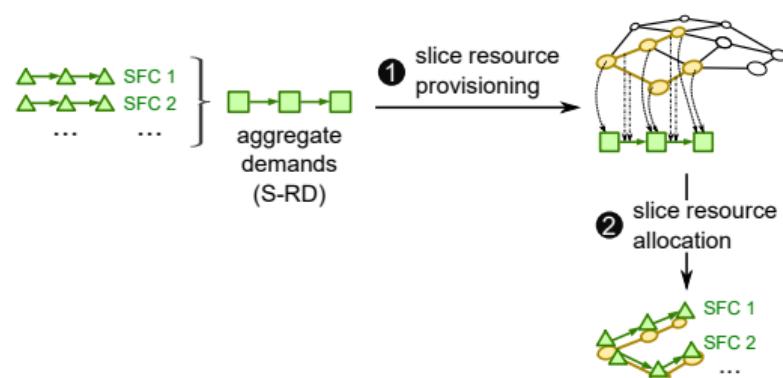


Often formulated as (Mixed) Integer Linear Programming and solved either directly or using heuristics

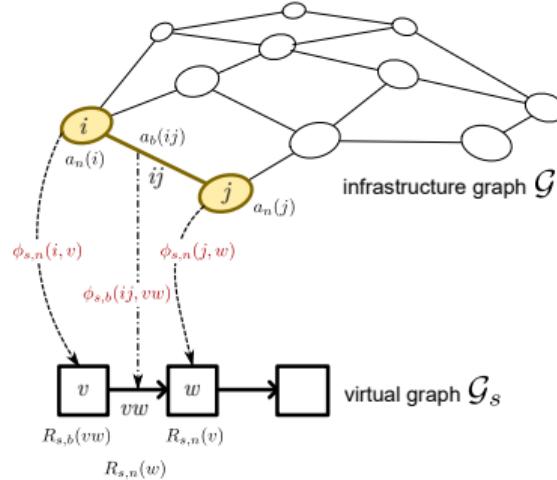
ILP: [Cohen2015], [Riera2016], [Vizarreta2017]

heuristics: column generation [Liu2017] , eigendecomposition [Mechtri2016]

Slice resource provisioning: (many-to-one mapping)



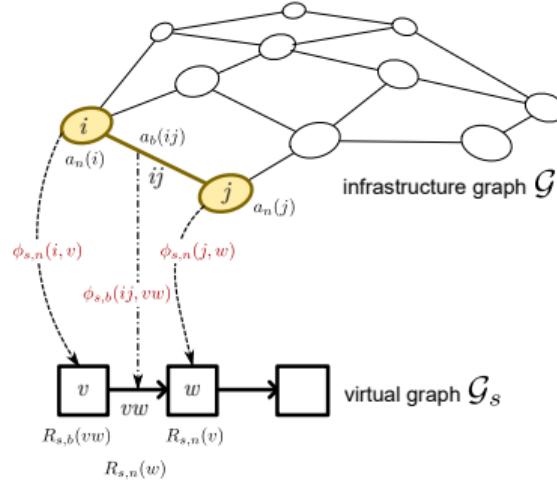
Related work: radio resource provisioning for slices:
[Sun2019], [Xiong2019]



- $\phi_{s,n}(i, v) \in [0, 1]$: Proportion of provisioned node resources
- $\phi_{s,b}(ij, vw) \in [0, 1]$: Proportion of provisioned link resources
- $\tilde{\phi}_{s,n}(i, v) \in \{0, 1\}$: Node mapping indicator

Objective function

$$\begin{aligned}
 c_{\text{total}}(\Phi) &= \sum_s c_{\text{total}}(\Phi_s) \\
 &= \sum_{s,i,v,n} a_n(i) \phi_{s,n}(i, v) c_n(i) && \text{(node resource leasing)} \\
 &+ \sum_{s,ij,vw} a_b(ij) \phi_{s,b}(ij, vw) c_b(ij) && \text{(link resource leasing)} \\
 &+ \sum_{s,i,v} \tilde{\phi}_{s,n}(i, v) c_d(i) && \text{(VNF deployment)}
 \end{aligned}$$



- $\phi_{s,n}(i, v) \in [0, 1]$: Proportion of provisioned node resources
- $\phi_{s,b}(ij, vw) \in [0, 1]$: Proportion of provisioned link resources
- $\tilde{\phi}_{s,n}(i, v) \in \{0, 1\}$: Node mapping indicator

Objective function

$$\begin{aligned}
 c_{\text{total}}(\Phi) &= \sum_s c_{\text{total}}(\Phi_s) \\
 &= \sum_{s,i,v,n} a_n(i) \phi_{s,n}(i, v) c_n(i) && \text{(node resource leasing)} \\
 &+ \sum_{s,ij,vw} a_b(ij) \phi_{s,b}(ij, vw) c_b(ij) && \text{(link resource leasing)} \\
 &+ \sum_{s,i,v} \tilde{\phi}_{s,n}(i, v) c_d(i) && \text{(VNF deployment)}
 \end{aligned}$$

(1) Network provisioning | Main constraints

Constraint 1: Slice resource requirements

$$\sum_i \underbrace{a_n(i)}_{\text{prov. rsrc from } i \text{ to } v} \phi_{s,n}(i, v) \geq \underbrace{R_n(v)}_{\text{S-RD at } v}, \forall s \in \mathcal{S}, v \in \mathcal{N}_s, n \in \Upsilon$$

$$\sum_{ij} \underbrace{a_b(ij)}_{\text{prov. rsrc from } ij \text{ to } vw} \phi_{s,b}(ij, vw) \geq \underbrace{R_b(vw)}_{\text{S-RD at } v}, \forall s \in \mathcal{S}, vw \in \mathcal{E}_s$$

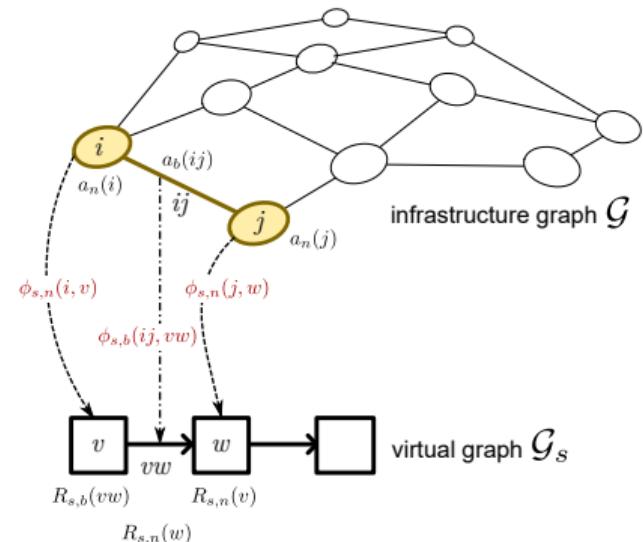
Constraint 2: Infrastructure resource limit

$$\sum_{s, v} \phi_{s,n}(i, v) \leq 1, \forall i \in \mathcal{N}, n \in \Upsilon,$$

$$\sum_{s, vw} \phi_{s,b}(ij, vw) \leq 1, \forall ij \in \mathcal{E}$$

Other constraints:

- proportionality of provisioned node resources (computing-memory-wireless)
- flow conservation: ensures flow continuity



(1) Network provisioning | Main constraints

Constraint 1: Slice resource requirements

$$\sum_i \underbrace{a_n(i)}_{\text{prov. rsrc from } i \text{ to } v} \phi_{s,n}(i, v) \geq \underbrace{R_n(v)}_{\text{S-RD at } v}, \forall s \in \mathcal{S}, v \in \mathcal{N}_s, n \in \Upsilon$$

$$\sum_{ij} \underbrace{a_b(ij)}_{\text{prov. rsrc from } ij \text{ to } vw} \phi_{s,b}(ij, vw) \geq \underbrace{R_b(vw)}_{\text{S-RD at } v}, \forall s \in \mathcal{S}, vw \in \mathcal{E}_s$$

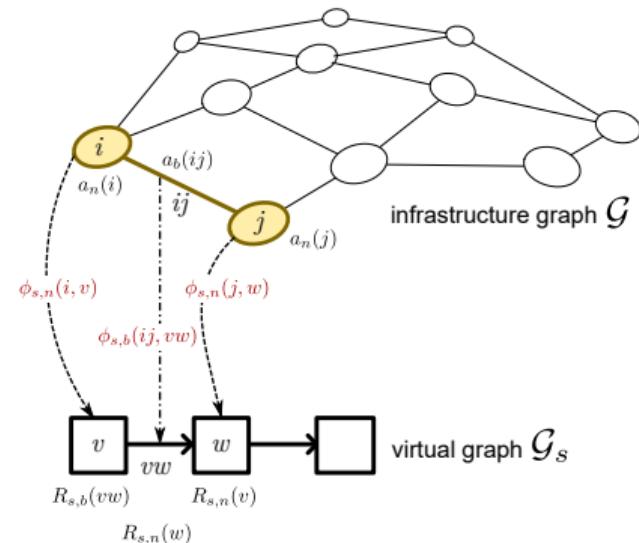
Constraint 2: Infrastructure resource limit

$$\sum_{s,v} \phi_{s,n}(i, v) \leq 1, \forall i \in \mathcal{N}, n \in \Upsilon,$$

$$\sum_{s,vw} \phi_{s,b}(ij, vw) \leq 1, \forall ij \in \mathcal{E}$$

Other constraints:

- proportionality of provisioned node resources (computing-memory-wireless)
- flow conservation: ensures flow continuity



(1) Network provisioning | Main constraints

Constraint 1: Slice resource requirements

$$\sum_i \underbrace{a_n(i)}_{\text{prov. rsrc from } i \text{ to } v} \phi_{s,n}(i, v) \geq \underbrace{R_n(v)}_{\text{S-RD at } v}, \forall s \in \mathcal{S}, v \in \mathcal{N}_s, n \in \Upsilon$$

$$\sum_{ij} \underbrace{a_b(ij)}_{\text{prov. rsrc from } ij \text{ to } vw} \phi_{s,b}(ij, vw) \geq \underbrace{R_b(vw)}_{\text{S-RD at } v}, \forall s \in \mathcal{S}, vw \in \mathcal{E}_s$$

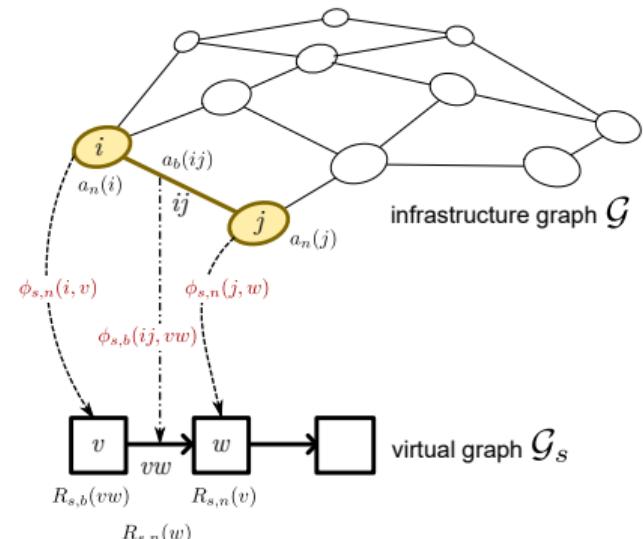
Constraint 2: Infrastructure resource limit

$$\sum_{s,v} \phi_{s,n}(i, v) \leq 1, \forall i \in \mathcal{N}, n \in \Upsilon,$$

$$\sum_{s,vw} \phi_{s,b}(ij, vw) \leq 1, \forall ij \in \mathcal{E}$$

Other constraints:

- proportionality of provisioned node resources (computing-memory-wireless)
- flow conservation: ensures flow continuity



MILP network resource provisioning

$$\text{minimize } c_{\text{total}}(\Phi) = \sum_{s \in \mathcal{S}} c_{\text{total}}(\Phi_s),$$

subject to

- (1) slice resource requirements
- (2) infrastructure resource limits
- (3) resource proportionality
- (4) flow conservation

Provisioning approaches:

- (1) Sequential approach (SP): Resources are provisioned slice-by-slice
- (2) Joint approach (JP): Provisioning performed for all slices simultaneously

MILP network resource provisioning

$$\text{minimize } c_{\text{total}}(\Phi) = \sum_{s \in \mathcal{S}} c_{\text{total}}(\Phi_s),$$

subject to

- (1) slice resource requirements
- (2) infrastructure resource limits
- (3) resource proportionality
- (4) flow conservation

Provisioning approaches:

- (1) Sequential approach (SP): Resources are provisioned slice-by-slice
- (2) Joint approach (JP): Provisioning performed for all slices simultaneously

MILP network resource provisioning

$$\text{minimize } c_{\text{total}}(\Phi) = \sum_{s \in S} c_{\text{total}}(\Phi_s),$$

subject to

- (1) slice resource requirements
- (2) infrastructure resource limits
- (3) resource proportionality
- (4) flow conservation

Provisioning approaches:

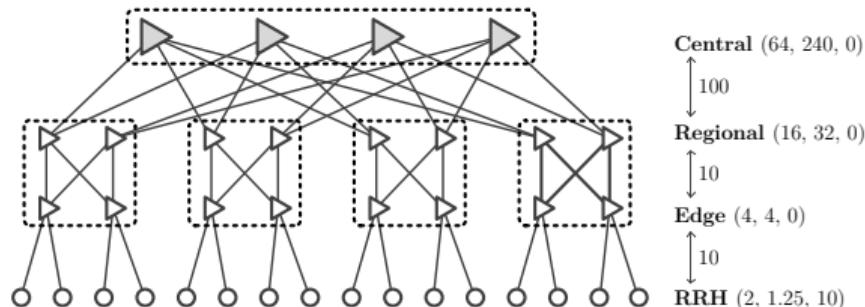
- (1) **Sequential** approach (**SP**): Resources are provisioned slice-by-slice
- (2) **Joint** approach (**JP**): Provisioning performed for all slices simultaneously

(1) Network provisioning | Evaluation

Infrastructure network

k-ary fat-tree topology

- Per-unit resource cost = 1 (whatever the type of resource)
- Deployment cost $c_d = 50$

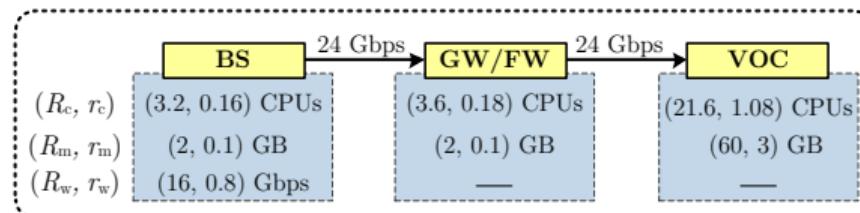


(1) Network provisioning | Evaluation

Slice resource demands

Slices: HD video streaming with *deterministic* S-RD

- Data rate at 4 Mbps for $N_s = 4000$ users.



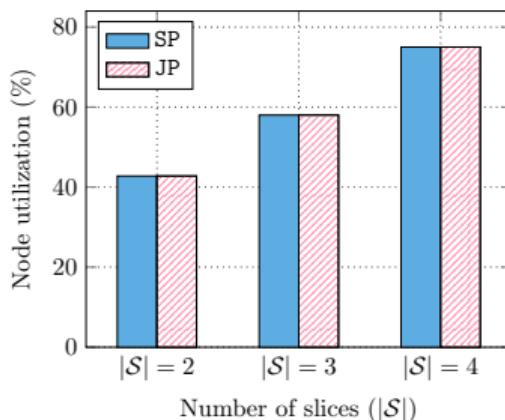
Simulation schemes

Comparison of 2 resource provisioning approaches:

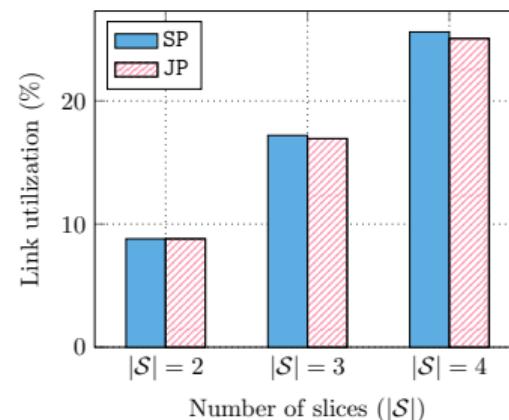
- Sequential approach (**SP**)
- Joint approach (**JP**)

Both schemes are evaluated using CPLEX MILP solver interfaced with MATLAB.

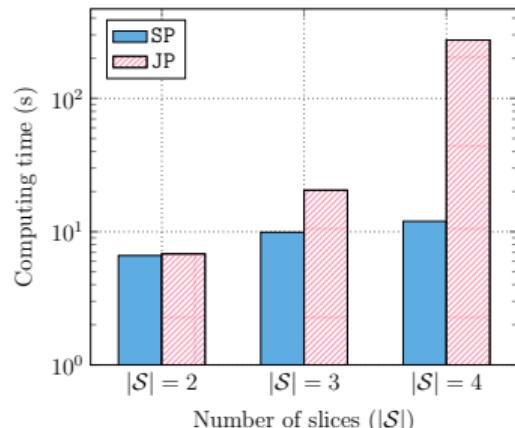
Sequential (SP) vs Joint (JP) approaches, with different $|\mathcal{S}|$



node utilization



link utilization

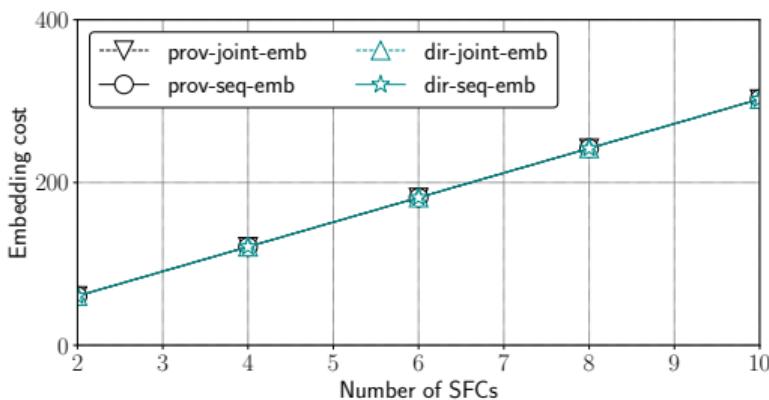


computing time

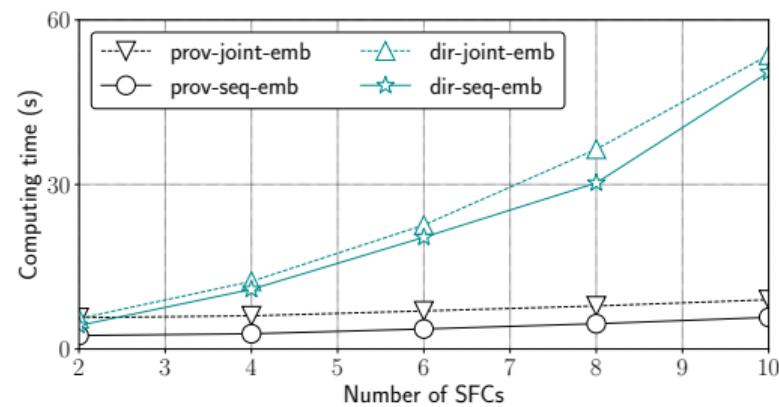
(1) Network provisioning | Evaluation

Results

Provisioned embedding vs Direct embedding (single slice)



Embedding cost



Computing time

Provisioned embedding: embedding is performed with **only a subset** of the infrastructure graph

- S-RD graph: aggregation of resource demands
- Slice resource provisioning: mapping of infrastructure graph to S-RD graph
- Two slice resource provisioning approaches:
 - ▶ Sequential approach (**SP**): more efficient in terms of computing time
 - ▶ Nevertheless: degraded node and link utilization compared to Joint approach (**JP**)
- Once provisioning completed, SFC embedding is **faster**

Contributions

Part I. Deterministic slice resource demands

1



network provisioning

2



coverage-aware

Part II. Random slice resource demands

3



uncertainty-aware

4



anticipated prov.



Related works

- dynamic slice/SFC deployment
 - ▶ [Liu'17]: deployment of newly arrived SFCs and readjustment of in-service SFCs
 - ▶ [Wang'19]: slice reconfiguration at two levels: SFCs within a slice at small time intervals, or multiple slices at larger time interval
- slice admission control:
 - ▶ [Bega'17], [Bega'20]: online SAC for radio resources at base stations
 - ▶ [Han'20]: SAC with dynamic requests, but not accounting for uncertainties

Related works

- dynamic slice/SFC deployment:
 - ▶ [Liu'17]: deployment of newly arrived SFCs and readjustment of in-service SFCs
 - ▶ [Wang'19]: slice reconfiguration at two levels: SFCs within a slice at small time intervals, or multiple slices at larger time interval
- slice admission control:
 - ▶ [Bega'17], [Bega'20]: online SAC for radio resources at base stations
 - ▶ [Han'20]: SAC with dynamic requests, but not accounting for uncertainties

This contribution:

Multiple slice requests from tenants to MNO

Uncertainties

- slice resource demands
- available infrastructure resource

Dynamic nature of slice provisioning requests

- arrive at anytime
- various activation times
- various lifetime
- resource needs vary during slice operation

This contribution:

Multiple slice requests from tenants to MNO

Uncertainties

- slice resource demands
- available infrastructure resource

Dynamic nature of slice provisioning requests

- arrive at anytime
- various activation times
- various lifetime
- resource needs vary during slice operation

This contribution:

Multiple slice requests from tenants to MNO

Uncertainties

- slice resource demands
- available infrastructure resource

Dynamic nature of slice provisioning requests

- arrive at anytime
- various activation times
- various lifetime
- resource needs vary during slice operation

This contribution:

Multiple slice requests from tenants to MNO

Uncertainties

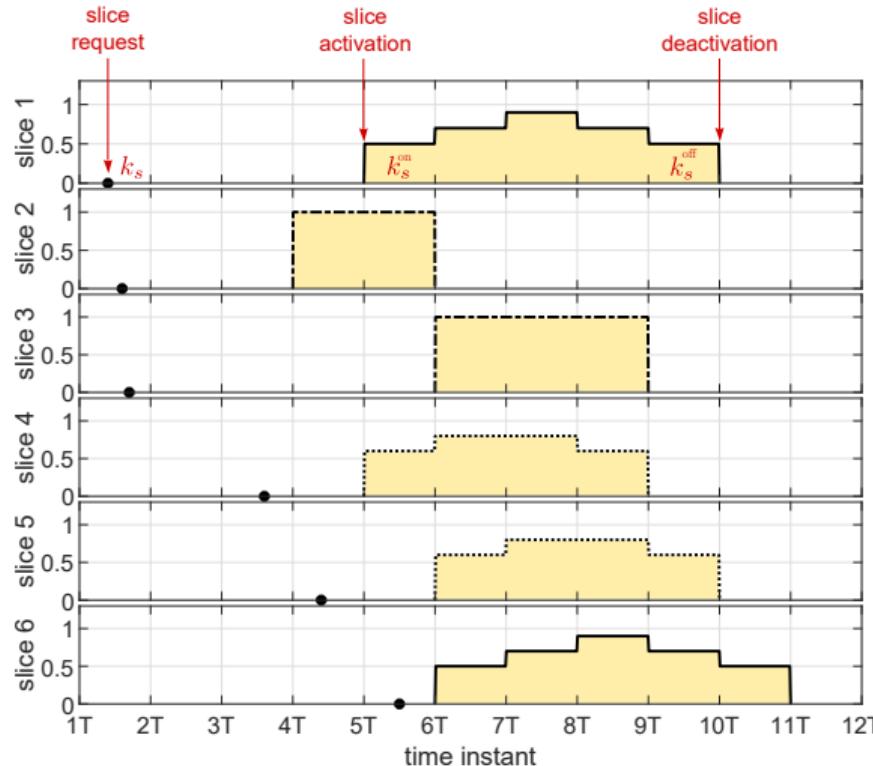
- slice resource demands
- available infrastructure resource

Dynamic nature of slice provisioning requests

- arrive at anytime
- various activation times
- various lifetime
- resource needs vary during slice operation

(4) Anticipated provisioning | System model

Slice provisioning requests



(4) Anticipated provisioning | Priority assignment

Slice classes: Premium and Standard

Initial priority assignment

$$P_{s,k} = \begin{cases} P_{\max}, & \text{for Premium slices,} \\ 0, & \text{for Standard slice, if } k_s^{\text{on}} > k + 1, \\ P_{\max} - 1 & \text{for Standard slice, if } k_s^{\text{on}} = k + 1. \end{cases}$$

Only slices with $P_{s,k} \geq P_{\text{thres}}$ are processed

$$P_{\text{thres}} = \alpha (P_{\max} - 1), \text{ with } \alpha \in [0, 1].$$

(4) Anticipated provisioning | Priority assignment

Slice classes: Premium and Standard

Initial priority assignment

$$P_{s,k} = \begin{cases} P_{\max}, & \text{for Premium slices,} \\ 0 & \text{for Standard slice, if } k_s^{\text{on}} > k + 1, \\ P_{\max} - 1 & \text{for Standard slice, if } k_s^{\text{on}} = k + 1. \end{cases}$$

Only slices with $P_{s,k} \geq P_{\text{thres}}$ are processed

$$P_{\text{thres}} = \alpha (P_{\max} - 1), \text{ with } \alpha \in [0, 1].$$

(4) Anticipated provisioning | Priority assignment

Slice classes: Premium and Standard
Initial priority assignment

$$P_{s,k} = \begin{cases} P_{\max}, & \text{for Premium slices,} \\ 0 & \text{for Standard slice, if } k_s^{\text{on}} > k + 1, \\ P_{\max} - 1 & \text{for Standard slice, if } k_s^{\text{on}} = k + 1. \end{cases}$$

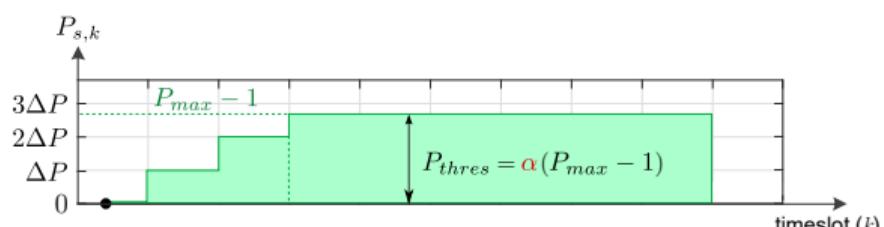
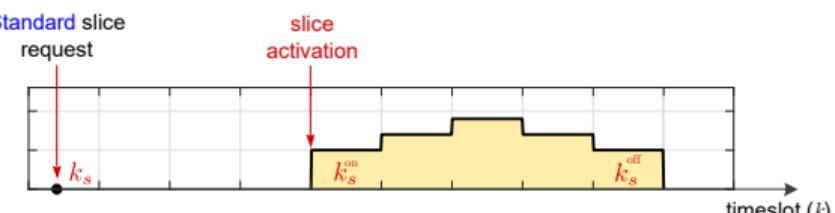
Only slices with $P_{s,k} \geq P_{\text{thres}}$ are processed

$$P_{\text{thres}} = \alpha(P_{\max} - 1), \text{ with } \alpha \in [0, 1],$$

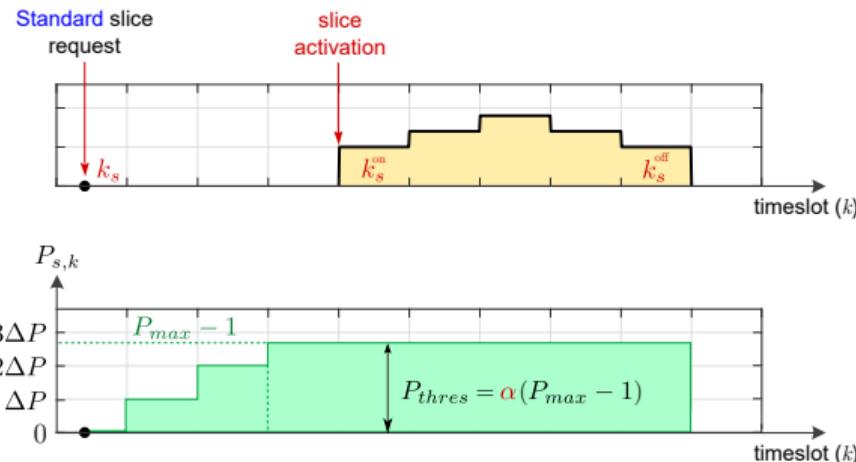
Priority update for unprocessed Standard slices in next time slots:

$$P_{s,k} = \begin{cases} \min \{P_{s,k-1} + \Delta P, P_{\max} - 1\} & \text{if } k_s^{\text{on}} > k + 1, \\ P_{\max} - 1 & \text{if } k_s^{\text{on}} = k + 1, \end{cases}$$

$\Delta P \geq 0$: priority increment.



(4) Anticipated provisioning | Priority assignment



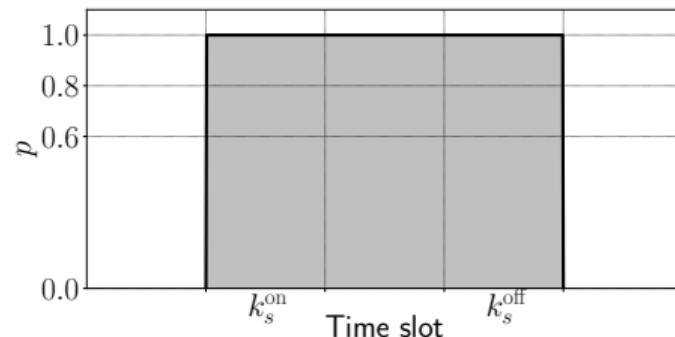
Parameterized priority level with α and ΔP

- $\alpha = 0$: Premium = Standard, $\forall \Delta P \rightarrow$ min benefit for Premium slices
- $\alpha = 1$, $\Delta P = 0$, Standard slices are delayed until the time slot preceding their activation \rightarrow max benefit for Premium slices

(4) Anticipated provisioning | S-RD

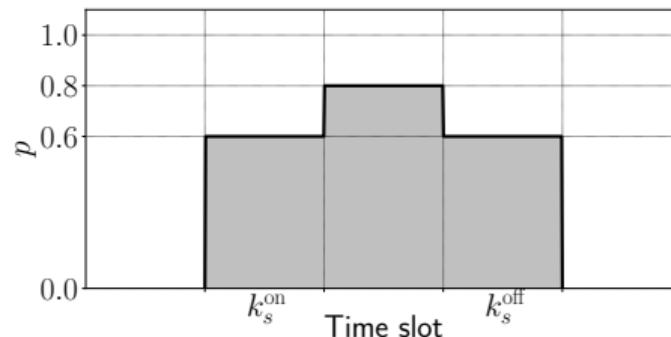
Slice resource demands:

- Number of users of slice s follows binomial distribution $\mathcal{B}(p_{s,k})$
- Temporal evolution of $p_{s,k}$



constant over a time interval

Service usage pattern

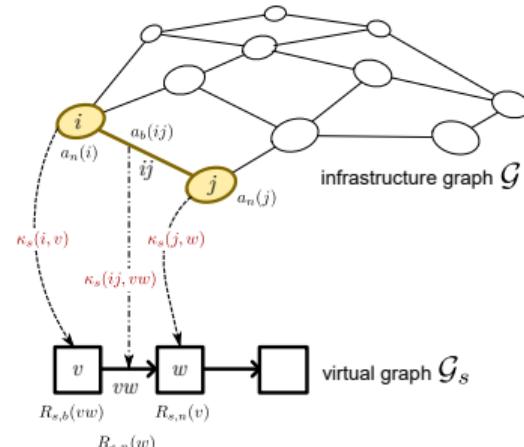


piece-wise constant

(4) Anticipated provisioning | Variables

Resource provisioning for slice s defined by mapping between \mathcal{G} and \mathcal{G}_s , for time slot ℓ

- $d_s \in \{0, 1\}$ indicates whether slice s is accepted or not
- $\kappa_{s,\ell}(i, v) \in \mathbb{N}$ determines number of VNFs that node i can provide resources for virtual node v
- $\kappa_{s,\ell}(ij, vw) \in \mathbb{N}$ determines bandwidth provisioned by link ij to support the traffic between virtual nodes of type v and w



Note:

Contributions 1 and 2 use Φ_s as main decision variables
Here: κ_s

(4) Anticipated provisioning | Constraints

Summary

Assignment $\kappa_{s,\ell}$ for slice s over time slot ℓ must satisfy

- infrastructure node and link capacity limits, w.r.t
 - ▶ background traffic
- service satisfaction probability, w.r.t
 - ▶ usage uncertainties
 - ▶ resource availability
- flow conservation constraint

(4) Anticipated provisioning | SSP constraint

Service Satisfaction Probability (SSP) of slice s at time slot k for a given assignment $\kappa_{s,\ell}$

$$p_{s,\ell}(\kappa_{s,\ell}, d_s) = \Pr \left\{ \begin{array}{l} \underbrace{\sum_i \kappa_{s,\ell}(i, v) r_{s,n}(v)}_{\text{prov. rsrc for } v} \geq \underbrace{d_s R_{s,n,\ell}(v)}_{\text{S-RD of } v}, \forall v, n, \\ \underbrace{\sum_{ij} \kappa_{s,\ell}(ij, vw) r_{s,b}(vw)}_{\text{prov. rsrc for } vw} \geq \underbrace{d_s R_{s,b,\ell}(vw)}_{\text{S-RD of } vw}, \forall vw \end{array} \right\} \geq \underbrace{p_s^{\text{sp}}}_{\text{required min SSP}}, \forall s \in \mathcal{R}_k, \ell \in \mathcal{K}_s$$

Relaxed constraints:

$$\sum_i \kappa_{s,\ell}(i, v) r_{s,n}(v) \geq \underbrace{d_s [\bar{R}_{s,n,\ell}(v) + \gamma_{s,\ell} \tilde{R}_{s,n,\ell}(v)]}_{\text{S-RD of } v}, \forall \ell > k, v \in \mathcal{N}_s, \forall n \in \Upsilon,$$

$$\sum_{ij} \kappa_{s,\ell}(ij, vw) r_{s,b}(vw) \geq \underbrace{d_s [\bar{R}_{s,b,\ell}(vw) + \gamma_{s,\ell} \tilde{R}_{s,b,\ell}(vw)]}_{\text{S-RD of } vw}, \forall \ell > k, vw \in \mathcal{E}_s.$$

\mathcal{R}_k : set of slices processed at time slot k

\bar{R}, \tilde{R} : mean and standard deviation of S-RD

$\gamma_{s,\ell}$: safety factor

(4) Anticipated provisioning | SSP constraint

Service Satisfaction Probability (SSP) of slice s at time slot k for a given assignment $\kappa_{s,\ell}$

$$p_{s,\ell}(\kappa_{s,\ell}, d_s) = \Pr \left\{ \begin{array}{l} \underbrace{\sum_i \kappa_{s,\ell}(i, v) r_{s,n}(v)}_{\text{prov. rsrc for } v} \geq \underbrace{d_s R_{s,n,\ell}(v)}_{\text{S-RD of } v}, \forall v, n, \\ \underbrace{\sum_{ij} \kappa_{s,\ell}(ij, vw) r_{s,b}(vw)}_{\text{prov. rsrc for } vw} \geq \underbrace{d_s R_{s,b,\ell}(vw)}_{\text{S-RD of } vw}, \forall vw \end{array} \right\} \geq \underbrace{p_s^{\text{sp}}}_{\text{required min SSP}}, \forall s \in \mathcal{R}_k, \ell \in \mathcal{K}_s$$

Relaxed constraints:

$$\sum_i \kappa_{s,\ell}(i, v) r_{s,n}(v) \geq \underbrace{d_s [\bar{R}_{s,n,\ell}(v) + \gamma_{s,\ell} \tilde{R}_{s,n,\ell}(v)]}_{\text{S-RD of } v}, \forall \ell > k, v \in \mathcal{N}_s, \forall n \in \Upsilon,$$

$$\sum_{ij} \kappa_{s,\ell}(ij, vw) r_{s,b}(vw) \geq \underbrace{d_s [\bar{R}_{s,b,\ell}(vw) + \gamma_{s,\ell} \tilde{R}_{s,b,\ell}(vw)]}_{\text{S-RD of } vw}, \forall \ell > k, vw \in \mathcal{E}_s.$$

\mathcal{R}_k : set of slices processed at time slot k

\bar{R}, \tilde{R} : mean and standard deviation of S-RD

$\gamma_{s,\ell}$: safety factor

(4) Anticipated provisioning | ImP constraint

Impact Probability (ImP) constraints on background services at infrastructure node i and link ij :

$$p_{n,\ell}^{\text{im}}(\kappa_{\mathcal{R}_k}, i) = \Pr \left\{ \underbrace{\sum_s \sum_v \kappa_{s,\ell}(i, v) r_{s,n}(v)}_{\text{resources prov. by } i} \geq a_n(i) - \underbrace{B_{n,\ell}(i)}_{\text{bkgd serv. at } i} \right\} \leq \underbrace{\bar{p}^{\text{im}}}_{\text{required max ImP}},$$
$$p_{b,\ell}^{\text{im}}(\kappa_{\mathcal{R}_k}, ij) = \Pr \left\{ \underbrace{\sum_s \sum_{vw} \kappa_{s,\ell}(ij, vw) r_{s,b}(vw)}_{\text{resources prov. by } ij} \geq a_b(ij) - \underbrace{B_{b,\ell}(ij)}_{\text{bkgd serv. at } ij} \right\} \leq \underbrace{\bar{p}^{\text{im}}}_{\text{required max ImP}}.$$

Relaxed constraints:

$$\sum_s \sum_v \kappa_{s,\ell}(i, v) r_{s,n}(v) \leq a_n(i) - [\bar{B}_{n,\ell}(i) + \gamma_{B,\ell} \tilde{B}_{n,\ell}(i)], \forall \ell > k, i \in \mathcal{N}, n \in \Upsilon,$$

$$\sum_s \sum_{vw} \kappa_{s,\ell}(ij, vw) r_{s,b}(vw) \leq a_b(ij) - \underbrace{[\bar{B}_{b,\ell}(ij) + \gamma_{B,\ell} \tilde{B}_{b,\ell}(ij)]}_{\text{background services}}, \forall \ell > k, ij \in \mathcal{E}.$$

\bar{B}, \tilde{B} : mean and standard deviation of resource consumption of background services

$\gamma_{B,\ell}$: safety factor

(4) Anticipated provisioning | ImP constraint

Impact Probability (ImP) constraints on background services at infrastructure node i and link ij :

$$p_{n,\ell}^{\text{im}}(\kappa_{\mathcal{R}_k}, i) = \Pr \left\{ \underbrace{\sum_s \sum_v \kappa_{s,\ell}(i, v) r_{s,n}(v)}_{\text{resources prov. by } i} \geq a_n(i) - \underbrace{B_{n,\ell}(i)}_{\text{bkgd serv. at } i} \right\} \leq \underbrace{\bar{p}^{\text{im}}}_{\text{required max ImP}},$$
$$p_{b,\ell}^{\text{im}}(\kappa_{\mathcal{R}_k}, ij) = \Pr \left\{ \underbrace{\sum_s \sum_{vw} \kappa_{s,\ell}(ij, vw) r_{s,b}(vw)}_{\text{resources prov. by } ij} \geq a_b(ij) - \underbrace{B_{b,\ell}(ij)}_{\text{bkgd serv. at } ij} \right\} \leq \underbrace{\bar{p}^{\text{im}}}_{\text{required max ImP}}.$$

Relaxed constraints:

$$\sum_s \sum_v \kappa_{s,\ell}(i, v) r_{s,n}(v) \leq a_n(i) - [\bar{B}_{n,\ell}(i) + \gamma_{B,\ell} \tilde{B}_{n,\ell}(i)], \forall \ell > k, i \in \mathcal{N}, n \in \Upsilon,$$

$$\sum_s \sum_{vw} \kappa_{s,\ell}(ij, vw) r_{s,b}(vw) \leq a_b(ij) - \underbrace{[\bar{B}_{b,\ell}(ij) + \gamma_{B,\ell} \tilde{B}_{b,\ell}(ij)]}_{\text{background services}}, \forall \ell > k, ij \in \mathcal{E}.$$

\bar{B}, \tilde{B} : mean and standard deviation of resource consumption of background services

$\gamma_{B,\ell}$: safety factor

(4) Anticipated provisioning | Objective function

Earnings of InP = costs charged to MNOs

Max-min problem: InP tries to

- maximize the nb of accepted slices
- minimize the cost for MNOs

Max-min anticipated provisioning problem

$$\max_{\mathbf{d}_{\mathcal{R}_k}} \min_{\boldsymbol{\kappa}_{\mathcal{R}_k}} E_k^{\text{InP}} (\mathbf{d}_{\mathcal{R}_k}, \boldsymbol{\kappa}_{\mathcal{R}_k}) = \sum_{s \in \mathcal{R}_k} \sum_{\ell \in \mathcal{K}_s} \left(\underbrace{C_r(\kappa_{s,\ell})}_{\text{resource}} + \underbrace{C_f(\kappa_{s,\ell})}_{\text{node disposal}} + \underbrace{C_a(\kappa_{s,\ell}, \kappa_{s,\ell-1})}_{\text{adaptation}} \right)$$

subject to

- service satisfaction probability constraints
- impact on background services constraints
- flow conservation constraint

Approaches

Sequential: solve problem for each slice s with initial $d_s = 1$

- $\min_{\kappa_{s_1}} E_k^{\text{InP}}(\kappa_{s_1}, \mathbf{1}), \min_{\kappa_{s_2}} E_k^{\text{InP}}(\kappa_{s_2}, \mathbf{1}), \dots$

Joint: start with $d_{\mathcal{R}_k} = 1$ ($d_s = 1, \forall s$), then refuse progressively low-prioritized slices if no solution found

- iteration 1: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 1, 1\})$
- iteration 2: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 1, \mathbf{0}\})$
- iteration 3: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, \mathbf{0}, \mathbf{0}\})$
- ... until a feasible solution found

Approaches

Sequential: solve problem for each slice s with initial $d_s = 1$

- $\min_{\kappa_{s_1}} E_k^{\text{InP}}(\kappa_{s_1}, 1), \min_{\kappa_{s_2}} E_k^{\text{InP}}(\kappa_{s_2}, 1), \dots$

Joint: start with $d_{\mathcal{R}_k} = 1$ ($d_s = 1, \forall s$), then refuse progressively low-prioritized slices if no solution found

- iteration 1: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 1, 1\})$
- iteration 2: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 1, 0\})$
- iteration 3: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 0, 0\})$
- ... until a feasible solution found

Approaches

Sequential: solve problem for each slice s with initial $d_s = 1$

- $\min_{\kappa_{s_1}} E_k^{\text{InP}}(\kappa_{s_1}, 1), \min_{\kappa_{s_2}} E_k^{\text{InP}}(\kappa_{s_2}, 1), \dots$

Joint: start with $d_{\mathcal{R}_k} = 1$ ($d_s = 1, \forall s$), then refuse progressively low-prioritized slices if no solution found

- iteration 1: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 1, 1\})$
- iteration 2: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 1, 0\})$
- iteration 3: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 0, 0\})$
- ... until a feasible solution found

Approaches

Sequential: solve problem for each slice s with initial $d_s = 1$

- $\min_{\kappa_{s_1}} E_k^{\text{InP}}(\kappa_{s_1}, 1), \min_{\kappa_{s_2}} E_k^{\text{InP}}(\kappa_{s_2}, 1), \dots$

Joint: start with $d_{\mathcal{R}_k} = 1$ ($d_s = 1, \forall s$), then refuse progressively low-prioritized slices if no solution found

- iteration 1: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 1, 1\})$
- iteration 2: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 1, 0\})$
- iteration 3: $\min_{\kappa_{\mathcal{R}_k}} E_k^{\text{InP}}(\kappa_{\mathcal{R}_k}, \{1, 1, \dots, 0, 0\})$
- ... until a feasible solution found

Slice model:

- Three types of slice
 - ▶ HD videos streaming, SSP constraint: $\underline{p}_s^{\text{sp}} = 0.99$
 - ▶ SD video streaming, SSP constraint: $\underline{p}_s^{\text{sp}} = 0.95$
 - ▶ video surveillance and traffic monitoring, SSP constraint: $\underline{p}_s^{\text{sp}} = 0.9$
- Slice type and demand pattern chosen uniformly at random.
- Number of provisioning request arrivals/time slot Pois(1)
- Arrival time uniformly distributed within each time slot
- Activation delay $\mathcal{U}(1, 6)$, lifetime $\mathcal{U}(1, 3)$

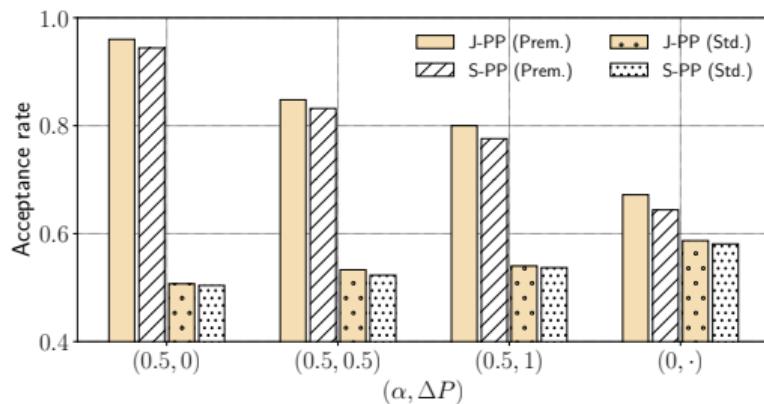
Background services:

- consume on average 20% of available infrastructure resources (with 5% std. dev.)
- ImP constraint: $\bar{p}^{\text{im}} = 0.1$

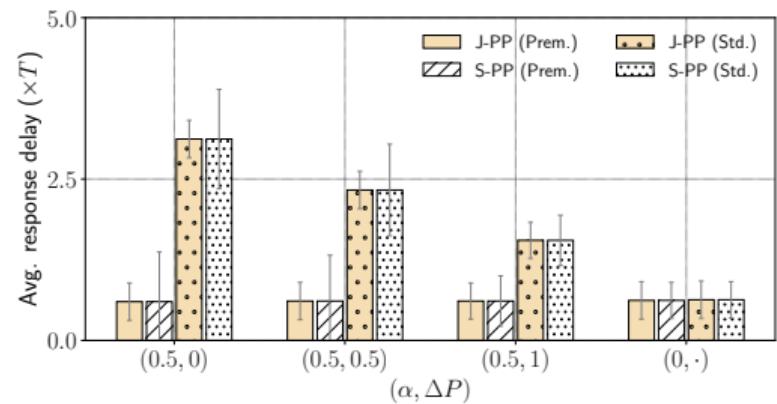
(4) Anticipated provisioning | Evaluation

Multiple slice provisioning

J-PP vs. S-PP with different processing strategies $(\alpha, \Delta P)$



acceptance rate

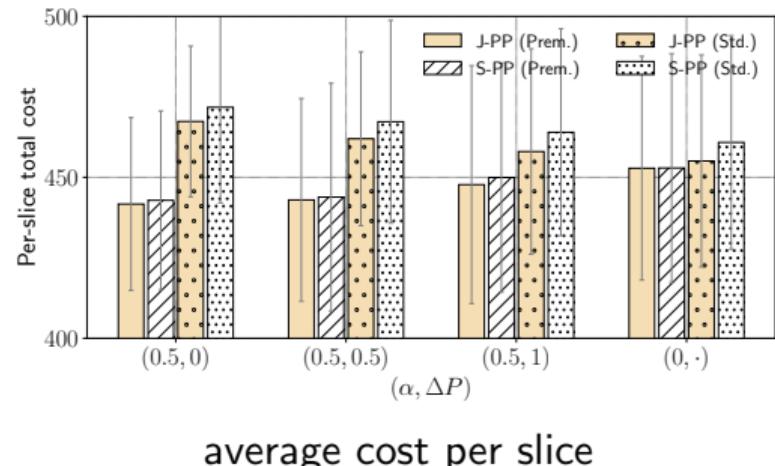
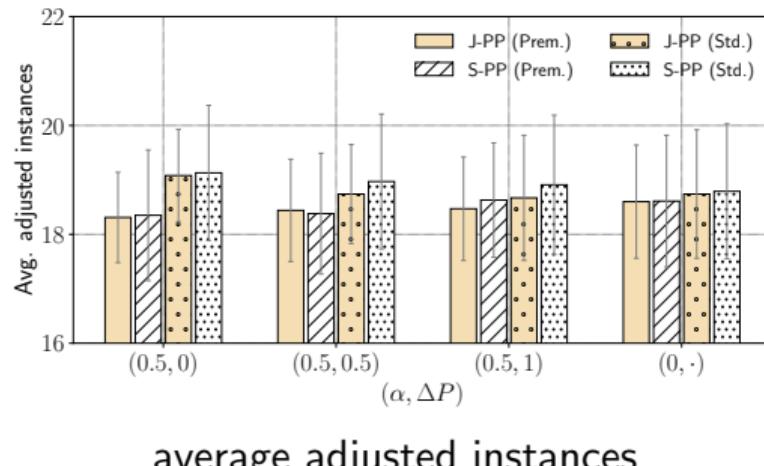


average response delay

(4) Anticipated provisioning | Evaluation

Multiple slice provisioning

J-PP vs. S-PP with different processing strategies ($\alpha, \Delta P$)

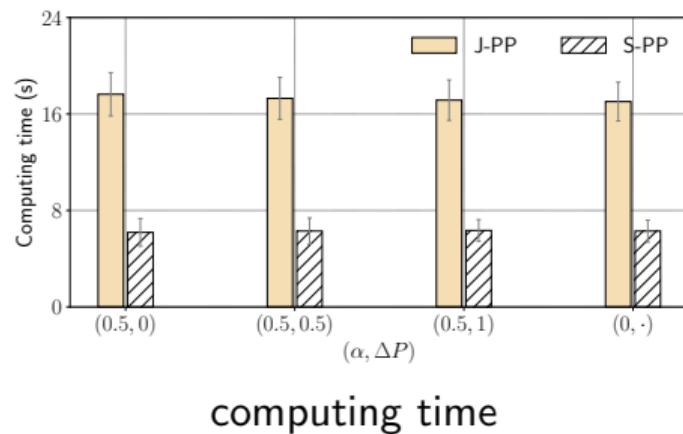


Objective: to minimize cost and nb of adjusted instances

(4) Anticipated provisioning | Evaluation

Multiple slice provisioning

J-PP vs. S-PP with different processing strategies $(\alpha, \Delta P)$



Two prioritized slice resource provisioning (PP) variants to

- **maximize** acceptance rate,
- **maximize** earnings of the InP + **minimize** costs charged to MNO

Advantages

- adapt to dynamic nature of slice provisioning requests;
- robustness to uncertainties.

1 System model

2 Main contributions

- Contribution 1: Network resource provisioning
- Contribution 4: Anticipated provisioning

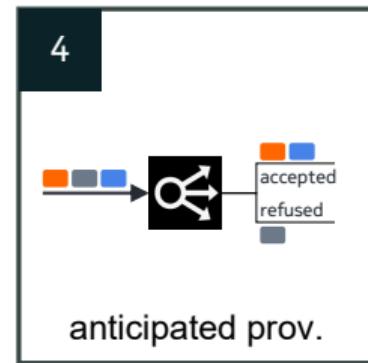
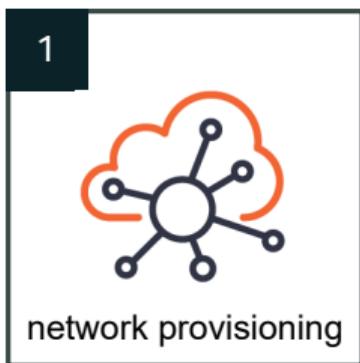
3 Conclusions and Perspective

- Contribution 2: Coverage-aware resource provisioning
- Contribution 3: Uncertainty-aware resource provisioning

Conclusions and Perspective

Contributions

Part I. Deterministic slice resource demands



Perspectives

- Additional constraints, e.g., latency, end-to-end error rate probability
- Efficient heuristics to scale up to more realistic setups
 - ▶ *eigen-decomposition* [Mechtri'16]
 - ▶ *column generation* [Nemhauser'12]
- Multiple domains network slicing (e.g., multi-InPs)
 - ▶ each InP finds its own resource provisioning solution
 - ▶ afterwards: coordination between InPs to obtain final provisioning solution
- Slot-by-slot provisioning: instead of considering all active time slots simultaneously, provision resources for the slice slot-by-slot → reduce time complexity

Perspectives

- Additional constraints, e.g., latency, end-to-end error rate probability
- Efficient heuristics to scale up to more realistic setups
 - ▶ *eigen-decomposition* [Mechtri'16]
 - ▶ *column generation* [Nemhauser'12]
- Multiple domains network slicing (e.g., multi-InPs)
 - ▶ each InP finds its own resource provisioning solution
 - ▶ afterwards: coordination between InPs to obtain final provisioning solution
- Slot-by-slot provisioning: instead of considering all active time slots simultaneously, provision resources for the slice slot-by-slot → reduce time complexity

Perspectives

- Additional constraints, e.g., latency, end-to-end error rate probability
- Efficient heuristics to scale up to more realistic setups
 - ▶ *eigen-decomposition* [Mechtri'16]
 - ▶ *column generation* [Nemhauser'12]
- Multiple domains network slicing (e.g., multi-InPs)
 - ▶ each InP finds its own resource provisioning solution
 - ▶ afterwards: coordination between InPs to obtain final provisioning solution
- Slot-by-slot provisioning: instead of considering all active time slots simultaneously, provision resources for the slice slot-by-slot → reduce time complexity

Journal papers

- (J1) Admission Control and Resource Provisioning for Prioritized Slice Requests with Uncertainties (IEEE TNSM 2021, submitted)
- (J2) Uncertainty-Aware Resource Provisioning for Network Slicing (IEEE TNSM 2021, published)
- (J3) Coverage-Aware Resource Provisioning Method for Network Slicing (IEEE/ACM TNET 2020, published)

Conference/Workshop Papers

- (C1) Foresighted Resource Provisioning for Network Slicing (IEEE HPSR 2021, to appear).
- (C2) Radio Resource Provisioning for Network Slicing with Coverage Constraints (IEEE ICC 2020, published)
- (C3) Aggregated Resource Provisioning for Network Slices (IEEE GLOBECOM 2018, published)

Project Deliverable

- (D1) Virtual Network Orchestration Framework and Algorithms (ANR MAESTRO-5G Project Deliverable D3.1, 2020)

Patents

- (P1) Method and Apparatus for Mapping Network Slices Onto Network Infrastructures With SLA Guarantee. WIPO Patent No. WO2020114608A1 (filed on Dec. 07, 2018 by Nokia Solutions and Networks)

BACK UP

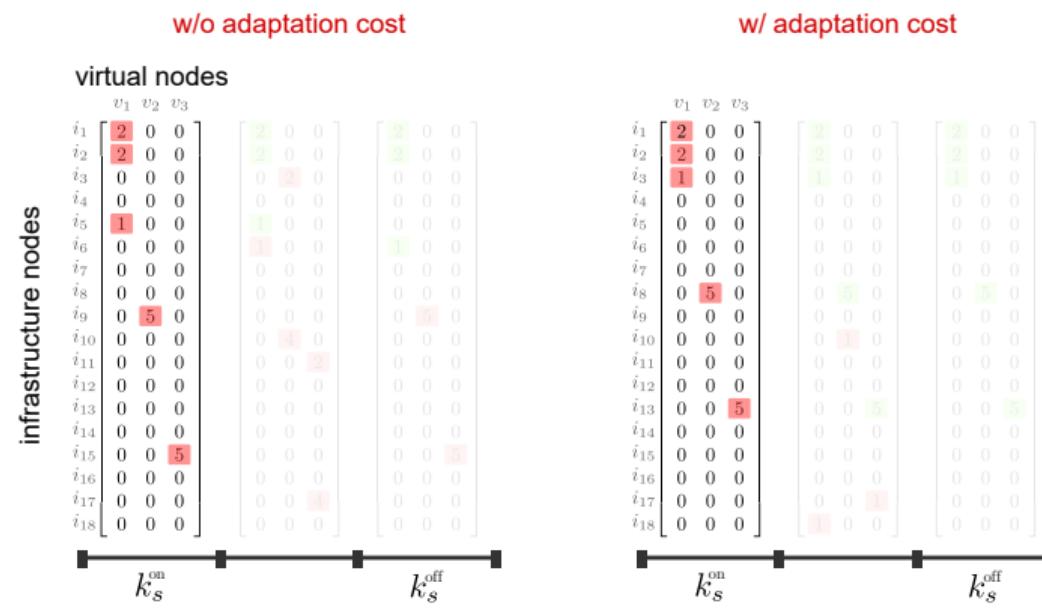
(4) Anticipated provisioning | Evaluation

Single slice provisioning

$$E_k^{\text{InP}} (\kappa_{\mathcal{R}_k}, d_{\mathcal{R}_k}) = \sum_{s \in \mathcal{R}_k} \sum_{\ell \in \mathcal{K}_s} \left(\underbrace{C_r(\kappa_{s,\ell})}_{\text{resource}} + \underbrace{C_f(\kappa_{s,\ell})}_{\text{node disposal}} + \underbrace{C_a(\kappa_{s,\ell}, \kappa_{s,\ell-1})}_{\text{adaptation}} \right)$$

Changes of $\kappa_{s,\ell}(i, v)$ along with ℓ :

- red: increase of $\kappa_{s,\ell}(i, v)$
- green: $\kappa_{s,\ell}(i, v)$ remains the same



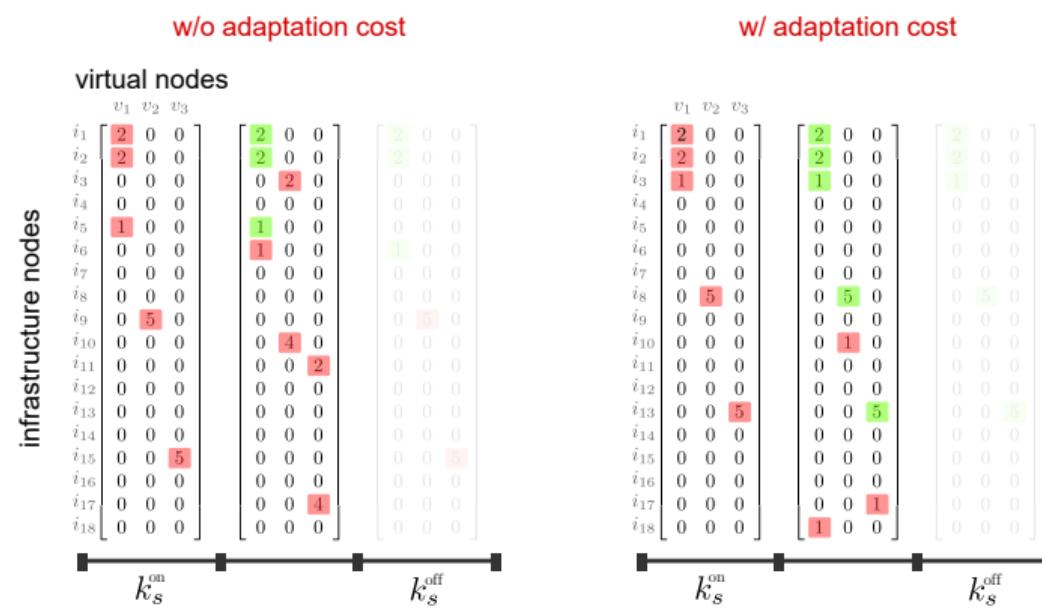
(4) Anticipated provisioning | Evaluation

Single slice provisioning

$$E_k^{\text{InP}} (\kappa_{\mathcal{R}_k}, d_{\mathcal{R}_k}) = \sum_{s \in \mathcal{R}_k} \sum_{\ell \in \mathcal{K}_s} \left(\underbrace{C_r(\kappa_{s,\ell})}_{\text{resource}} + \underbrace{C_f(\kappa_{s,\ell})}_{\text{node disposal}} + \underbrace{C_a(\kappa_{s,\ell}, \kappa_{s,\ell-1})}_{\text{adaptation}} \right)$$

Changes of $\kappa_{s,\ell}(i, v)$ along with ℓ :

- red: increase of $\kappa_{s,\ell}(i, v)$
- green: $\kappa_{s,\ell}(i, v)$ remains the same



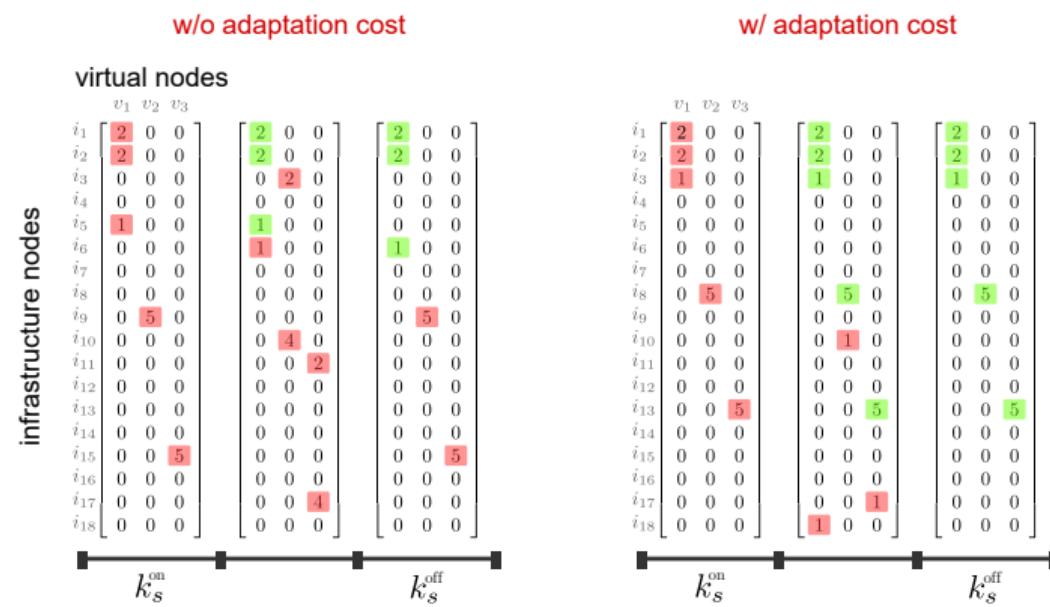
(4) Anticipated provisioning | Evaluation

Single slice provisioning

$$E_k^{\text{InP}} (\kappa_{\mathcal{R}_k}, d_{\mathcal{R}_k}) = \sum_{s \in \mathcal{R}_k} \sum_{\ell \in \mathcal{K}_s} \left(\underbrace{C_r(\kappa_{s,\ell})}_{\text{resource}} + \underbrace{C_f(\kappa_{s,\ell})}_{\text{node disposal}} + \underbrace{C_a(\kappa_{s,\ell}, \kappa_{s,\ell-1})}_{\text{adaptation}} \right)$$

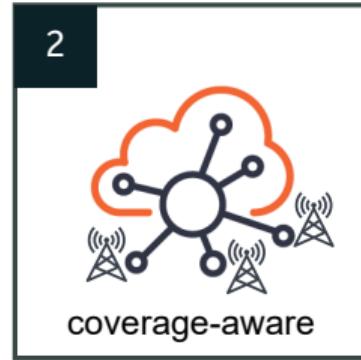
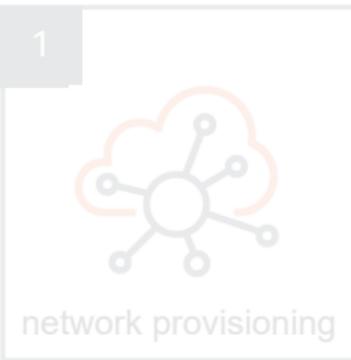
Changes of $\kappa_{s,\ell}(i, v)$ along with ℓ :

- red: increase of $\kappa_{s,\ell}(i, v)$
- green: $\kappa_{s,\ell}(i, v)$ remains the same

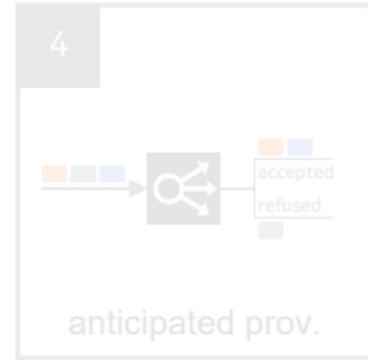


Contributions

Part I. Deterministic slice resource demands



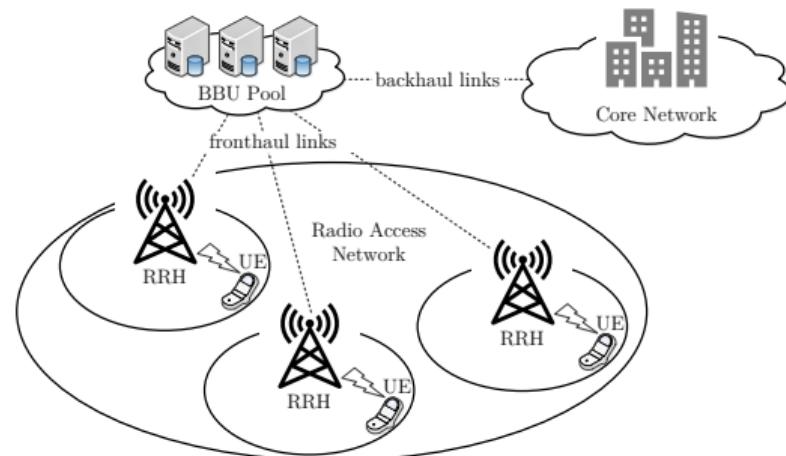
Part II. Random slice resource demands



(2) Coverage-aware | Problem statement

RP and NP problems

- **RP:** Radio resource Provisioning: radio
- **NP:** (other) Network resource Provisioning: computing, storage, and bandwidth



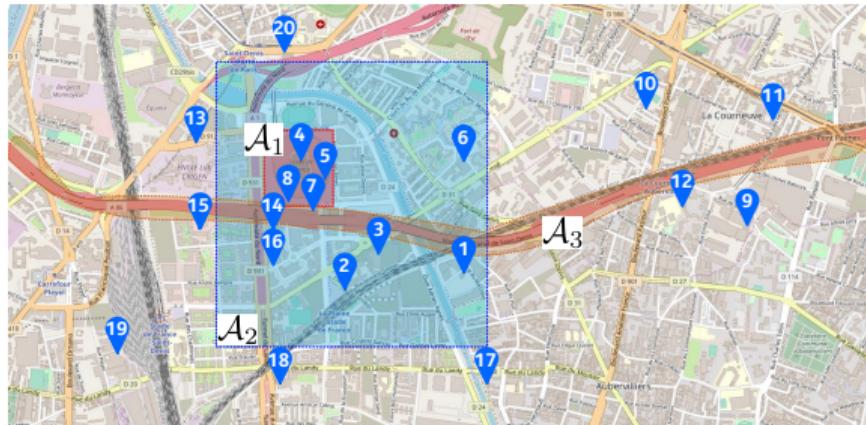
Cloud Radio Access Network (CRAN) general architecture

- RRH: Remote Radio Head
- BBU: Base Band Unit

(2) Coverage-aware | Radio provisioning (RP)

Problem statement

- SM-SLA involves:
 - ▶ **coverage constraints**: geographical distribution slice users. This distribution is described by the user density function $\rho_s(x)$, with $x \in \mathcal{A}$.
 - ▶ **supported service type**
 - ▶ **targeted QoS** such as a minimum average data rate $U_{s,u}$ and $U_{s,d}$ for the wireless uplink and downlink traffic for **each user**

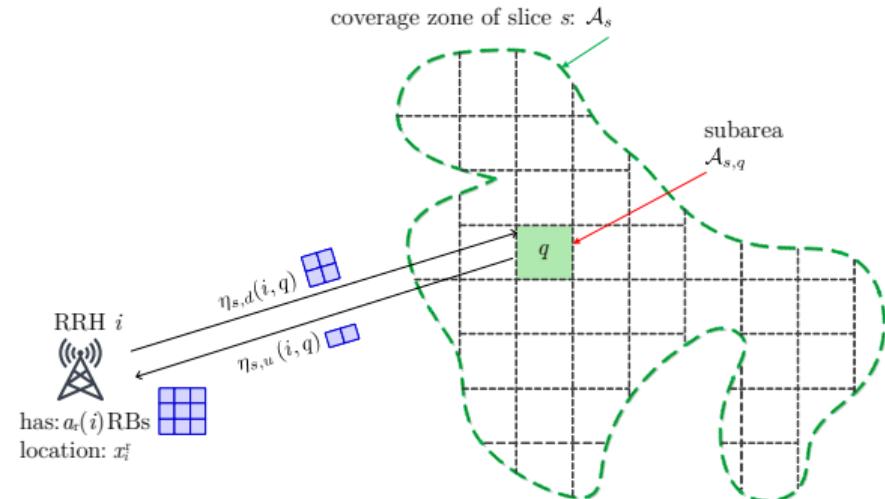


\mathcal{A} : the geographical area under study,
 \mathcal{A}_s : the subarea over which service s has to be made available

(2) Coverage-aware | Radio provisioning (RP)

Proposed approach for RP

- For slice s , InP has to provide a minimum data rate ($U_{s,u}$, $U_{s,d}$) to each user spread over \mathcal{A}_s with density $\rho_s(x)$
- Each RRH has a fixed amount $a_r(i)$ of resource blocks (RB) per time unit



- \mathcal{A}_s is partitioned into Q_s convex subareas $\mathcal{A}_{s,q}$,
 $q \in \mathcal{Q}_s = \{1, \dots, Q_s\}$
- Instead of allocating RBs to users, RRH nodes allocate RBs to subareas

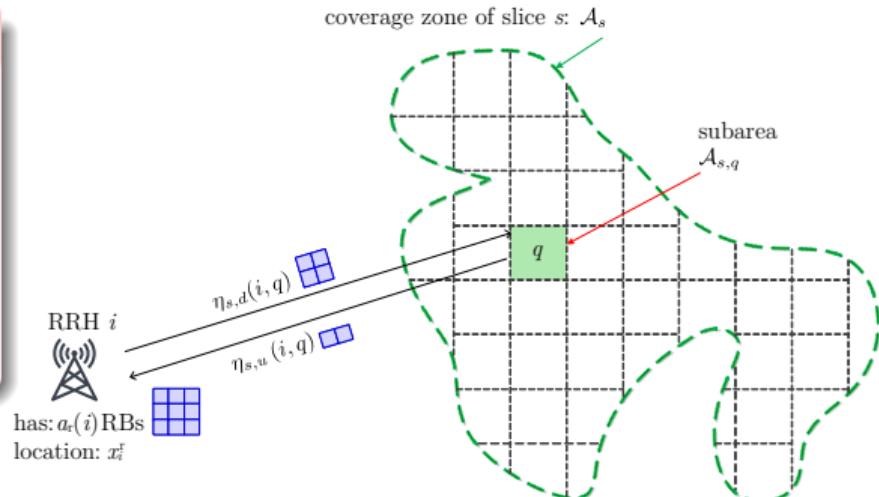
(2) Coverage-aware | Radio provisioning (RP)

RP cost function

Cost function of RP problem

$$\begin{aligned} c_{\text{rr}}(\boldsymbol{\eta}) = & \sum_{\sigma \in \mathcal{S}} \sum_{i \in \mathcal{N}_r} c_f(i) \tilde{\eta}_s(i) \\ & + \sum_{s \in \mathcal{S}} \sum_{i \in \mathcal{N}_r} \sum_{q \in \mathcal{Q}_s} [c_r(i) - \lambda b_u(x_i^r, \mathcal{A}_{s,q})] a_r(i) \eta_{s,u}(i, q) \\ & + \sum_{s \in \mathcal{S}} \sum_{i \in \mathcal{N}_r} \sum_{q \in \mathcal{Q}_s} [c_r(i) - \lambda b_d(x_i^r, \mathcal{A}_{s,q})] a_r(i) \eta_{s,d}(i, q) \end{aligned}$$

- $\eta_s(i, q) \in [0, 1]$: proportion of RBs that RRH i provisions to subarea q
- $\lambda b(x_i^r, \mathcal{A}_{s,q}) a_r(i) \eta_s(i, q)$: rate-related discount to penalty the use of far RRHs



(2) Coverage-aware | Radio provisioning (RP)

RP main constraints

(1) **RRH capacity limit:** $\forall i \in \mathcal{N}_r$,

$$\sum_{s \in \mathcal{S}} \sum_{q \in \mathcal{Q}_s} (\eta_{s,u}(i, q) + \eta_{s,d}(i, q)) \leq 1,$$

(2) **Rate requirement per user:** $\forall s \in \mathcal{S}, q \in \mathcal{Q}_s$,

$$\sum_{i \in \mathcal{N}_r} \eta_{s,u}(i, q) a_r(i) b_u(x_i^r, \mathcal{A}_{s,q}) \geq U_{s,u} \int_{\mathcal{A}_{s,q}} \rho_s(x) dx,$$

$$\sum_{i \in \mathcal{N}_r} \eta_{s,d}(i, q) a_r(i) b_d(x_i^r, \mathcal{A}_{s,q}) \geq U_{s,d} \int_{\mathcal{A}_{s,q}} \rho_s(x) dx,$$

(3) **Rate requirement per slice:** $\forall s \in \mathcal{S}$,

$$\sum_{q \in \mathcal{Q}_s} \sum_{i \in \mathcal{N}_r} \eta_{s,u}(i, q) a_r(i) b_u(x_i^r, \mathcal{A}_{s,q}) \geq R_{s,u}(v_r),$$

$$\sum_{q \in \mathcal{Q}_s} \sum_{i \in \mathcal{N}_r} \eta_{s,d}(i, q) a_r(i) b_d(x_i^r, \mathcal{A}_{s,q}) \geq R_{s,d}(v_r),$$

(4) **UL-DL proportionality:** $\forall q \in \mathcal{Q}_s$,

$$\frac{\eta_{s,u}(i, q) a_r(i) b_u(x_i^r, \mathcal{A}_{s,q})}{R_{s,u}(v_r)} = \frac{\eta_{s,d}(i, q) a_r(i) b_d(x_i^r, \mathcal{A}_{s,q})}{R_{s,d}(v_r)}.$$

(2) Coverage-aware | Single-step vs two-step provisioning

- Global provisioning problem (single-step):

$$\begin{aligned} \min \quad & c_{\text{tot}}(\boldsymbol{\eta}, \Phi) = c_{\text{rr}}(\boldsymbol{\eta}) + c_{\text{wr}}(\Phi) \\ \text{s.t.} \quad & \text{NP and RP constraints} \end{aligned} \quad (1)$$

- When the size of Φ and $\boldsymbol{\eta}$ increases, the problem may become intractable

→ Two-step provisioning algorithm: **CARP** (**Coverage-Aware Resource Provisioning**), which minimizes both terms of (1) separately.

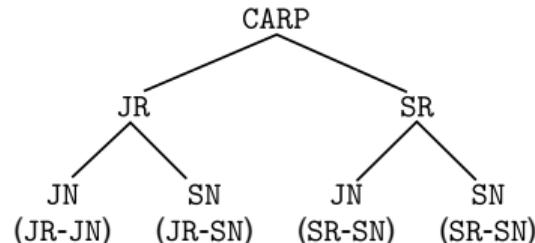


Figure: Four variants of CARP.

Infrastructure

Infrastructure model:

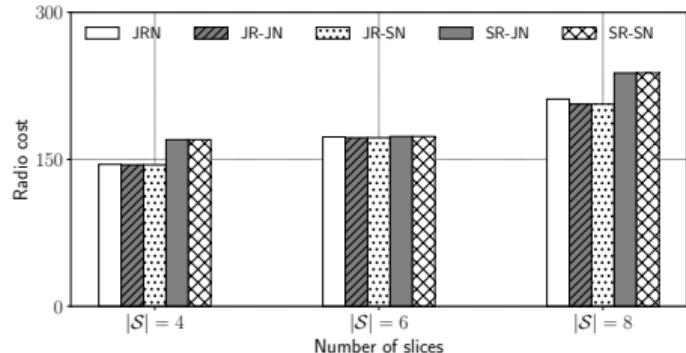
- k -ary fat tree topology
 - ▶ Leasing cost for each resource unit = 1
 - ▶ For coverage constraints, we consider a $1.43 \text{ km} \times 4.95 \text{ km}$ area around the Stade de France in Seine-Saint-Denis

Slices: three types:

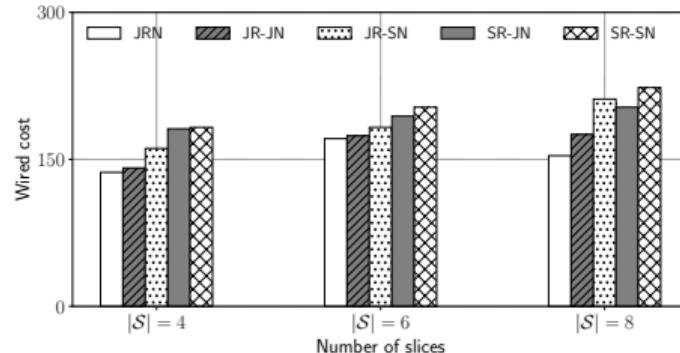
- HD video streaming
- SD video streaming
- Video surveillance and traffic monitoring

(2) Coverage-aware | Evaluation

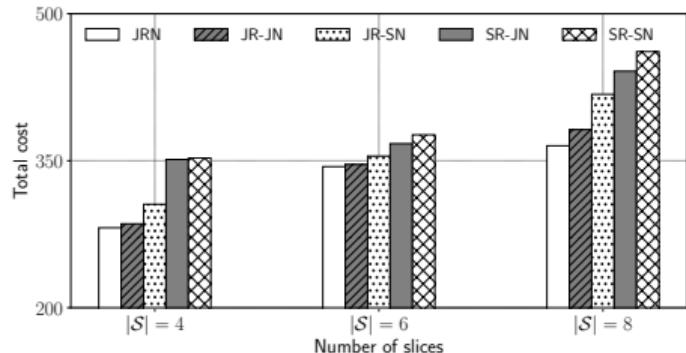
Performance of JRN vs 4 CARP variants.



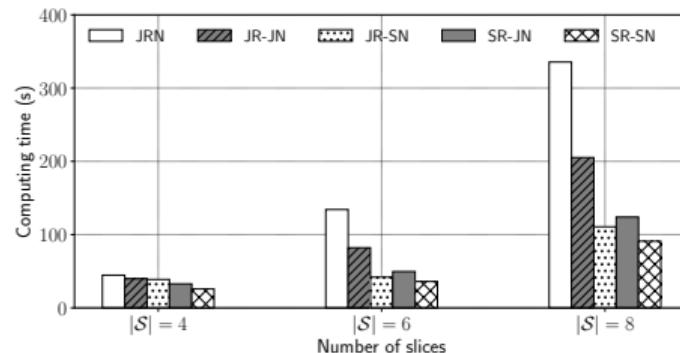
(a) Radio cost



(b) Wired cost



(c) Total cost



(d) Computing time

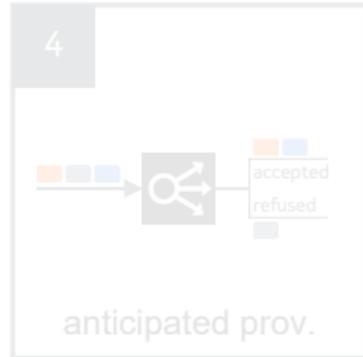
- Extension to Contribution 1: Slice resource provisioning accounting coverage constraints:
 - ▶ provide a minimum data rate for users in the geographical areas covered by slices
- Two-step approach to decompose the global optimization problem (JRN)
- JRN vs 4 variants of CARP:
 - ▶ least provisioning cost
 - ▶ but requires a much larger time complexity

Contributions

Part I. Deterministic slice resource demands



Part II. Random slice resource demands

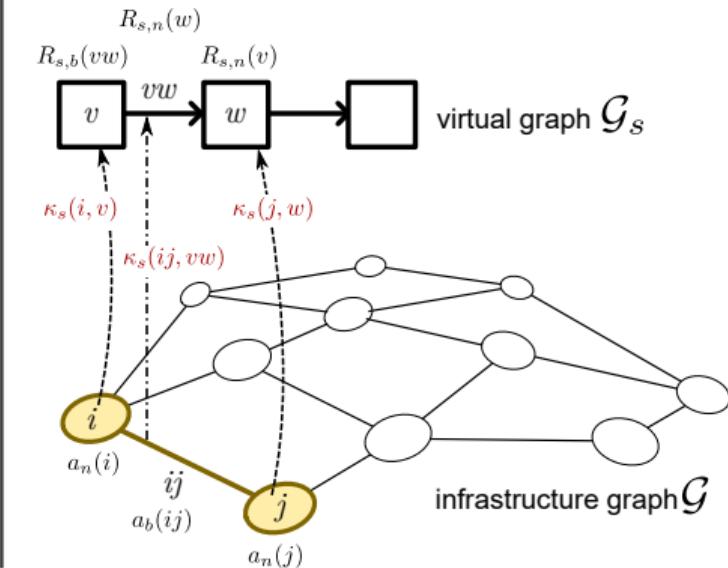


- Uncertainties may lead to slice QoS < expected level
 - ▶ S-RDs: traffic dynamics, e.g., flow arrival/departure
 - ▶ available infrastructure resources
- This contribution:
 - ▶ provisions infrastructure resources for network slices, while being **robust to a partly unknown number of users with a random usage** of the slice resources.
 - ▶ guarantees a predefined **service satisfaction probability (SSP)**
 - ▶ moreover, provisioning is performed so as to **limit its impact on background services**

(3) Uncertainty-aware | Variables

Resource provisioning for slice s

- $\kappa_s(i, v) \in \mathbb{N}$, $\kappa_s(ij, vw) \in \mathbb{N}$
- $\tilde{\kappa}_s(i, v) \in \{0, 1\}$ node mapping indicator function.
- $\tilde{\kappa}(i) \in \{0, 1\}$ indicator function for use of infrastructure node i .



(3) Uncertainty-aware | Static vs uncertainty-aware constraints

Slice resource demand constraints

Static constraints: S-RD $R_s = (R_{s,n}(v), R_{s,b}(vw))_{n \in \Upsilon, (v,vw) \in \mathcal{G}_s}^\top$ is deterministic

$$\sum_{i \in \mathcal{N}} \kappa_s(i, v) r_{s,n}(v) \geq R_{s,n}(v), \forall s \in \mathcal{S}, v \in \mathcal{N}_s, n \in \Upsilon$$

$$\sum_{ij \in \mathcal{E}} \kappa_s(ij, vw) r_{s,b}(vw) \geq R_{s,b}(vw), \forall s \in \mathcal{S}, vw \in \mathcal{E}_s$$

Uncertainty-aware constraints: S-RD R_s is random \rightarrow Service Satisfaction Probability (SSP) constraint:

$$p_s(\kappa_s) = \Pr \left\{ \begin{array}{l} \sum_{i \in \mathcal{N}} \kappa_s(i, v) r_{s,n}(v) \geq R_{s,n}(v), \forall v, n, \\ \sum_{ij \in \mathcal{E}} \kappa_s(ij, vw) r_{s,b}(vw) \geq R_{s,b}(vw), \forall vw \end{array} \right\} \geq \underline{p}_s, \forall s \in \mathcal{S}.$$

\underline{p}_s : minimum SSP for slice s .

(3) Uncertainty-aware | Static vs uncertainty-aware constraints

Slice resource demand constraints

Static constraints: S-RD $R_s = (R_{s,n}(v), R_{s,b}(vw))_{n \in \Upsilon, (v,vw) \in \mathcal{G}_s}^\top$ is deterministic

$$\sum_{i \in \mathcal{N}} \kappa_s(i, v) r_{s,n}(v) \geq R_{s,n}(v), \forall s \in \mathcal{S}, v \in \mathcal{N}_s, n \in \Upsilon$$

$$\sum_{ij \in \mathcal{E}} \kappa_s(ij, vw) r_{s,b}(vw) \geq R_{s,b}(vw), \forall s \in \mathcal{S}, vw \in \mathcal{E}_s$$

Uncertainty-aware constraints: S-RD R_s is random \rightarrow Service Satisfaction Probability (SSP) constraint:

$$p_s(\kappa_s) = \Pr \left\{ \begin{array}{l} \sum_{i \in \mathcal{N}} \kappa_s(i, v) r_{s,n}(v) \geq R_{s,n}(v), \forall v, n, \\ \sum_{ij \in \mathcal{E}} \kappa_s(ij, vw) r_{s,b}(vw) \geq R_{s,b}(vw), \forall vw \end{array} \right\} \geq \underline{p}_s, \forall s \in \mathcal{S}.$$

\underline{p}_s : minimum SSP for slice s .

Slice resource demand constraints

Relaxed SSP constraints:

$$\sum_{i \in \mathcal{N}} \kappa_s(i, v) r_{s,n}(v) \geq \bar{R}_{s,n}(v) + \gamma_s \tilde{R}_{s,n}(v), \forall s \in \mathcal{S}, v \in \mathcal{N}_s, n \in \Upsilon$$

$$\sum_{ij \in \mathcal{E}} \kappa_s(ij, vw) r_{s,b}(vw) \geq \bar{R}_{s,b}(vw) + \gamma_s \tilde{R}_{s,b}(vw), \forall s \in \mathcal{S}, vw \in \mathcal{E}_s$$

The SSP becomes

$$p_s(\gamma_s) = \Pr \left\{ \begin{array}{l} \hat{R}_{s,n}(v, \gamma_s) \geq R_{s,n}(v), \forall s \in \mathcal{S}, v \in \mathcal{N}_s, n \in \Upsilon \\ \hat{R}_{s,b}(vw, \gamma_s) \geq R_{s,b}(vw), \forall s \in \mathcal{S}, vw \in \mathcal{E}_s \end{array} \right\}$$

A sufficiently large γ_s would satisfy the SSP constraint.

Slice resource demand constraints

Finding γ_s : The smallest value of γ_s such that $p_s(\gamma_s) \geq \underline{p}_s$

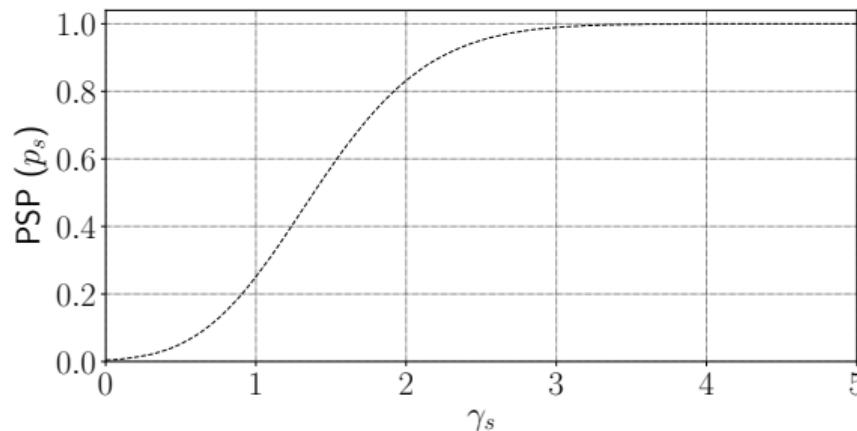


Figure: Evolution of p_s as function of γ_s .

γ_s can then be found by applying a bisection search¹.

¹R. L. Burden and J. Douglas Faires, *Numerical Analysis*, 9th ed. Brooks/Cole, Cengage Learning, 2011.

(3) Uncertainty-aware | Static vs uncertainty-aware constraints

Impact on background service (ImP) constraints

Static constraints: Infrastructure resources ($a_n(i)$, $a_b(ij)$) are deterministic

$$\sum_{s \in S} \sum_{v \in \mathcal{N}_s} \kappa_s(i, v) r_{s,n}(v) \leq a_n(i), \forall i \in \mathcal{N}, n \in \Upsilon,$$

$$\sum_{s \in S} \sum_{vw \in \mathcal{E}_s} \kappa_s(ij, vw) r_{s,b}(vw) \leq a_b(ij), \forall ij \in \mathcal{E},$$

Uncertainty-aware constraints: ($a_n(i)$, $a_b(ij)$) are random \rightarrow **Impact Probabilities (ImPs)** constraints:

$$p_n^{\text{im}}(i) \leq \bar{p}^{\text{im}}, \forall i \in \mathcal{N}, n \in \Upsilon,$$

$$p_b^{\text{im}}(ij) \leq \bar{p}^{\text{im}}, \forall ij \in \mathcal{E},$$

\bar{p}^{im} : minimum ImP on background services.

Relaxed ImP constraints:

$$\sum_{s,v} \kappa_s(i, v) r_{s,n}(v) \leq a_n(i) - [\bar{B}_n(i) + \gamma_B \tilde{B}_n(i)], \forall n \in \Upsilon, i \in \mathcal{N},$$

$$\sum_{s,vw} \kappa_s(ij, vw) r_{s,b}(vw) \leq a_b(ij) - [\bar{B}_b(ij) + \gamma_B \tilde{B}_b(ij)], \forall ij \in \mathcal{E},$$

\rightarrow Optimality: $\gamma_B = \Phi^{-1}(1 - \bar{p}^{\text{im}})$ guarantees $p_n^{\text{im}}(i) \leq \bar{p}^{\text{im}}$, $\forall i \in \mathcal{N}, n \in \Upsilon$, and $p_b^{\text{im}}(ij) \leq \bar{p}^{\text{im}}$, $\forall ij \in \mathcal{E}$.

Uncertainty-aware provisioning for multiple slices

$$\underset{\{d, \kappa\} = \{d_s, \kappa_s\}_{s \in \mathcal{S}}}{\text{maximize}} \quad \sum_{s \in \mathcal{S}} (l_s d_s - C_s(\kappa_s)),$$

$$\text{subject to } \sum_{i \in \mathcal{N}} \kappa_s(i, v) r_{s,n}(v) \geq \widehat{R}_{s,n}(v, \gamma_s) d_s, \forall s \in \mathcal{S}, v \in \mathcal{N}_s, n \in \Upsilon$$

$$\sum_{ij \in \mathcal{E}} \kappa_s(ij, vw) r_{s,b}(vw) \geq \widehat{R}_{s,b}(vw, \gamma_s) d_s, \forall s \in \mathcal{S}, vw \in \mathcal{E}_s,$$

$$\sum_{s \in \mathcal{S}} \sum_{v \in \mathcal{N}_s} \kappa_s(i, v) r_{s,n}(v) \leq a_n(i) - \widehat{B}_n(i, \gamma_B), \forall i \in \mathcal{N}, n \in \Upsilon,$$

$$\sum_{s \in \mathcal{S}} \sum_{vw \in \mathcal{E}_s} \kappa_s(ij, vw) r_{s,b}(vw) \leq a_b(ij) - \widehat{B}_b(ij, \gamma_B), \forall ij \in \mathcal{E}.$$

flow conservation constraint

(3) Uncertainty-aware | Evaluation

Setup

Table: Provisioning approaches.

	Jointly	Sequentially
background services not taken into account	JP	SP
background services accounted	JP-B	SP-B

- Background services: consume on average **20%** of available infrastructure resources (**5%** std. dev.)
- Types of slices:
 - (1) HD video streaming at 4 Mbps. $N_s \sim \mathcal{B}(300, 0.9)$, $I_s = 900$, $\underline{p}_s = 0.99$
 - (2) SD video streaming at 2 Mbps. $N_s \sim \mathcal{B}(1000, 0.8)$, $I_s = 1000$, $\underline{p}_s = 0.95$
 - (3) Video surveillance and traffic monitoring at 1 Mbps. $N_s = 50$, $I_s = 800$, $\underline{p}_s = 0.9$

(3) Uncertainty-aware | Evaluation

Provisioning of a single slice

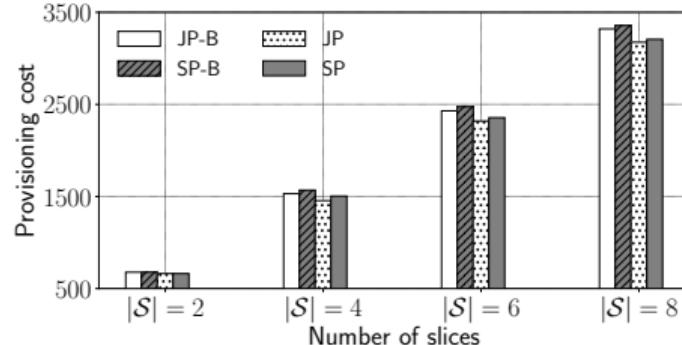
- Provisioning of a single slice of Type 1, with $\underline{p}_s = 0.99$ and $\bar{p}^{\text{im}} = 0.1$
- SP-B vs and SP (w/ and w/o background traffic):
 - ▶ SP has a **lower provisioning cost**, and yields a **higher earning**, but has a **higher impact** on background services

Table: SP-B vs and SP (w/ and w/o background traffic).

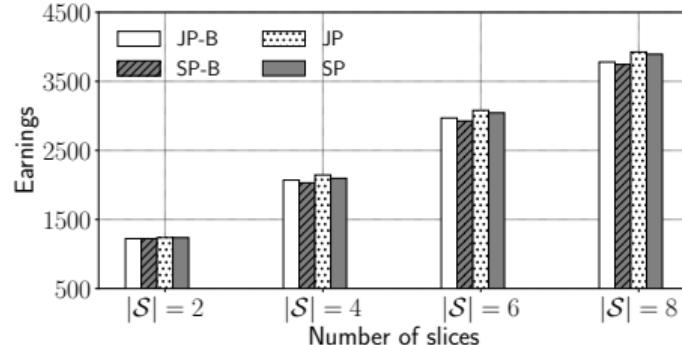
Criteria	SP-B	SP
Maximal p^{im}	1.26e-4	0.58
Provisioning cost	332	326
Total earnings	568	574
#impacted nodes	0	1
#impacted links	0	0

(3) Uncertainty-aware | Evaluation

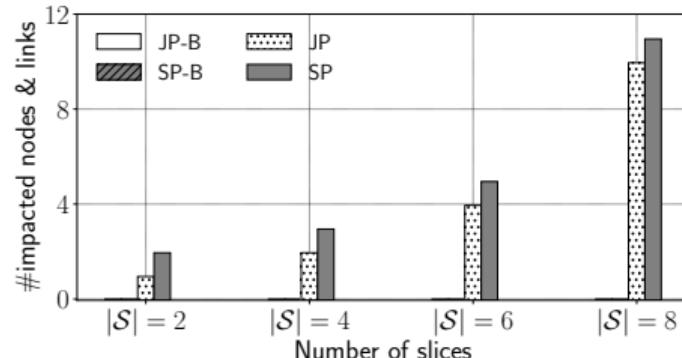
Provisioning of several slices of different types



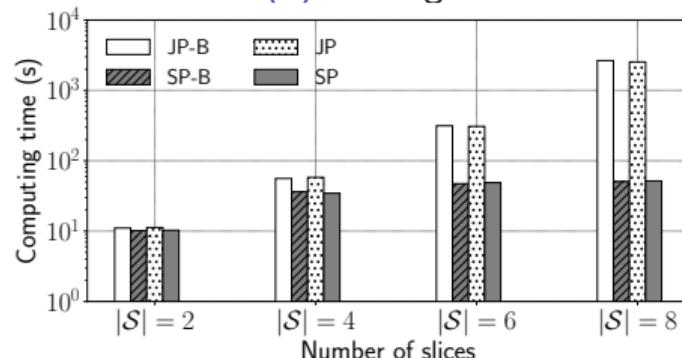
(a) Provisioning cost



(b) Earnings



(c) #impacted nodes & links



(d) Computing time

(3) Uncertainty-aware | Evaluation

Benefits of the uncertainty-aware slice resource provisioning

- **UPE:** uncertainty-aware provisioning and embedding solution
- **DPE:** deterministic provisioning and embedding solution
- S-RD: nb of users associated to the slice follows a binomial distribution $\mathcal{B}(m, p)$
 - ▶ m is fixed to 300
 - ▶ p varies from 0.4 to 0.9

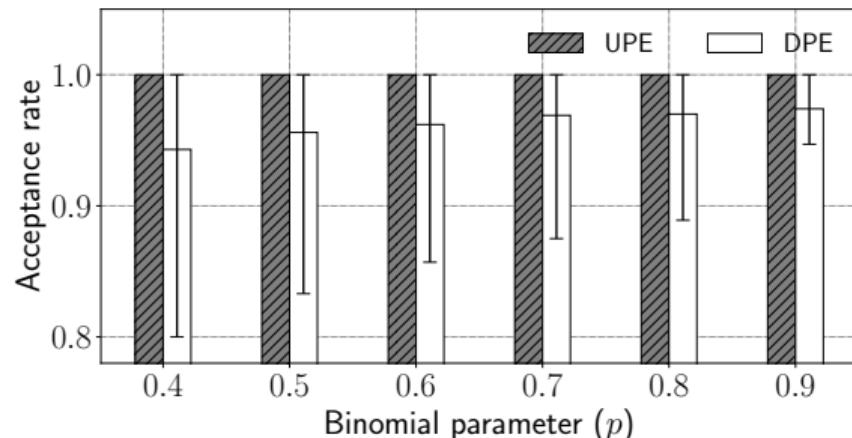


Figure: UPE vs DPE in terms of SFC acceptance rate.

- Resource provisioning method for network slicing that maximizes
 - ▶ acceptance rate,
 - ▶ earnings of the InP
- Advantages
 - ▶ robustness to uncertainties
 - ▶ keep the impact on the background services under a threshold imposed by the InP.

References

-  Cohen, R., Lewin-Eytan, L., Naor, J. S., and Raz, D. (2015).
Near Optimal Placement of Virtual Network Functions.
In *Proc. IEEE INFOCOM*, pages 1346–1354.
-  GSM Alliance (2017).
An Introduction to Network Slicing.
White Paper.
-  Liang, C. and Yu, F. R. (2014).
Wireless Network Virtualization: A Survey, Some Research Issues and Challenges.
IEEE Commun. Surveys Tuts., pages 1–24.
-  Liu, J., Lu, W., Zhou, F., Lu, P., and Zhu, Z. (2017).
On Dynamic Service Function Chain Deployment and Readjustment.
IEEE Trans. Netw. Service Manag., 14(3):543–553.
-  Mechtri, M., Ghribi, C., and Zeghlache, D. (2016).
A Scalable Algorithm for the Placement of Service Function Chains.
IEEE Trans. Netw. Service Manag., 13(3):533–546.
-  Nemhauser, G. (2012).
Column Generation for Linear and Integer Programming.
Optimization Stories, I:65–73.
-  Riera, J. F., Batalle, J., Bonnet, J., Dias, M., McGrath, M., Petralia, G., Liberati, F., Giuseppi, A., Pietrabissa, A., Ceselli, A., Petrini, A., Trubian, M., Papadimitrou, P., Dietrich, D., Ramos, A., Melian, J., Xilouris, G., Kourtis, A., Kourtis, T., and Markakis, E. K. (2016).
TeNOR: Steps towards an orchestration platform for multi-PoP NFV deployment.
In *IEEE NetSoft Conference and Workshops: Software-Defined Infrastructure for Networks, Clouds, IoT and Services*, pages 243–250.
-  Riggio, R., Bradai, A., Harutyunyan, D., Rasheed, T., and Ahmed, T. (2016).