

Walchand College of Engineering, Sangli
Computer Science & Engineering
Third Year
Course: Design and analysis of algorithm Lab
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Week 6 Assignment

Greedy method

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Prn: 2020BTECS00051

To apply Greedy method to solve problems of

1) Job sequencing with deadlines

1.A) Generate table of feasible, processing sequencing, profit.

1.B) What is the solution generated by the function JS when $n=7$, $(p_1, p_2, \dots, p_7) = (3, 5, 20, 18, 1, 6, 30)$ and $(d_1, d_2, d_3, \dots, d_7) = (1, 3, 4, 3, 2, 1, 2)$?

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Assignment 6.
DAA Lab

1.A) Generate table of feasible, processing sequencing, Profit

1.B) What is solution generated by function JS when $n=7$, $(p_1, p_2, \dots, p_7) = (3, 5, 20, 18, 1, 6, 30)$ and $(d_1, d_2, \dots, d_7) = (1, 3, 4, 3, 2, 1, 2)$?

→ Given.

	T_1	T_2	T_3	T_4	T_5	T_6	T_7
Profit	3	5	20	18	1	6	30
Deadline	1	3	4	3	2	1	2

Sort jobs as per decreasing order of profit

	T_7	T_3	T_4	T_6	T_2	T_1	T_5
Profit	30	20	18	6	5	3	1
Deadline	2	4	3	1	3	1	2

Maximum deadline is 4.
Therefore create 4 slots.
Now allocate Jobs to highest slot, starting from job with highest profit.

Select Job 7 - allocate to slot 1
slot 2 Job T_6 T_7 T_4 T_3

Select Job 3 - allocate to slot 4
slot 4

Select Job 4 - allocate to slot 3
Select Job 6 - allocate to slot 1

Total profit = $30 + 20 + 18 + 6 = 74$.

1.C) **Input:** Five Jobs with following deadlines and profits

JobID	Deadline	Profit
a	2	100
b	1	19
c	2	27
d	1	25
e	3	15

Output: Following is maximum profit sequence of jobs:

c, a, e

1.D) Study and implement Disjoint set algorithm to reduce time complexity of JS from $O(n^2)$ to nearly $O(n)$

Brute force:

Algorithm:

- a. Sort all jobs in decreasing order of profit.
- b. Iterate on jobs in decreasing order of profit. For each job, do the following :
 - i. Find a time slot i , such that slot is empty and $i < \text{deadline}$ and i is greatest. Put the job in this slot and mark this slot filled.
 - ii. If no such i exists, then ignore the job.

Code:

```
#include<bits/stdc++.h>

using namespace std;

struct job{
    char jobId;
    int deadline;
    int profit;
};

bool compare(job a,job b){
    return (a.profit > b.profit);
}

int main(){
    job jobs[] =
    {'a',2,100},{'b',1,19},{'c',2,27},{'d',1,25},{'e',3,15}};
    int n=sizeof(jobs)/sizeof(jobs[0]);
```

```

// sort according to increasing order of
sort(jobs,jobs+n,compare);

int result[n]={0};
bool slot[n]={false};

for(int i=0;i<n;i++){
    for(int j=min(jobs[i].deadline,n)-1;j>=0;j--){
        if(slot[j]==false){
            result[j]=i;
            slot[j]=true;
            break;
        }
    }
}

cout<<"The jobs that can be performed within deadline
to maximize the profit :"<<endl;
for(int i=0;i<n;i++){
    if(slot[i]){
        cout<<jobs[result[i]].jobId<< " ";
    }
}
}

```

Complexity Analysis:

Time complexity: $O(n^2)$

Space Complexity: $O(n)$

Output:

```

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Windows PowerShell
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PS C:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ> cd "c:\Use
f ($?) { g++ A6Q1a.cpp -o A6Q1a } ; if ($?) { .\A6Q1a }
The jobs that can be performed within deadline to maximize the profit :
c a e
PS C:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ>

```

Disjoint set algorithm:**Algorithm:**

1. Sort all jobs in decreasing order of profit.
2. Initialize the result sequence as first job in sorted jobs.
3. Do following for remaining n-1 jobs
 If the current job can fit in the current result sequence without missing the deadline, add current job to the result. Else ignore the current job.

Code:

```
#include<bits/stdc++.h>

using namespace std;

struct Job{
    char id;
    int deadLine, profit;
};

struct DisjointSet{
    int *parent;

    DisjointSet(int n){
        parent = new int[n+1];

        for (int i = 0; i <= n; i++)
            parent[i] = i;
    }

    int find(int s){
        if (s == parent[s])
            return s;
        return parent[s] = find(parent[s]);
    }

    void merge(int u, int v){
        parent[v] = u;
    }
};

bool cmp(Job a, Job b){
```

```

        return (a.profit > b.profit);
    }

int main()
{
    Job arr[] = { { 'a', 2, 100 }, { 'b', 1, 19 }, { 'c',
2, 27 }, { 'd', 1, 25 }, { 'e', 3, 15 } };
    int n = sizeof(arr) / sizeof(arr[0]);

    sort(arr, arr + n, cmp);

    int maxDeadline = INT_MIN;
    for (int i = 0; i < n; i++){
        maxDeadline = max(maxDeadline, arr[i].deadLine);
    }

    DisjointSet ds(maxDeadline);

    cout<<"The jobs that can be performed within deadline
to maximize the profit :"<<endl;
    for (int i = 0; i < n; i++){
        int availableSlot = ds.find(arr[i].deadLine);

        if (availableSlot > 0){
            ds.merge(ds.find(availableSlot -
1),availableSlot);
            cout << arr[i].id << " ";
        }
    }
    return 0;
}

```

Complexity Analysis:

Time complexity: $O(n \log d)$

n - total number of jobs

d - maximum possible deadline.

Space Complexity: $O(d)$

Output:

```
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PS C:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ> cd "c:\Use
f ($?) { g++ A6Q1b.cpp -o A6Q1b } ; if ($?) { .\A6Q1b }
The jobs that can be performed within deadline to maximize the profit :
a c e
PS C:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ>
```

2) To implement Fractional Knapsack problem 3 objects (n=3).

$(w_1, w_2, w_3) = (18, 15, 10)$

$(p_1, p_2, p_3) = (25, 24, 15)$

$M=20$

With strategy

a) Largest-profit strategy

Algorithm:

1. Sort the vector according to decreasing order of profit
2. Traverse the sorted vector till capacity of knapsack is not full
3. For every iteration check if the weight of current object is less than or equal to knapsack capacity, if yes then include it.
4. Else take the fractional part according to need.

Code:

```
#include<bits/stdc++.h>
using namespace std;

int main(){
    // {profit,weight}
    vector<pair<int,int>> v={{25,18},{24,15},{15,10}};
    int capacity=20;
    float ans=0;

    // Sort according to decreasing order of profit
    sort(v.begin(),v.end(),greater<>());

    int i=0;
    while(capacity!=0){
        // if capacity of knapsack is more than or
        equal to current weight then add it in knapsack
        if(capacity>=v[i].second){
            ans+=v[i].first;
```

```

        capacity-=v[i].second;
    }else{ //take the fractional part of it
        float x=(float)v[i].first/v[i].second;
        float y=x*capacity;
        ans+=y;
        capacity-=x;
    }
    i++;
}

cout<<"The maximum profit is "<<ans<<endl;
}

```

Complexity Analysis:

Time Complexity: $O(n \log n)$

Space Complexity: $O(1)$

Output:

```

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PS C:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ> cd "c:\Users\tr
f ($?) { g++ A6Q2a.cpp -o A6Q2a } ; if ($?) { .\A6Q2a }
The maximum profit is 28.2
PS C:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ>

```

b) Smallest-weight strategy

Algorithm:

1. Sort the vector according to increasing order of weight
2. Traverse the sorted vector till capacity of knapsack is not full
3. For every iteration check if the weight of current object is less than or equal to knapsack capacity, if yes then include it.
4. Else take the fractional part according to need.

Code:

```

#include<bits/stdc++.h>
using namespace std;

int main(){
    // {weight,profit}
    vector<pair<int,int>> v={{18,25},{15,24},{10,15}};

```

```

int capacity=20;
float ans=0;

// Sort according to increasing of weight
sort(v.begin(),v.end());

int i=0;
while(capacity!=0){
    // if capacity of knapsack is more than or
    equal to current weight then add it in knapsack
    if(capacity>=v[i].first){
        ans+=v[i].second;
        capacity-=v[i].first;
        // cout<<ans<<" "<<capacity<<endl;
    }else{ // take the fractional part of it
        int c=capacity;
        float x=(float)capacity/v[i].first;
        float y=x*v[i].second;

        ans+=y;
        capacity-=c;
    }
    i++;
}

cout<<"The maximum profit is "<<ans<<endl;
}

```

Complexity Analysis:

Time Complexity: $O(n \log n)$

Space Complexity: $O(1)$

Output:

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```
PS C:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ> cd "c:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ"
f ($?) { g++ A6Q2b.cpp -o A6Q2b } ; if ($?) { .\A6Q2b }
```

The maximum profit is 31

```
PS C:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ>
```


c) Largest profit-weight ratio strategy

Algorithm:

1. Sort the vector according to decreasing order of profit/weight ratio
2. Traverse the sorted vector till capacity of knapsack is not full
3. For every iteration check if the weight of current object is less than or equal to knapsack capacity, if yes then include it.
4. Else take the fractional part according to need.

Code:

```
#include<bits/stdc++.h>
using namespace std;

bool sorted(pair<int,int>& a,pair<int,int>& b){
    float d1=(float)a.first/a.second;
    float d2=(float)b.first/b.second;

    if(d1 > d2){
        return true;
    }
    return false;
}

int main(){
    // {profit,weight}
    vector<pair<int,int>> v={{25,18},{24,15},{15,10}};
    int capacity=20;
    float ans=0;

    // Sort according to decreasing order of
profit/weight ratio
    sort(v.begin(),v.end(),sorted);

    int i=0;
    while(capacity!=0){
        // if capacity of knapsack is more than or
equal to current weight then add it in knapsack
        if(capacity>=v[i].second){
            ans+=v[i].first;
            capacity-=v[i].second;
        }else{ //take fractional part of it
```

```

        int c=capacity;
        float x=(float)capacity/v[i].second;
        float y=x*v[i].first;

        ans+=y;
        capacity-=c;
    }
    i++;
}

cout<<"The maximum profit is "<<ans<<endl;
}

```

Complexity Analysis:

Time Complexity: $O(n \log n)$

Space Complexity: $O(1)$

Output:

```

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PS C:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ>
f ($?) { g++ A6Q2c.cpp -o A6Q2c } ; if ($?) { .\A6Q2c }
The maximum profit is 31.5
PS C:\Users\trupti patil\OneDrive\Desktop\ACADEMICS\SEM5\DAA\ExpQ>

```