

1. Fibonacci Series (Recursive & Non-Recursive)

Aim:

To write recursive and non-recursive Python programs to generate Fibonacci numbers and analyze time and space complexity.

Theory:

The Fibonacci sequence is a series where each term is the sum of the previous two terms: $F(n) = F(n-1) + F(n-2)$. Base cases are $F(0)=0$ and $F(1)=1$.

Approaches:

1. Recursive: Uses function calls repeatedly.
2. Iterative: Uses loop and variables to calculate terms efficiently.

Complexity:

Recursive → Time: $O(2^n)$, Space: $O(n)$

Iterative → Time: $O(n)$, Space: $O(1)$

Sample Output:

Fibonacci Series: 0 1 1 2 3 5 8 ...

Viva Questions:

1. What is the Fibonacci series?
→ It is a series where each term is the sum of the previous two terms.
2. What is the difference between recursion and iteration?
→ Recursion calls itself; iteration uses loops.
3. Why is recursive Fibonacci inefficient?
→ It repeats computations, giving $O(2^n)$ time.
4. What are base cases in recursion?
→ $F(0)=0$, $F(1)=1$.
5. Which method is faster?
→ Iterative, because it uses $O(n)$ time and $O(1)$ space.

2. Huffman Encoding using Greedy Strategy

Aim:

To implement Huffman Encoding for data compression using the greedy method.

Theory:

Huffman coding is a lossless compression algorithm assigning shorter codes to frequent symbols and longer codes to rare ones.

Steps:

1. Count frequency of each symbol.
2. Build a min-heap.
3. Merge two smallest nodes repeatedly.
4. Assign binary codes (0-left, 1-right).

Complexity:

Time: $O(n \log n)$

Space: $O(n)$

Sample Output:

Character	Frequency	Code
A	5	0
B	2	10
C	1	11

Viva Questions:

1. What is Huffman encoding?
→ A greedy-based compression technique.
2. Why is it called greedy?
→ It always combines the smallest frequency nodes first.
3. Is Huffman coding lossless?
→ Yes, it preserves original data.
4. What is the time complexity?
→ $O(n \log n)$.
5. What type of tree is used?
→ Binary tree.

3. Fractional Knapsack using Greedy Method

Aim:

To solve the fractional knapsack problem using a greedy approach to maximize total profit.

Theory:

In this problem, items can be broken into fractions. The greedy approach selects items based on highest value-to-weight ratio.

Steps:

1. Input items, weights, and values.
2. Compute value/weight ratio.
3. Sort by ratio.
4. Pick items or fractions until full.

Complexity:

Time: $O(n \log n)$

Space: $O(n)$

Sample Output:

Maximum value in knapsack = 240.0

Viva Questions:

1. What is the greedy choice here?
→ Selecting the item with the highest value-to-weight ratio.
2. What is the difference between fractional and 0/1 knapsack?
→ Fractional allows partial items.
3. Time complexity?
→ $O(n \log n)$.
4. Can greedy solve 0/1 knapsack?
→ No, it fails there.
5. Why is greedy optimal here?
→ Because it satisfies the greedy-choice property.

4. 0/1 Knapsack using Dynamic Programming**Aim:**

To solve the 0/1 Knapsack problem using dynamic programming for optimal profit.

Theory:

Items are indivisible. DP builds a table $dp[i][w]$ storing max profit using first i items and weight limit w .

Recurrence Relation:

$dp[i][w] = \max(\text{value}[i-1] + dp[i-1][w-\text{weight}[i-1]], dp[i-1][w])$ if $\text{weight}[i-1] \leq w$

Complexity:

Time: $O(nW)$
Space: $O(nW)$

Sample Output:

Maximum Profit = 220

Viva Questions:

1. What does 0/1 mean?
→ Either take or leave the item.
2. Why DP and not greedy?
→ Greedy fails for 0/1 case.
3. What is the time complexity?
→ $O(nW)$.
4. What is stored in DP table?
→ Maximum profit for subproblems.
5. What principle is used?
→ Principle of optimality.

5. N-Queens Problem using Backtracking

Aim:

To place N queens on an $N \times N$ chessboard so that no two queens attack each other.

Theory:

N-Queens is a constraint satisfaction problem solved using backtracking. Queens must not share row, column, or diagonal.

Steps:

1. Place first queen in a safe position.
2. Try placing next queen.
3. If unsafe, backtrack and try another position.
4. Repeat until all queens are placed.

Complexity:

Time: $O(N!)$

Space: $O(N^2)$

Sample Output:

One possible solution (for $N=4$):

```
0 1 0 0
0 0 0 1
1 0 0 0
0 0 1 0
```

Viva Questions:

1. What is the N-Queens problem?
→ To place N queens so none attack each other.
2. Which method is used?
→ Backtracking.
3. Time complexity?
→ $O(N!)$.
4. How is a safe position checked?
→ Check same row, column, and diagonals.
5. What is backtracking?
→ Trying all possibilities and undoing invalid moves.