

# Outline

- 1 Interpolating Polynomial Error Bound
- 2 Example: 2nd Lagrange Interpolating Polynomial Error Bound
- 3 Example: Interpolating Polynomial Error for Tabulated Data

# Use of the Interpolating Polynomial Error Bound

## Example: Tabulated Data

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- What step size  $h$  will ensure that linear interpolation gives an absolute error of at most  $10^{-6}$  for all  $x$  in  $[0, 1]$ ?

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Let  $x_0, x_1, \dots$  be the numbers at which  $f$  is evaluated,  $x$  be in  $[0, 1]$ , and suppose  $j$  satisfies  $x_j \leq x \leq x_{j+1}$ .

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Let  $x_0, x_1, \dots$  be the numbers at which  $f$  is evaluated,  $x$  be in  $[0, 1]$ , and suppose  $j$  satisfies  $x_j \leq x \leq x_{j+1}$ . The error bound theorem ▶ Theorem implies that the error in linear interpolation is

$$|f(x) - P(x)| = \left| \frac{f^{(2)}(\xi)}{2!} (x - x_j)(x - x_{j+1}) \right| = \frac{|f^{(2)}(\xi)|}{2} |(x - x_j)| |(x - x_{j+1})|$$

# Use of the Interpolating Polynomial Error Bound

## Solution (1/3)

The step size is  $h$ , so  $x_j = jh$ ,  $x_{j+1} = (j+1)h$ , and

$$|f(x) - P(x)| \leq \frac{|f^{(2)}(\xi)|}{2!} |(x - jh)(x - (j+1)h)|.$$



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Hence

$$\begin{aligned} |f(x) - P(x)| &\leq \frac{\max_{\xi \in [0,1]} e^{\xi}}{2} \max_{x_j \leq x \leq x_{j+1}} |(x - jh)(x - (j+1)h)| \\ &\leq \frac{e}{2} \max_{x_j \leq x \leq x_{j+1}} |(x - jh)(x - (j+1)h)|. \end{aligned}$$

# Use of the Interpolating Polynomial Error Bound

## Solution (2/3)

Consider the function  $g(x) = (x - jh)(x - (j + 1)h)$ , for  $jh \leq x \leq (j + 1)h$ .

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$$g'(x) = (x - (j + 1)h) + (x - jh) = 2 \left( x - jh - \frac{h}{2} \right),$$

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the only critical point for  $g$  is at  $x = jh + \frac{h}{2}$ , with

$$g \left( jh + \frac{h}{2} \right) = \left( \frac{h}{2} \right)^2 = \frac{h^2}{4}$$

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Since  $g(jh) = 0$  and  $g((j + 1)h) = 0$ , the maximum value of  $|g'(x)|$  in  $[jh, (j + 1)h]$  must occur at the critical point.

# Use of the Interpolating Polynomial Error Bound

## Solution (3/3)

This implies that

$$|f(x) - P(x)| \leq \frac{e}{2} \max_{x_j \leq x \leq x_{j+1}} |g(x)| \leq \frac{e}{2} \cdot \frac{h^2}{4} = \frac{eh^2}{8}.$$

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Consequently, to ensure that the the error in linear interpolation is bounded by  $10^{-6}$ , it is sufficient for  $h$  to be chosen so that

$$\frac{eh^2}{8} \leq 10^{-6}. \quad \text{This implies that} \quad h < 1.72 \times 10^{-3}.$$

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Because  $n = \frac{(1-0)}{h}$  must be an integer, a reasonable choice for the step size is  $h = 0.001$ .



Questions?

# Reference Material

# Generalized Rolle's Theorem

Suppose  $f \in C[a, b]$  is  $n$  times differentiable on  $(a, b)$ . If

$$f(x) = 0$$

at the  $n + 1$  distinct numbers  $a \leq x_0 < x_1 < \dots < x_n \leq b$ , then a number  $c$  in  $(x_0, x_n)$ , and hence in  $(a, b)$ , exists with

$$f^{(n)}(c) = 0$$

[◀ Return to Error Bound Theorem](#)

# The Lagrange Polynomial: Theoretical Error Bound

Suppose  $x_0, x_1, \dots, x_n$  are distinct numbers in the interval  $[a, b]$  and  $f \in C^{n+1}[a, b]$ . Then, for each  $x$  in  $[a, b]$ , a number  $\xi(x)$  (generally unknown) between  $x_0, x_1, \dots, x_n$ , and hence in  $(a, b)$ , exists with

$$f(x) = P(x) + \frac{f^{(n+1)}(\xi(x))}{(n+1)!} (x - x_0)(x - x_1) \cdots (x - x_n)$$

where  $P(x)$  is the interpolating polynomial given by

$$P(x) = f(x_0)L_{n,0}(x) + \cdots + f(x_n)L_{n,n}(x) = \sum_{k=0}^n f(x_k)L_{n,k}(x)$$

[◀ Return to Second Lagrange Interpolating Polynomial Example](#)

[◀ Return to Tabulated data example with  \$f\(x\) = e^x\$](#)

## A Recursive Method

We want to construct a  $n^{\text{th}}$ -degree polynomial that interpolates  $(n+1)$  points,  $(x_i, y_i)$ ,  $i=0, 1, \dots, n$ .

Let  $p_i(x)$  the  $i^{\text{th}}$ -degree polynomial that interpolates  $(x_k, y_k)$ ,  $k=0, 1, \dots, i$ . We have

$$p_0(x) = y_0$$

$$p_1(x) = y_0 + \frac{x-x_0}{x_1-x_0} (y_1 - y_0)$$

Suppose we have already obtained  $P_k(x)$ ,  $k \leq i$ .  
We can construct  $P_{i+1}(x)$  as follows

$$P_{i+1}(x) = P_i(x) + \frac{\pi_i(x)}{\pi_i(x_{i+1})} (y_{i+1} - P_i(x_{i+1}))$$

where

$$\pi_i(x) = \prod_{k=0}^i (x - x_k)$$

It can be easily verified that

(i)  $P_{n+1}(x_k) = y_k, \quad k=0, 1, \dots, n+1$

(ii)  $P_{n+1}(x)$  is an  $(n+1)^{\text{st}}$ -degree polynomial.

(iii) It is easy, by induction, to prove.

INPUT  $n, \{(x_i, y_i)\}_{i=0}^n, x$   
 $P = y_0$

For  $i = 0, 1, \dots, n-1$  repeat

$\Pi = 1$

For  $j = 0, 1, \dots, i$

$$\Pi = \Pi \times (x - x_j) / (x_{i+1} - x_j)$$

End

$$P = P + \Pi \times (y_{i+1} - P_i(x_{i+1}))$$

End

OUTPUT  $P$

STOP



INPUT  $n, \{(x_i, y_i)\}_{i=0}^n$

For  $i=0, 1, \dots, n$  repeat

$$P_0(x_i) = y_0$$

End

For  $i=0, 1, \dots, n-1$  repeat

For  $j=i+1, \dots, n$  repeat

$$\Pi = 1$$

For  $k=0, \dots, i$  repeat

$$\Pi = \Pi \times (x_j - x_k) / (x_{i+1} - x_k)$$

End

$$P_{i+1}(x_j) = P_i(x_j) + \Pi \times (y_{i+1} - P_i(x_{i+1}))$$

End

End

OUTPUT  $P_i(x_{i+1}), i=0, 1, \dots, n-1$

STOP

(3)

$P_j(x_i)$	$x_0$	$x_1$	$x_2$	$x_3$	...			$x_{n-1}$	$x_n$
$P_0$		$P_0(x_1)$	$P_0(x_2)$	$P_0(x_3)$	...			$P_0(x_{n-1})$	$P_0(x_n)$
$P_1$			$P_1(x_2)$	$P_1(x_3)$	...			$P_1(x_{n-1})$	$P_1(x_n)$
$P_2$				$P_2(x_3)$	...			$P_2(x_{n-1})$	$P_2(x_n)$
					...				
					...				
					...				
					...				
					...				
					...				
$P_{n-2}$					...			$P_{n-2}(x_{n-1})$	$P_{n-2}(x_n)$
$P_{n-1}$					...				$P_{n-1}(x_n)$

# Interpolation & Polynomial Approximation

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## Hermite Interpolation I

Numerical Analysis (9th Edition)

R L Burden & J D Faires

Beamer Presentation Slides

prepared by

John Carroll

Dublin City University

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## 1 Osculating & Hermite Polynomials

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# Hermite Interpolation: Osculating Polynomials

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- **Osculating polynomials** generalize both the Taylor polynomials and the Lagrange polynomials.

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- Suppose that we are given  $n + 1$  distinct numbers  $x_0, x_1, \dots, x_n$  in  $[a, b]$  and nonnegative integers  $m_0, m_1, \dots, m_n$ , and  $m = \max\{m_0, m_1, \dots, m_n\}$ .

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- The osculating polynomial approximating a function  $f \in C^m[a, b]$  at  $x_i$ , for each  $i = 0, \dots, n$ , is the polynomial of **least degree** that has the same values as the function  $f$  and all its derivatives of order less than or equal to  $m_i$  at each  $x_i$ .

# Hermite Interpolation: Osculating Polynomials

## Osculating Polynomials (Cont'd)

The degree of this osculating polynomial is at most

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The degree of this osculating polynomial is at most

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because the number of conditions to be satisfied is  $\sum_{i=0}^n m_i + (n + 1)$ , and a polynomial of degree  $M$  has  $M + 1$  coefficients that can be used to satisfy these conditions.

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- The **osculating polynomial** approximating  $f$  is the polynomial  $P(x)$  of least degree such that

$$\frac{d^k P(x_i)}{dx^k} = \frac{d^k f(x_i)}{dx^k}$$

for each  $i = 0, 1, \dots, n$  and  $k = 0, 1, \dots, m_i$ .

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The osculating polynomial approximating  $f$  is the  $m_0$ th Taylor polynomial for  $f$  at  $x_0$  when  $n = 0$  and the  $n$ th Lagrange polynomial interpolating  $f$  on  $x_0, x_1, \dots, x_n$  when  $m_i = 0$  for each  $i$ .

# Hermite Polynomials

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- For a given function  $f$ , these polynomials agree with  $f$  at  $x_0, x_1, \dots, x_n$ .

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## Hermite Polynomials

- The case when  $m_i = 1$ , for each  $i = 0, 1, \dots, n$ , gives the **Hermite polynomials**.
- For a given function  $f$ , these polynomials agree with  $f$  at  $x_0, x_1, \dots, x_n$ .
- In addition, since their first derivatives agree with those of  $f$ , they have the same “shape” as the function at  $(x_i, f(x_i))$  in the sense that the **tangent lines** to the polynomial and the function agree.

# Outline

- 1 Osculating & Hermite Polynomials
- 2 The Precise Form of the Hermite Polynomials**
- 3 Example: Constructing the Hermite Polynomial using Lagrange Polynomials

# Precise Form of the Hermite Polynomials

## Theorem

If  $f \in C^1[a, b]$  and  $x_0, \dots, x_n \in [a, b]$  are distinct, the unique polynomial of least degree agreeing with  $f$  and  $f'$  at  $x_0, \dots, x_n$  is the Hermite polynomial of degree at most  $2n + 1$  given by

$$H_{2n+1}(x) = \sum_{j=0}^n f(x_j) H_{n,j}(x) + \sum_{j=0}^n f'(x_j) \hat{H}_{n,j}(x)$$

where, for  $L_{n,j}(x)$  denoting the  $j$ th Lagrange coefficient polynomial of degree  $n$ , we have

$$H_{n,j}(x) = [1 - 2(x - x_j)L'_{n,j}(x_j)]L_{n,j}^2(x) \quad \text{and} \quad \hat{H}_{n,j}(x) = (x - x_j)L_{n,j}^2(x)$$

Continued on the next slide ...



# Precise Form of the Hermite Polynomials

$$H_{2n+1}(x) = \sum_{j=0}^n f(x_j) H_{n,j}(x) + \sum_{j=0}^n f'(x_j) \hat{H}_{n,j}(x)$$

## Theorem (Cont'd)

Moreover, if  $f \in C^{2n+2}[a, b]$ , then

$$f(x) = H_{2n+1}(x) + \frac{(x - x_0)^2 \dots (x - x_n)^2}{(2n+2)!} f^{(2n+2)}(\xi(x))$$

for some (generally unknown)  $\xi(x)$  in the interval  $(a, b)$ .

# Precise Form of the Hermite Polynomials

## Proof (1/4)

First recall that

$$L_{n,j}(x_i) = \begin{cases} 0, & \text{if } i \neq j \\ 1, & \text{if } i = j \end{cases}$$

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Hence when  $i \neq j$ ,

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whereas, for each  $i$ ,

$$\begin{aligned} H_{n,i}(x_i) &= [1 - 2(x_i - x_i)L'_{n,i}(x_i)] \cdot 1 = 1 \\ \text{and} \quad \hat{H}_{n,i}(x_i) &= (x_i - x_i) \cdot 1^2 = 0 \end{aligned}$$

# Precise Form of the Hermite Polynomials

## Proof (2/4)

- As a consequence

$$H_{2n+1}(x_i) = \sum_{\substack{j=0 \\ j \neq i}}^n f(x_j) \cdot 0 + f(x_i) \cdot 1 + \sum_{j=0}^n f'(x_j) \cdot 0 = f(x_i)$$

so  $H_{2n+1}$  agrees with  $f$  at  $x_0, x_1, \dots, x_n$ .

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so  $H_{2n+1}$  agrees with  $f$  at  $x_0, x_1, \dots, x_n$ .

- To show the agreement of  $H'_{2n+1}$  with  $f'$  at the nodes, first note that  $L_{n,j}(x)$  is a factor of  $H'_{n,j}(x)$ , so  $H'_{n,j}(x_i) = 0$  when  $i \neq j$ .

# Precise Form of the Hermite Polynomials

## Proof (3/4)

In addition, when  $i = j$  we have  $L_{n,i}(x_i) = 1$ , so

$$\begin{aligned}H'_{n,i}(x_i) &= -2L'_{n,i}(x_i) \cdot L_{n,i}^2(x_i) + [1 - 2(x_i - x_i)L'_{n,i}(x_i)]2L_{n,i}(x_i)L'_{n,i}(x_i) \\&= -2L'_{n,i}(x_i) + 2L'_{n,i}(x_i) = 0\end{aligned}$$

Hence,  $H'_{n,j}(x_i) = 0$  for all  $i$  and  $j$ .

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Hence,  $H'_{n,j}(x_i) = 0$  for all  $i$  and  $j$ .

Finally,

$$\begin{aligned} \hat{H}'_{n,j}(x_i) &= L_{n,j}^2(x_i) + (x_i - x_j)2L_{n,j}(x_i)L'_{n,j}(x_i) \\ &= L_{n,j}(x_i)[L_{n,j}(x_i) + 2(x_i - x_j)L'_{n,j}(x_i)] \end{aligned}$$

so  $\hat{H}'_{n,j}(x_i) = 0$  if  $i \neq j$  and  $\hat{H}'_{n,i}(x_i) = 1$ .



# Precise Form of the Hermite Polynomials

## Proof (4/4)

Combining these facts, we have

$$H'_{2n+1}(x_i) = \sum_{j=0}^n f(x_j) \cdot 0 + \sum_{\substack{j=0 \\ j \neq i}}^n f'(x_j) \cdot 0 + f'(x_i) \cdot 1 = f'(x_i)$$

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Therefore,  $H_{2n+1}$  agrees with  $f$  and  $H'_{2n+1}$  with  $f'$  at  $x_0, x_1, \dots, x_n$ .

# Outline

- 1 Osculating & Hermite Polynomials
- 2 The Precise Form of the Hermite Polynomials
- 3 Example: Constructing the Hermite Polynomial using Lagrange Polynomials

# Constructing the Hermite Polynomial

## Example: Constructing $H_5(x)$

Use the Hermite polynomial that agrees with the data listed in the following table to find an approximation to  $f(1.5)$ .

$k$	$x_k$	$f(x_k)$	$f'(x_k)$
0	1.3	0.6200860	-0.5220232
1	1.6	0.4554022	-0.5698959
2	1.9	0.2818186	-0.5811571

# Constructing the Hermite Polynomial $H_5(x)$

## Solution (1/5)

We first compute the Lagrange polynomials and their derivatives.

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$$L_{2,0}(x) = \frac{(x - x_1)(x - x_2)}{(x_0 - x_1)(x_0 - x_2)} = \frac{50}{9}x^2 - \frac{175}{9}x + \frac{152}{9}$$

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$$L'_{2,1}(x) = \frac{-200}{9}x + \frac{320}{9}$$

# Constructing the Hermite Polynomial $H_5(x)$

## Solution (2/5)

and, finally

$$L_{2,2} = \frac{(x - x_0)(x - x_1)}{(x_2 - x_0)(x_2 - x_1)} = \frac{50}{9}x^2 - \frac{145}{9}x + \frac{104}{9}$$

# Constructing the Hermite Polynomial $H_5(x)$

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# Constructing the Hermite Polynomial $H_5(x)$

## Solution (3/5)

The polynomials  $H_{2,j}(x)$  and  $\hat{H}_{2,j}(x)$  are then

$$H_{2,0}(x) = [1 - 2(x - 1.3)(-5)] \left( \frac{50}{9}x^2 - \frac{175}{9}x + \frac{152}{9} \right)^2$$

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$$H_{2,1}(x) = 1 \cdot \left( \frac{-100}{9}x^2 + \frac{320}{9}x - \frac{247}{9} \right)^2$$

$$H_{2,2}(x) = 10(2 - x) \left( \frac{50}{9}x^2 - \frac{145}{9}x + \frac{104}{9} \right)^2$$

# Constructing the Hermite Polynomial $H_5(x)$

## Solution (4/5)

$$\hat{H}_{2,0}(x) = (x - 1.3) \left( \frac{50}{9}x^2 - \frac{175}{9}x + \frac{152}{9} \right)^2$$



# Constructing the Hermite Polynomial $H_5(x)$

## Solution (4/5)

$$\begin{aligned}\hat{H}_{2,0}(x) &= (x - 1.3) \left( \frac{50}{9}x^2 - \frac{175}{9}x + \frac{152}{9} \right)^2 \\ \hat{H}_{2,1}(x) &= (x - 1.6) \left( \frac{-100}{9}x^2 + \frac{320}{9}x - \frac{247}{9} \right)^2\end{aligned}$$

# Constructing the Hermite Polynomial $H_5(x)$

## Solution (4/5)

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$$\hat{H}_{2,1}(x) = (x - 1.6) \left( \frac{-100}{9}x^2 + \frac{320}{9}x - \frac{247}{9} \right)^2$$

$$\hat{H}_{2,2}(x) = (x - 1.9) \left( \frac{50}{9}x^2 - \frac{145}{9}x + \frac{104}{9} \right)^2$$

# Constructing the Hermite Polynomial $H_5(x)$

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$$\begin{aligned} H_5(x) = & 0.6200860H_{2,0}(x) + 0.4554022H_{2,1}(x) + 0.2818186H_{2,2}(x) \\ & - 0.5220232\hat{H}_{2,0}(x) - 0.5698959\hat{H}_{2,1}(x) \\ & - 0.5811571\hat{H}_{2,2}(x) \end{aligned}$$

so that

$$\begin{aligned} H_5(1.5) = & 0.6200860 \left( \frac{4}{27} \right) + 0.4554022 \left( \frac{64}{81} \right) + 0.2818186 \left( \frac{5}{81} \right) \\ & - 0.5220232 \left( \frac{4}{405} \right) - 0.5698959 \left( \frac{-32}{405} \right) \\ & - 0.5811571 \left( \frac{-2}{405} \right) = 0.5118277 \end{aligned}$$

a result that is accurate to the places listed.

# Constructing the Hermite Polynomial

## Observation

Although the theorem provides a complete description of the Hermite polynomials, it is clear from this example that the need to determine and evaluate the Lagrange polynomials and their derivatives makes the procedure tedious even for small values of  $n$ .

# Constructing the Hermite Polynomial

## Observation

Although the theorem provides a complete description of the Hermite polynomials, it is clear from this example that the need to determine and evaluate the Lagrange polynomials and their derivatives makes the procedure tedious even for small values of  $n$ .

## Remedy

We will turn to an alternative method for generating Hermite approximations that has as its basis the Newton interpolatory divided-difference formula at  $x_0, x_1, \dots, x_n$ , that is,

$$P_n(x) = f[x_0] + \sum_{k=1}^n f[x_0, x_1, \dots, x_k](x - x_0) \cdots (x - x_{k-1})$$

Questions?