

1

The right side of the root node = 0 since $\min(y, 3, 1, 6) = y = 0$. This means X needs to be ≥ 0 since the root node is maximizing. However, the root of the left side is ≤ 3 because of one of its children = 3. The right side of the left side can therefore be pruned, since it is $\max(5, 6) = 6$ which is > 3 . Since the parent of x is maximizing, x needs to be \geq its sibling which is $\min(2, 1, 7) = 1$. This means that x needs to be ≤ 3 , and ≥ 1 .

2

Since $x=0$, it can be pruned as its sibling is greater ($\min(2, 1, 7) = 1$). The parent of x (1) is smaller than its siblings ($\min(3, 1, \max(5, 6)) = 1$). The left side of the tree is therefore = 1. A will be $\max(6, \min(3, 1)) = \max(6, 1) = 6$. Y therefore needs to be = 6 in order to become the same value as the right side. We now have 1 on the left, and 6 on the right side. $\max(1, 6) = 6$. Y is now equal to the root node (6).

3

Pruning A (=6) means that we need a better (smaller, since the parent of A is minimizing) value earlier on (assuming earlier means more left, if traversal starts on the left side) in the tree. This means that the left side of the tree must be $\geq y$. This is because it will then not be necessary to look for smaller values than y inside A, since the root node will then have a better (higher) value on its left side.

X needs to be $\geq y$ when $y = [1, 3]$.