

# OS Lab2

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tags : **OS**

## 1.

### main thread

#### step1

我先做parse，policy 跟 priority 我先儲存成一個string，等等會再分解

```
for(int i=1 ; i<argc ; i+=2){
    if(strcmp(argv[i],"-n")==0){
        thread_nums = atoi(argv[i+1]);
    }else if(strcmp(argv[i],"-t")==0){
        char *eptr;
        time_wait = strtod(argv[i+1],&eptr);
    }else if(strcmp(argv[i],"-s")==0){
        policy = argv[i+1];
    }else if(strcmp(argv[i],"-p")==0){
        priority = argv[i+1];
    }
}
```

宣告policys為二維陣列，儲存各個thread的policy為FIFO or Normal

宣告prioritys為一維陣列，儲存各個thread的priority

用strtok去parse剛剛儲存的string

```
int count = 0;
char policys[thread_nums][32];
int prioritys[thread_nums];

char *token = strtok(policy, ",");

while(token!=NULL){
    strcpy(policys[count++], token);
}
```

```

        token = strtok(NULL, ",");
    }

    count=0;
    token = strtok(priority, ",");
    while(token!=NULL){
        prioritys[count++] = atoi(token);
        token = strtok(NULL, ",");
    }

```

## step2

宣告各個thread以及需要的info和attribute

```

pthread_t t[thread_nums];
thread_info_t info[thread_nums];
pthread_attr_t attr[thread_nums]; //attr
cpu_set_t cpuset; //affinity
struct sched_param param[thread_nums]; //priority

```

## step3

set cpu affinity, 宣告使用第0個cpu

```

CPU_ZERO(&cpuset);
CPU_SET(0, &cpuset);

```

## step4

init barrier

```

pthread_barrier_init(&barrier, NULL, thread_nums);

```

set info to each thread, 然後判斷 sched\_policy 是 SCHED\_OTHER or SCHED\_FIFO

```

info[i].thread_id = (pthread_t)i;
info[i].thread_num = thread_nums;
info[i].sched_policy = (strcmp(policys[i], "NORMAL")==0)? SCHED_OTHER : SCHED_FIFO;
info[i].sched_priority = prioritys[i];

```

init attribute, and set affinity to thread

```
pthread_attr_init(&attr[i]);
if(pthread_attr_setaffinity_np(&attr[i], sizeof(cpu_set_t), &cpuset) != 0){
    printf("affinity error\\n");
}
```

set policy and priority, if policy is NORMAL, do not give priority to it

```
if(pthread_attr_setschedpolicy(&attr[i], info[i].sched_policy)){
    printf("policy error\\n");
}

if(info[i].sched_policy == SCHED_FIFO){
    param[i].sched_priority = priority[i];
    if(pthread_attr_setschedparam(&attr[i], &param[i]) != 0){
        printf("param error\\n");
    }
}
```

不要讓thread繼承main thread的性質，所以要下額外的指令，宣告完後create thread

```
if(pthread_attr_setinheritsched(&attr[i], PTHREAD_EXPLICIT_SCHED) != 0){
    printf("inherit error \\n");
    return -1;
}
pthread_create(&t[i], &attr[i], thread_func, (void *)&info[i]);
```

## step5

等待所有的thread做完，且destroy掉barrier

```
for(int i=0 ; i<thread_nums ; i++){
    pthread_join(t[i], NULL);
}
pthread_barrier_destroy(&barrier);
return 0;
```

## Thread function

## step1

設置barrier，等待所有的thread create完再執行

```
pthread_barrier_wait(&barrier);
```

## step2

thread do the task

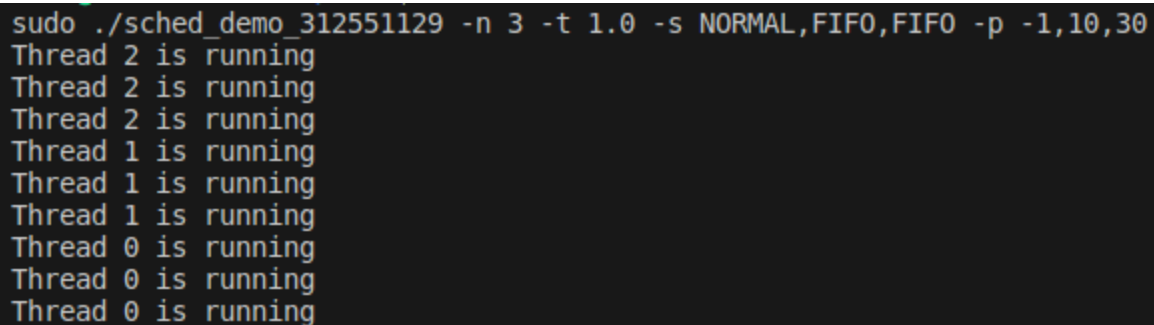
```
start = clock();
while(1){
    current = clock();
    double elapse = ((double)(current - start)) / CLOCKS_PER_SEC;
    if(elapse >= time_wait){
        break;
    }
}
```

## step3

跑完則exit

```
thread_exit(NULL);
```

## 2

A terminal window with a black background and white text. The first line is a command: 'sudo ./sched\_demo\_312551129 -n 3 -t 1.0 -s NORMAL,FIFO,FIFO -p -1,10,30'. The subsequent lines show the execution of three threads: 'Thread 2 is running' (three times), 'Thread 1 is running' (three times), and 'Thread 0 is running' (three times). This sequence demonstrates that threads with higher priority (Thread 2 with priority 30) execute first, followed by threads with lower priority (Thread 1 with priority 10, then Thread 0 with priority -1).

```
sudo ./sched_demo_312551129 -n 3 -t 1.0 -s NORMAL,FIFO,FIFO -p -1,10,30
Thread 2 is running
Thread 2 is running
Thread 2 is running
Thread 1 is running
Thread 1 is running
Thread 1 is running
Thread 0 is running
Thread 0 is running
Thread 0 is running
```

因為FIFO是real time，所以thread 1 2會先做，而thread 2的priority為30，比thread 1的priority 10還大，所以會比它更早做。

### 3

```
sudo ./sched_demo_312551129 -n 4 -t 0.5 -s NORMAL,FIFO,NORMAL,FIFO -p -1,10,-1,30
Thread 3 is running
Thread 3 is running
Thread 3 is running
Thread 1 is running
Thread 1 is running
Thread 1 is running
Thread 0 is running
Thread 2 is running
Thread 0 is running
Thread 2 is running
Thread 0 is running
Thread 2 is running
```

因為thread1,3都是FIFO所以執行順序會比較快，且thread3的priority比thread1的priority高，所以會先做thread3再來thread1，而thread0跟2都是NORMAL，所以會輪流做，因此會輪流輸出。

### 4

先用start紀錄現在的時間，用一個無限迴圈去跑，每次紀錄current time，如果start-current > time\_wait，則中止迴圈，這樣就不會讓它context switch。

```
start = clock();
while(1){
    current = clock();
    double elapse = ((double)(current - start)) / CLOCKS_PER_SEC;
    if(elapse >= time_wait){
        break;
    }
}
```