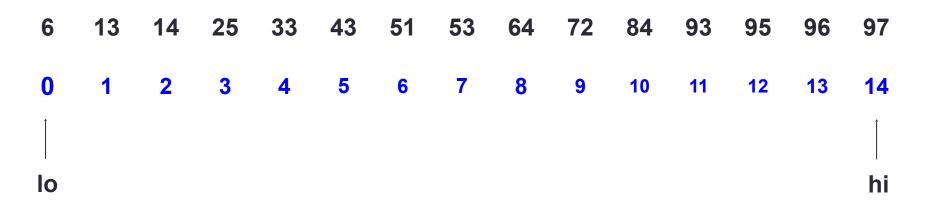
## BINARY SEARCH TREES

Problem Solving with Computers-II

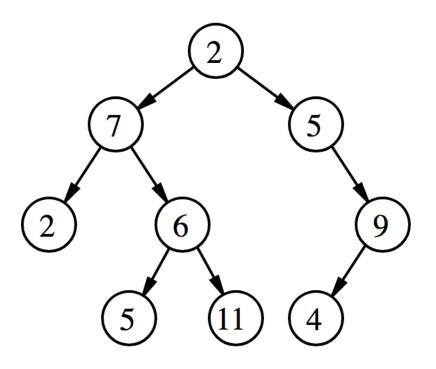


### Binary Search

- Binary search. Given value and sorted array a[], find index i such that a[i] = value, or report that no such index exists.
- Invariant. Algorithm maintains a [lo] ≤ value ≤ a [hi].
- Ex. Binary search for 33.



#### Trees



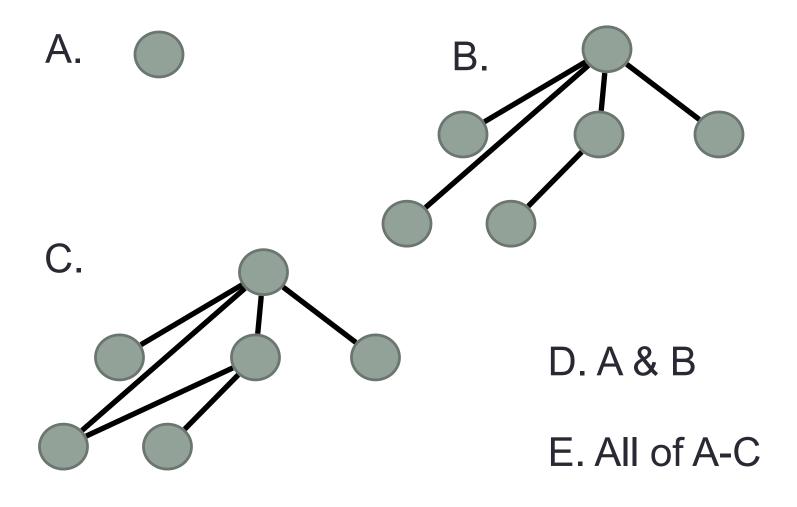
A tree has following general properties:

- One node is distinguished as a **root**;
- Every node (exclude a root) is connected by a directed edge *from* exactly one other node;

A direction is: *parent -> children* 

• Leaf node: Node that has no children

## Which of the following is/are a tree?



### Binary Search Trees

What are the operations supported?

What are the running times of these operations?

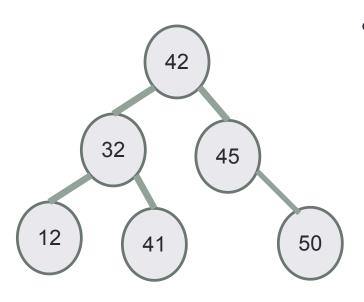
How do you implement the BST i.e. operations supported by it?

#### Operations supported by Sorted arrays and Binary Search Trees (BST)

Operations	
Min	
Max	
Successor	
Predecessor	
Search	
Insert	
Delete	
Print elements in order	

Example keys: 42, 32, 45, 12, 41, 50

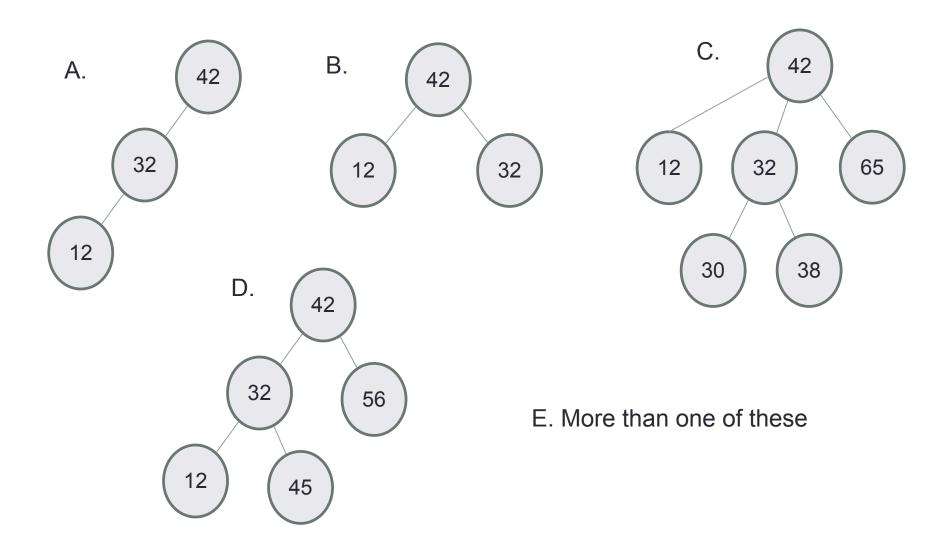
### Binary Search Tree – What is it?



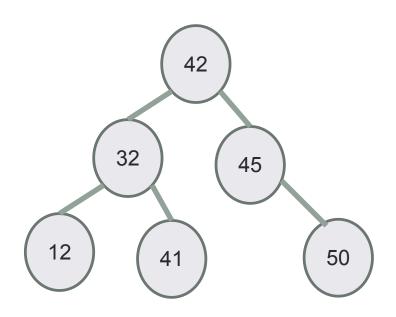
- Each node:
  - stores a key (k)
  - has a pointer to left child, right child and parent (optional)
  - Satisfies the Search Tree Property

For any node, Keys in node's left subtree < Node's key Node's key < Keys in node's right subtree

### Which of the following is/are a binary search tree?



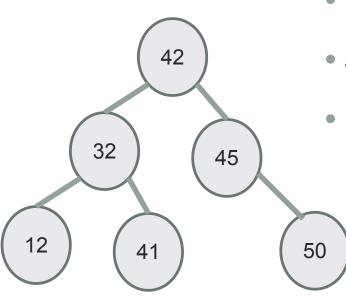
#### BSTs allow efficient search!



- Start at the root;
- Trace down a path by comparing **k** with the key of the current node x:
  - If the keys are equal: we have found the key
  - If k < key[x] search in the left subtree of x
  - If k > key[x] search in the right subtree of x

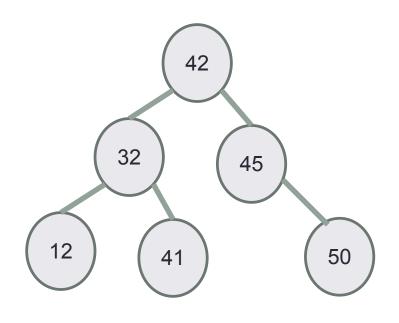


#### Insert



- Insert 40
- Search for the key
- Insert at the spot you expected to find it

#### Min/Max



# Which of the following described the algorithm to find the maximum value in the BST?

- A. Follow right child pointers from the root, until a node with no right child is encountered, return that node's key
- B. Follow left child pointers from the root, until a node with no left child is encountered, return that node's key
- C. Traverse to the last level in the tree and traverse the tree left to right, return the key of the last node in the last level.

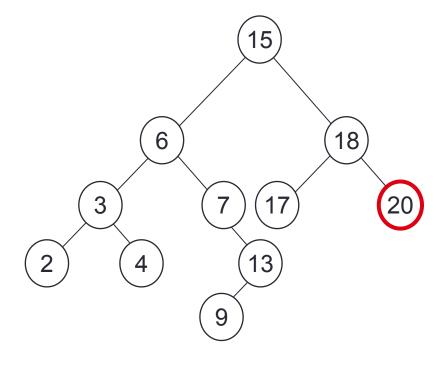
### Define the BST ADT

<b>Operations</b>	
Search	42
Insert	
Min	
Max	$\begin{pmatrix} 32 \end{pmatrix} \begin{pmatrix} 45 \end{pmatrix}$
Successor	
Predecessor	
Delete	$ \left(\begin{array}{c} 12 \right) \left(\begin{array}{c} 41 \right) $
Print elements in order	

```
class BSTNode {
public:
  BSTNode* left;
  BSTNode* right;
  BSTNode* parent;
  int const data;
  BSTNode(int d) : data(d) {
    left = right = parent = nullptr;
```

#### Max: find the maximum key value in a BST

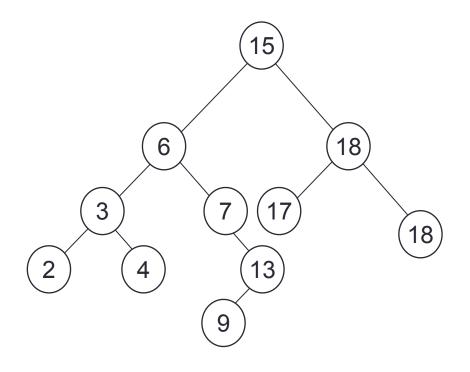
Alg: int BST::max()



Maximum = 20

#### Min: find the minimum key value in a BST

```
Alg: int BST::min() {
Start at the root.
Follow
        child
pointers from the root, until
a node with no left child is
encountered.
Return the key of that node
```

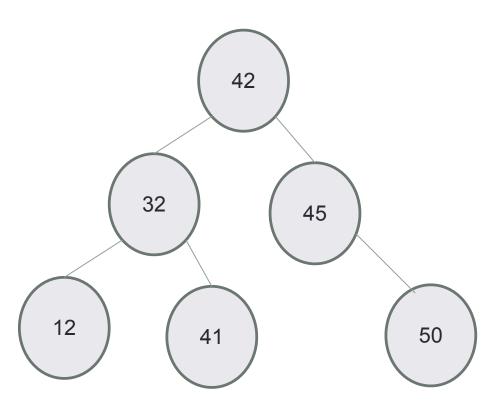


Min = ?

### Traversing down the tree

• Suppose n is a pointer to the root. What is the output of the following code:

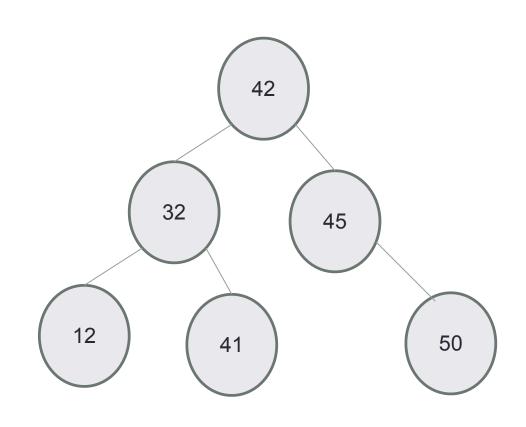
```
n = n->left;
n = n->right;
cout<<n->data<<endl;
 A. 42
 B. 32
 C. 12
 D. 41
 E. Segfault
```



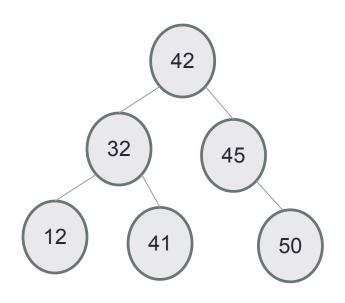
### Traversing up the tree

- Suppose n is a pointer to the node with value 50.
- What is the output of the following code:

```
n = n-parent;
  = n->parent;
n = n->left;
cout<<n->data<<endl;
 A. 42
 B. 32
 C. 12
 D. 45
 E. Segfault
```



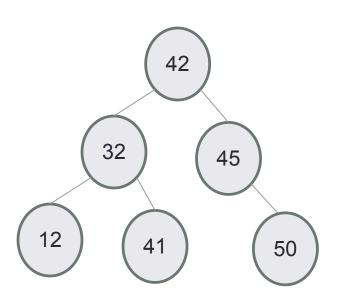
### In order traversal: print elements in sorted order



Algorithm Inorder(tree)

- 1. Traverse the left subtree, i.e., call Inorder(left-subtree)
- 2. Visit the root.
- 3. Traverse the right subtree, i.e., call Inorder(right-subtree)

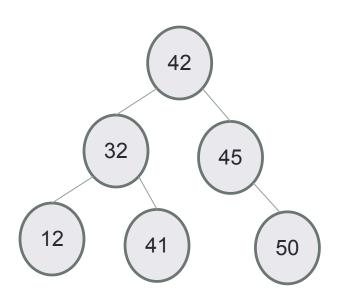
### Pre-order traversal: nice way to linearize your tree!



Algorithm Preorder(tree)

- 1. Visit the root.
- 2. Traverse the left subtree, i.e., call Preorder(left-subtree)
- 3. Traverse the right subtree, i.e., call Preorder(right-subtree)

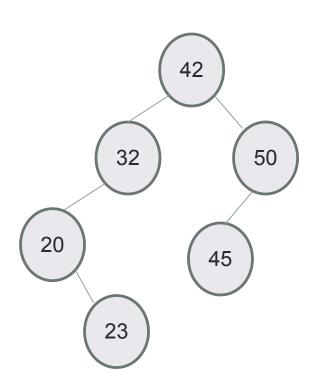
#### Post-order traversal: use to recursively clear the tree!



Algorithm Postorder(tree)

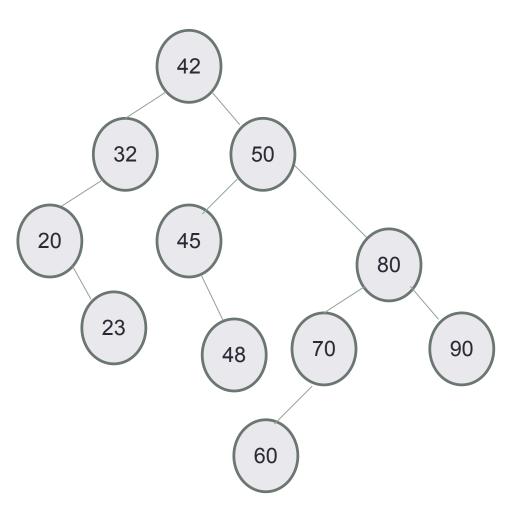
- 1. Traverse the left subtree, i.e., call Postorder(left-subtree)
- 2. Traverse the right subtree, i.e., call Postorder(right-subtree)
- 3. Visit the root.

#### Predecessor: Next smallest element



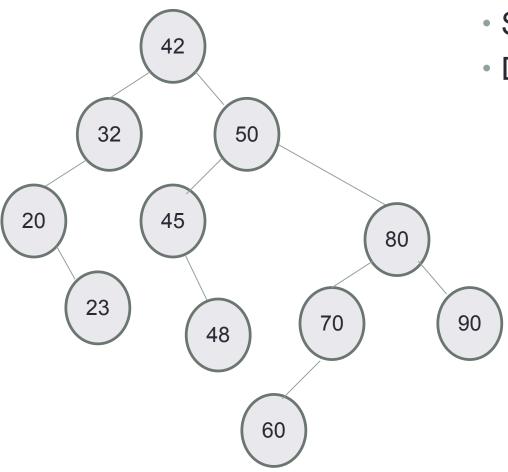
- What is the predecessor of 32?
- What is the predecessor of 45?

# Successor: Next largest element



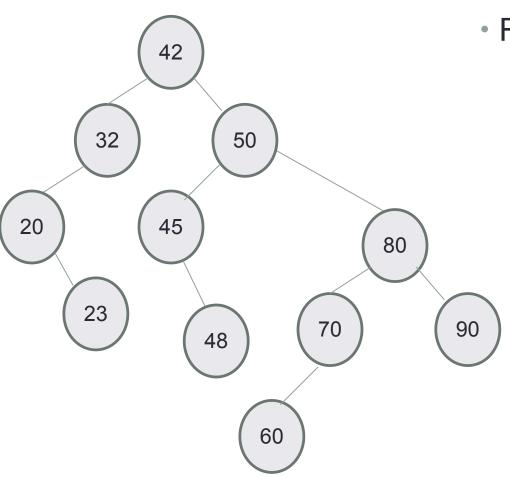
- What is the successor of 45?
- What is the successor of 50?
- What is the successor of 60?

#### Delete: Case 1 - Node is a leaf node



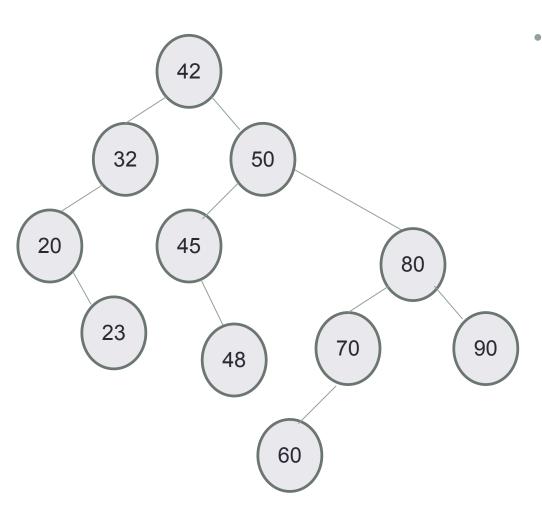
- Set parent's (left/right) child pointer to null
- Delete the node

# Delete: Case 2 - Node has only one child



Replace the node by its only child

#### Delete: Case 3 - Node has two children



 Can we still replace the node by one of its children? Why or Why not?