

₁ ALGEBRAIC NUMBER THEORY,
₂ A COMPUTATIONAL APPROACH

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Chapter 1

Basic Commutative Algebra

The commutative algebra in this chapter provides a foundation for understanding the more refined number-theoretic structures associated to number fields.

First we prove the structure theorem for finitely generated abelian groups. Then we establish the standard properties of Noetherian rings and modules, including a proof of the Hilbert basis theorem. We also observe that finitely generated abelian groups are Noetherian \mathbb{Z} -modules. After establishing properties of Noetherian rings, we consider rings of algebraic integers and discuss some of their properties.

1.1 Finitely Generated Abelian Groups

Finitely generated abelian groups arise all over algebraic number theory. For example, they will appear in this book as class groups, unit groups, and the underlying additive groups of rings of integers, and as Mordell-Weil groups of elliptic curves.

In this section, we prove the structure theorem for finitely generated abelian groups, since it will be crucial for much of what we will do later.

Let $\mathbb{Z} = \{0, \pm 1, \pm 2, \dots\}$ denote the ring of (rational) integers, and for each positive integer n , let $\mathbb{Z}/n\mathbb{Z}$ denote the ring of integers modulo n , which is a cyclic abelian group of order n under addition.

Definition 1.1.1 (Finitely Generated). A group G is *finitely generated* if there exists $g_1, \dots, g_n \in G$ such that every element of G can be expressed as a finite product (or sum, if we write G additively) of positive or negative powers of the g_i .

31 For example, the group \mathbb{Z} is finitely generated, since it is generated by 1.

Theorem 1.1.2 (Structure Theorem for Finitely Generated Abelian Groups).
 Let G be a finitely generated abelian group. Then there is an isomorphism

$$G \approx (\mathbb{Z}/n_1\mathbb{Z}) \oplus (\mathbb{Z}/n_2\mathbb{Z}) \oplus \cdots \oplus (\mathbb{Z}/n_s\mathbb{Z}) \oplus \mathbb{Z}^r,$$

32 where $r, s \geq 0$, $n_i > 1$ for all i , and $n_1 \mid n_2 \mid \cdots \mid n_s$. Furthermore, the n_i
 33 and r are uniquely determined by G .

34 **Exercise 1.1.3.** Quick! Guess how many abelian groups there are of order
 35 less than 12. Use Theorem 1.1.2 to classify all abelian groups of order less
 36 than 12. How many do you think there are? How many are there?

37 We will prove the theorem as follows. We first remark that any subgroup
 38 of a finitely generated free abelian group is finitely generated. Then we see
 39 how to represent finitely generated abelian groups as quotients of finite rank
 40 free abelian groups, and how to reinterpret such a presentation in terms of
 41 matrices over the integers. Next we describe how to use row and column
 42 operations over the integers to show that every matrix over the integers is
 43 equivalent to one in a canonical diagonal form, called the Smith normal form.
 44 We obtain a proof of the theorem by reinterpreting the in terms of groups.
 45 Finally, we observe that the representation in the theorem is necessarily
 46 unique.

47 **Proposition 1.1.4.** If H is a subgroup of a finitely generated abelian group
 48 G , then H is finitely generated.

49 The key reason that this is true is that G is a finitely generated module
 50 over the principal ideal domain \mathbb{Z} . We defer the proof of Proposition 1.1.4 to
 51 Section 1.2, where we will give a complete proof of a beautiful generalization
 52 in the context of Noetherian rings (the Hilbert basis theorem).

53 **Corollary 1.1.5.** Suppose G is a finitely generated abelian group. Then
 54 there are finitely generated free abelian groups F_1 and F_2 and there is a
 55 homomorphism $\psi : F_2 \rightarrow F_1$ such that $G \approx F_1/\psi(F_2)$.

56 *Proof.* Let x_1, \dots, x_m be generators for G . Let $F_1 = \mathbb{Z}^m$ and let $\varphi : F_1 \rightarrow G$
 57 be the homomorphism that sends the i th generator $(0, 0, \dots, 1, \dots, 0)$ of \mathbb{Z}^m
 58 to x_i . Then φ is surjective, and by Proposition 1.1.4 the kernel $\ker(\varphi)$ of
 59 φ is a finitely generated abelian group. Suppose there are n generators for
 60 $\ker(\varphi)$, let $F_2 = \mathbb{Z}^n$ and fix a surjective homomorphism $\psi : F_2 \rightarrow \ker(\varphi)$.
 61 Then $F_1/\psi(F_2)$ is isomorphic to G . \square

An *sequence* of homomorphisms of abelian groups

$$H \xrightarrow{f} G \xrightarrow{g} K$$

is exact if $\text{im}(f) = \ker(g)$. For longer sequences, exactness means every three consecutive terms with two arrows are exact. Given a finitely generated abelian group G , Corollary 1.1.5 provides an exact sequence

$$F_2 \xrightarrow{\psi} F_1 \rightarrow G \rightarrow 0.$$

Suppose G is a nonzero finitely generated abelian group. By the corollary, there are free abelian groups F_1 and F_2 and there is a homomorphism $\psi : F_2 \rightarrow F_1$ such that $G \approx F_1/\psi(F_2)$. Upon choosing a basis for F_1 and F_2 , we obtain isomorphisms $F_1 \approx \mathbb{Z}^n$ and $F_2 \approx \mathbb{Z}^m$ for integers n and m . Just as in linear algebra, we view $\psi : F_2 \rightarrow F_1$ as being given by left multiplication by the $n \times m$ matrix A whose columns are the images of the generators of F_2 in \mathbb{Z}^n . We visualize this as follows:

$$\mathbb{Z}^m \xrightarrow{A} \mathbb{Z}^n \rightarrow G \rightarrow 0$$

62 The *cokernel* of the homomorphism defined by A is the quotient of \mathbb{Z}^n
 63 by the image of A (i.e., the \mathbb{Z} -span of the columns of A), and this cokernel
 64 is isomorphic to G .

65 The following proposition implies that we may choose a bases for F_1 and
 66 F_2 such that the matrix of A only has nonzero entries along the diagonal,
 67 so that the structure of the cokernel of A is trivial to understand.

68 **Proposition 1.1.6** (Smith normal form). *Suppose A is an $n \times m$ integer*
 69 *matrix. Then there exist invertible integer matrices P and Q such that*
 70 *$A' = PAQ$ only has nonzero entries along the diagonal, and these entries*
 71 *are $n_1, n_2, \dots, n_s, 0, \dots, 0$, where $s \geq 0$, $n_i \geq 1$ for all i , and $n_1 \mid n_2 \mid \dots \mid n_s$.*

72 *Example 1.1.7.* An example of a matrix in Smith normal form is

$$A = \begin{pmatrix} 2 & 0 & 0 & 0 \\ 0 & 6 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}.$$

73 *Remark 1.1.8.* Note that the matrices P and Q are invertible as integer
 74 matrices, so $\det(P)$ and $\det(Q)$ are ± 1 . In particular $\det A' = \pm \det A$. We
 75 will see in the proof of Theorem 1.1.2 that A' is uniquely determined by A .

76 **Definition 1.1.9.** The matrix A' in Proposition 1.1.6 is called the *Smith*
 77 *normal form* of A .

78 *Proof of Proposition 1.1.6.* The matrix P will be a product of matrices that
 79 define elementary row operations and Q will be a product corresponding to
 80 elementary column operations. The elementary row and column operations
 81 over \mathbb{Z} are as follows:

82 **Add multiple:** Add an integer multiple of one row to another (or a multi-
 83 ple of one column to another).

84 **Swap:** Interchange two rows or two columns.

85 **Rescale:** Multiply a row by -1 .

86 Each of these operations is given by left or right multiplying by an invertible
 87 matrix E with integer entries, where E is the result of applying the given
 88 operation to the identity matrix, and E is invertible because each operation
 89 can be reversed using another row or column operation over the integers.

90 To see that the proposition must be true, assume $A \neq 0$ and perform
 91 the following steps (compare [Art91, pg. 459]):

- 92 1. By permuting rows and columns, move a nonzero entry of A with
 93 smallest absolute value to the upper left corner of A . Now “attempt”
 94 (as explained in detail below) to make all other entries in the first row
 95 and column 0 by adding multiples of the top row or first column to
 96 other rows or columns, as follows:

97 Suppose a_{i1} is a nonzero entry in the first column, with
 98 $i > 1$. Using the division algorithm, write $a_{i1} = a_{11}q + r$,
 99 with $0 \leq r < a_{11}$. Now add $-q$ times the first row to the
 100 i th row. If $r > 0$, then go to step 1 (so that an entry with
 101 absolute value at most r is the upper left corner).

102 If at any point this operation produces a nonzero entry in the matrix
 103 with absolute value smaller than $|a_{11}|$, start the process over by per-
 104 muting rows and columns to move that entry to the upper left corner
 105 of A . Since the integers $|a_{11}|$ are a decreasing sequence of positive
 106 integers, we will not have to move an entry to the upper left corner
 107 infinitely often, so when this step is done the upper left entry of the
 108 matrix is nonzero, and all entries in the first row and column are 0.

109 2. We may now assume that a_{11} is the only nonzero entry in the first
 110 row and column. If some entry a_{ij} of A is not divisible by a_{11} , add
 111 the column of A containing a_{ij} to the first column, thus producing an
 112 entry in the first column that is nonzero. When we perform step 2,
 113 the remainder r will be greater than 0. Permuting rows and columns
 114 results in a smaller $|a_{11}|$. Since $|a_{11}|$ can only shrink finitely many
 115 times, eventually we will get to a point where every a_{ij} is divisible by
 116 a_{11} . If a_{11} is negative, multiple the first row by -1 .

117 After performing the above operations, the first row and column of A are
 118 zero except for a_{11} which is positive and divides all other entries of A . We
 119 repeat the above steps for the matrix B obtained from A by deleting the first
 120 row and column. The upper left entry of the resulting matrix will be divisible
 121 by a_{11} , since every entry of B is. Repeating the argument inductively proves
 122 the proposition. \square

123 *Example 1.1.10.* The matrix $\begin{pmatrix} -2 & 2 \\ -3 & 4 \end{pmatrix}$ has Smith normal form $\begin{pmatrix} 1 & 0 \\ 0 & 2 \end{pmatrix}$, and
 124 the matrix $\begin{pmatrix} 1 & 4 & 9 \\ 16 & 25 & 36 \\ 49 & 64 & 81 \end{pmatrix}$ has Smith normal form $\begin{pmatrix} 1 & 0 & 0 \\ 0 & 3 & 0 \\ 0 & 0 & 72 \end{pmatrix}$. As a double
 125 check, note that the determinants of a matrix and its Smith normal form
 126 match, up to sign. This is because

$$\det(PAQ) = \det(P) \det(A) \det(Q) = \pm \det(A).$$

127 We compute each of the above Smith forms using **Sage**, along with the
 128 corresponding transformation matrices. To do this we use the **Sage** com-
 129 mand **matrix**, which takes as input the base ring, the number of rows, and
 130 the entries. The output of **matrix** is a matrix object which has the method
 131 **smith_form**.

132 First the 2×2 matrix.

```
A = matrix(ZZ, 2, [-2,2, -3,4])
S, P, Q = A.smith_form(); S
```

```
[1 0]
[0 2]
```

```
P*A*Q
```

```
[1 0]
[0 2]
```

```
P
```

```
[0 1]
[1 0]
```

```
Q
```

```
[1 -4]
[1 -3]
```

133

134 Next the 3×3 matrix.

```
A = matrix(ZZ, 3, [1,4,9, 16,25,36, 49,64,81])
S, P, Q = A.smith_form(); S
```

```
[ 1  0  0]
[ 0  3  0]
[ 0  0 72]
```

```
P*A*Q
```

```
[ 1  0  0]
[ 0  3  0]
[ 0  0 72]
```

```
P
```

```
[ 0  0  1]
[ 0  1 -1]
[ 1 -20 -17]
```

```
Q
```

```
[ 47  74  93]
[-79 -125 -156]
[ 34  54  67]
```

135

136 For one more example, we compute the Smith form of a 3×3 matrix of
 137 rank 2:

```
m = matrix(ZZ, 3, [2..10]); m
```

```
[ 2  3  4]
[ 5  6  7]
[ 8  9 10]
```

```
m.smith_form()[0]
```

```
[1 0 0]
[0 3 0]
[0 0 0]
```

139 *Proof of Theorem 1.1.2.* Suppose G is a finitely generated abelian group,
 140 which we may assume is nonzero. As in the paragraph before Proposition
 141 1.1.6, we use Corollary 1.1.5 to write G as the cokernel of an $n \times$
 142 m integer matrix A . By Proposition 1.1.6 there are isomorphisms $Q :$
 143 $\mathbb{Z}^m \rightarrow \mathbb{Z}^m$ and $P : \mathbb{Z}^n \rightarrow \mathbb{Z}^n$ such that $A' = PAQ$ has diagonal entries
 144 $n_1, n_2, \dots, n_s, 0, \dots, 0$, where $n_1 > 1$ and $n_1 \mid n_2 \mid \dots \mid n_s$. Then G is
 145 isomorphic to the cokernel of the diagonal matrix A' , so

$$G \cong (\mathbb{Z}/n_1\mathbb{Z}) \oplus (\mathbb{Z}/n_2\mathbb{Z}) \oplus \dots \oplus (\mathbb{Z}/n_s\mathbb{Z}) \oplus \mathbb{Z}^r, \quad (1.1)$$

146 as claimed. The n_i are determined by G , because n_i is the smallest positive
 147 integer n such that nG requires at most $s + r - i$ generators. We see from
 148 the representation (1.1) of G as a product that n_i has this property and that
 149 no smaller positive integer does. \square

150 **Exercise 1.1.11.** Recall Smith normal form defined in Proposition 1.1.6.
 151 With only minor modifications, then the proposition and proof will work
 152 over any principle ideal domain. Find and apply these modifications then

153 find the Smith normal form of the matrix $\begin{pmatrix} 1 & 2 & 3 \\ 0 & 1+i & 2 \\ 0 & 1 & 5 \end{pmatrix}$.

154 [*Hint:* You can use **Sage** to verify your answer. However, you will need
 155 to make explicitly construct the Gaussian integers in order to input the
 156 matrix. You can do this by the following code.]

```
K.<i> = QuadraticField(-1)
R = K.maximal_order()
M = matrix(R, 3, [1,2,3,0,1+i,2,0,1,5]); show(M)
#show(M.smith_form()[0]) #uncomment for the answer
```

158 **Exercise 1.1.12.** Let $A = \begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{pmatrix}$.

- 159 (a) Find the Smith normal form of A .
- 160 (b) Prove that the cokernel of the map $\mathbb{Z}^3 \rightarrow \mathbb{Z}^3$ given by multiplication
161 by A is isomorphic to $\mathbb{Z}/3\mathbb{Z} \oplus \mathbb{Z}$.

162 1.2 Noetherian Rings and Modules

163 A module M over a commutative ring R with unit element is much like
164 a vector space, but with more subtle structure. In this book, most of the
165 modules we encounter will be noetherian, which is a generalization of the
166 “finite dimensional” property of vector spaces. This section is about prop-
167 erties of noetherian modules (and rings), which are crucial to much of this
168 book. We thus give complete proofs of these properties, so you will have a
169 solid foundation on which to learn algebraic number theory.

170 We first define noetherian rings and modules, then introduce several
171 equivalent characterizations of them. We prove that when the base ring is
172 noetherian, a module is finitely generated if and only if it is noetherian.
173 Next we define short exact sequences, and prove that the middle module
174 in a sequence is noetherian if and only if the first and last modules are
175 noetherian. Finally, we prove the Hilbert basis theorem, which asserts that
176 adjoining finitely many elements to a noetherian ring results in a noetherian
177 ring.

178 Let R be a commutative ring with unity. An R -module is an additive
179 abelian group M equipped with a map $R \times M \rightarrow M$ such that for all $r, r' \in R$
180 and all $m, m' \in M$ we have $(rr')m = r(r'm)$, $(r + r')m = rm + r'm$,
181 $r(m + m') = rm + rm'$, and $1m = m$. A *submodule* of M is a subgroup of
182 M that is preserved by the action of R . For example, R is a module over
183 itself, and any ideal I in R is an R -submodule of R .

184 *Example 1.2.1.* Abelian groups are the same as \mathbb{Z} -modules, and vector spaces
185 over a field K are the same as K -modules.

186 An R -module M is finitely generated if there are elements $m_1, \dots, m_n \in$
187 M such that every element of M is an R -linear combination of the m_i . The
188 noetherian property is stronger than just being finitely generated:

189 **Definition 1.2.2** (Noetherian). An R -module M is *noetherian* if every sub-
190 module of M is finitely generated. A ring R is *noetherian* if R is noetherian
191 as a module over itself, i.e., if every ideal of R is finitely generated.

Any submodule M' of a noetherian module M is also noetherian. Indeed, if every submodule of M is finitely generated then so is every submodule of M' , since submodules of M' are also submodules of M .

Example 1.2.3. Let $R = M = \mathbb{Q}[x_1, x_2, \dots]$ be a polynomial ring over \mathbb{Q} in infinitely many indeterminants x_i . Then M is finitely generated as an R -module (!), since it is generated by 1. Consider the submodule $I = (x_1, x_2, \dots)$ of polynomials with 0 constant term, and suppose it is generated by polynomials f_1, \dots, f_n . Let x_i be an indeterminant that does not appear in any f_j , and suppose there are $h_k \in R$ such that $\sum_{k=1}^n h_k f_k = x_i$. Setting $x_i = 1$ and all other $x_j = 0$ on both sides of this equation and using that the f_k all vanish (they have 0 constant term), yields $0 = 1$, a contradiction. We conclude that the ideal I is not finitely generated, hence M is not a noetherian R -module, despite being finitely generated.

Definition 1.2.4 (Ascending chain condition). An R -module M satisfies the *ascending chain condition* if every sequence $M_1 \subset M_2 \subset M_3 \subset \dots$ of submodules of M eventually stabilizes, i.e., there is some n such that $M_n = M_{n+1} = M_{n+2} = \dots$.

We will use the notion of maximal element below. If \mathcal{X} is a set of subsets of a set S , ordered by inclusion, then a *maximal element* $A \in \mathcal{X}$ is a set such that no superset of A is contained in \mathcal{X} . Note that \mathcal{X} may contain many different maximal elements.

Proposition 1.2.5. *If M is an R -module, then the following are equivalent:*

1. M is noetherian,
2. M satisfies the ascending chain condition, and
3. Every nonempty set of submodules of M contains at least one maximal element.

Proof.

(1 \implies 2): Suppose $M_1 \subset M_2 \subset \dots$ is a sequence of submodules of M . Then $M_\infty = \cup_{n=1}^\infty M_n$ is a submodule of M . Since M is noetherian and M_∞ is a submodule of M , there is a finite set a_1, \dots, a_m of generators for M_∞ . Each a_i must be contained in some M_j , so there is an n such that $a_1, \dots, a_m \in M_n$. But then $M_k = M_n$ for all $k \geq n$, which proves that the chain of M_i stabilizes, so the ascending chain condition holds for M .

(2 \implies 3) : Suppose 3 were false, so there exists a nonempty set S of submodules of M that does not contain a maximal element. We will use S to construct an infinite ascending chain of submodules of M that does not stabilize. Note that S is infinite, otherwise it would contain a maximal element. Let M_1 be any element of S . Then there is an M_2 in S that strictly contains M_1 , otherwise S would contain the maximal element M_1 . Continuing inductively in this way we find an M_3 in S that properly contains M_2 , etc., and we produce an infinite ascending chain of submodules of M , which contradicts the ascending chain condition.

(3 \implies 1) : Suppose 1 is false, so there is a submodule M' of M that is not finitely generated. We will show that the set S of all finitely generated submodules of M' does not have a maximal element, which will be a contradiction. Suppose S does have a maximal element L . Since L is finitely generated and $L \subset M'$, and M' is not finitely generated, there is an $a \in M'$ such that $a \notin L$. Then $L' = L + Ra$ is an element of S that strictly contains the presumed maximal element L , a contradiction.

□

Definition 1.2.6 (Module Homomorphism). A *homomorphism* of R -modules $\varphi : M \rightarrow N$ is an abelian group homomorphism such that for any $r \in R$ and $m \in M$ we have $\varphi(rm) = r\varphi(m)$. A sequence

$$L \xrightarrow{f} M \xrightarrow{g} N,$$

where f and g are homomorphisms of R -modules, is *exact* if $\text{im}(f) = \ker(g)$. A *short exact sequence* of R -modules is a sequence

$$0 \rightarrow L \xrightarrow{f} M \xrightarrow{g} N \rightarrow 0$$

that is exact at each point, i.e., f is injective, g is surjective, and $\text{im}(f) = \ker(g)$.

Example 1.2.7. The sequence

$$0 \rightarrow \mathbb{Z} \xrightarrow{2} \mathbb{Z} \rightarrow \mathbb{Z}/2\mathbb{Z} \rightarrow 0$$

is an exact sequence, where the first map sends 1 to 2, and the second is the natural quotient map.

Lemma 1.2.8. *If*

$$0 \rightarrow L \xrightarrow{f} M \xrightarrow{g} N \rightarrow 0$$

248 *is a short exact sequence of R -modules, then M is noetherian if and only if*
 249 *both L and N are noetherian.*

250 *Proof.* First suppose that M is noetherian. Then L is a submodule of M ,
 251 so L is noetherian. Let N' be a submodule of N ; then the inverse image of
 252 N' in M is a submodule of M , so it is finitely generated, hence its image N'
 253 is also finitely generated. Thus N is noetherian as well.

254 Next assume nothing about M , but suppose that both L and N are
 255 noetherian. Suppose M' is a submodule of M ; then $M_0 = f(L) \cap M'$ is
 256 isomorphic to a submodule of the noetherian module L , so M_0 is generated
 257 by finitely many elements a_1, \dots, a_n . The quotient M'/M_0 is isomorphic
 258 (via g) to a submodule of the noetherian module N , so M'/M_0 is generated
 259 by finitely many elements b_1, \dots, b_m . For each $i \leq m$, let c_i be a lift of b_i to
 260 M' , modulo M_0 . Then the elements $a_1, \dots, a_n, c_1, \dots, c_m$ generate M' , for
 261 if $x \in M'$, then there is some element $y \in M_0$ such that $x - y$ is an R -linear
 262 combination of the c_i , and y is an R -linear combination of the a_i . \square

263 **Proposition 1.2.9.** *Suppose R is a noetherian ring. Then an R -module M*
 264 *is noetherian if and only if it is finitely generated.*

265 *Proof.* If M is noetherian then every submodule of M is finitely generated
 266 so M itself is finitely generated. Conversely, suppose M is finitely generated,
 267 say by elements a_1, \dots, a_n . Then there is a surjective homomorphism from
 268 $R^n = R \oplus \dots \oplus R$ to M that sends $(0, \dots, 0, 1, 0, \dots, 0)$ (1 in the i th factor)
 269 to a_i . Using Lemma 1.2.8 and exact sequences of R -modules such as $0 \rightarrow$
 270 $R \rightarrow R \oplus R \rightarrow R \rightarrow 0$, we see inductively that R^n is noetherian. Again by
 271 Lemma 1.2.8, homomorphic images of noetherian modules are noetherian,
 272 so M is noetherian. \square

273 **Lemma 1.2.10.** *Suppose $\varphi : R \rightarrow S$ is a surjective homomorphism of rings*
 274 *and R is noetherian. Then S is noetherian.*

Proof. The kernel of φ is an ideal I in R , and we have an exact sequence

$$0 \rightarrow I \rightarrow R \rightarrow S \rightarrow 0$$

275 with R noetherian. This is an exact sequence of R -modules, where S has
 276 the R -module structure induced from φ (if $r \in R$ and $s \in S$, then we define
 277 $rs = \varphi(r)s$). By Lemma 1.2.8, it follows that S is a noetherian R -modules.

Suppose J is an ideal of S . Since J is an R -submodule of S , if we view J as an R -module, then J is finitely generated. Since R acts on J through S , the R -generators of J are also S -generators of J , so J is finitely generated as an ideal. Thus S is noetherian. \square

Theorem 1.2.11 (Hilbert Basis Theorem). *If R is a noetherian ring and S is finitely generated as a ring over R , then S is noetherian. In particular, for any n the polynomial ring $R[x_1, \dots, x_n]$ and any of its quotients are noetherian.*

Proof. Assume first that we have already shown that for any n the polynomial ring $R[x_1, \dots, x_n]$ is noetherian. Suppose S is finitely generated as a ring over R , so there are generators s_1, \dots, s_n for S . Then the map $x_i \mapsto s_i$ extends uniquely to a surjective homomorphism $\pi : R[x_1, \dots, x_n] \twoheadrightarrow S$, and Lemma 1.2.10 implies that S is noetherian.

The rings $R[x_1, \dots, x_n]$ and $(R[x_1, \dots, x_{n-1}])[x_n]$ are isomorphic, so it suffices to prove that if R is noetherian then $R[x]$ is also noetherian. (Our proof follows [Art91, §12.5].) Thus suppose I is an ideal of $R[x]$ and that R is noetherian. We will show that I is finitely generated.

Let A be the set of leading coefficients of polynomials in I . (The leading coefficient of a polynomial is the coefficient of the highest degree monomial, or 0 if the polynomial is 0; thus $3x^7 + 5x^2 - 4$ has leading coefficient 3.) We will first show that A is an ideal of R . Suppose $a, b \in A$ are nonzero with $a + b \neq 0$. Then there are polynomials f and g in I with leading coefficients a and b . If $\deg(f) \leq \deg(g)$, then $a + b$ is the leading coefficient of $x^{\deg(g)-\deg(f)}f + g$, so $a + b \in A$; the argument when $\deg(f) > \deg(g)$ is analogous. Suppose $r \in R$ and $a \in A$ with $ra \neq 0$. Then ra is the leading coefficient of rf , so $ra \in A$. Thus A is an ideal in R .

Since R is noetherian and A is an ideal of R , there exist nonzero $a_1, \dots, a_n \in A$ that generate A as an ideal. Since A is the set of leading coefficients of elements of I , and the a_j are in A , we can choose for each $j \leq n$ an element $f_j \in I$ with leading coefficient a_j . By multiplying the f_j by some power of x , we may assume that the f_j all have the same degree $d \geq 1$.

Let $S_{<d}$ be the set of elements of I that have degree strictly less than d . This set is closed under addition and under multiplication by elements of R , so $S_{<d}$ is a module over R . The module $S_{<d}$ is the submodule of the R -module of polynomials of degree less than n , which is noetherian by Proposition 1.2.9 because it is generated by $1, x, \dots, x^{n-1}$. Thus $S_{<d}$ is finitely generated, and we may choose generators h_1, \dots, h_m for $S_{<d}$.

We finish by proving using induction on the degree that every $g \in I$ is an $R[x]$ -linear combination of $f_1, \dots, f_n, h_1, \dots, h_m$. If $g \in I$ has degree 0, then

317 $g \in S_{<d}$, since $d \geq 1$, so g is a linear combination of h_1, \dots, h_m . Next suppose
 318 $g \in I$ has degree e , and that we have proven the statement for all elements
 319 of I of degree $< e$. If $e \leq d$, then $g \in S_{<d}$, so g is in the $R[x]$ -ideal generated
 320 by h_1, \dots, h_m . Next suppose that $e \geq d$. Then the leading coefficient b
 321 of g lies in the ideal A of leading coefficients of elements of I , so there exist
 322 $r_i \in R$ such that $b = r_1 a_1 + \dots + r_n a_n$. Since f_i has leading coefficient a_i , the
 323 difference $g - x^{e-d} r_i f_i$ has degree less than the degree e of g . By induction
 324 $g - x^{e-d} r_i f_i$ is an $R[x]$ linear combination of $f_1, \dots, f_n, h_1, \dots, h_m$, so g is
 325 also an $R[x]$ linear combination of $f_1, \dots, f_n, h_1, \dots, h_m$. Since each f_i and
 326 h_j lies in I , it follows that I is generated by $f_1, \dots, f_n, h_1, \dots, h_m$, so I is
 327 finitely generated, as required. \square

328 1.2.1 The Ring \mathbb{Z} is Noetherian

329 The ring \mathbb{Z} is noetherian since every ideal of \mathbb{Z} is generated by one element.

330 **Proposition 1.2.12.** *Every ideal of the ring \mathbb{Z} is principal.*

331 *Proof.* Suppose I is a nonzero ideal in \mathbb{Z} . Let d be the least positive element
 332 of I . Suppose that $a \in I$ is any nonzero element of I . Using the division
 333 algorithm, we write $a = dq + r$, where q is an integer and $0 \leq r < d$. We have
 334 $r = a - dq \in I$ and $r < d$, so our assumption that d is minimal implies that
 335 $r = 0$, hence $a = dq$ is in the ideal generated by d . Thus I is the principal
 336 ideal generated by d . \square

337 *Example 1.2.13.* Let $I = (12, 18)$ be the ideal of \mathbb{Z} generated by 12 and 18.
 338 If $n = 12a + 18b \in I$, with $a, b \in \mathbb{Z}$, then $6 \mid n$, since $6 \mid 12$ and $6 \mid 18$. Also,
 339 $6 = 18 - 12 \in I$, so $I = (6)$.

340 The ring \mathbb{Z} in Sage is ZZ, which is Noetherian.

```

ZZ.is_noetherian()

```

341 **True**

342 We create the ideal I in Sage as follows, and note that it is principal:

```

I = ideal(12,18); I

```

Principal ideal (6) of Integer Ring

```

I.is_principal()

```

True

344 We could also create I as follows:

```
345      ZZ.ideal(12,18)
```

```
345      | Principal ideal (6) of Integer Ring
```

346 Propositions 1.2.9 and 1.2.12 together imply that any finitely generated
 347 abelian group is noetherian. This means that subgroups of finitely generated
 348 abelian groups are finitely generated, which provides the missing step in our
 349 proof of the structure theorem for finitely generated abelian groups.

350 **Exercise 1.2.14.** There is another way to show every principle ideal domain
 351 (for example \mathbb{Z}) is noetherian (contrast to the proof in Section 1.2.1). Let
 352 R be a PID and (a) an arbitrary ideal. Use the facts that $(b) \supseteq (a)$ if and
 353 only if $b \mid a$ and that R is a UFD to show that any ascending chain of ideals
 354 starting with (a) must stabilize.

355 1.3 Rings of Algebraic Integers

356 In this section we introduce the central objects of this book, which are the
 357 rings of algebraic integers. These are noetherian rings with an enormous
 358 amount of structure. We also introduce a function field analogue of these
 359 rings.

360 An *algebraic number* is a root of some nonzero polynomial $f(x) \in \mathbb{Q}[x]$.
 361 For example, $\sqrt{2}$ and $\sqrt{5}$ are both algebraic numbers, being roots of $x^2 - 2$
 362 and $x^2 - 5$, respectively. But is $\sqrt{2} + \sqrt{5}$ necessarily the root of some
 363 polynomial in $\mathbb{Q}[x]$? This isn't quite so obvious.

364 **Proposition 1.3.1.** *An element α of a field extension of \mathbb{Q} is an algebraic*
 365 *number if and only if the ring $\mathbb{Q}[\alpha]$ generated by α is finite dimensional as*
 366 *a \mathbb{Q} vector space.*

367 *Proof.* Suppose α is an algebraic number, so there is a nonzero polynomial
 368 $f(x) \in \mathbb{Q}[x]$, so that $f(\alpha) = 0$. The equation $f(\alpha) = 0$ implies that $\alpha^{\deg(f)}$
 369 can be written in terms of smaller powers of α , so $\mathbb{Q}[\alpha]$ is spanned by the
 370 finitely many numbers $1, \alpha, \dots, \alpha^{\deg(f)-1}$, hence finite dimensional. Con-
 371 versely, suppose $\mathbb{Q}[\alpha]$ is finite dimensional. Then for some $n \geq 1$, we have
 372 that α^n is in the \mathbb{Q} -vector space spanned by $1, \alpha, \dots, \alpha^{n-1}$. Thus α satisfies
 373 a polynomial $f(x) \in \mathbb{Q}[x]$ of degree n . \square

374 **Proposition 1.3.2.** *Suppose K is a field and $\alpha, \beta \in K$ are two algebraic*
 375 *numbers. Then $\alpha\beta$ and $\alpha + \beta$ are also algebraic numbers.*

376 *Proof.* Let $n = \dim_{\mathbb{Q}} \mathbb{Q}[\alpha]$ and $n = \dim_{\mathbb{Q}} \mathbb{Q}[\beta]$. The subring $\mathbb{Q}[\alpha, \beta] \subset K$ is
 377 a \mathbb{Q} -vector space that is spanned by the numbers $\alpha^i \beta^j$, where $0 \leq i < n$ and
 378 $0 \leq j < m$. Thus $\mathbb{Q}[\alpha, \beta]$ is finite dimensional, and since $\alpha + \beta$ and $\alpha\beta$ are
 379 both in $\mathbb{Q}[\alpha, \beta]$, we conclude by Proposition 1.3.1 that both are algebraic
 380 numbers. \square

381 Suppose C is a field extension of \mathbb{Q} such that every polynomial $f(x) \in$
 382 $\mathbb{Q}[x]$ factors completely in C . The algebraic closure $\overline{\mathbb{Q}}$ of \mathbb{Q} inside C is the
 383 field generated by all roots in C of polynomials in $\mathbb{Q}[x]$. The fundamental
 384 theorem of algebra tells us that $C = \mathbb{C}$ is one choice of field C as above.
 385 There are other fields C , e.g., constructed using p -adic numbers. One can
 386 show that any two choices of $\overline{\mathbb{Q}}$ are isomorphic; however, there will be *many*
 387 isomorphisms between them.

388 **Definition 1.3.3** (Algebraic Integer). An element $\alpha \in \overline{\mathbb{Q}}$ is an *algebraic*
 389 *integer* if it is a root of some monic polynomial with coefficients in \mathbb{Z} .

390 For example, $\sqrt{2}$ is an algebraic integer, since it is a root of the monic
 391 integral polynomial $x^2 - 2$. As we will see below, $1/2$ is not an algebraic
 392 integer.

393 The following two propositions are analogous to Propositions 1.3.1–1.3.2
 394 above, with the proofs replacing basic facts about vector spaces with facts
 395 we proved above about noetherian rings and modules.

396 **Proposition 1.3.4.** *An element $\alpha \in \overline{\mathbb{Q}}$ is an algebraic integer if and only*
 397 *if $\mathbb{Z}[\alpha]$ is finitely generated as a \mathbb{Z} -module.*

398 *Proof.* Suppose α is integral and let $f \in \mathbb{Z}[x]$ be a monic integral poly-
 399 nomial such that $f(\alpha) = 0$. Then, as a \mathbb{Z} -module, $\mathbb{Z}[\alpha]$ is generated by
 400 $1, \alpha, \alpha^2, \dots, \alpha^{d-1}$, where d is the degree of f . Conversely, suppose $\alpha \in \overline{\mathbb{Q}}$
 401 is such that $\mathbb{Z}[\alpha]$ is finitely generated as a module over \mathbb{Z} , say by elements
 402 $f_1(\alpha), \dots, f_n(\alpha)$. Let d be any integer bigger than the degrees of all f_i .
 403 Then there exist integers a_i such that $\alpha^d = \sum_{i=1}^n a_i f_i(\alpha)$, hence α satis-
 404 fies the monic polynomial $x^d - \sum_{i=1}^n a_i f_i(x) \in \mathbb{Z}[x]$, so α is an algebraic
 405 integer. \square

406 The proof of the following proposition uses repeatedly that any submod-
 407 ule of a finitely generated \mathbb{Z} -module is finitely generated, which uses that \mathbb{Z}
 408 is noetherian and that finitely generated modules over a noetherian ring are
 409 noetherian.

410 **Proposition 1.3.5.** *Suppose K is a field and $\alpha, \beta \in K$ are two algebraic*
 411 *integers. Then $\alpha\beta$ and $\alpha + \beta$ are also algebraic integers.*

412 *Proof.* Let m, n be the degrees of monic integral polynomials that have α, β
 413 as roots, respectively. Then we can write α^m in terms of smaller powers of
 414 α and likewise for β^n , so the elements $\alpha^i \beta^j$ for $0 \leq i < m$ and $0 \leq j < n$
 415 span the \mathbb{Z} -module $\mathbb{Z}[\alpha, \beta]$. Since $\mathbb{Z}[\alpha + \beta]$ is a submodule of the finitely-
 416 generated \mathbb{Z} -module $\mathbb{Z}[\alpha, \beta]$, it is finitely generated, so $\alpha + \beta$ is integral.
 417 Likewise, $\mathbb{Z}[\alpha\beta]$ is a submodule of $\mathbb{Z}[\alpha, \beta]$, so it is also finitely generated,
 418 and $\alpha\beta$ is integral. \square

419 1.3.1 Minimal Polynomials

420 **Definition 1.3.6** (Minimal Polynomial). The *minimal polynomial* of $\alpha \in \overline{\mathbb{Q}}$
 421 is the monic polynomial $f \in \mathbb{Q}[x]$ of least positive degree such that $f(\alpha) = 0$.

422 It is a consequence of Lemma 1.3.9 below that “the” minimal polynomial
 423 of α is unique. The minimal polynomial of $1/2$ is $x - 1/2$, and the minimal
 424 polynomial of $\sqrt[3]{2}$ is $x^3 - 2$.

425 *Example 1.3.7.* We compute the minimal polynomial of $(\sqrt[3]{2})^2 + 3$. in terms
 426 of $\sqrt[4]{2}$:

this is confusing,
 sometimes easier
 to use numberfield
 to construct ele-
 ments rather than
 typing `(sqrt(2) +
 3).minpoly()`

```
427 K.<a> = NumberField(x^4 - 2)
      a^4
```

```
| 2
```

```
(a^2 + 3).minpoly()
```

```
| x^2 - 6*x + 7
```

428 **Exercise 1.3.8.** Find the minimal polynomial of $\sqrt{2} + \sqrt{3}$ by hand. Check
 429 your result with **Sage**.

430 **Lemma 1.3.9.** Suppose $\alpha \in \overline{\mathbb{Q}}$. Then the minimal polynomial of α divides
 431 any polynomial h such that $h(\alpha) = 0$.

Proof. Let f be a choice of minimal polynomial of α , as in Definition 1.3.6,
 and let h be a polynomial with $h(\alpha) = 0$. Use the division algorithm to
 write $h = qf + r$, where $0 \leq \deg(r) < \deg(f)$. We have

$$r(\alpha) = h(\alpha) - q(\alpha)f(\alpha) = 0,$$

432 so α is a root of r . However, f is a polynomial of least positive degree with
 433 root α , so $r = 0$. \square

434 **Exercise 1.3.10.** Show that the minimal polynomial of an algebraic number
 435 $\alpha \in \overline{\mathbb{Q}}$ is unique.

436 **Lemma 1.3.11.** Suppose $\alpha \in \overline{\mathbb{Q}}$. Then α is an algebraic integer if and only
 437 if the minimal polynomial f of α has coefficients in \mathbb{Z} .

438 *Proof.* First suppose that the minimal polynomial f of α has coefficients in
 439 \mathbb{Z} . Since $f \in \mathbb{Z}[x]$ is monic (by definition) and $f(\alpha) = 0$, we see immediately
 440 that α is an algebraic integer.

441 Now suppose that α is an algebraic integer. Then there is some nonzero
 442 monic $g \in \mathbb{Z}[x]$ such that $g(\alpha) = 0$. By Lemma 1.3.9, we have $g = fh$,
 443 for some $h \in \mathbb{Q}[x]$, and h is monic because f and g are. If $f \notin \mathbb{Z}[x]$, then
 444 some prime p divides the denominator of some coefficient of f . Let p^i be
 445 the largest power of p that divides some denominator of some coefficient f ,
 446 and likewise let p^j be the largest power of p that divides some denominator
 447 of a coefficient of h . Then $p^{i+j}g = (p^i f)(p^j h)$, and if we reduce both sides
 448 modulo p , then the left hand side is 0 but the right hand side is a product
 449 of two nonzero polynomials in $\mathbb{F}_p[x]$, hence nonzero, a contradiction. \square

450 **Exercise 1.3.12.** Which of the following numbers are algebraic integers?

- 451 (a) The number $(1 + \sqrt{5})/2$.
 452 (b) The number $(2 + \sqrt{5})/2$.
 453 (c) The value of the infinite sum $\sum_{n=1}^{\infty} 1/n^2$.
 454 (d) The number $\alpha/3$, where α is a root of $x^4 + 54x + 243$.

455 **Example 1.3.13.** We compute some minimal polynomials in Sage. The min-
 456 imal polynomial of $1/2$:

```
(1/2).minpoly()
```

```
x - 1/2
```

We construct a root a of $x^2 - 2$ and compute its minimal polynomial:

457

```
K.<a> = NumberField(x^2 - 2)
a^2 - 2
```

```
0
```

```
a.minpoly()
```

```
x^2 - 2
```

make sure this is bold

make sure we use big
K for number fields

458 Finally we compute the minimal polynomial of $\alpha = \sqrt{2}/2 + 3$, which is not
 459 integral, hence Proposition 1.3.4 implies that α is not an algebraic integer:

```
(a/2 + 3).minpoly()
```

460

```
| x^2 - 6*x + 17/2
```

The only elements of \mathbb{Q} that are algebraic integers are the usual integers \mathbb{Z} , since $\mathbb{Z}[1/d]$ is not finitely generated as a \mathbb{Z} -module. Watch out since there are elements of \mathbb{Q} that seem to *appear* to have denominators when written down, but are still algebraic integers. This is an artifact of how we write them down, e.g., if we wrote our integers as a multiple of $\alpha = 2$, then we would write 1 as $\alpha/2$. For example,

$$\alpha = \frac{1 + \sqrt{5}}{2}$$

461 is an algebraic integer, since it is a root of the monic integral polynomial
 462 $x^2 - x - 1$. We verify this using **Sage** below, though of course this is easy
 463 to do by hand (you should try much more complicated examples in **Sage**).

```
k.<a> = QuadraticField(5)
a^2
```

464

```
| 5
```

```
alpha = (1 + a)/2
alpha.minpoly()
```

```
| x^2 - x - 1
```

```
alpha.is_integral()
```

```
| True
```

465 Since $\sqrt{5}$ can be expressed in terms of radicals, we can also compute this
 466 minimal polynomial using the symbolic functionality in Sage.

```
alpha = (1+sqrt(5))/2
alpha.minpoly()
```

```
| x^2 - x - 1
```

467 Here is a more complicated example using a similar approach:

```
alpha = sqrt(2) + 3^(1/4)
alpha.minpoly()
```

```
| x^8 - 8*x^6 + 18*x^4 - 104*x^2 + 1
```

468 *Example 1.3.14.* We illustrate an example of a sum and product of two
 469 algebraic integers being an algebraic integer. We first make the relative
 470 number field obtained by adjoining a root of $x^3 - 5$ to the field $\mathbb{Q}(\sqrt{2})$:

```
k.<a, b> = NumberField([x^2 - 2, x^3 - 5])
k
```

471

```
| Number Field in a with defining polynomial x^2 + -2 over its base field
```

472 Here a and b are roots of $x^2 - 2$ and $x^3 - 5$, respectively.

```
a^2
```

```
| 2
```

473

```
b^3
```

```
| 5
```

474 We compute the minimal polynomial of the sum and product of $\sqrt[3]{5}$ and
 475 $\sqrt{2}$. The command `absolute_minpoly` gives the minimal polynomial of the
 476 element over the rational numbers \mathbb{Q} .

```
(a+b).absolute_minpoly()
```

```
| x^6 - 6*x^4 - 10*x^3 + 12*x^2 - 60*x + 17
```

477

```
(a*b).absolute_minpoly()
```

```
| x^6 - 200
```

The minimal polynomial of the product is $\sqrt[3]{5}\sqrt{2}$ is trivial to compute by hand. In light of the Cayley-Hamilton theorem, we can compute the minimal polynomial of $\alpha = \sqrt[3]{5} + \sqrt{2}$ by hand by computing the determinant of the matrix given by left multiplication by α on the basis

$$1, \sqrt{2}, \sqrt[3]{5}, \sqrt[3]{5}\sqrt{2}, \sqrt[3]{5}^2, \sqrt[3]{5}^2\sqrt{2}.$$

478 This is a general method which works well for computers. However it can
479 also be done using simple algebra.

480 The following is an alternative, more symbolic way to compute the min-
481 imal polynomials above, though it is not provably correct. We compute α
482 to 100 bits precision (via the `n` command), then use the LLL algorithm (via
483 the `algdep` command) to heuristically find a linear relation between the first
484 6 powers of α (see Section 1.5 below for more about LLL).

```

a = 5^(1/3); b = sqrt(2)
c = a+b; c

5^(1/3) + sqrt(2)

(a+b).n(100).algdep(6)

x^6 - 6*x^4 - 10*x^3 + 12*x^2 - 60*x + 17

(a*b).n(100).algdep(6)

x^6 - 200
```

is this example too
long?

487 **Exercise 1.3.15.** Compute the minimal polynomial of $\alpha = \sqrt[3]{5} + \sqrt{2}$ by
488 hand without finding the determinate of a 6×6 matrix.

489 [Hint: Let $a^2 = 2$, $b^3 = 5$, and $x = a + b$. Then $(x - a)^3 = b^3 = 5$. Now
490 simplify and use the fact that $a^2 = 2$.]

491 **Exercise 1.3.16.** Let $\alpha = \sqrt{2} + \frac{1+\sqrt{5}}{2}$.

492 (a) Is α an algebraic integer?

493 (b) Explicitly write down the minimal polynomial of α as an element of
494 $\mathbb{Q}[x]$.

1.3.2 Number fields, rings of integers, and orders

Definition 1.3.17 (Number field). A *number field* is a field K that contains the rational numbers \mathbb{Q} such that the degree $[K : \mathbb{Q}] = \dim_{\mathbb{Q}}(K)$ is finite.

If K is a number field, then by the primitive element theorem there is an $\alpha \in K$ so that $K = \mathbb{Q}(\alpha)$. Let $f(x) \in \mathbb{Q}[x]$ be the minimal polynomial of α . Fix a choice of algebraic closure $\overline{\mathbb{Q}}$ of \mathbb{Q} . Associated to each of the $\deg(f)$ roots $\alpha' \in \overline{\mathbb{Q}}$ of f , we obtain a field embedding $K \hookrightarrow \overline{\mathbb{Q}}$ that sends α to α' . Thus any number field can be embedded in $[K : \mathbb{Q}] = \deg(f)$ distinct ways in $\overline{\mathbb{Q}}$.

Definition 1.3.18 (Ring of Integers). The *ring of integers* of a number field K is the ring

$$\mathcal{O}_K = \{x \in K : x \text{ is an algebraic integer}\}.$$

One of the most basic facts about \mathcal{O}_K is that it is indeed a ring. This fact is important enough to be stated as a separate theorem.

Theorem 1.3.19. *Let K be a number field. The ring of integers \mathcal{O}_K is a ring.*

Proof. This follows directly from Proposition 1.3.5. □

Example 1.3.20. The field \mathbb{Q} of rational numbers is a number field of degree 1, and the ring of integers of \mathbb{Q} is \mathbb{Z} . The field $K = \mathbb{Q}(i)$ of Gaussian integers has degree 2 and $\mathcal{O}_K = \mathbb{Z}[i]$.

Example 1.3.21. The golden ratio $\varphi = (1 + \sqrt{5})/2$ is in the quadratic number field $K = \mathbb{Q}(\sqrt{5}) = \mathbb{Q}(\varphi)$; notice that φ satisfies $x^2 - x - 1$, so $\varphi \in \mathcal{O}_K$. To see that $\mathcal{O}_K = \mathbb{Z}[\varphi]$ directly, we proceed as follows. By Proposition 1.3.4, the algebraic integers K are exactly the elements $a + b\sqrt{5} \in K$, with $a, b \in \mathbb{Q}$ that have integral minimal polynomial. The matrix of $a + b\sqrt{5}$ with respect to the basis $1, \sqrt{5}$ for K is $m = \begin{pmatrix} a & 5b \\ b & a \end{pmatrix}$. The characteristic polynomial of m is $f = (x - a)^2 - 5b^2 = x^2 - 2ax + a^2 - 5b^2$, which is in $\mathbb{Z}[x]$ if and only if $2a \in \mathbb{Z}$ and $a^2 - 5b^2 \in \mathbb{Z}$. Thus $a = a'/2$ with $a' \in \mathbb{Z}$, and $(a'/2)^2 - 5b^2 \in \mathbb{Z}$, so $5b^2 \in \frac{1}{4}\mathbb{Z}$, so $b \in \frac{1}{2}\mathbb{Z}$ as well. If a has a denominator of 2, then b must also have a denominator of 2 to ensure that the difference $a^2 - 5b^2$ is an integer. This proves that $\mathcal{O}_K = \mathbb{Z}[\varphi]$.

Example 1.3.22. The ring of integers of $K = \mathbb{Q}(\sqrt[3]{9})$ is $\mathbb{Z}[\sqrt[3]{3}]$, where $\sqrt[3]{3} = \frac{1}{3}(\sqrt[3]{9})^2 \notin \mathbb{Z}[\sqrt[3]{9}]$. As we will see, in general the problem of computing \mathcal{O}_K given K may be very hard, since it requires factoring a certain potentially large integer.

make this better

527 **Exercise 1.3.23.** From basic definitions, find the rings of integers of the
 528 fields $\mathbb{Q}(\sqrt{11})$ and $\mathbb{Q}(\sqrt{-6})$.

529 **Definition 1.3.24** (Order). An *order* in \mathcal{O}_K is any subring R of \mathcal{O}_K such
 530 that the quotient \mathcal{O}_K/R of abelian groups is finite. (By definition R must
 531 contain 1 because it is a ring.)

532 **Exercise 1.3.25.** Let R be a subring of \mathcal{O}_K . Show that R is an order of
 533 \mathcal{O}_K if and only if R contains a spanning set for K as a vector space over \mathbb{Q} .

534 **Exercise 1.3.26.** Let K be a number field of degree n . Suppose $\{\alpha_1, \dots, \alpha_n\}$
 535 is a \mathbb{Z} -independent set of algebraic integers. Is $\mathbb{Z}\alpha_1 + \dots + \mathbb{Z}\alpha_n$ an ideal of
 536 \mathcal{O}_K ?

537 As noted above, $\mathbb{Z}[i]$ is the ring of integers of $\mathbb{Q}(i)$. For every nonzero
 538 integer n , the subring $\mathbb{Z} + ni\mathbb{Z}$ of $\mathbb{Z}[i]$ is an order. The subring \mathbb{Z} of $\mathbb{Z}[i]$ is not
 539 an order, because \mathbb{Z} does not have finite index in $\mathbb{Z}[i]$. Also the subgroup
 540 $2\mathbb{Z} + i\mathbb{Z}$ of $\mathbb{Z}[i]$ is not an order because it is not a ring.

541 **Exercise 1.3.27.** Let K be a quadratic extension of \mathbb{Q} and R be any order
 542 in \mathcal{O}_K . Show that \mathcal{O}_K/R is cyclic as an abelian group and that there is a
 543 bijection between orders of \mathcal{O}_K containing R and divisors of $[\mathcal{O}_K : R]$.

544 *Remark 1.3.28.* Exercise 1.3.27 is used in elliptic curve cryptography to
 545 measure the number of isogenies; for example, see [KKM11, §11.2].

546 **Exercise 1.3.29.** Let K be a number field of degree n . Suppose $\{\alpha_1, \dots, \alpha_n\}$
 547 is a \mathbb{Z} -independent set of algebraic integers. Is $\mathbb{Z}\alpha_1 + \dots + \mathbb{Z}\alpha_n$ an ideal of
 548 \mathcal{O}_K ?

find a good place for
this

549 We define the number field $\mathbb{Q}(i)$ and compute its ring of integers.

```
K.<i> = NumberField(x^2 + 1)
OK = K.ring_of_integers(); OK
```

550

```
Order with module basis 1, i in Number Field in i with
defining polynomial x^2 + 1
```

551 Next we compute the order $\mathbb{Z} + 3i\mathbb{Z}$.

```
O3 = K.order(3*i); O3
```

```
Order with module basis 1, 3*i in Number Field in i with
defining polynomial x^2 + 1
```

552

```
O3.gens()
```

```
[1, 3*i]
```


553 We test whether certain elements are in the order.

```
5 + 9*i in O3
```

```
| True
```

554

```
1 + 2*i in O3
```

```
| False
```

555 We will frequently consider orders because they are often much easier
 556 to write down explicitly than \mathcal{O}_K . For example, if $K = \mathbb{Q}(\alpha)$ and α is an
 557 algebraic integer, then $\mathbb{Z}[\alpha]$ is an order in \mathcal{O}_K , but frequently $\mathbb{Z}[\alpha] \neq \mathcal{O}_K$.

558 *Example 1.3.30.* In this example $[\mathcal{O}_K : \mathbb{Z}[a]] = 2197$. First we define the
 559 number field $K = \mathbb{Q}(a)$ where a is a root of $x^3 - 15x^2 - 94x - 3674$, then
 560 we compute the order $\mathbb{Z}[a]$ generated by a .

```
K.<a> = NumberField(x^3 - 15*x^2 - 94*x - 3674)
Oa = K.order(a); Oa
```

561

```
| Order with module basis 1, a, a^2 in Number Field in a with defining
  polynomial x^3 - 15*x^2 - 94*x - 3674
```

```
Oa.basis()
```

```
| [1, a, a^2]
```

562 Next we compute a \mathbb{Z} -basis for the maximal order \mathcal{O}_K of K , and compute
 563 that the index of $\mathbb{Z}[a]$ in \mathcal{O}_K is $2197 = 13^3$.

```
OK = K.maximal_order()
OK.basis()
```

564

```
| [25/169*a^2 + 10/169*a + 1/169, 5/13*a^2 + 1/13*a, a^2]
```

```
Oa.index_in(OK)
```

```
| 2197
```

565 **Lemma 1.3.31.** *Let \mathcal{O}_K be the ring of integers of a number field. Then*
 566 *$\mathcal{O}_K \cap \mathbb{Q} = \mathbb{Z}$ and $\mathbb{Q}\mathcal{O}_K = K$.*

567 *Proof.* Suppose $\alpha \in \mathcal{O}_K \cap \mathbb{Q}$ with $\alpha = a/b \in \mathbb{Q}$ in lowest terms and $b > 0$.
 568 Since α is integral, $\mathbb{Z}[a/b]$ is finitely generated as a module, so $b = 1$.

To prove that $\mathbb{Q}\mathcal{O}_K = K$, suppose $\alpha \in K$, and let $f(x) \in \mathbb{Q}[x]$ be the minimal monic polynomial of α . For any positive integer d , the minimal monic polynomial of $d\alpha$ is $d^{\deg(f)}f(x/d)$, i.e., the polynomial obtained from $f(x)$ by multiplying the coefficient of $x^{\deg(f)}$ by 1, multiplying the coefficient of $x^{\deg(f)-1}$ by d , multiplying the coefficient of $x^{\deg(f)-2}$ by d^2 , etc. If d is the least common multiple of the denominators of the coefficients of f , then the minimal monic polynomial of $d\alpha$ has integer coefficients, so $d\alpha$ is integral and $d\alpha \in \mathcal{O}_K$. This proves that $\mathbb{Q}\mathcal{O}_K = K$. \square

Exercise 1.3.32. Which of the following rings are orders in the given number field, i.e. orders in the ring of integers of the given number field.

- (a) The ring $R = \mathbb{Z}[i]$ in the number field $\mathbb{Q}(i)$.
- (b) The ring $R = \mathbb{Z}[i/2]$ in the number field $\mathbb{Q}(i)$.
- (c) The ring $R = \mathbb{Z}[17i]$ in the number field $\mathbb{Q}(i)$.
- (d) The ring $R = \mathbb{Z}[i]$ in the number field $\mathbb{Q}(\sqrt[4]{-1})$.

Exercise 1.3.33. Find the ring of integers of $\mathbb{Q}(\alpha)$, where $\alpha^5 + 7\alpha + 1 = 0$ using a computer.

1.3.3 Function fields

Let k be any field. We can also make the same definitions, but with \mathbb{Q} replaced by the field $k(t)$ of rational functions in an indeterminate t , and \mathbb{Z} replaced by $k[t]$. The analogue of a number field is called a *function field*; it is a finite algebraic extension field K of $k(t)$. Elements of K have a unique minimal polynomial as above, and the ring of integers of K consists of those elements whose monic minimal polynomial has coefficients in the polynomial ring $k[t]$.

Geometrically, if $F(x, t) = 0$ is an affine equation that defines (via projective closure) a nonsingular projective curve C , then $K = k(t)[x]/(F(x, t))$ is a function field. We view the field K as the field of all rational functions on the projective closure of the curve C . The ring of integers \mathcal{O}_K is the subring of rational functions that have no poles on the affine curve $F(x, t) = 0$, though they may have poles at infinity, i.e., at the extra points we introduce when passing to the projective closure C . The algebraic arguments we gave above prove that \mathcal{O}_K is a ring. This is also geometrically intuitive, since the sum and product of two functions with no poles also have no poles.

Exercise 1.3.34. Let $k = \mathbb{F}_p$ be the finite field with p elements where p is some prime. Find all automorphisms of $k(t)$. Note that an automorphism is completely characterized by its value on t . How many such automorphisms are there?

[*Hint:* For some people, it is easier to think about the equivalent question: What rational functions $f \in k(t)$ is the map $k(t) \rightarrow k(t)$ given by $t \mapsto f(t)$ an automorphism?]

1.4 Norms and Traces

In this section we develop some basic properties of norms, traces, and discriminants, and give more properties of rings of integers in the general context of Dedekind domains.

Before discussing norms and traces we introduce some notation for field extensions. If $K \subset L$ are number fields, we let $[L : K]$ denote the dimension of L viewed as a K -vector space. If K is a number field and $a \in \overline{\mathbb{Q}}$, let $K(a)$ be the extension of K generated by a , which is the smallest number field that contains both K and a . If $a \in \overline{\mathbb{Q}}$ then a has a minimal polynomial $f(x) \in \mathbb{Q}[x]$, and the *Galois conjugates* of a are the roots of f . These are called the Galois conjugates because they are the orbit of a under the action of $\text{Gal}(\overline{\mathbb{Q}}/\mathbb{Q})$.

Example 1.4.1. The element $\sqrt{2}$ has minimal polynomial $x^2 - 2$ and the Galois conjugates of $\sqrt{2}$ are $\sqrt{2}$ and $-\sqrt{2}$. The cube root $\sqrt[3]{2}$ has minimal polynomial $x^3 - 2$ and three Galois conjugates $\sqrt[3]{2}, \zeta_3 \sqrt[3]{2}, \zeta_3^2 \sqrt[3]{2}$, where ζ_3 is a cube root of unity, e.g. $\zeta_3 = e^{2\pi i/3}$.

We can create the extension $\mathbb{Q}(\zeta_3)(\sqrt[3]{2})$ in Sage in this way:

```
L.<cuberooroot2> = CyclotomicField(3).extension(x^3 - 2)
cuberooroot2^3
```

626

2

Then we list the Galois conjugates of $\sqrt[3]{2}$.

```
cuberooroot2.galois_conjugates(L)

[cuberooroot2, (-zeta3 - 1)*cuberooroot2, zeta3*cuberooroot2]
```

628

Note that $\zeta_3^2 = -\zeta_3 - 1$:

629

```

630
zeta3 = L.base_field().0
zeta3^2

- zeta3 - 1

```

use alpha instead of zeta3

Suppose $K \subset L$ is an inclusion of number fields and let $a \in L$. Then left multiplication by a defines a K -linear transformation $\ell_a : L \rightarrow L$. (The transformation ℓ_a is K -linear because L is commutative.)

Example 1.4.2. Let $K = \mathbb{Q}$ and $L = \mathbb{Q}(\sqrt{5})$. Then $B = \{1, \sqrt{5}\}$ is a basis for $\mathbb{Q}(\sqrt{5})$ as a \mathbb{Q} -vector space. So we can identify $\mathbb{Q}(\sqrt{5})$ with \mathbb{Q}^2 by

$$a + b\sqrt{5} \leftrightarrow \begin{pmatrix} a \\ b \end{pmatrix}$$

Let $\alpha = 7 + 3\sqrt{5}$. The matrix for ℓ_α with respect to the basis B is

$$\ell_\alpha = \begin{pmatrix} 7 & 15 \\ 3 & 7 \end{pmatrix}.$$

The following is an example of how to translate from the language of algebraic numbers to the language of linear algebra:

$$\alpha(2 + \sqrt{5}) + (3 + 5\sqrt{5}) \leftrightarrow \begin{pmatrix} 7 & 15 \\ 3 & 7 \end{pmatrix} \begin{pmatrix} 2 \\ 1 \end{pmatrix} + \begin{pmatrix} 3 \\ 5 \end{pmatrix}.$$

Definition 1.4.3 (Norm and Trace). The *norm* and *trace* of a from L to K are

$$\text{Norm}_{L/K}(a) = \det(\ell_a) \quad \text{and} \quad \text{Trace}_{L/K}(a) = \text{Trace}(\ell_a).$$

Example 1.4.4. Continuing example 1.4.2, we can compute

$$\text{Norm}_{\mathbb{Q}(\sqrt{5})/\mathbb{Q}}(7 + 3\sqrt{5}) = \det \begin{pmatrix} 7 & 15 \\ 3 & 7 \end{pmatrix} = 49 - 45 = 4$$

and

$$\text{Trace}_{\mathbb{Q}(\sqrt{5})/\mathbb{Q}}(7 + 3\sqrt{5}) = \text{Trace} \begin{pmatrix} 7 & 15 \\ 3 & 7 \end{pmatrix} = 14.$$

We know from linear algebra that determinants are multiplicative and traces are additive, so for $a, b \in L$ we have

$$\text{Norm}_{L/K}(ab) = \text{Norm}_{L/K}(a) \cdot \text{Norm}_{L/K}(b)$$

and

$$\text{Trace}_{L/K}(a + b) = \text{Trace}_{L/K}(a) + \text{Trace}_{L/K}(b).$$

641 Note that if $f \in \mathbb{Q}[x]$ is the characteristic polynomial of ℓ_a , then the
 642 constant term of f is $(-1)^{\deg(f)} \det(\ell_a)$, and the coefficient of $x^{\deg(f)-1}$ is
 643 $-\text{Trace}(\ell_a)$.

644 **Proposition 1.4.5.** *Let $a \in L$ and let $\sigma_1, \dots, \sigma_d$, where $d = [L : K]$, be the
 645 distinct field embeddings $L \hookrightarrow \overline{\mathbb{Q}}$ that fix every element of K . Then*

$$\text{Norm}_{L/K}(a) = \prod_{i=1}^d \sigma_i(a) \quad \text{and} \quad \text{Trace}_{L/K}(a) = \sum_{i=1}^d \sigma_i(a).$$

646 *Proof.* We prove the proposition by computing the characteristic polynomial
 647 of a . Let $f \in K[x]$ be the minimal polynomial of a over K , and note that f
 648 has distinct roots and is irreducible, since it is the polynomial in $K[x]$ of
 649 least degree that is satisfied by a and K has characteristic 0. Since f is
 650 irreducible, we have $K(a) \cong K[x]/(f)$, so $[K(a) : K] = \deg(f)$. Also a
 651 satisfies a polynomial if and only if ℓ_a does, so the characteristic polynomial
 652 of ℓ_a acting on $K(a)$ is f . Let b_1, \dots, b_n be a basis for L over $K(a)$ and
 653 note that $1, \dots, a^m$ is a basis for $K(a)/K$, where $m = \deg(f) - 1$. Then
 654 $a^i b_j$ is a basis for L over K , and left multiplication by a acts the same way
 655 on the span of $b_j, ab_j, \dots, a^m b_j$ as on the span of $b_k, ab_k, \dots, a^m b_k$, for any
 656 pair $j, k \leq n$. Thus the matrix of ℓ_a on L is a block direct sum of copies
 657 of the matrix of ℓ_a acting on $K(a)$, so the characteristic polynomial of ℓ_a
 658 on L is $f^{[L:K(a)]}$. The proposition follows because the roots of $f^{[L:K(a)]}$ are
 659 exactly the images $\sigma_i(a)$, with multiplicity $[L : K(a)]$, since each embedding
 660 of $K(a)$ into $\overline{\mathbb{Q}}$ extends in exactly $[L : K(a)]$ ways to L . \square

661 **Warning 1.4.6.** It is important in Proposition 1.4.5 that the product and
 662 sum be over *all* the images $\sigma_i(a)$, not over just the distinct images. For
 663 example, if $a = 1 \in L$, then $\text{Trace}_{L/K}(a) = [L : K]$, whereas the sum of the
 664 distinct conjugates of a is 1.

665 *Remark 1.4.7.* Let $K \subset L$ be an extension of number fields. If $\alpha \in \mathcal{O}_L$, then
 666 the formula of Proposition 1.4.5 implies that the norm and trace down to K
 667 of α is an element of \mathcal{O}_K , because the sum and product of algebraic integers
 668 is an algebraic integer.

669 *Example 1.4.8.* Continuing example 1.4.2, let $\alpha = 7 + 3\sqrt{5}$. The images
 670 of α in the embeddings $\mathbb{Q}(\sqrt{5}) \rightarrow \mathbb{C}$ are $7 + 3\sqrt{5}$ and $7 - 3\sqrt{5}$. So using
 671 Proposition 1.4.5 we can compute

$$\text{Norm}_{\mathbb{Q}(\sqrt{5})/\mathbb{Q}}(7 + 3\sqrt{5}) = (7 + 3\sqrt{5})(7 - 3\sqrt{5}) = 4$$

672 and

$$\text{Trace}_{\mathbb{Q}(\sqrt{5})/\mathbb{Q}}(7 + 3\sqrt{5}) = (7 + 3\sqrt{5}) + (7 - 3\sqrt{5}) = 14.$$

673 The following corollary asserts that the norm and trace behave well in
674 towers.

Corollary 1.4.9. *Suppose $K \subset L \subset M$ is a tower of number fields, and let $a \in M$. Then*

$$\text{Norm}_{M/K}(a) = \text{Norm}_{L/K}(\text{Norm}_{M/L}(a)) \quad \text{and} \quad \text{Trace}_{M/K}(a) = \text{Trace}_{L/K}(\text{Trace}_{M/L}(a)).$$

675 *Proof.* The proof uses that every embedding $L \hookrightarrow \overline{\mathbb{Q}}$ extends in exactly
676 $[M : L]$ way to an embedding $M \hookrightarrow \overline{\mathbb{Q}}$. This is clear if we view M as
677 $L[x]/(h(x))$ for some irreducible polynomial $h(x) \in L[x]$ of degree $[M : L]$,
678 and note that the extensions of $L \hookrightarrow \overline{\mathbb{Q}}$ to M correspond to the roots of h ,
679 of which there are $\deg(h)$, since $\overline{\mathbb{Q}}$ is algebraically closed.

680 For the first equation, both sides are the product of $\sigma_i(a)$, where σ_i
681 runs through the embeddings of M into $\overline{\mathbb{Q}}$ that fix K . To see this, suppose
682 $\sigma : L \rightarrow \overline{\mathbb{Q}}$ fixes K . If σ' is an extension of σ to M , and τ_1, \dots, τ_d are
683 the embeddings of M into $\overline{\mathbb{Q}}$ that fix L , then $\sigma'\tau_1, \dots, \sigma'\tau_d$ are exactly the
684 extensions of σ to M . For the second statement, both sides are the sum of
685 the $\sigma_i(a)$. \square

686 **Proposition 1.4.10.** *Let K be a number field. The ring of integers \mathcal{O}_K is*
687 *a lattice in K , i.e., $\mathbb{Q}\mathcal{O}_K = K$ and \mathcal{O}_K is an abelian group of rank $[K : \mathbb{Q}]$.*

Proof. We saw in Lemma 1.3.31 that $\mathbb{Q}\mathcal{O}_K = K$. Thus there exists a basis
 a_1, \dots, a_n for K , where each a_i is in \mathcal{O}_K . Suppose that as $x = \sum_{i=1}^n c_i a_i \in$
 \mathcal{O}_K varies over all elements of \mathcal{O}_K the denominators of the coefficients c_i
are not all uniformly bounded. Then subtracting off integer multiples of the
 a_i , we see that as $x = \sum_{i=1}^n c_i a_i \in \mathcal{O}_K$ varies over elements of \mathcal{O}_K with c_i
between 0 and 1, the denominators of the c_i are also arbitrarily large. This
implies that there are infinitely many elements of \mathcal{O}_K in the bounded subset

$$S = \{c_1 a_1 + \dots + c_n a_n : c_i \in \mathbb{Q}, 0 \leq c_i \leq 1\} \subset K.$$

688 Thus for any $\varepsilon > 0$, there are elements $a, b \in \mathcal{O}_K$ such that the coefficients
689 of $a - b$ are all less than ε (otherwise the elements of \mathcal{O}_K would all be a
690 “distance” of least ε from each other, so only finitely many of them would
691 fit in S).

692 As mentioned above, the norms of elements of \mathcal{O}_K are integers. Since the
693 norm of an element is the determinant of left multiplication by that element,
694 the norm is a homogenous polynomial of degree n in the indeterminate
695 coefficients c_i , which is 0 only on the element 0, so the constant term of
696 this polynomial is 0. If the c_i get arbitrarily small for elements of \mathcal{O}_K , then

the values of the norm polynomial get arbitrarily small, which would imply that there are elements of \mathcal{O}_K with positive norm too small to be in \mathbb{Z} , a contradiction. So the set S contains only finitely many elements of \mathcal{O}_K . Thus the denominators of the c_i are bounded, so for some d , we have that \mathcal{O}_K has finite index in $A = \frac{1}{d}\mathbb{Z}a_1 + \cdots + \frac{1}{d}\mathbb{Z}a_n$. Since A is isomorphic to \mathbb{Z}^n , it follows from the structure theorem for finitely generated abelian groups that \mathcal{O}_K is isomorphic as a \mathbb{Z} -module to \mathbb{Z}^n , as claimed. \square

Corollary 1.4.11. *The ring of integers \mathcal{O}_K of a number field is noetherian.*

Proof. By Proposition 1.4.10, the ring \mathcal{O}_K is finitely generated as a module over \mathbb{Z} , so it is certainly finitely generated as a ring over \mathbb{Z} . By Theorem 1.2.11, \mathcal{O}_K is noetherian. \square

1.5 Recognizing Algebraic Numbers using LLL

Suppose we somehow compute a decimal approximation α to some rational number $\beta \in \mathbb{Q}$ and from this wish to recover β . For concreteness, say

$$\beta = \frac{22}{389} = 0.05655526992287917737789203084832904884318766066838046 \dots$$

and we compute

$$\alpha = 0.056555.$$

Now suppose given only α that you would like to recover β . A standard technique is to use continued fractions, which yields a sequence of good rational approximations for α ; by truncating right before a surprisingly big partial quotient (the 23 in the continued fraction \mathfrak{v}), we obtain β :

```

v = continued_fraction(0.056555); v

[0, 17, 1, 2, 6, 1, 23, 1, 1, 1, 1, 1, 2]

convergents([0, 17, 1, 2, 6, 1])

[0, 1/17, 1/18, 3/53, 19/336, 22/389]
```

Generalizing this, suppose next that somehow you numerically approximate an algebraic number, e.g., by evaluating a special function and get a decimal approximation $\alpha \in \mathbb{C}$ to an algebraic number $\beta \in \overline{\mathbb{Q}}$. For concreteness, suppose $\beta = \frac{1}{3} + \sqrt[4]{3}$:

```
N(1/3 + 3^(1/4), digits=50)
```

718

```
1.64940734628582579415255223513033238849340192353916
```

719 Now suppose you very much want to find the (rescaled) minimal polynomial
 720 $f(x) \in \mathbb{Z}[x]$ of β just given this numerical approximation α . This is of great
 721 value even without proof, since often in practice once you know a potential
 722 minimal polynomial you can verify that it is in fact right. Exactly this
 723 situation arises in the explicit construction of class fields (a more advanced
 724 topic in number theory) and in the construction of Heegner points on elliptic
 725 curves. As we will see, the LLL algorithm provides a polynomial time way
 726 to solve this problem, assuming α has been computed to sufficient precision.

727 1.5.1 LLL Reduced Basis

Given a basis b_1, \dots, b_n for \mathbb{R}^n , the *Gramm-Schmidt orthogonalization* process produces an orthogonal basis b_1^*, \dots, b_n^* for \mathbb{R}^n as follows. Define inductively

$$b_i^* = b_i - \sum_{j < i} \mu_{i,j} b_j^*$$

where

$$\mu_{i,j} = \frac{b_i \cdot b_j^*}{b_j^* \cdot b_j^*}.$$

728 *Example 1.5.1.* We compute the Gramm-Schmidt orthogonal basis of the
 729 rows of a matrix. Note that no square roots are introduced in the process;
 730 there would be square roots if we constructed an orthonormal basis.

```
A = matrix(ZZ, 2, [1,2, 3,4]); A
```

731

```
[1 2]
[3 4]
```

```
Bstar, mu = A.gramm_schmidt()
```

732 The rows of the matrix B^* are obtained from the rows of A by the Gramm-
 733 Schmidt procedure.

734

```

Bstar
[ 1 2]
[ 4/5 -2/5]

```

```

mu
[ 0 0]
[11/5 0]

```

735

A lattice $L \subset \mathbb{R}^n$ is a subgroup that is free of rank n such that $\mathbb{R}L = \mathbb{R}^n$.

Definition 1.5.2 (LLL-reduced basis). The basis b_1, \dots, b_n for a lattice $L \subset \mathbb{R}^n$ is *LLL reduced* if for all i, j ,

$$|\mu_{i,j}| \leq \frac{1}{2}$$

and for each $i \geq 2$,

$$|b_i^*|^2 \geq \left(\frac{3}{4} - \mu_{i,i-1}^2 \right) |b_{i-1}^*|^2$$

For example, the basis $b_1 = (1, 2)$, $b_2 = (3, 4)$ for a lattice L is *not* LLL reduced because $b_1^* = b_1$ and

$$\mu_{2,1} = \frac{b_2 \cdot b_1^*}{b_1^* \cdot b_1^*} = \frac{11}{5} > \frac{1}{2}.$$

However, the basis $b_1 = (1, 0)$, $b_2 = (0, 2)$ for L is LLL reduced, since

$$\mu_{2,1} = \frac{b_2 \cdot b_1^*}{b_1^* \cdot b_1^*} = 0,$$

and

$$2^2 \geq (3/4) \cdot 1^2.$$

736

```

A = matrix(ZZ, 2, [1, 2, 3, 4])
A.LLL()

```

```

[1 0]
[0 2]

```

1.5.2 What LLL really means

The following theorem is not too difficult to prove.

Let b_1, \dots, b_n be an LLL reduced basis for a lattice $L \subset \mathbb{R}^n$. Let $d(L)$ denote the absolute value of the determinant of any matrix whose rows are basis for L . Then the vectors b_i are “nearly orthogonal” and “short” in the sense of the following theorem:

Theorem 1.5.3. *We have*

1. $d(L) \leq \prod_{i=1}^n |b_i| \leq 2^{n(n-1)/4} d(L).$

2. *For $1 \leq j \leq i \leq n$, we have*

$$|b_j| \leq 2^{(i-1)/2} |b_i^*|.$$

3. *The vector b_1 is very short in the sense that*

$$|b_1| \leq 2^{(n-1)/4} d(L)^{1/n}$$

and for every nonzero $x \in L$ we have

$$|b_1| \leq 2^{(n-1)/2} |x|.$$

4. *More generally, for any linearly independent $x_1, \dots, x_t \in L$, we have*

$$|b_j| \leq 2^{(n-1)/2} \max(|x_1|, \dots, |x_t|)$$

for $1 \leq j \leq t$.

Perhaps the most amazing thing about the idea of an LLL reduced basis is that there is an algorithm (in fact many) that given a basis for a lattice L produce an LLL reduced basis for L , and do so *quickly*, i.e., in polynomial time in the number of digits of the input. The current optimal implementation (and practically optimal algorithms) for computing LLL reduced basis are due to Damien Stehle, and are included standard in Magma in **Sage**. Stehle’s code is amazing – it can LLL reduce a random lattice in \mathbb{R}^n for $n < 1000$ in a matter of minutes!

```

A = random_matrix(ZZ, 200)
t = cputime()
B = A.LLL()
cputime(t)      # random output

```

3.0494159999999999

755 There is even a very fast variant of Stehle's implementation that computes
 756 a basis for L that is very likely LLL reduced but may in rare cases fail to
 757 be LLL reduced.

```

t = cputime()
B = A.LLL(algorithm="fpLLL:fast") # not tested
cputime(t) # random output

```

758 0.968426999999999837

759 1.5.3 Applying LLL

760 The LLL definition and algorithm has many application in number theory,
 761 e.g., to cracking lattice-based cryptosystems, to enumerating all short vec-
 762 tors in a lattice, to finding relations between decimal approximations to
 763 complex numbers, to very fast univariate polynomial factorization in $\mathbb{Z}[x]$
 764 and more generally in $K[x]$ where K is a number fields, and to computation
 765 of kernels and images of integer matrices. LLL can also be used to solve
 766 the problem of recognizing algebraic numbers mentioned at the beginning
 767 of Section 1.5.

768 Suppose as above that α is a decimal approximation to some algebraic
 769 number β , and to for simplicity assume that $\alpha \in \mathbb{R}$ (the general case of
 770 $\alpha \in \mathbb{C}$ is described in [Coh93]). We finish by explaining how to use LLL to
 771 find a polynomial $f(x) \in \mathbb{Z}[x]$ such that $f(\alpha)$ and its coefficients are small,
 772 hence has a shot at being the minimal polynomial of β .

773 Given a real number decimal approximation α , an integer d (the degree),
 774 and an integer K (a function of the precision to which α is known), the
 775 following steps produce a polynomial $f(x) \in \mathbb{Z}[x]$ of degree at most d such
 776 that $f(\alpha)$ is small.

- 777 1. Form the lattice in \mathbb{R}^{d+2} with basis the rows of the matrix A whose
 778 first $(d+1) \times (d+1)$ part is the identity matrix, and whose last column
 779 has entries

$$K, \lfloor K\alpha \rfloor, \lfloor K\alpha^2 \rfloor, \dots, \lfloor K\alpha^d \rfloor. \quad (1.2)$$

780 (Note this matrix is $(d+1) \times (d+2)$ so the lattice is not of full rank
 781 in \mathbb{R}^{d+2} , which isn't a problem, since the LLL definition also makes
 782 sense for fewer vectors.)

- 783 2. Compute an LLL reduced basis for the \mathbb{Z} -span of the rows of A , and
 784 let B be the corresponding matrix. Let $b_1 = (a_0, a_1, \dots, a_{d+1})$ be the

785 first row of B and notice that B is obtained from A by left multipli-
 786 cation by an invertible integer matrix. Thus a_0, \dots, a_d are the linear
 787 combination of the (1.2) that equals a_{d+1} . Moreover, since B is LLL
 788 reduced we expect that a_{d+1} is relatively small.

789 3. Output $f(x) = a_0 + a_1x + \dots + a_dx^d$. We have that $f(\alpha) \sim a_{d+1}/K$,
 790 which is small. Thus $f(x)$ may be a very good candidate for the
 791 minimal polynomial of β (the algebraic number we are approximating),
 792 assuming d was chosen minimally and α was computed to sufficient
 793 precision.

794 The following is a complete implementation of the above algorithm in
 795 Sage:

```

def myalgdep(a, d, K=10^6):
    aa = [floor(K*a^i) for i in range(d+1)]
    A = identity_matrix(ZZ, d+1)
    B = matrix(ZZ, d+1, 1, aa)
796 A = A.augment(B)
    L = A.LLL()
    v = L[0][: -1].list()
    return ZZ['x'](v)
  
```

797 Here is an example of using it:

```

R.<x> = RDF[]
f = 2*x^3 - 3*x^2 + 10*x - 4
a = f.roots()[0][0]; a
798 myalgdep(a, 3, 10^6) # not tested
  
```

```

2*x^3 - 3*x^2 + 10*x - 4
  
```

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