

₁ ALGEBRAIC NUMBER THEORY,
₂ A COMPUTATIONAL APPROACH

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₄ May 15, 2019

Chapter 1

Introduction

1.1 Mathematical background

In addition to general mathematical maturity, this book assumes you have the following background:

- Basics of finite group theory
- Commutative rings, ideals, quotient rings
- Some elementary number theory
- Basic Galois theory of fields
- Point set topology
- Basics of topological rings, groups, and measure theory

For example, if you have never worked with finite groups before, you should read another book first. If you haven't seen much elementary ring theory, there is still hope, but you will have to do some additional reading and exercises. We will briefly review the basics of the Galois theory of number fields.

Some of the homework problems involve using a computer, but there are examples which you can build on. We will not assume that you have a programming background or know much about algorithms. Most of the book uses **Sage** (<http://sagemath.org>), which is free open source mathematical software. The following is an example **Sage** session:

```

2 + 2

```

```

4

```

```

k.<a> = NumberField(x^2 + 1); k

```

```

Number Field in a with defining polynomial x^2 + 1

```

1.2 What is algebraic number theory?

A number field K is a finite degree extension of the rational numbers \mathbb{Q} . The primitive element theorem from field theory asserts that every such extension can be represented as the set of all polynomials of degree less than $d = [K : \mathbb{Q}] = \dim_{\mathbb{Q}} K$ in a single root α of some polynomial with coefficients in \mathbb{Q} :

$$K = \mathbb{Q}(\alpha) = \left\{ \sum_{n=0}^{d-1} a_n \alpha^n : a_n \in \mathbb{Q} \right\}.$$

Note that $\mathbb{Q}(\alpha)$ is non-canonically isomorphic to $\mathbb{Q}[x]/(f)$, where f is the minimal polynomial of α over \mathbb{Q} . The homomorphism $\mathbb{Q}[x] \rightarrow \mathbb{Q}(\alpha)$ that sends x to α has kernel (f) , hence it induces an isomorphism between $\mathbb{Q}[x]/(f)$ and $\mathbb{Q}(\alpha)$. It is not canonical, since $\mathbb{Q}(\alpha)$ could have nontrivial automorphisms. For example, if $\alpha = \sqrt{2}$, then $\mathbb{Q}(\sqrt{2})$ is isomorphic as a field to $\mathbb{Q}(-\sqrt{2})$ via $\sqrt{2} \mapsto -\sqrt{2}$. There are two isomorphisms $\mathbb{Q}[x]/(x^2 - 2) \rightarrow \mathbb{Q}(\sqrt{2})$.

Algebraic number theory is the study of number fields, their rings of integers, and related objects (e.g., function fields, elliptic curves, etc.). To gain a deeper understanding of these concepts, one uses techniques from (mostly commutative) algebra and finite group theory. The main objects that we study in this book are number fields, rings of integers of number fields, unit groups, ideal class groups, norms, traces, discriminants, prime ideals, Hilbert and other class fields and associated reciprocity laws, zeta and L -functions, and algorithms for computing with each of the above.

1.2.1 Topics in this book

These are some of the main topics that are discussed in this book:

- Rings of integers of number fields
- Unique factorization of nonzero ideals in Dedekind domains

- Structure of the group of units of the ring of integers
- Fractional ideals and class groups
- Decomposition and inertia groups, Frobenius elements
- Ramification
- Discriminant and different
- Quadratic and biquadratic fields
- Cyclotomic fields (and applications)
- How to use **Sage** to compute with many of the above objects

We will also touch on elliptic curves and L -functions. However we will not do anything nontrivial with these subjects.

1.3 Some applications of algebraic number theory

The following examples illustrate some of the power, depth, and importance of algebraic number theory.

Integer factorization: The number field sieve is the asymptotically fastest known algorithm for factoring general large integers (that don't have too special of a form). On December 12, 2009, the number field sieve was used to factor the RSA-768 challenge, which is a 232 digit number that is a product of two primes:

```
rsa768 = 12301866845301177551304949583849627207728535695\
95334792197322452151726400507263657518745202199786469389\
95647494277406384592519255732630345373154826850791702612\
21429134616704292143116022212404792747377940806653514195\
97459856902143413
n = 3347807169895689878604416984821269081770479498371376\
85689124313889828837938780022876147116525317430877378144\
67999489
m = 3674604366679959042824463379962795263227915816434308\
76426760322838157396665112792333734171433968102700927987\
36308917
n*m == rsa768
```

True

This record integer factorization cracked a certain 768-bit public key cryptosystem (see [KAF⁺10]), thus establishing a lower bound on one's choice of key size:

```

\ $ man ssh-keygen    # in ubuntu-12.04
...
    -b bits
69                      Specifies the number of bits in the key to
                        create. For RSA keys, the minimum size is
                        768 bits ...

```

70 **Primality testing:** Agrawal and his students Saxena and Kayal found in
71 2002 the first ever deterministic polynomial-time (in the number of
72 digits) primality test [AKS04]. Their methods involve arithmetic in
73 quotients of $(\mathbb{Z}/n\mathbb{Z})[x]$, which are best understood in the context of
74 algebraic number theory.

75 **Deeper point of view:** Some questions in number theory are best viewed
76 from the point of view of algebraic number theory such as:

- 77 • Pell's Equation $x^2 - dy^2 = 1$ can be reinterpreted in terms of units
78 in real quadratic fields, which leads to a study of unit groups of
79 number fields.
- 80 • Integer factorization is a special case of factoring nonzero ideals
81 in rings of integers of number fields.
- 82 • The Riemann hypothesis about the zeros of $\zeta(s)$ generalizes to
83 zeta functions of number fields.
- 84 • Reinterpreting Gauss's quadratic reciprocity law in terms of the
85 arithmetic of cyclotomic fields $\mathbb{Q}(e^{2\pi i/n})$ leads to class field theory,
86 which in turn leads to the Langlands program.

87 **Fermat's Last Theorem:** This classical theorem says $x^n + y^n = z^n$ has no
88 solutions with x, y, z, n all positive integers and $n \geq 3$. Wiles's proof of
89 Fermat's Last Theorem uses methods from algebraic number theory
90 extensively, in addition to many other deep techniques. Attempts
91 to prove Fermat's Last Theorem long ago were hugely influential in
92 the development of algebraic number theory by Dedekind, Hilbert,
93 Kummer, Kronecker, and others.

94 **Arithmetic geometry:** This is a huge field that studies solutions to poly-
95 nomial equations that lie in arithmetically interesting rings, such as
96 the integers or number fields. A major triumph of arithmetic geometry
97 is Faltings's proof of Mordell's Conjecture.

98 **Theorem 1.3.1** (Faltings). *Let X be a nonsingular plane algebraic*
99 *curve over a number field K . Assume that the manifold $X(\mathbb{C})$ of com-*
100 *plex solutions to X has genus at least 2 (i.e., $X(\mathbb{C})$ is topologically a*

101 *donut with at least two holes). Then the set $X(K)$ of points on X with*
102 *coordinates in K is finite.*

103 For example, Theorem 1.3.1 implies that for any $n \geq 4$ and any number
104 field K , there are only finitely many solutions in \bar{K} to $x^n + y^n = 1$.

105 A major open problem in arithmetic geometry is the *Birch and Swinnerton-*
106 *Dyer conjecture*. An *elliptic curve* E is an algebraic curve with at least
107 one point with coordinates in K such that the set of complex points
108 $E(\mathbb{C})$ is a topological torus (i.e., $E(\mathbb{C})$ is topologically a donut with
109 one hole). The Birch and Swinnerton-Dyer conjecture gives a criterion
110 for whether or not $E(K)$ is infinite in terms of analytic properties of
111 the L -function $L(E, s)$. See [http://www.claymath.org/millennium/](http://www.claymath.org/millennium/Birch_and_Swinnerton-Dyer_Conjecture/)
112 [Birch_and_Swinnerton-Dyer_Conjecture/](http://www.claymath.org/millennium/Birch_and_Swinnerton-Dyer_Conjecture/).

Chapter 2

Basic Commutative Algebra

The commutative algebra in this chapter provides a foundation for understanding the more refined number-theoretic structures associated to number fields.

First we prove the structure theorem for finitely generated abelian groups. Then we establish the standard properties of Noetherian rings and modules, including a proof of the Hilbert basis theorem. We also observe that finitely generated abelian groups are Noetherian \mathbb{Z} -modules. After establishing properties of Noetherian rings, we consider rings of algebraic integers and discuss some of their properties.

2.1 Finitely Generated Abelian Groups

Finitely generated abelian groups arise all over algebraic number theory. For example, they will appear in this book as class groups, unit groups, and the underlying additive groups of rings of integers, and as Mordell-Weil groups of elliptic curves.

In this section, we prove the structure theorem for finitely generated abelian groups, since it will be crucial for much of what we will do later.

Let $\mathbb{Z} = \{0, \pm 1, \pm 2, \dots\}$ denote the ring of (rational) integers, and for each positive integer n , let $\mathbb{Z}/n\mathbb{Z}$ denote the ring of integers modulo n , which is a cyclic abelian group of order n under addition.

Definition 2.1.1 (Finitely Generated). A group G is *finitely generated* if there exists $g_1, \dots, g_n \in G$ such that every element of G can be expressed as a finite product (or sum, if we write G additively) of positive or negative powers of the g_i .

139 For example, the group \mathbb{Z} is finitely generated, since it is generated by 1.

Theorem 2.1.2 (Structure Theorem for Finitely Generated Abelian Groups).
 Let G be a finitely generated abelian group. Then there is an isomorphism

$$G \approx (\mathbb{Z}/n_1\mathbb{Z}) \oplus (\mathbb{Z}/n_2\mathbb{Z}) \oplus \cdots \oplus (\mathbb{Z}/n_s\mathbb{Z}) \oplus \mathbb{Z}^r,$$

140 where $r, s \geq 0$, $n_i > 1$ for all i , and $n_1 \mid n_2 \mid \cdots \mid n_s$. Furthermore, the n_i
 141 and r are uniquely determined by G .

142 **Exercise 2.1.3.** Quick! Guess how many abelian groups there are of order
 143 less than 12. Use Theorem 2.1.2 to classify all abelian groups of order less
 144 than 12. How many do you think there are? How many are there?

145 We will prove the theorem as follows. We first remark that any subgroup
 146 of a finitely generated free abelian group is finitely generated. Then we see
 147 how to represent finitely generated abelian groups as quotients of finite rank
 148 free abelian groups, and how to reinterpret such a presentation in terms of
 149 matrices over the integers. Next we describe how to use row and column
 150 operations over the integers to show that every matrix over the integers is
 151 equivalent to one in a canonical diagonal form, called the Smith normal form.
 152 We obtain a proof of the theorem by reinterpreting the in terms of groups.
 153 Finally, we observe that the representation in the theorem is necessarily
 154 unique.

155 **Proposition 2.1.4.** If H is a subgroup of a finitely generated abelian group
 156 G , then H is finitely generated.

157 The key reason that this is true is that G is a finitely generated module
 158 over the principal ideal domain \mathbb{Z} . We defer the proof of Proposition 2.1.4 to
 159 Section 2.2, where we will give a complete proof of a beautiful generalization
 160 in the context of Noetherian rings (the Hilbert basis theorem).

161 **Corollary 2.1.5.** Suppose G is a finitely generated abelian group. Then
 162 there are finitely generated free abelian groups F_1 and F_2 and there is a
 163 homomorphism $\psi : F_2 \rightarrow F_1$ such that $G \approx F_1/\psi(F_2)$.

164 *Proof.* Let x_1, \dots, x_m be generators for G . Let $F_1 = \mathbb{Z}^m$ and let $\varphi : F_1 \rightarrow G$
 165 be the homomorphism that sends the i th generator $(0, 0, \dots, 1, \dots, 0)$ of \mathbb{Z}^m
 166 to x_i . Then φ is surjective, and by Proposition 2.1.4 the kernel $\ker(\varphi)$ of
 167 φ is a finitely generated abelian group. Suppose there are n generators for
 168 $\ker(\varphi)$, let $F_2 = \mathbb{Z}^n$ and fix a surjective homomorphism $\psi : F_2 \rightarrow \ker(\varphi)$.
 169 Then $F_1/\psi(F_2)$ is isomorphic to G . \square

An *sequence* of homomorphisms of abelian groups

$$H \xrightarrow{f} G \xrightarrow{g} K$$

is exact if $\text{im}(f) = \ker(g)$. For longer sequences, exactness means every three consecutive terms with two arrows are exact. Given a finitely generated abelian group G , Corollary 2.1.5 provides an exact sequence

$$F_2 \xrightarrow{\psi} F_1 \rightarrow G \rightarrow 0.$$

Suppose G is a nonzero finitely generated abelian group. By the corollary, there are free abelian groups F_1 and F_2 and there is a homomorphism $\psi : F_2 \rightarrow F_1$ such that $G \approx F_1/\psi(F_2)$. Upon choosing a basis for F_1 and F_2 , we obtain isomorphisms $F_1 \approx \mathbb{Z}^n$ and $F_2 \approx \mathbb{Z}^m$ for integers n and m . Just as in linear algebra, we view $\psi : F_2 \rightarrow F_1$ as being given by left multiplication by the $n \times m$ matrix A whose columns are the images of the generators of F_2 in \mathbb{Z}^n . We visualize this as follows:

$$\mathbb{Z}^m \xrightarrow{A} \mathbb{Z}^n \rightarrow G \rightarrow 0$$

170 The *cokernel* of the homomorphism defined by A is the quotient of \mathbb{Z}^n
 171 by the image of A (i.e., the \mathbb{Z} -span of the columns of A), and this cokernel
 172 is isomorphic to G .

173 The following proposition implies that we may choose a bases for F_1 and
 174 F_2 such that the matrix of A only has nonzero entries along the diagonal,
 175 so that the structure of the cokernel of A is trivial to understand.

176 **Proposition 2.1.6** (Smith normal form). *Suppose A is an $n \times m$ integer*
 177 *matrix. Then there exist invertible integer matrices P and Q such that*
 178 *$A' = PAQ$ only has nonzero entries along the diagonal, and these entries*
 179 *are $n_1, n_2, \dots, n_s, 0, \dots, 0$, where $s \geq 0$, $n_i \geq 1$ for all i , and $n_1 \mid n_2 \mid \dots \mid n_s$.*

180

INSERT EXAMPLE

181 **Remark 2.1.7.** Note that the matrices P and Q are invertible as integer
 182 matrices, so $\det(P)$ and $\det(Q)$ are ± 1 . In particular $\det A' = \pm \det A$.

183 **Definition 2.1.8.** The matrix A' in Proposition 2.1.6 is called the *Smith*
 184 *normal form* of A .

We will see in the proof of Theorem 2.1.2 that A' is uniquely determined by A . An example of a matrix in Smith normal form is

$$A = \begin{pmatrix} 2 & 0 & 0 & 0 \\ 0 & 6 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}.$$

185 *Proof of Proposition 2.1.6.* The matrix P will be a product of matrices that
 186 define elementary row operations and Q will be a product corresponding to
 187 elementary column operations. The elementary row and column operations
 188 over \mathbb{Z} are as follows:

189 **Add multiple:** Add an integer multiple of one row to another (or a multi-
 190 ple of one column to another).

191 **Swap:** Interchange two rows or two columns.

192 **Rescale:** Multiply a row by -1 .

193 Each of these operations is given by left or right multiplying by an invertible
 194 matrix E with integer entries, where E is the result of applying the given
 195 operation to the identity matrix, and E is invertible because each operation
 196 can be reversed using another row or column operation over the integers.

197 To see that the proposition must be true, assume $A \neq 0$ and perform
 198 the following steps (compare [Art91, pg. 459]):

- 199 1. By permuting rows and columns, move a nonzero entry of A with
 200 smallest absolute value to the upper left corner of A . Now “attempt”
 201 (as explained in detail below) to make all other entries in the first row
 202 and column 0 by adding multiples of the top row or first column to
 203 other rows or columns, as follows:

204 Suppose a_{i1} is a nonzero entry in the first column, with
 205 $i > 1$. Using the division algorithm, write $a_{i1} = a_{11}q + r$,
 206 with $0 \leq r < a_{11}$. Now add $-q$ times the first row to the
 207 i th row. If $r > 0$, then go to step 1 (so that an entry with
 208 absolute value at most r is the upper left corner).

209 If at any point this operation produces a nonzero entry in the matrix
 210 with absolute value smaller than $|a_{11}|$, start the process over by per-
 211 muting rows and columns to move that entry to the upper left corner
 212 of A . Since the integers $|a_{11}|$ are a decreasing sequence of positive
 213 integers, we will not have to move an entry to the upper left corner
 214 infinitely often, so when this step is done the upper left entry of the
 215 matrix is nonzero, and all entries in the first row and column are 0.

- 216 2. We may now assume that a_{11} is the only nonzero entry in the first
 217 row and column. If some entry a_{ij} of A is not divisible by a_{11} , add
 218 the column of A containing a_{ij} to the first column, thus producing an
 219 entry in the first column that is nonzero. When we perform step 2,

220 the remainder r will be greater than 0. Permuting rows and columns
 221 results in a smaller $|a_{11}|$. Since $|a_{11}|$ can only shrink finitely many
 222 times, eventually we will get to a point where every a_{ij} is divisible by
 223 a_{11} . If a_{11} is negative, multiple the first row by -1 .

224 After performing the above operations, the first row and column of A are
 225 zero except for a_{11} which is positive and divides all other entries of A . We
 226 repeat the above steps for the matrix B obtained from A by deleting the first
 227 row and column. The upper left entry of the resulting matrix will be divisible
 228 by a_{11} , since every entry of B is. Repeating the argument inductively proves
 229 the proposition. \square

Example 2.1.9. The matrix $\begin{pmatrix} -2 & 2 \\ -3 & 4 \end{pmatrix}$ has Smith normal form $\begin{pmatrix} 1 & 0 \\ 0 & 2 \end{pmatrix}$, and
 the matrix $\begin{pmatrix} 1 & 4 & 9 \\ 16 & 25 & 36 \\ 49 & 64 & 81 \end{pmatrix}$ has Smith normal form $\begin{pmatrix} 1 & 0 & 0 \\ 0 & 3 & 0 \\ 0 & 0 & 72 \end{pmatrix}$. As a double
 check, note that the determinants of a matrix and its Smith normal form
 match, up to sign. This is because

$$\det(PAQ) = \det(P) \det(A) \det(Q) = \pm \det(A).$$

230 We compute each of the above Smith forms using **Sage**, along with the
 231 corresponding transformation matrices. To do this we use the **Sage** com-
 232 mand **matrix**, which takes as input the base ring, the number of rows, and
 233 the entries. The output of **matrix** is a matrix object which has the method
 234 **smith_form**.

235 First the 2×2 matrix.

```
A = matrix(ZZ, 2, [-2,2, -3,4])
S, P, Q = A.smith_form(); S
```

```
[1 0]
[0 2]
```

```
P*A*Q
```

```
[1 0]
[0 2]
```

```
P
```

```
[0 1]
[1 0]
```

```
Q
```

```
[1 -4]
[1 -3]
```

236

237 Next the 3×3 matrix.

```
A = matrix(ZZ, 3, [1,4,9, 16,25,36, 49,64,81])
S, P, Q = A.smith_form(); S
```

```
[ 1  0  0]
[ 0  3  0]
[ 0  0 72]
```

```
P*A*Q
```

```
[ 1  0  0]
[ 0  3  0]
[ 0  0 72]
```

```
P
```

```
[ 0  0  1]
[ 0  1 -1]
[ 1 -20 -17]
```

```
Q
```

```
[ 47  74  93]
[-79 -125 -156]
[ 34  54  67]
```

238

239 For one more example, we compute the Smith form of a 3×3 matrix of
 240 rank 2:

```
m = matrix(ZZ, 3, [2..10]); m
```

```
[ 2  3  4]
[ 5  6  7]
[ 8  9 10]
```

```
m.smith_form()[0]
```

```
[1 0 0]
[0 3 0]
[0 0 0]
```

242 *Proof of Theorem 2.1.2.* Suppose G is a finitely generated abelian group,
 243 which we may assume is nonzero. As in the paragraph before Proposition
 244 2.1.6, we use Corollary 2.1.5 to write G as the cokernel of an $n \times$
 245 m integer matrix A . By Proposition 2.1.6 there are isomorphisms $Q :$
 246 $\mathbb{Z}^m \rightarrow \mathbb{Z}^m$ and $P : \mathbb{Z}^n \rightarrow \mathbb{Z}^n$ such that $A' = PAQ$ has diagonal entries
 247 $n_1, n_2, \dots, n_s, 0, \dots, 0$, where $n_1 > 1$ and $n_1 \mid n_2 \mid \dots \mid n_s$. Then G is
 248 isomorphic to the cokernel of the diagonal matrix A' , so

$$G \cong (\mathbb{Z}/n_1\mathbb{Z}) \oplus (\mathbb{Z}/n_2\mathbb{Z}) \oplus \dots \oplus (\mathbb{Z}/n_s\mathbb{Z}) \oplus \mathbb{Z}^r, \quad (2.1)$$

249 as claimed. The n_i are determined by G , because n_i is the smallest positive
 250 integer n such that nG requires at most $s + r - i$ generators. We see from
 251 the representation (2.1) of G as a product that n_i has this property and that
 252 no smaller positive integer does. \square

253 **Exercise 2.1.10.** Recall Smith normal form defined in Proposition 2.1.6.
 254 With only minor modifications, then the proposition and proof will work
 255 over any principle ideal domain. Find and apply these modifications then

256 find the Smith normal form of the matrix $\begin{pmatrix} 1 & 2 & 3 \\ 0 & 1+i & 2 \\ 0 & 1 & 5 \end{pmatrix}$.

257 [*Hint:* You can use **Sage** to verify your answer. However, you will need
 258 to make explicitly construct the Gaussian integers in order to input the
 259 matrix. You can do this by the following code.]

```
K.<i> = QuadraticField(-1)
R = K.maximal_order()
M = matrix(R, 3, [1,2,3,0,1+i,2,0,1,5]); show(M)
#show(M.smith_form()[0]) #uncomment for the answer
```

261 **Exercise 2.1.11.** Let $A = \begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{pmatrix}$.

262 (a) Find the Smith normal form of A .

263 (b) Prove that the cokernel of the map $\mathbb{Z}^3 \rightarrow \mathbb{Z}^3$ given by multiplication
264 by A is isomorphic to $\mathbb{Z}/3\mathbb{Z} \oplus \mathbb{Z}$.

265 2.2 Noetherian Rings and Modules

266 A module M over a commutative ring R with unit element is much like
267 a vector space, but with more subtle structure. In this book, most of the
268 modules we encounter will be noetherian, which is a generalization of the
269 “finite dimensional” property of vector spaces. This section is about prop-
270 erties of noetherian modules (and rings), which are crucial to much of this
271 book. We thus give complete proofs of these properties, so you will have a
272 solid foundation on which to learn algebraic number theory.

273 We first define noetherian rings and modules, then introduce several
274 equivalent characterizations of them. We prove that when the base ring is
275 noetherian, a module is finitely generated if and only if it is noetherian.
276 Next we define short exact sequences, and prove that the middle module
277 in a sequence is noetherian if and only if the first and last modules are
278 noetherian. Finally, we prove the Hilbert basis theorem, which asserts that
279 adjoining finitely many elements to a noetherian ring results in a noetherian
280 ring.

281 Let R be a commutative ring with unity. An R -module is an additive
282 abelian group M equipped with a map $R \times M \rightarrow M$ such that for all $r, r' \in R$
283 and all $m, m' \in M$ we have $(rr')m = r(r'm)$, $(r + r')m = rm + r'm$,
284 $r(m + m') = rm + rm'$, and $1m = m$. A *submodule* of M is a subgroup of
285 M that is preserved by the action of R . For example, R is a module over
286 itself, and any ideal I in R is an R -submodule of R .

287 *Example 2.2.1.* Abelian groups are the same as \mathbb{Z} -modules, and vector spaces
288 over a field K are the same as K -modules.

289 An R -module M is finitely generated if there are elements $m_1, \dots, m_n \in$
290 M such that every element of M is an R -linear combination of the m_i . The
291 noetherian property is stronger than just being finitely generated:

292 **Definition 2.2.2** (Noetherian). An R -module M is *noetherian* if every sub-
293 module of M is finitely generated. A ring R is *noetherian* if R is noetherian
294 as a module over itself, i.e., if every ideal of R is finitely generated.

Any submodule M' of a noetherian module M is also noetherian. Indeed, if every submodule of M is finitely generated then so is every submodule of M' , since submodules of M' are also submodules of M .

Example 2.2.3. Let $R = M = \mathbb{Q}[x_1, x_2, \dots]$ be a polynomial ring over \mathbb{Q} in infinitely many indeterminants x_i . Then M is finitely generated as an R -module (!), since it is generated by 1. Consider the submodule $I = (x_1, x_2, \dots)$ of polynomials with 0 constant term, and suppose it is generated by polynomials f_1, \dots, f_n . Let x_i be an indeterminant that does not appear in any f_j , and suppose there are $h_k \in R$ such that $\sum_{k=1}^n h_k f_k = x_i$. Setting $x_i = 1$ and all other $x_j = 0$ on both sides of this equation and using that the f_k all vanish (they have 0 constant term), yields $0 = 1$, a contradiction. We conclude that the ideal I is not finitely generated, hence M is not a noetherian R -module, despite being finitely generated.

Definition 2.2.4 (Ascending chain condition). An R -module M satisfies the *ascending chain condition* if every sequence $M_1 \subset M_2 \subset M_3 \subset \dots$ of submodules of M eventually stabilizes, i.e., there is some n such that $M_n = M_{n+1} = M_{n+2} = \dots$.

We will use the notion of maximal element below. If \mathcal{X} is a set of subsets of a set S , ordered by inclusion, then a *maximal element* $A \in \mathcal{X}$ is a set such that no superset of A is contained in \mathcal{X} . Note that \mathcal{X} may contain many different maximal elements.

Proposition 2.2.5. *If M is an R -module, then the following are equivalent:*

1. M is noetherian,
2. M satisfies the ascending chain condition, and
3. Every nonempty set of submodules of M contains at least one maximal element.

Proof.

(1 \Rightarrow 2) : Suppose $M_1 \subset M_2 \subset \dots$ is a sequence of submodules of M . Then $M_\infty = \cup_{n=1}^\infty M_n$ is a submodule of M . Since M is noetherian and M_∞ is a submodule of M , there is a finite set a_1, \dots, a_m of generators for M_∞ . Each a_i must be contained in some M_j , so there is an n such that $a_1, \dots, a_m \in M_n$. But then $M_k = M_n$ for all $k \geq n$, which proves that the chain of M_i stabilizes, so the ascending chain condition holds for M .

(2 \Rightarrow 3) : Suppose 3 were false, so there exists a nonempty set S of submodules of M that does not contain a maximal element. We will use S to construct an infinite ascending chain of submodules of M that does not stabilize. Note that S is infinite, otherwise it would contain a maximal element. Let M_1 be any element of S . Then there is an M_2 in S that strictly contains M_1 , otherwise S would contain the maximal element M_1 . Continuing inductively in this way we find an M_3 in S that properly contains M_2 , etc., and we produce an infinite ascending chain of submodules of M , which contradicts the ascending chain condition.

(3 \Rightarrow 1) : Suppose 1 is false, so there is a submodule M' of M that is not finitely generated. We will show that the set S of all finitely generated submodules of M' does not have a maximal element, which will be a contradiction. Suppose S does have a maximal element L . Since L is finitely generated and $L \subset M'$, and M' is not finitely generated, there is an $a \in M'$ such that $a \notin L$. Then $L' = L + Ra$ is an element of S that strictly contains the presumed maximal element L , a contradiction.

□

Definition 2.2.6. A *homomorphism* of R -modules $\varphi : M \rightarrow N$ is an abelian group homomorphism such that for any $r \in R$ and $m \in M$ we have $\varphi(rm) = r\varphi(m)$. A sequence

$$L \xrightarrow{f} M \xrightarrow{g} N,$$

where f and g are homomorphisms of R -modules, is *exact* if $\text{im}(f) = \ker(g)$. A *short exact sequence* of R -modules is a sequence

$$0 \rightarrow L \xrightarrow{f} M \xrightarrow{g} N \rightarrow 0$$

that is exact at each point, i.e., f is injective, g is surjective, and $\text{im}(f) = \ker(g)$.

Example 2.2.7. The sequence

$$0 \rightarrow \mathbb{Z} \xrightarrow{2} \mathbb{Z} \rightarrow \mathbb{Z}/2\mathbb{Z} \rightarrow 0$$

is an exact sequence, where the first map sends 1 to 2, and the second is the natural quotient map.

Lemma 2.2.8. *If*

$$0 \rightarrow L \xrightarrow{f} M \xrightarrow{g} N \rightarrow 0$$

351 *is a short exact sequence of R -modules, then M is noetherian if and only if*
 352 *both L and N are noetherian.*

353 *Proof.* First suppose that M is noetherian. Then L is a submodule of M ,
 354 so L is noetherian. Let N' be a submodule of N ; then the inverse image of
 355 N' in M is a submodule of M , so it is finitely generated, hence its image N'
 356 is also finitely generated. Thus N is noetherian as well.

357 Next assume nothing about M , but suppose that both L and N are
 358 noetherian. Suppose M' is a submodule of M ; then $M_0 = f(L) \cap M'$ is
 359 isomorphic to a submodule of the noetherian module L , so M_0 is generated
 360 by finitely many elements a_1, \dots, a_n . The quotient M'/M_0 is isomorphic
 361 (via g) to a submodule of the noetherian module N , so M'/M_0 is generated
 362 by finitely many elements b_1, \dots, b_m . For each $i \leq m$, let c_i be a lift of b_i to
 363 M' , modulo M_0 . Then the elements $a_1, \dots, a_n, c_1, \dots, c_m$ generate M' , for
 364 if $x \in M'$, then there is some element $y \in M_0$ such that $x - y$ is an R -linear
 365 combination of the c_i , and y is an R -linear combination of the a_i . \square

366 **Proposition 2.2.9.** *Suppose R is a noetherian ring. Then an R -module M*
 367 *is noetherian if and only if it is finitely generated.*

368 *Proof.* If M is noetherian then every submodule of M is finitely generated
 369 so M itself is finitely generated. Conversely, suppose M is finitely generated,
 370 say by elements a_1, \dots, a_n . Then there is a surjective homomorphism from
 371 $R^n = R \oplus \dots \oplus R$ to M that sends $(0, \dots, 0, 1, 0, \dots, 0)$ (1 in the i th factor)
 372 to a_i . Using Lemma 2.2.8 and exact sequences of R -modules such as $0 \rightarrow$
 373 $R \rightarrow R \oplus R \rightarrow R \rightarrow 0$, we see inductively that R^n is noetherian. Again by
 374 Lemma 2.2.8, homomorphic images of noetherian modules are noetherian,
 375 so M is noetherian. \square

376 **Lemma 2.2.10.** *Suppose $\varphi : R \rightarrow S$ is a surjective homomorphism of rings*
 377 *and R is noetherian. Then S is noetherian.*

Proof. The kernel of φ is an ideal I in R , and we have an exact sequence

$$0 \rightarrow I \rightarrow R \rightarrow S \rightarrow 0$$

378 with R noetherian. This is an exact sequence of R -modules, where S has
 379 the R -module structure induced from φ (if $r \in R$ and $s \in S$, then we define
 380 $rs = \varphi(r)s$). By Lemma 2.2.8, it follows that S is a noetherian R -modules.

381 Suppose J is an ideal of S . Since J is an R -submodule of S , if we view J
 382 as an R -module, then J is finitely generated. Since R acts on J through S ,
 383 the R -generators of J are also S -generators of J , so J is finitely generated
 384 as an ideal. Thus S is noetherian. \square

385 **Theorem 2.2.11** (Hilbert Basis Theorem). *If R is a noetherian ring and*
 386 *S is finitely generated as a ring over R , then S is noetherian. In particular,*
 387 *for any n the polynomial ring $R[x_1, \dots, x_n]$ and any of its quotients are*
 388 *noetherian.*

389 *Proof.* Assume first that we have already shown that for any n the polyno-
 390 mial ring $R[x_1, \dots, x_n]$ is noetherian. Suppose S is finitely generated as a
 391 ring over R , so there are generators s_1, \dots, s_n for S . Then the map $x_i \mapsto s_i$
 392 extends uniquely to a surjective homomorphism $\pi : R[x_1, \dots, x_n] \twoheadrightarrow S$, and
 393 Lemma 2.2.10 implies that S is noetherian.

394 The rings $R[x_1, \dots, x_n]$ and $(R[x_1, \dots, x_{n-1}])[x_n]$ are isomorphic, so it
 395 suffices to prove that if R is noetherian then $R[x]$ is also noetherian. (Our
 396 proof follows [Art91, §12.5].) Thus suppose I is an ideal of $R[x]$ and that R
 397 is noetherian. We will show that I is finitely generated.

398 Let A be the set of leading coefficients of polynomials in I . (The leading
 399 coefficient of a polynomial is the coefficient of the highest degree monomial,
 400 or 0 if the polynomial is 0; thus $3x^7 + 5x^2 - 4$ has leading coefficient 3.)
 401 We will first show that A is an ideal of R . Suppose $a, b \in A$ are nonzero
 402 with $a + b \neq 0$. Then there are polynomials f and g in I with leading
 403 coefficients a and b . If $\deg(f) \leq \deg(g)$, then $a + b$ is the leading coefficient
 404 of $x^{\deg(g)-\deg(f)}f + g$, so $a + b \in A$; the argument when $\deg(f) > \deg(g)$ is
 405 analogous. Suppose $r \in R$ and $a \in A$ with $ra \neq 0$. Then ra is the leading
 406 coefficient of rf , so $ra \in A$. Thus A is an ideal in R .

407 Since R is noetherian and A is an ideal of R , there exist nonzero $a_1, \dots, a_n \in$
 408 A that generate A as an ideal. Since A is the set of leading coefficients of
 409 elements of I , and the a_j are in A , we can choose for each $j \leq n$ an element
 410 $f_j \in I$ with leading coefficient a_j . By multiplying the f_j by some power of x ,
 411 we may assume that the f_j all have the same degree $d \geq 1$.

412 Let $S_{<d}$ be the set of elements of I that have degree strictly less than d .
 413 This set is closed under addition and under multiplication by elements of R ,
 414 so $S_{<d}$ is a module over R . The module $S_{<d}$ is the submodule of the R -
 415 module of polynomials of degree less than n , which is noetherian by Propo-
 416 sition 2.2.9 because it is generated by $1, x, \dots, x^{n-1}$. Thus $S_{<d}$ is finitely
 417 generated, and we may choose generators h_1, \dots, h_m for $S_{<d}$.

418 We finish by proving using induction on the degree that every $g \in I$ is an
 419 $R[x]$ -linear combination of $f_1, \dots, f_n, h_1, \dots, h_m$. If $g \in I$ has degree 0, then

420 $g \in S_{<d}$, since $d \geq 1$, so g is a linear combination of h_1, \dots, h_m . Next suppose
 421 $g \in I$ has degree e , and that we have proven the statement for all elements
 422 of I of degree $< e$. If $e \leq d$, then $g \in S_{<d}$, so g is in the $R[x]$ -ideal generated
 423 by h_1, \dots, h_m . Next suppose that $e \geq d$. Then the leading coefficient b
 424 of g lies in the ideal A of leading coefficients of elements of I , so there exist
 425 $r_i \in R$ such that $b = r_1 a_1 + \dots + r_n a_n$. Since f_i has leading coefficient a_i , the
 426 difference $g - x^{e-d} r_i f_i$ has degree less than the degree e of g . By induction
 427 $g - x^{e-d} r_i f_i$ is an $R[x]$ linear combination of $f_1, \dots, f_n, h_1, \dots, h_m$, so g is
 428 also an $R[x]$ linear combination of $f_1, \dots, f_n, h_1, \dots, h_m$. Since each f_i and
 429 h_j lies in I , it follows that I is generated by $f_1, \dots, f_n, h_1, \dots, h_m$, so I is
 430 finitely generated, as required. \square

431 2.2.1 The Ring \mathbb{Z} is Noetherian

432 The ring \mathbb{Z} is noetherian since every ideal of \mathbb{Z} is generated by one element.

433 **Proposition 2.2.12.** *Every ideal of the ring \mathbb{Z} is principal.*

434 *Proof.* Suppose I is a nonzero ideal in \mathbb{Z} . Let d be the least positive element
 435 of I . Suppose that $a \in I$ is any nonzero element of I . Using the division
 436 algorithm, we write $a = dq + r$, where q is an integer and $0 \leq r < d$. We have
 437 $r = a - dq \in I$ and $r < d$, so our assumption that d is minimal implies that
 438 $r = 0$, hence $a = dq$ is in the ideal generated by d . Thus I is the principal
 439 ideal generated by d . \square

440 *Example 2.2.13.* Let $I = (12, 18)$ be the ideal of \mathbb{Z} generated by 12 and 18.
 441 If $n = 12a + 18b \in I$, with $a, b \in \mathbb{Z}$, then $6 \mid n$, since $6 \mid 12$ and $6 \mid 18$. Also,
 442 $6 = 18 - 12 \in I$, so $I = (6)$.

443 The ring \mathbb{Z} in Sage is ZZ, which is Noetherian.

```
ZZ.is_noetherian()
```

444 **True**

445 We create the ideal I in Sage as follows, and note that it is principal:

```
I = ideal(12,18); I
```

Principal ideal (6) of Integer Ring

```
I.is_principal()
```

True

447 We could also create I as follows:

```
448      ZZ.ideal(12,18)
```

```
448      Principal ideal (6) of Integer Ring
```

449 Propositions 2.2.9 and 2.2.12 together imply that any finitely generated
 450 abelian group is noetherian. This means that subgroups of finitely generated
 451 abelian groups are finitely generated, which provides the missing step in our
 452 proof of the structure theorem for finitely generated abelian groups.

453 **Exercise 2.2.14.** There is another way to show every principle ideal domain
 454 (for example \mathbb{Z}) is noetherian (contrast to the proof in Section 2.2.1). Let
 455 R be a PID and (a) an arbitrary ideal. Use the facts that $(b) \supseteq (a)$ if and
 456 only if $b \mid a$ and R is a UFD to show that ascending chain of ideals starting
 457 with (a) must stabilize.

458 2.3 Rings of Algebraic Integers

459 In this section we introduce the central objects of this book, which are the
 460 rings of algebraic integers. These are noetherian rings with an enormous
 461 amount of structure. We also introduce a function field analogue of these
 462 rings.

463 An *algebraic number* is a root of some nonzero polynomial $f(x) \in \mathbb{Q}[x]$.
 464 For example, $\sqrt{2}$ and $\sqrt{5}$ are both algebraic numbers, being roots of $x^2 - 2$
 465 and $x^2 - 5$, respectively. But is $\sqrt{2} + \sqrt{5}$ necessarily the root of some
 466 polynomial in $\mathbb{Q}[x]$? This isn't quite so obvious.

467 **Proposition 2.3.1.** *An element α of a field extension of \mathbb{Q} is an algebraic*
 468 *number if and only if the ring $\mathbb{Q}[\alpha]$ generated by α is finite dimensional as*
 469 *a \mathbb{Q} vector space.*

470 *Proof.* Suppose α is an algebraic number, so there is a nonzero polynomial
 471 $f(x) \in \mathbb{Q}[x]$, so that $f(\alpha) = 0$. The equation $f(\alpha) = 0$ implies that $\alpha^{\deg(f)}$
 472 can be written in terms of smaller powers of α , so $\mathbb{Q}[\alpha]$ is spanned by the
 473 finitely many numbers $1, \alpha, \dots, \alpha^{\deg(f)-1}$, hence finite dimensional. Con-
 474 versely, suppose $\mathbb{Q}[\alpha]$ is finite dimensional. Then for some $n \geq 1$, we have
 475 that α^n is in the \mathbb{Q} -vector space spanned by $1, \alpha, \dots, \alpha^{n-1}$. Thus α satisfies
 476 a polynomial $f(x) \in \mathbb{Q}[x]$ of degree n . \square

477 **Proposition 2.3.2.** *Suppose K is a field and $\alpha, \beta \in K$ are two algebraic*
 478 *numbers. Then $\alpha\beta$ and $\alpha + \beta$ are also algebraic numbers.*

479 *Proof.* Let $n = \dim_{\mathbb{Q}} \mathbb{Q}[\alpha]$ and $n = \dim_{\mathbb{Q}} \mathbb{Q}[\beta]$. The subring $\mathbb{Q}[\alpha, \beta] \subset K$ is
 480 a \mathbb{Q} -vector space that is spanned by the numbers $\alpha^i \beta^j$, where $0 \leq i < n$ and
 481 $0 \leq j < m$. Thus $\mathbb{Q}[\alpha, \beta]$ is finite dimensional, and since $\alpha + \beta$ and $\alpha\beta$ are
 482 both in $\mathbb{Q}[\alpha, \beta]$, we conclude by Proposition 2.3.1 that both are algebraic
 483 numbers. \square

484 Suppose C is a field extension of \mathbb{Q} such that every polynomial $f(x) \in$
 485 $\mathbb{Q}[x]$ factors completely in C . The algebraic closure $\overline{\mathbb{Q}}$ of \mathbb{Q} inside C is the
 486 field generated by all roots in C of polynomials in $\mathbb{Q}[x]$. The fundamental
 487 theorem of algebra tells us that $C = \mathbb{C}$ is one choice of field C as above.
 488 There are other fields C , e.g., constructed using p -adic numbers. One can
 489 show that any two choices of $\overline{\mathbb{Q}}$ are isomorphic; however, there will be *many*
 490 isomorphisms between them.

491 **Definition 2.3.3** (Algebraic Integer). An element $\alpha \in \overline{\mathbb{Q}}$ is an *algebraic*
 492 *integer* if it is a root of some monic polynomial with coefficients in \mathbb{Z} .

493 For example, $\sqrt{2}$ is an algebraic integer, since it is a root of the monic
 494 integral polynomial $x^2 - 2$. As we will see below, $1/2$ is not an algebraic
 495 integer.

496 The following two propositions are analogous to Propositions 2.3.1–2.3.2
 497 above, with the proofs replacing basic facts about vector spaces with facts
 498 we proved above about noetherian rings and modules.

499 **Proposition 2.3.4.** *An element $\alpha \in \overline{\mathbb{Q}}$ is an algebraic integer if and only*
 500 *if $\mathbb{Z}[\alpha]$ is finitely generated as a \mathbb{Z} -module.*

501 *Proof.* Suppose α is integral and let $f \in \mathbb{Z}[x]$ be a monic integral poly-
 502 nomial such that $f(\alpha) = 0$. Then, as a \mathbb{Z} -module, $\mathbb{Z}[\alpha]$ is generated by
 503 $1, \alpha, \alpha^2, \dots, \alpha^{d-1}$, where d is the degree of f . Conversely, suppose $\alpha \in \overline{\mathbb{Q}}$
 504 is such that $\mathbb{Z}[\alpha]$ is finitely generated as a module over \mathbb{Z} , say by elements
 505 $f_1(\alpha), \dots, f_n(\alpha)$. Let d be any integer bigger than the degrees of all f_i .
 506 Then there exist integers a_i such that $\alpha^d = \sum_{i=1}^n a_i f_i(\alpha)$, hence α satis-
 507 fies the monic polynomial $x^d - \sum_{i=1}^n a_i f_i(x) \in \mathbb{Z}[x]$, so α is an algebraic
 508 integer. \square

509 The proof of the following proposition uses repeatedly that any submod-
 510 ule of a finitely generated \mathbb{Z} -module is finitely generated, which uses that \mathbb{Z}
 511 is noetherian and that finitely generated modules over a noetherian ring are
 512 noetherian.

513 **Proposition 2.3.5.** *Suppose K is a field and $\alpha, \beta \in K$ are two algebraic*
 514 *integers. Then $\alpha\beta$ and $\alpha + \beta$ are also algebraic integers.*

515 *Proof.* Let m, n be the degrees of monic integral polynomials that have α, β
 516 as roots, respectively. Then we can write α^m in terms of smaller powers of
 517 α and likewise for β^n , so the elements $\alpha^i \beta^j$ for $0 \leq i < m$ and $0 \leq j < n$
 518 span the \mathbb{Z} -module $\mathbb{Z}[\alpha, \beta]$. Since $\mathbb{Z}[\alpha + \beta]$ is a submodule of the finitely-
 519 generated \mathbb{Z} -module $\mathbb{Z}[\alpha, \beta]$, it is finitely generated, so $\alpha + \beta$ is integral.
 520 Likewise, $\mathbb{Z}[\alpha\beta]$ is a submodule of $\mathbb{Z}[\alpha, \beta]$, so it is also finitely generated,
 521 and $\alpha\beta$ is integral. \square

522 2.3.1 Minimal Polynomials

523 **Definition 2.3.6** (Minimal Polynomial). The *minimal polynomial* of $\alpha \in \overline{\mathbb{Q}}$
 524 is the monic polynomial $f \in \mathbb{Q}[x]$ of least positive degree such that $f(\alpha) = 0$.

525 It is a consequence of Lemma 2.3.9 below that “the” minimal polynomial
 526 of α is unique. The minimal polynomial of $1/2$ is $x - 1/2$, and the minimal
 527 polynomial of $\sqrt[3]{2}$ is $x^3 - 2$.

528 *Example 2.3.7.* We compute the minimal polynomial of $(\sqrt[3]{2})^2 + 3$. in terms
 529 of $\sqrt[4]{2}$:

this is confusing,
 sometimes easier
 to use numberfield
 to construct ele-
 ments rather than
 typing `(sqrt(2) +
 3).minpoly()`

```
530 K.<a> = NumberField(x^4 - 2)
      a^4
```

```
| 2
```

```
(a^2 + 3).minpoly()
```

```
| x^2 - 6*x + 7
```

531 **Exercise 2.3.8.** Find the minimal polynomial of $\sqrt{2} + \sqrt{3}$ by hand. Check
 532 your result with **Sage**.

533 **Lemma 2.3.9.** Suppose $\alpha \in \overline{\mathbb{Q}}$. Then the minimal polynomial of α divides
 534 any polynomial h such that $h(\alpha) = 0$.

Proof. Let f be a choice of minimal polynomial of α , as in Definition 2.3.6,
 and let h be a polynomial with $h(\alpha) = 0$. Use the division algorithm to
 write $h = qf + r$, where $0 \leq \deg(r) < \deg(f)$. We have

$$r(\alpha) = h(\alpha) - q(\alpha)f(\alpha) = 0,$$

535 so α is a root of r . However, f is a polynomial of least positive degree with
 536 root α , so $r = 0$. \square

537 **Exercise 2.3.10.** Show that the minimal polynomial of an algebraic number
 538 $\alpha \in \overline{\mathbb{Q}}$ is unique.

539 **Lemma 2.3.11.** Suppose $\alpha \in \overline{\mathbb{Q}}$. Then α is an algebraic integer if and only
 540 if the minimal polynomial f of α has coefficients in \mathbb{Z} .

541 *Proof.* First suppose that the minimal polynomial f of α has coefficients in
 542 \mathbb{Z} . Since $f \in \mathbb{Z}[x]$ is monic (by definition) and $f(\alpha) = 0$, we see immediately
 543 that α is an algebraic integer.

544 Now suppose that α is an algebraic integer. Then there is some nonzero
 545 monic $g \in \mathbb{Z}[x]$ such that $g(\alpha) = 0$. By Lemma 2.3.9, we have $g = fh$,
 546 for some $h \in \mathbb{Q}[x]$, and h is monic because f and g are. If $f \notin \mathbb{Z}[x]$, then
 547 some prime p divides the denominator of some coefficient of f . Let p^i be
 548 the largest power of p that divides some denominator of some coefficient f ,
 549 and likewise let p^j be the largest power of p that divides some denominator
 550 of a coefficient of h . Then $p^{i+j}g = (p^i f)(p^j h)$, and if we reduce both sides
 551 modulo p , then the left hand side is 0 but the right hand side is a product
 552 of two nonzero polynomials in $\mathbb{F}_p[x]$, hence nonzero, a contradiction. \square

553 **Exercise 2.3.12.** Which of the following numbers are algebraic integers?

- 554 (a) The number $(1 + \sqrt{5})/2$.
 555 (b) The number $(2 + \sqrt{5})/2$.
 556 (c) The value of the infinite sum $\sum_{n=1}^{\infty} 1/n^2$.
 557 (d) The number $\alpha/3$, where α is a root of $x^4 + 54x + 243$.

558 **Example 2.3.13.** We compute some minimal polynomials in Sage. The min-
 559 imal polynomial of $1/2$:

```
(1/2).minpoly()
```

```
| x - 1/2
```

We construct a root a of $x^2 - 2$ and compute its minimal polynomial:

560

```
K.<a> = NumberField(x^2 - 2)
a^2 - 2
```

```
| 0
```

```
a.minpoly()
```

```
| x^2 - 2
```

make sure this is bold

make sure we use big
K for number fields

561 Finally we compute the minimal polynomial of $\alpha = \sqrt{2}/2 + 3$, which is not
 562 integral, hence Proposition 2.3.4 implies that α is not an algebraic integer:

```
(a/2 + 3).minpoly()
```

563

```
| x^2 - 6*x + 17/2
```

The only elements of \mathbb{Q} that are algebraic integers are the usual integers \mathbb{Z} , since $\mathbb{Z}[1/d]$ is not finitely generated as a \mathbb{Z} -module. Watch out since there are elements of \mathbb{Q} that seem to *appear* to have denominators when written down, but are still algebraic integers. This is an artifact of how we write them down, e.g., if we wrote our integers as a multiple of $\alpha = 2$, then we would write 1 as $\alpha/2$. For example,

$$\alpha = \frac{1 + \sqrt{5}}{2}$$

564 is an algebraic integer, since it is a root of the monic integral polynomial
 565 $x^2 - x - 1$. We verify this using **Sage** below, though of course this is easy
 566 to do by hand (you should try much more complicated examples in **Sage**).

```
k.<a> = QuadraticField(5)
a^2
```

567

```
| 5
```

```
alpha = (1 + a)/2
alpha.minpoly()
```

```
| x^2 - x - 1
```

```
alpha.is_integral()
```

```
| True
```

568 Since $\sqrt{5}$ can be expressed in terms of radicals, we can also compute this
 569 minimal polynomial using the symbolic functionality in Sage.

```
alpha = (1+sqrt(5))/2
alpha.minpoly()
```

```
| x^2 - x - 1
```

570 Here is a more complicated example using a similar approach:

```
alpha = sqrt(2) + 3^(1/4)
alpha.minpoly()
```

```
| x^8 - 8*x^6 + 18*x^4 - 104*x^2 + 1
```

571 *Example 2.3.14.* We illustrate an example of a sum and product of two
 572 algebraic integers being an algebraic integer. We first make the relative
 573 number field obtained by adjoining a root of $x^3 - 5$ to the field $\mathbb{Q}(\sqrt{2})$:

```
k.<a, b> = NumberField([x^2 - 2, x^3 - 5])
k
```

574

```
| Number Field in a with defining polynomial x^2 + -2 over its base field
```

575 Here a and b are roots of $x^2 - 2$ and $x^3 - 5$, respectively.

```
a^2
```

```
| 2
```

576

```
b^3
```

```
| 5
```

577 We compute the minimal polynomial of the sum and product of $\sqrt[3]{5}$ and
 578 $\sqrt{2}$. The command `absolute_minpoly` gives the minimal polynomial of the
 579 element over the rational numbers \mathbb{Q} .

```
(a+b).absolute_minpoly()
```

```
| x^6 - 6*x^4 - 10*x^3 + 12*x^2 - 60*x + 17
```

580

```
(a*b).absolute_minpoly()
```

```
| x^6 - 200
```

The minimal polynomial of the product is $\sqrt[3]{5}\sqrt{2}$ is trivial to compute by hand. In light of the Cayley-Hamilton theorem, we can compute the minimal polynomial of $\alpha = \sqrt[3]{5} + \sqrt{2}$ by hand by computing the determinant of the matrix given by left multiplication by α on the basis

$$1, \sqrt{2}, \sqrt[3]{5}, \sqrt[3]{5}\sqrt{2}, \sqrt[3]{5}^2, \sqrt[3]{5}^2\sqrt{2}.$$

581 This is a general method which works well for computers. However it can
582 also be done using simple algebra.

583 The following is an alternative, more symbolic way to compute the min-
584 imal polynomials above, though it is not provably correct. We compute α
585 to 100 bits precision (via the `n` command), then use the LLL algorithm (via
586 the `algdep` command) to heuristically find a linear relation between the first
587 6 powers of α (see Section 2.5 below for more about LLL).

```

a = 5^(1/3); b = sqrt(2)
c = a+b; c

5^(1/3) + sqrt(2)

(a+b).n(100).algdep(6)

x^6 - 6*x^4 - 10*x^3 + 12*x^2 - 60*x + 17

(a*b).n(100).algdep(6)

x^6 - 200

```

is this example too
long?

590 **Exercise 2.3.15.** Compute the minimal polynomial of $\alpha = \sqrt[3]{5} + \sqrt{2}$ by
591 hand without finding the determinate of a 6×6 matrix.

592 [Hint: Let $a^2 = 2$, $b^3 = 5$, and $x = a + b$. Then $(x - a)^3 = b^3 = 5$. Now
593 simplify and use the fact that $a^2 = 2$.]

594 **Exercise 2.3.16.** Let $\alpha = \sqrt{2} + \frac{1+\sqrt{5}}{2}$.

595 (a) Is α an algebraic integer?

596 (b) Explicitly write down the minimal polynomial of α as an element of
597 $\mathbb{Q}[x]$.

2.3.2 Number fields, rings of integers, and orders

Definition 2.3.17 (Number field). A *number field* is a field K that contains the rational numbers \mathbb{Q} such that the degree $[K : \mathbb{Q}] = \dim_{\mathbb{Q}}(K)$ is finite.

If K is a number field, then by the primitive element theorem there is an $\alpha \in K$ so that $K = \mathbb{Q}(\alpha)$. Let $f(x) \in \mathbb{Q}[x]$ be the minimal polynomial of α . Fix a choice of algebraic closure $\overline{\mathbb{Q}}$ of \mathbb{Q} . Associated to each of the $\deg(f)$ roots $\alpha' \in \overline{\mathbb{Q}}$ of f , we obtain a field embedding $K \hookrightarrow \overline{\mathbb{Q}}$ that sends α to α' . Thus any number field can be embedded in $[K : \mathbb{Q}] = \deg(f)$ distinct ways in $\overline{\mathbb{Q}}$.

Definition 2.3.18 (Ring of Integers). The *ring of integers* of a number field K is the ring

$$\mathcal{O}_K = \{x \in K : x \text{ is an algebraic integer}\}.$$

One of the most basic facts about \mathcal{O}_K is that it is indeed a ring. This fact is important enough to be stated as a separate theorem.

Theorem 2.3.19. *Let K be a number field. The ring of integers \mathcal{O}_K is a ring.*

Proof. This follows directly from Proposition 2.3.5. □

Example 2.3.20. The field \mathbb{Q} of rational numbers is a number field of degree 1, and the ring of integers of \mathbb{Q} is \mathbb{Z} . The field $K = \mathbb{Q}(i)$ of Gaussian integers has degree 2 and $\mathcal{O}_K = \mathbb{Z}[i]$.

Example 2.3.21. The golden ratio $\varphi = (1 + \sqrt{5})/2$ is in the quadratic number field $K = \mathbb{Q}(\sqrt{5}) = \mathbb{Q}(\varphi)$; notice that φ satisfies $x^2 - x - 1$, so $\varphi \in \mathcal{O}_K$. To see that $\mathcal{O}_K = \mathbb{Z}[\varphi]$ directly, we proceed as follows. By Proposition 2.3.4, the algebraic integers K are exactly the elements $a + b\sqrt{5} \in K$, with $a, b \in \mathbb{Q}$ that have integral minimal polynomial. The matrix of $a + b\sqrt{5}$ with respect to the basis $1, \sqrt{5}$ for K is $m = \begin{pmatrix} a & 5b \\ b & a \end{pmatrix}$. The characteristic polynomial of m is $f = (x - a)^2 - 5b^2 = x^2 - 2ax + a^2 - 5b^2$, which is in $\mathbb{Z}[x]$ if and only if $2a \in \mathbb{Z}$ and $a^2 - 5b^2 \in \mathbb{Z}$. Thus $a = a'/2$ with $a' \in \mathbb{Z}$, and $(a'/2)^2 - 5b^2 \in \mathbb{Z}$, so $5b^2 \in \frac{1}{4}\mathbb{Z}$, so $b \in \frac{1}{2}\mathbb{Z}$ as well. If a has a denominator of 2, then b must also have a denominator of 2 to ensure that the difference $a^2 - 5b^2$ is an integer. This proves that $\mathcal{O}_K = \mathbb{Z}[\varphi]$.

Example 2.3.22. The ring of integers of $K = \mathbb{Q}(\sqrt[3]{9})$ is $\mathbb{Z}[\sqrt[3]{3}]$, where $\sqrt[3]{3} = \frac{1}{3}(\sqrt[3]{9})^2 \notin \mathbb{Z}[\sqrt[3]{9}]$. As we will see, in general the problem of computing \mathcal{O}_K given K may be very hard, since it requires factoring a certain potentially large integer.

630 **Exercise 2.3.23.** From basic definitions, find the rings of integers of the
 631 fields $\mathbb{Q}(\sqrt{11})$ and $\mathbb{Q}(\sqrt{-6})$.

632 **Definition 2.3.24** (Order). An *order* in \mathcal{O}_K is any subring R of \mathcal{O}_K such
 633 that the quotient \mathcal{O}_K/R of abelian groups is finite. (By definition R must
 634 contain 1 because it is a ring.)

635 **Exercise 2.3.25.** Let R be a subring of \mathcal{O}_K . Show that R is an order of
 636 \mathcal{O}_K if and only if R contains a spanning set for K as a vector space over \mathbb{Q} .

637 As noted above, $\mathbb{Z}[i]$ is the ring of integers of $\mathbb{Q}(i)$. For every nonzero
 638 integer n , the subring $\mathbb{Z} + ni\mathbb{Z}$ of $\mathbb{Z}[i]$ is an order. The subring \mathbb{Z} of $\mathbb{Z}[i]$ is not
 639 an order, because \mathbb{Z} does not have finite index in $\mathbb{Z}[i]$. Also the subgroup
 640 $2\mathbb{Z} + i\mathbb{Z}$ of $\mathbb{Z}[i]$ is not an order because it is not a ring.

641 **Exercise 2.3.26.** Let K be a quadratic extension of \mathbb{Q} and R be any order
 642 in \mathcal{O}_K . Show that \mathcal{O}_K/R is cyclic as an abelian group and that there is a
 643 bijection between orders of \mathcal{O}_K containing R and divisors of $[\mathcal{O}_K : R]$.

644 *Remark 2.3.27.* Exercise 2.3.26 is used in elliptic curve cryptography to
 645 measure the number of isogenies; for example, see [KKM11, §11.2].

646 **Exercise 2.3.28.** Let K be a number field of degree n . Suppose $\{\alpha_1, \dots, \alpha_n\}$
 647 is a \mathbb{Z} -independent set of algebraic integers. Is $\mathbb{Z}\alpha_1 + \dots + \mathbb{Z}\alpha_n$ an ideal of
 648 \mathcal{O}_K ?

find a good place for
this

649 We define the number field $\mathbb{Q}(i)$ and compute its ring of integers.

```
K.<i> = NumberField(x^2 + 1)
OK = K.ring_of_integers(); OK
```

650

```
Order with module basis 1, i in Number Field in i with
defining polynomial x^2 + 1
```

651 Next we compute the order $\mathbb{Z} + 3i\mathbb{Z}$.

```
O3 = K.order(3*i); O3
```

```
Order with module basis 1, 3*i in Number Field in i with
defining polynomial x^2 + 1
```

652

```
O3.gens()
```

```
[1, 3*i]
```

653 We test whether certain elements are in the order.

```
5 + 9*i in O3
```

```
True
```

654

```
1 + 2*i in O3
```

```
False
```

655 We will frequently consider orders because they are often much easier
 656 to write down explicitly than \mathcal{O}_K . For example, if $K = \mathbb{Q}(\alpha)$ and α is an
 657 algebraic integer, then $\mathbb{Z}[\alpha]$ is an order in \mathcal{O}_K , but frequently $\mathbb{Z}[\alpha] \neq \mathcal{O}_K$.

658 *Example 2.3.29.* In this example $[\mathcal{O}_K : \mathbb{Z}[a]] = 2197$. First we define the
 659 number field $K = \mathbb{Q}(a)$ where a is a root of $x^3 - 15x^2 - 94x - 3674$, then
 660 we compute the order $\mathbb{Z}[a]$ generated by a .

```
K.<a> = NumberField(x^3 - 15*x^2 - 94*x - 3674)
Oa = K.order(a); Oa
```

```
Order with module basis 1, a, a^2 in Number Field in a with defining
polynomial x^3 - 15*x^2 - 94*x - 3674
```

661

```
Oa.basis()
```

```
[1, a, a^2]
```

662 Next we compute a \mathbb{Z} -basis for the maximal order \mathcal{O}_K of K , and compute
 663 that the index of $\mathbb{Z}[a]$ in \mathcal{O}_K is $2197 = 13^3$.

```
OK = K.maximal_order()
OK.basis()
```

```
[25/169*a^2 + 10/169*a + 1/169, 5/13*a^2 + 1/13*a, a^2]
```

664

```
Oa.index_in(OK)
```

```
2197
```

665 **Lemma 2.3.30.** *Let \mathcal{O}_K be the ring of integers of a number field. Then*
 666 *$\mathcal{O}_K \cap \mathbb{Q} = \mathbb{Z}$ and $\mathbb{Q}\mathcal{O}_K = K$.*

667 *Proof.* Suppose $\alpha \in \mathcal{O}_K \cap \mathbb{Q}$ with $\alpha = a/b \in \mathbb{Q}$ in lowest terms and $b > 0$.
 668 Since α is integral, $\mathbb{Z}[a/b]$ is finitely generated as a module, so $b = 1$.

To prove that $\mathbb{Q}\mathcal{O}_K = K$, suppose $\alpha \in K$, and let $f(x) \in \mathbb{Q}[x]$ be the minimal monic polynomial of α . For any positive integer d , the minimal monic polynomial of $d\alpha$ is $d^{\deg(f)} f(x/d)$, i.e., the polynomial obtained from $f(x)$ by multiplying the coefficient of $x^{\deg(f)}$ by 1, multiplying the coefficient of $x^{\deg(f)-1}$ by d , multiplying the coefficient of $x^{\deg(f)-2}$ by d^2 , etc. If d is the least common multiple of the denominators of the coefficients of f , then the minimal monic polynomial of $d\alpha$ has integer coefficients, so $d\alpha$ is integral and $d\alpha \in \mathcal{O}_K$. This proves that $\mathbb{Q}\mathcal{O}_K = K$. \square

Exercise 2.3.31. Which of the following rings are orders in the given number field, i.e. orders in the ring of integers of the given number field.

- (a) The ring $R = \mathbb{Z}[i]$ in the number field $\mathbb{Q}(i)$.
- (b) The ring $R = \mathbb{Z}[i/2]$ in the number field $\mathbb{Q}(i)$.
- (c) The ring $R = \mathbb{Z}[17i]$ in the number field $\mathbb{Q}(i)$.
- (d) The ring $R = \mathbb{Z}[i]$ in the number field $\mathbb{Q}(\sqrt[4]{-1})$.

Exercise 2.3.32. Find the ring of integers of $\mathbb{Q}(a)$, where $a^5 + 7a + 1 = 0$ using a computer.

2.3.3 Function fields

Let k be any field. We can also make the same definitions, but with \mathbb{Q} replaced by the field $k(t)$ of rational functions in an indeterminate t , and \mathbb{Z} replaced by $k[t]$. The analogue of a number field is called a *function field*; it is a finite algebraic extension field K of $k(t)$. Elements of K have a unique minimal polynomial as above, and the ring of integers of K consists of those elements whose monic minimal polynomial has coefficients in the polynomial ring $k[t]$.

Geometrically, if $F(x, t) = 0$ is an affine equation that defines (via projective closure) a nonsingular projective curve C , then $K = k(t)[x]/(F(x, t))$ is a function field. We view the field K as the field of all rational functions on the projective closure of the curve C . The ring of integers \mathcal{O}_K is the subring of rational functions that have no poles on the affine curve $F(x, t) = 0$, though they may have poles at infinity, i.e., at the extra points we introduce when passing to the projective closure C . The algebraic arguments we gave above prove that \mathcal{O}_K is a ring. This is also geometrically intuitive, since the sum and product of two functions with no poles also have no poles.

Exercise 2.3.33. Let $k = \mathbb{F}_p$ be the finite field with p elements where p is some prime. Find all automorphisms of $k(t)$. Note that an automorphism is completely characterized by its value on t . How many such automorphisms are there?

[*Hint:* For some people, it is easier to think about the equivalent question: What rational functions $f \in k(t)$ is the map $k(t) \rightarrow k(t)$ given by $t \mapsto f(t)$ an automorphism?]

2.4 Norms and Traces

In this section we develop some basic properties of norms, traces, and discriminants, and give more properties of rings of integers in the general context of Dedekind domains.

Before discussing norms and traces we introduce some notation for field extensions. If $K \subset L$ are number fields, we let $[L : K]$ denote the dimension of L viewed as a K -vector space. If K is a number field and $a \in \overline{\mathbb{Q}}$, let $K(a)$ be the extension of K generated by a , which is the smallest number field that contains both K and a . If $a \in \overline{\mathbb{Q}}$ then a has a minimal polynomial $f(x) \in \mathbb{Q}[x]$, and the *Galois conjugates* of a are the roots of f . These are called the Galois conjugates because they are the orbit of a under the action of $\text{Gal}(\overline{\mathbb{Q}}/\mathbb{Q})$.

Example 2.4.1. The element $\sqrt{2}$ has minimal polynomial $x^2 - 2$ and the Galois conjugates of $\sqrt{2}$ are $\sqrt{2}$ and $-\sqrt{2}$. The cube root $\sqrt[3]{2}$ has minimal polynomial $x^3 - 2$ and three Galois conjugates $\sqrt[3]{2}, \zeta_3 \sqrt[3]{2}, \zeta_3^2 \sqrt[3]{2}$, where ζ_3 is a cube root of unity.

We create the extension $\mathbb{Q}(\zeta_3)(\sqrt[3]{2})$ in Sage.

```
L.<cuberooroot2> = CyclotomicField(3).extension(x^3 - 2)
cuberooroot2^3
```

2

Then we list the Galois conjugates of $\sqrt[3]{2}$.

```
cuberooroot2.galois_conjugates(L)
```

[cuberooroot2, (-zeta3 - 1)*cuberooroot2, zeta3*cuberooroot2]

Note that $\zeta_3^2 = -\zeta_3 - 1$:

```

zeta3 = L.base_field().0
zeta3^2
730
- zeta3 - 1

```

731 Suppose $K \subset L$ is an inclusion of number fields and let $a \in L$. Then
 732 left multiplication by a defines a K -linear transformation $\ell_a : L \rightarrow L$. (The
 733 transformation ℓ_a is K -linear because L is commutative.)

Definition 2.4.2 (Norm and Trace). The *norm* and *trace* of a from L to K are

$$\text{Norm}_{L/K}(a) = \det(\ell_a) \quad \text{and} \quad \text{Trace}_{L/K}(a) = \text{Trace}(\ell_a).$$

We know from linear algebra that determinants are multiplicative and traces are additive, so for $a, b \in L$ we have

$$\text{Norm}_{L/K}(ab) = \text{Norm}_{L/K}(a) \cdot \text{Norm}_{L/K}(b)$$

and

$$\text{Trace}_{L/K}(a + b) = \text{Trace}_{L/K}(a) + \text{Trace}_{L/K}(b).$$

734 Note that if $f \in \mathbb{Q}[x]$ is the characteristic polynomial of ℓ_a , then the
 735 constant term of f is $(-1)^{\deg(f)} \det(\ell_a)$, and the coefficient of $x^{\deg(f)-1}$ is
 736 $-\text{Trace}(\ell_a)$.

Proposition 2.4.3. Let $a \in L$ and let $\sigma_1, \dots, \sigma_d$, where $d = [L : K]$, be the distinct field embeddings $L \hookrightarrow \overline{\mathbb{Q}}$ that fix every element of K . Then

$$\text{Norm}_{L/K}(a) = \prod_{i=1}^d \sigma_i(a) \quad \text{and} \quad \text{Trace}_{L/K}(a) = \sum_{i=1}^d \sigma_i(a).$$

737 *Proof.* We prove the proposition by computing the characteristic polynomial
 738 of a . Let $f \in K[x]$ be the minimal polynomial of a over K , and note that f
 739 has distinct roots and is irreducible, since it is the polynomial in $K[x]$ of
 740 least degree that is satisfied by a and K has characteristic 0. Since f is
 741 irreducible, we have $K(a) \cong K[x]/(f)$, so $[K(a) : K] = \deg(f)$. Also a
 742 satisfies a polynomial if and only if ℓ_a does, so the characteristic polynomial
 743 of ℓ_a acting on $K(a)$ is f . Let b_1, \dots, b_n be a basis for L over $K(a)$ and
 744 note that $1, \dots, a^m$ is a basis for $K(a)/K$, where $m = \deg(f) - 1$. Then
 745 $a^i b_j$ is a basis for L over K , and left multiplication by a acts the same way
 746 on the span of $b_j, ab_j, \dots, a^m b_j$ as on the span of $b_k, ab_k, \dots, a^m b_k$, for any
 747 pair $j, k \leq n$. Thus the matrix of ℓ_a on L is a block direct sum of copies

of the matrix of ℓ_a acting on $K(a)$, so the characteristic polynomial of ℓ_a on L is $f^{[L:K(a)]}$. The proposition follows because the roots of $f^{[L:K(a)]}$ are exactly the images $\sigma_i(a)$, with multiplicity $[L : K(a)]$, since each embedding of $K(a)$ into $\overline{\mathbb{Q}}$ extends in exactly $[L : K(a)]$ ways to L . \square

Warning 2.4.4. It is important in Proposition 2.4.3 that the product and sum be over *all* the images $\sigma_i(a)$, not over just the distinct images. For example, if $a = 1 \in L$, then $\text{Trace}_{L/K}(a) = [L : K]$, whereas the sum of the distinct conjugates of a is 1.

Remark 2.4.5. Let $K \subset L$ be an extension of number fields. If $\alpha \in \mathcal{O}_L$, then the formula of Proposition 2.4.3 implies that the norm and trace down to K of α is an element of \mathcal{O}_K , because the sum and product of algebraic integers is an algebraic integer.

The following corollary asserts that the norm and trace behave well in towers.

Corollary 2.4.6. *Suppose $K \subset L \subset M$ is a tower of number fields, and let $a \in M$. Then*

$$\text{Norm}_{M/K}(a) = \text{Norm}_{L/K}(\text{Norm}_{M/L}(a)) \quad \text{and} \quad \text{Trace}_{M/K}(a) = \text{Trace}_{L/K}(\text{Trace}_{M/L}(a)).$$

Proof. The proof uses that every embedding $L \hookrightarrow \overline{\mathbb{Q}}$ extends in exactly $[M : L]$ way to an embedding $M \hookrightarrow \overline{\mathbb{Q}}$. This is clear if we view M as $L[x]/(h(x))$ for some irreducible polynomial $h(x) \in L[x]$ of degree $[M : L]$, and note that the extensions of $L \hookrightarrow \overline{\mathbb{Q}}$ to M correspond to the roots of h , of which there are $\deg(h)$, since $\overline{\mathbb{Q}}$ is algebraically closed.

For the first equation, both sides are the product of $\sigma_i(a)$, where σ_i runs through the embeddings of M into $\overline{\mathbb{Q}}$ that fix K . To see this, suppose $\sigma : L \rightarrow \overline{\mathbb{Q}}$ fixes K . If σ' is an extension of σ to M , and τ_1, \dots, τ_d are the embeddings of M into $\overline{\mathbb{Q}}$ that fix L , then $\sigma'\tau_1, \dots, \sigma'\tau_d$ are exactly the extensions of σ to M . For the second statement, both sides are the sum of the $\sigma_i(a)$. \square

Proposition 2.4.7. *Let K be a number field. The ring of integers \mathcal{O}_K is a lattice in K , i.e., $\mathbb{Q}\mathcal{O}_K = K$ and \mathcal{O}_K is an abelian group of rank $[K : \mathbb{Q}]$.*

Proof. We saw in Lemma 2.3.30 that $\mathbb{Q}\mathcal{O}_K = K$. Thus there exists a basis a_1, \dots, a_n for K , where each a_i is in \mathcal{O}_K . Suppose that as $x = \sum_{i=1}^n c_i a_i \in \mathcal{O}_K$ varies over all elements of \mathcal{O}_K the denominators of the coefficients c_i are not all uniformly bounded. Then subtracting off integer multiples of the a_i , we see that as $x = \sum_{i=1}^n c_i a_i \in \mathcal{O}_K$ varies over elements of \mathcal{O}_K with c_i

between 0 and 1, the denominators of the c_i are also arbitrarily large. This implies that there are infinitely many elements of \mathcal{O}_K in the bounded subset

$$S = \{c_1 a_1 + \cdots + c_n a_n : c_i \in \mathbb{Q}, 0 \leq c_i \leq 1\} \subset K.$$

Thus for any $\varepsilon > 0$, there are elements $a, b \in \mathcal{O}_K$ such that the coefficients of $a - b$ are all less than ε (otherwise the elements of \mathcal{O}_K would all be a “distance” of least ε from each other, so only finitely many of them would fit in S).

As mentioned above, the norms of elements of \mathcal{O}_K are integers. Since the norm of an element is the determinant of left multiplication by that element, the norm is a homogenous polynomial of degree n in the indeterminate coefficients c_i , which is 0 only on the element 0, so the constant term of this polynomial is 0. If the c_i get arbitrarily small for elements of \mathcal{O}_K , then the values of the norm polynomial get arbitrarily small, which would imply that there are elements of \mathcal{O}_K with positive norm too small to be in \mathbb{Z} , a contradiction. So the set S contains only finitely many elements of \mathcal{O}_K . Thus the denominators of the c_i are bounded, so for some d , we have that \mathcal{O}_K has finite index in $A = \frac{1}{d}\mathbb{Z}a_1 + \cdots + \frac{1}{d}\mathbb{Z}a_n$. Since A is isomorphic to \mathbb{Z}^n , it follows from the structure theorem for finitely generated abelian groups that \mathcal{O}_K is isomorphic as a \mathbb{Z} -module to \mathbb{Z}^n , as claimed. \square

Corollary 2.4.8. *The ring of integers \mathcal{O}_K of a number field is noetherian.*

Proof. By Proposition 2.4.7, the ring \mathcal{O}_K is finitely generated as a module over \mathbb{Z} , so it is certainly finitely generated as a ring over \mathbb{Z} . By Theorem 2.2.11, \mathcal{O}_K is noetherian. \square

2.5 Recognizing Algebraic Numbers using LLL

Suppose we somehow compute a decimal approximation α to some rational number $\beta \in \mathbb{Q}$ and from this wish to recover β . For concreteness, say

$$\beta = \frac{22}{389} = 0.05655526992287917737789203084832904884318766066838046 \dots$$

and we compute

$$\alpha = 0.056555.$$

Now suppose given only α that you would like to recover β . A standard technique is to use continued fractions, which yields a sequence of good rational approximations for α ; by truncating right before a surprisingly big partial quotient (the 23 in the continued fraction v), we obtain β :

```

v = continued_fraction(0.056555); v

```

```

[0, 17, 1, 2, 6, 1, 23, 1, 1, 1, 1, 1, 2]

```

```

convergents([0, 17, 1, 2, 6, 1])

```

```

[0, 1/17, 1/18, 3/53, 19/336, 22/389]

```

Generalizing this, suppose next that somehow you numerically approximate an algebraic number, e.g., by evaluating a special function and get a decimal approximation $\alpha \in \mathbb{C}$ to an algebraic number $\beta \in \overline{\mathbb{Q}}$. For concreteness, suppose $\beta = \frac{1}{3} + \sqrt[4]{3}$:

```

N(1/3 + 3^(1/4), digits=50)

```

```

1.64940734628582579415255223513033238849340192353916

```

Now suppose you very much want to find the (rescaled) minimal polynomial $f(x) \in \mathbb{Z}[x]$ of β just given this numerical approximation α . This is of great value even without proof, since often in practice once you know a potential minimal polynomial you can verify that it is in fact right. Exactly this situation arises in the explicit construction of class fields (a more advanced topic in number theory) and in the construction of Heegner points on elliptic curves. As we will see, the LLL algorithm provides a polynomial time way to solve this problem, assuming α has been computed to sufficient precision.

2.5.1 LLL Reduced Basis

Given a basis b_1, \dots, b_n for \mathbb{R}^n , the *Gram-Schmidt orthogonalization* process produces an orthogonal basis b_1^*, \dots, b_n^* for \mathbb{R}^n as follows. Define inductively

$$b_i^* = b_i - \sum_{j < i} \mu_{i,j} b_j^*$$

where

$$\mu_{i,j} = \frac{b_i \cdot b_j^*}{b_j^* \cdot b_j^*}.$$

Example 2.5.1. We compute the Gram-Schmidt orthogonal basis of the rows of a matrix. Note that no square roots are introduced in the process; there would be square roots if we constructed an orthonormal basis.

```
A = matrix(ZZ, 2, [1,2, 3,4]); A
```

```

818  [1 2]
      [3 4]

```

```
Bstar, mu = A.gramm_schmidt()
```

819 The rows of the matrix B^* are obtained from the rows of A by the Gramm-
820 Schmidt procedure.

```
Bstar
```

```

821  [ 1 2]
      [ 4/5 -2/5]

```

```
mu
```

```

      [ 0 0]
      [11/5 0]

```

822 A lattice $L \subset \mathbb{R}^n$ is a subgroup that is free of rank n such that $\mathbb{R}L = \mathbb{R}^n$.

Definition 2.5.2 (LLL-reduced basis). The basis b_1, \dots, b_n for a lattice $L \subset \mathbb{R}^n$ is *LLL reduced* if for all i, j ,

$$|\mu_{i,j}| \leq \frac{1}{2}$$

and for each $i \geq 2$,

$$|b_i^*|^2 \geq \left(\frac{3}{4} - \mu_{i,i-1}^2 \right) |b_{i-1}^*|^2$$

For example, the basis $b_1 = (1, 2)$, $b_2 = (3, 4)$ for a lattice L is *not* LLL reduced because $b_1^* = b_1$ and

$$\mu_{2,1} = \frac{b_2 \cdot b_1^*}{b_1^* \cdot b_1^*} = \frac{11}{5} > \frac{1}{2}.$$

However, the basis $b_1 = (1, 0)$, $b_2 = (0, 2)$ for L is LLL reduced, since

$$\mu_{2,1} = \frac{b_2 \cdot b_1^*}{b_1^* \cdot b_1^*} = 0,$$

and

$$2^2 \geq (3/4) \cdot 1^2.$$

```
A = matrix(ZZ, 2, [1,2, 3,4])
A.LLL()
```

823

```

[1 0]
[0 2]
```

824 **2.5.2 What LLL really means**

825 The following theorem is not too difficult to prove.

826 Let b_1, \dots, b_n be an LLL reduced basis for a lattice $L \subset \mathbb{R}^n$. Let $d(L)$
 827 denote the absolute value of the determinant of any matrix whose rows are
 828 basis for L . Then the vectors b_i are “nearly orthogonal” and “short” in the
 829 sense of the following theorem:

830 **Theorem 2.5.3.** *We have*831 1. $d(L) \leq \prod_{i=1}^n |b_i| \leq 2^{n(n-1)/4} d(L)$.2. For $1 \leq j \leq i \leq n$, we have

$$|b_j| \leq 2^{(i-1)/2} |b_i^*|.$$

3. The vector b_1 is very short in the sense that

$$|b_1| \leq 2^{(n-1)/4} d(L)^{1/n}$$

and for every nonzero $x \in L$ we have

$$|b_1| \leq 2^{(n-1)/2} |x|.$$

4. More generally, for any linearly independent $x_1, \dots, x_t \in L$, we have

$$|b_j| \leq 2^{(n-1)/2} \max(|x_1|, \dots, |x_t|)$$

832 for $1 \leq j \leq t$.

833 Perhaps the most amazing thing about the idea of an LLL reduced basis
 834 is that there is an algorithm (in fact many) that given a basis for a lattice L
 835 produce an LLL reduced basis for L , and do so *quickly*, i.e., in polynomial
 836 time in the number of digits of the input. The current optimal implementa-
 837 tion (and practically optimal algorithms) for computing LLL reduced basis
 838 are due to Damien Stehle, and are included standard in Magma in **Sage**.
 839 Stehle’s code is amazing – it can LLL reduce a random lattice in \mathbb{R}^n for
 840 $n < 1000$ in a matter of minutes!

```

A = random_matrix(ZZ, 200)
t = cputime()
B = A.LLL()
841 cputime(t)      # random output

```

```

| 3.0494159999999999

```

842 There is even a very fast variant of Stehle's implementation that computes
843 a basis for L that is very likely LLL reduced but may in rare cases fail to
844 be LLL reduced.

```

t = cputime()
B = A.LLL(algorithm="fpLLL:fast") # not tested
845 cputime(t)      # random output

```

```

| 0.968426999999999837

```

846 2.5.3 Applying LLL

847 The LLL definition and algorithm has many application in number theory,
848 e.g., to cracking lattice-based cryptosystems, to enumerating all short vec-
849 tors in a lattice, to finding relations between decimal approximations to
850 complex numbers, to very fast univariate polynomial factorization in $\mathbb{Z}[x]$
851 and more generally in $K[x]$ where K is a number fields, and to computation
852 of kernels and images of integer matrices. LLL can also be used to solve
853 the problem of recognizing algebraic numbers mentioned at the beginning
854 of Section 2.5.

855 Suppose as above that α is a decimal approximation to some algebraic
856 number β , and to for simplicity assume that $\alpha \in \mathbb{R}$ (the general case of
857 $\alpha \in \mathbb{C}$ is described in [Coh93]). We finish by explaining how to use LLL to
858 find a polynomial $f(x) \in \mathbb{Z}[x]$ such that $f(\alpha)$ and its coefficients are small,
859 hence has a shot at being the minimal polynomial of β .

860 Given a real number decimal approximation α , an integer d (the degree),
861 and an integer K (a function of the precision to which α is known), the
862 following steps produce a polynomial $f(x) \in \mathbb{Z}[x]$ of degree at most d such
863 that $f(\alpha)$ is small.

- 864 1. Form the lattice in \mathbb{R}^{d+2} with basis the rows of the matrix A whose
865 first $(d+1) \times (d+1)$ part is the identity matrix, and whose last column
866 has entries

$$K, \lfloor K\alpha \rfloor, \lfloor K\alpha^2 \rfloor, \dots, \lfloor K\alpha^d \rfloor. \quad (2.2)$$

867 (Note this matrix is $(d+1) \times (d+2)$ so the lattice is not of full rank
 868 in \mathbb{R}^{d+2} , which isn't a problem, since the LLL definition also makes
 869 sense for fewer vectors.)

870 2. Compute an LLL reduced basis for the \mathbb{Z} -span of the rows of A , and
 871 let B be the corresponding matrix. Let $b_1 = (a_0, a_1, \dots, a_{d+1})$ be the
 872 first row of B and notice that B is obtained from A by left multipli-
 873 cation by an invertible integer matrix. Thus a_0, \dots, a_d are the linear
 874 combination of the (2.2) that equals a_{d+1} . Moreover, since B is LLL
 875 reduced we expect that a_{d+1} is relatively small.

876 3. Output $f(x) = a_0 + a_1x + \dots + a_dx^d$. We have that $f(\alpha) \sim a_{d+1}/K$,
 877 which is small. Thus $f(x)$ may be a very good candidate for the
 878 minimal polynomial of β (the algebraic number we are approximating),
 879 assuming d was chosen minimally and α was computed to sufficient
 880 precision.

881 The following is a complete implementation of the above algorithm in
 882 Sage:

```

def myalgdep(a, d, K=10^6):
    aa = [floor(K*a^i) for i in range(d+1)]
    A = identity_matrix(ZZ, d+1)
    B = matrix(ZZ, d+1, 1, aa)
883 A = A.augment(B)
    L = A.LLL()
    v = L[0][:-1].list()
    return ZZ['x'](v)

```

884 Here is an example of using it:

```

R.<x> = RDF[]
f = 2*x^3 - 3*x^2 + 10*x - 4
a = f.roots()[0][0]; a
885 myalgdep(a, 3, 10^6)          # not tested

```

2*x^3 - 3*x^2 + 10*x - 4

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