

Today.

Polynomials.

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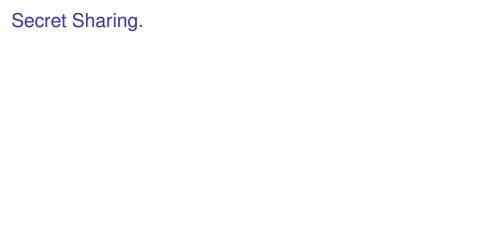
Secret Sharing.

Today.

Polynomials.

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Correcting for loss or even corruption.



Share secret among \boldsymbol{n} people.

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Secrecy: Any k-1 knows nothing.

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The idea of the day.

Two points make a line. Lots of lines go through one point.

A polynomial

$$P(x) = a_d x^d + a_{d-1} x^{d-1} \cdots + a_0.$$

is specified by **coefficients** $a_d, \dots a_0$.

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Polynomials over reals: $a_1, \ldots, a_d \in \Re$, use $x \in \Re$.

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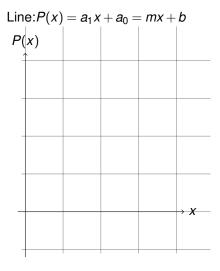
Polynomials P(x) with arithmetic modulo p: ¹ $a_i \in \{0, ..., p-1\}$ and

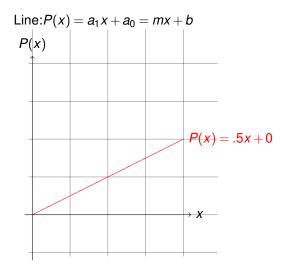
$$P(x) = a_d x^d + a_{d-1} x^{d-1} \cdots + a_0 \pmod{p},$$
 for $x \in \{0, \dots, p-1\}.$

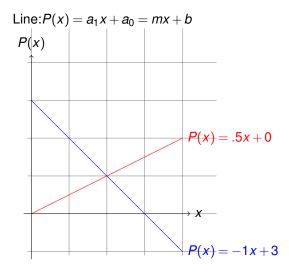
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Parabola: $P(x) = a_2x^2 + a_1x + a_0$

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Parabola: $P(x) = a_2x^2 + a_1x + a_0 = ax^2 + bx + c$

Line:
$$P(x) = a_1 x + a_0 = mx + b$$

$$P(x) = 0.5x^2 - x + 0.1$$

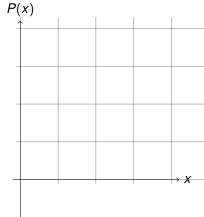
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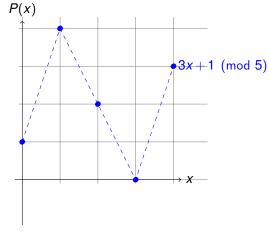
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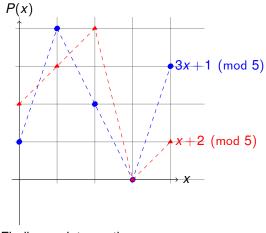
$$P(x) = 0.5x^2 - x + 0.1$$

$$P(x) = -.3x^2 + 1x + .1$$

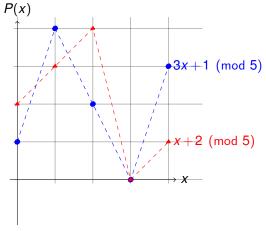
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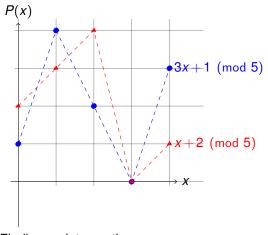
Finding an intersection.
$$x+2 \equiv 3x+1 \pmod{5}$$
 $\implies 2x \equiv 1 \pmod{5}$



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3 is multiplicative inverse of 2 modulo 5.



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$$x+2\equiv 3x+1\pmod{5}$$
 $\implies 2x\equiv 1\pmod{5}$ $\implies x\equiv 3\pmod{5}$ 3 is multiplicative inverse of 2 modulo 5. Good when modulus is prime!!

Fact: Exactly 1 degree $\leq d$ polynomial contains d+1 points. ²

²Points with different *x* values.

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Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime p contains d+1 pts.

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Poll.

Two points determine a line. What facts below tell you this?

Say points are $(x_1, y_1), (x_2, y_2)$ **.**

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- (A) Line is y = mx + b.
- (B) Plug in a point gives an equation: $y_1 = mx_1 + b$
- (C) The unknowns are *m* and *b*.
- (D) If equations have unique solution, done.

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All true.

Flow Poll.

Why solution? Why unique?

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- (A) Solution cuz: $m = (y_2 y_1)/(x_2 x_1), b = y_1 m(x_1)$
- (B) Unique cuz, only one line goes through two points.
- (C) Try: $(m'x + b') (mx + b) = (m' m)x + (b b') = ax + c \neq 0$.
- (D) Either $ax_1 + c \neq 0$ or $ax_2 + c \neq 0$.
- (E) Contradiction.

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Flow poll. (All true. (B) is not a proof, it is restatement.)

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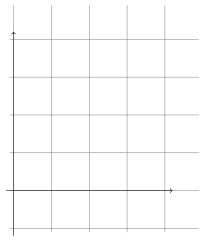
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- (C) $a_0 = m$
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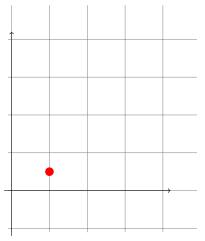
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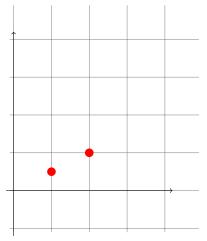
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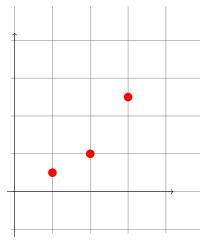
Fact: Exactly 1 degree $\leq d$ polynomial contains d+1 points. ³



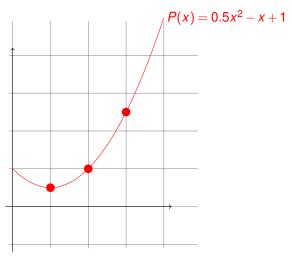
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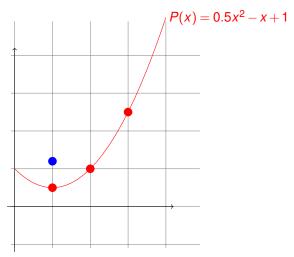
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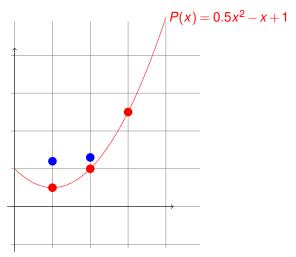
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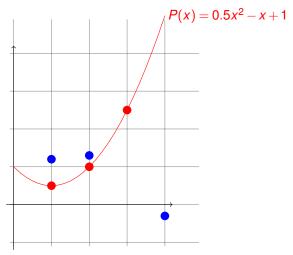
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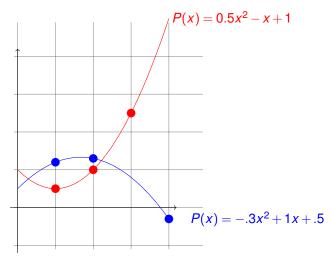
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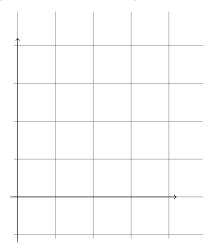


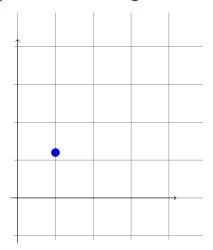
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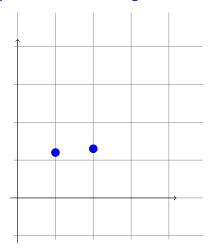


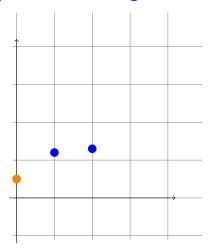
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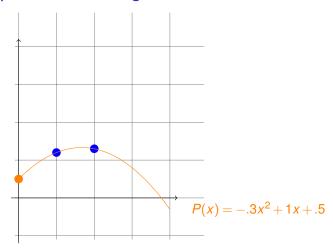
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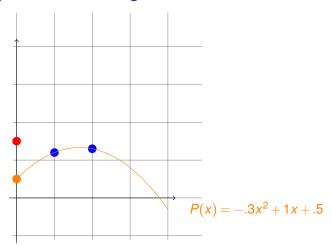


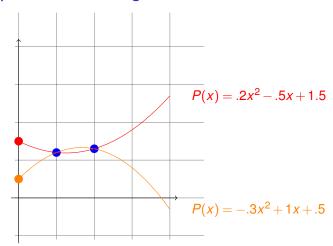


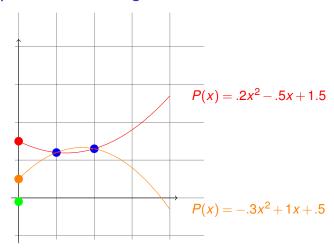


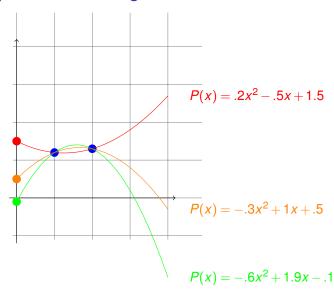


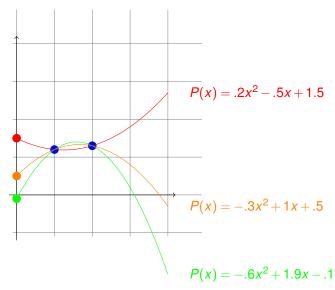












Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime p contains d+1 pts.

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Knowing k pts \implies only one P(x) \implies evaluate P(0).

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Poll:example.

The polynomial from the scheme: $P(x) = 2x^2 + 1x + 3 \pmod{5}$. What is true for the secret sharing scheme using P(x)?

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- (A) The secret is "2".
- (B) The secret is "3".
- (C) A share could be (1,5) cuz P(1) = 5
- (D) A share could be (2,4)
- (E) A share could be (0,3)

$$P(1) = m(1) + b \equiv m + b$$

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For a line, $a_1x + a_0 = mx + b$ contains points (1,3) and (2,4).

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Backsolve: $b \equiv 2 \pmod{5}$.

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And the line is...

$$x+2 \mod 5$$
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Quadratic

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So polynomial is $2x^2 + 1x + 4 \pmod{5}$

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Will this always work?

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Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime p contains d+1 pts.

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Multiplicative inverses due to gcd(x,p) = 1, for all $x \in \{1,...,p-1\}$

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For set of *x*-values, x_1, \ldots, x_{d+1} .

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Will $y_2\Delta_2(x)$ contain (x_2,y_2) ?

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See the idea?

For set of *x*-values, x_1, \ldots, x_{d+1} .

$$\Delta_i(x) = \begin{cases} 1, & \text{if } x = x_i. \\ 0, & \text{if } x = x_j \text{ for } j \neq i. \\ ?, & \text{otherwise.} \end{cases}$$
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Mark what's true.

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- (D) $\Delta_1(x_3) = 1$
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Find $\Delta_1(x)$ polynomial contains (1,1); (2,0); (3,0).

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Put the delta functions together.

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A parabola (degree 2), can intersect y = 0 at only two x's.

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Intuitively, a field is a set with operations corresponding to addition, multiplication, and division.

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(Almost) the same as what is missing: one P(i).



Runtime.

Runtime: polynomial in k, n, and $\log p$.

- 1. Evaluate degree k-1 polynomial n times using $\log p$ -bit numbers.
- 2. Reconstruct secret by solving system of *k* equations using $\log p$ -bit arithmetic.

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Infinite number for reals, rationals, complex numbers!

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Secret Sharing:

k points on degree k-1 polynomial is great! Can hand out n points on polynomial as shares.