1 Polynomial Practice

(a) If f and g are non-zero real polynomials, how many roots do the following polynomials have at least? How many can they have at most? (Your answer may depend on the degrees of f and g.)

(i) $f+g$

(ii) $f \cdot g$

(iii) f/g, assuming that f/g is a polynomial

(b) Now let f and g be polynomials over $\text{GF}(p)$.

(i) We say a polynomial $f = 0$ if $\forall x, f(x) = 0$. If $f \cdot g = 0$, is it true that either $f = 0$ or $g = 0$?

(ii) How many f of degree *exactly* $d < p$ are there such that $f(0) = a$ for some fixed $a \in \{0, 1, \ldots, p-1\}$?

(c) Find a polynomial f over $\text{GF}(5)$ that satisfies $f(0) = 1, f(2) = 2, f(4) = 0$. How many such polynomials are there?

Solution:

(a) (i) It could be that $f+g$ has no roots at all (example: $f(x) = 2x^2 - 1$ and $g(x) = -x^2 + 2$), so the minimum number is 0. However, if the highest degree of $f+g$ is odd, then it has to cross the x-axis at least once, meaning that the minimum number of roots for odd degree polynomials is 1. On the other hand, $f+g$ is a polynomial of degree at most $m = \max(\deg f, \deg g)$, so it can have at most m roots. The one exception to this expression is if $f = -g$. In that case, $f+g = 0$, so the polynomial has an infinite number of roots!

(ii) A product is zero if and only if one of its factors vanishes. So if $f(x) \cdot g(x) = 0$ for some x, then either x is a root of f or it is a root of g, which gives a maximum of $\deg f + \deg g$ possibilities. Again, there may not be any roots if neither f nor g have any roots (example: $f(x) = g(x) = x^2 + 1$).

(iii) If f/g is a polynomial, then it must be of degree $d = \deg f - \deg g$ and so there are at most d roots. Once more, it may not have any roots, e.g. if $f(x) = g(x)(x^2+1)$, $f/g = x^2 + 1$ has no root.

(b) (i) No.

Example 1: $x^{p-1}-1$ and x are both non-zero polynomials on $GF(p)$ for any p. x has a root at 0, and by Little Fermat, $x^{p-1}-1$ has a root at all non-zero points in $GF(p)$. So, their product x^p-x must have a zero on all points in $GF(p)$.

Example 2: To satisfy $f \cdot g = 0$, all we need is $(\forall x \in S, f(x) = 0 \vee g(x) = 0)$ where $S = \{0,\dots,p-1\}$. We may see that this is not equivalent to $(\forall x \in S, f(x) = 0)) \vee (\forall x \in S, g(x) = 0)$.

To construct a concrete example, let $p = 2$ and we enforce $f(0) = 1, f(1) = 0$ (e.g. $f(x) = 1-x$), and $g(0) = 0, g(1) = 1$ (e.g. $g(x) = x$). Then $f \cdot g = 0$ but neither f nor g is the zero polynomial.

(ii) We know that in general each of the $d+1$ coefficients of $f(x) = \sum_{k=0}^{d} c_k x^k$ can take any of p values. However, the conditions $f(0)$ and $\deg f = d$ impose constraints on the constant coefficient $f(0) = c_0 = a$ and the top coefficient $x_d \neq 0$. Hence we are left with $(p-1)\cdot p^{d-1}$ possibilities.

(c) We know by part (b) that any polynomial over GF(5) can be of degree at most 4. A polynomial of degree ≤ 4 is determined by 5 points (x_i, y_i). We have assigned three, which leaves $5^2 = 25$ possibilities. To find a specific polynomial, we use Lagrange interpolation:

$$\Delta_0(x) = 2(x-2)(x-4) \qquad \Delta_2(x) = x(x-4) \qquad \Delta_4(x) = 2x(x-2),$$

and so $f(x) = \Delta_0(x) + 2\Delta_2(x) = 4x^2 + 1$.

2 Lagrange Interpolation in Finite Fields

Find a unique polynomial $p(x)$ of degree at most 3 that passes through points $(-1,3)$, $(0,1)$, $(1,2)$, and $(2,0)$ in modulo 5 arithmetic using the Lagrange interpolation.

(a) Find $p_{-1}(x)$ where $p_{-1}(0) \equiv p_{-1}(1) \equiv p_{-1}(2) \equiv 0 \pmod 5$ and $p_{-1}(-1) \equiv 1 \pmod 5$.

(b) Find $p_0(x)$ where $p_0(-1) \equiv p_0(1) \equiv p_0(2) \equiv 0 \pmod 5$ and $p_0(0) \equiv 1 \pmod 5$.

(c) Find $p_1(x)$ where $p_1(-1) \equiv p_1(0) \equiv p_1(2) \equiv 0 \pmod 5$ and $p_1(1) \equiv 1 \pmod 5$.

(d) Find $p_2(x)$ where $p_2(-1) \equiv p_2(0) \equiv p_2(1) \equiv 0 \pmod 5$ and $p_2(2) \equiv 1 \pmod 5$.

(e) Construct $p(x)$ using a linear combination of $p_{-1}(x)$, $p_0(x)$, $p_1(x)$ and $p_2(x)$.

Solution:

(a)

$$\begin{aligned} p_{-1}(x) &\equiv x(x-1)(x-2)\big((-1)(-1-1)(-1-2)\big)^{-1} \equiv x(x-1)(x-2)(-6)^{-1} \\ &\equiv 4x(x-1)(x-2) \equiv x(x-1)(x-2)(-6)^{-1} \equiv 4x(x-1)(x-2) \pmod 5 \end{aligned}$$

(b)

$$\begin{aligned} p_0(x) &\equiv (x+1)(x-1)(x-2)\big((1)(-1)(-2)\big)^{-1} \equiv 3(x+1)(x-1)(x-2) \\ &\equiv (x+1)(x-1)(x-2)(2)^{-1} \equiv 3(x+1)(x-1)(x-2) \pmod 5 \end{aligned}$$

(c)

$$\begin{aligned} p_1(x) &\equiv (x+1)(x)(x-2)\big((2)(1)(-1)\big)^{-1} \\ &\equiv 2(x+1)(x)(x-2) \equiv (x+1)(x)(x-2)(-2)^{-1} \equiv 2(x+1)(x)(x-2) \pmod 5 \end{aligned}$$

(d) $p_2(x) \equiv (x+1)(x)(x-1)(6)^{-1} \equiv (x+1)(x)(x-2) \pmod 5$.

(e) We don't need $p_2(x)$.
$p(x) \equiv 3 \cdot p_{-1}(x) + 1 \cdot p_0(x) + 2 \cdot p_1(x) + 0 \cdot p_2(x) \equiv 4x^3 + 4x^2 + 3x + 1 \pmod 5$.

3 Secrets in the United Nations

A vault in the United Nations can be opened with a secret combination $s \in \mathbb{Z}$. In only two situations should this vault be opened: (i) all 193 member countries must agree, or (ii) at least 55 countries, plus the U.N. Secretary-General, must agree.

(a) Propose a scheme that gives private information to the Secretary-General and all 193 member countries so that the secret combination s can only be recovered under either one of the two specified conditions.

(b) The General Assembly of the UN decides to add an extra level of security: each of the 193 member countries has a delegation of 12 representatives, all of whom must agree in order for that country to help open the vault. Propose a scheme that adds this new feature. The scheme should give private information to the Secretary-General and to each representative of each country.

Solution:

(a) Create a polynomial of degree 192 and give each country one point. Give the Secretary General $193 - 55 = 138$ points, so that if she collaborates with 55 countries, they will have a total of 193 points and can reconstruct the polynomial. Without the Secretary-General, the polynomial can still be recovered if all 193 countries come together. (We do all our work in GF(p) where $p \geq d+1$).

Alternatively, we could have one scheme for condition (i) and another for (ii). The first condition is the secret-sharing setup we discussed in the notes, so a single polynomial of degree 192 suffices, with each country receiving one point, and evaluation at zero returning the combination s. For the second condition, create a polynomial f of degree 1 with $f(0) = s$, and give $f(1)$ to the Secretary-General. Now create a second polynomial g of degree 54, with $g(0) = f(2)$, and give one point of g to each country. This way any 55 countries can recover $g(0) = f(2)$, and then can consult with the Secretary-General to recover $s = f(0)$ from $f(1)$ and $f(2)$.

(b) We'll layer an *additional* round of secret-sharing onto the scheme from part (a). If t_i is the key given to the ith country, produce a degree-11 polynomial f_i so that $f_i(0) = t_i$, and give one point of f_i to each of the 12 delegates. Do the same for each country (using different f_i each time, of course).

4 To The Moon!

A secret number s is required to launch a rocket, and Alice distributed the values $(1, p(1)), (2, p(2)), \ldots, (n+1, p(n+1))$ of a degree n polynomial p to a group of \$GME holders $\text{Bob}_1, \ldots, \text{Bob}_{n+1}$. As usual, she chose p such that $p(0) = s$. Bob_1 through Bob_{n+1} now gather to jointly discover the secret. However, Bob_1 is secretly a partner at Melvin Capital and already knows s, and wants to sabotage $\text{Bob}_2, \ldots, \text{Bob}_{n+1}$, making them believe that the secret is in fact some fixed $s' \neq s$. How could he achieve this? In other words, what value should he report (in terms variables known in the problem, such as s', s or y_1) in order to make the others believe that the secret is s'?

Solution:

We know that in order to discover s, the Bobs would compute

$$s = y_1 \Delta_1(0) + \sum_{k=2}^{n+1} y_k \Delta_k(0), \tag{1}$$

where $y_i = p(i)$. Bob_1 now wants to change his value y_1 to some y_1', so that

$$s' = y_1' \Delta_1(0) + \sum_{k=2}^{n+1} y_k \Delta_k(0). \tag{2}$$

Subtracting Equation 1 from 2 and solving for y_1', we see that

$$y_1' = (\Delta_1(0))^{-1} \left(s' - s\right) + y_1,$$

where $(\Delta_1(0))^{-1}$ exists, because $\deg \Delta_1(x) = n$ with its n roots at $2, \ldots, n+1$ (so $\Delta_1(0) \neq 0$).