6 Algorithmic Journeys with Concepts

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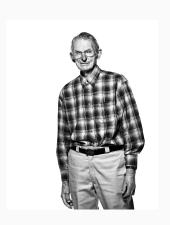
Rails Reactor / Giphy

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The Software Industry is Not Industrialized

Software components (routines), to be widely applicable to different machines and users, should be available in families arranged according to precision, robustness, generality and time-space performance.



A Familiar Example. Douglas McIlroy about sin

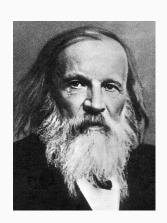
Dimensions along which we wish to have variablity:

- 1. precision, for which perhaps ten different approximating functionsmight suffice
- 2. floating vs fixed computation
- 3. argument ranges $[0, \pi/2]$, [0, 2pi], also $[-\pi/2, pi/2]$, $[-\pi, \pi]$, [-big, +big]
- robustness ranging from no argument validation through signaling of complete loss of significance, to signaling of specified range violations

Periodic Table



Figure 1: A Russian periodic table based on Dmitri Mendeleyev's original table of 1869.



Species Plantarum

Lists every species of plant known at the time, classified into genera. It is the first work to consistently apply binomial names and was the starting point for the naming of plants.



Elements

- 1. Definitions
- 2. Postulates
- 3. Common notions



Common Notions

- 1. Things which are equal to the same thing are also equal to one other.
- 2. If equals be added to equals, the wholes are equal.
- 3. If equals be subructed from equals, the remainders are equal.
- 4. Things which coincide with one another are equal to one another.
- 5. The whole is greater than the part.

Basic idea

The essence of generic programming lies in the idea of concepts. A concept is a way of describing a family of related object types.

Natural	Mathematics	Programming	Programming
Science			Examples
genus	theory	concept	Integral, Character
species	model	type or class	uint8_t, char
individual	element	instance	01000001(65, 'A')

Definitions

- 1. Datum
- 2. Value
- 3. Value type
- 4. Object
- 5. Object type

Datum

Definition A datum is a sequence of bits.

Example 01000001 is an example of a datum.

Value

Definition

A value is a datum together with its interpretation.

Example

The datum 01000001 might have the interpretation of the integer 65, or the character "A".

Explanation

Every **value** must be associated with a **datum** in memory; there is no way to refer to disembodied **values** in modern programming languages.

Value type

Definition A **value type** is a set of values sharing a common interpretation.

Object

Definition

An **object** is a collection of bits in memory that contain a **value** of a given **value type**.

Explanation

An object is immutable if the value never changes, and mutable otherwise. An object is unrestricted if it can contain any value of its value type.

Object type

DefinitionAn **object type** is a uniform method of storing and retrieving **values** of a given **value type** from a particular **object** when given its address.

Programming with concepts

Regular type

Operation

- 1. Copy construction
- 2. Assignment
- 3. Equality
- 4. Destruction

Semantic

$$\forall a \ \forall b \ \forall c : T \ a(b) \implies (b = c \implies a = c)$$
 $\forall a \ \forall b \ \forall c : a \leftarrow b \implies (b = c \implies a = c)$
 $\forall f \in Regular Function : a = b \implies f(a) = f(b)$

Semiregular type

Operation

- 1. Copy construction
- 2. Assignment
- 3. Destruction

Semantic

$$\forall a \ \forall b \ \forall c : T \ a(b) \implies (b = c \implies a = c)$$

 $\forall a \ \forall b \ \forall c : a \leftarrow b \implies (b = c \implies a = c)$
 $\forall f \in Regular Function : a = b \implies f(a) = f(b)$

```
\begin{tabular}{ll} Readable(T) &\triangleq Regular(T) \land \\ & ValueType : Readable \to Regular \land \\ & source : T \to ValueType(T) \land \\ \end{tabular}
```

```
Writable(T) \triangleq Regular(T) \land
ValueType : Writable \rightarrow Regular \land
(\forall x \in T)(\forall v \in ValueType(T)) sink(x) \leftarrow v
is \ a \ well - formed \ statement
source : T \rightarrow ValueType(T) \land
```

```
 \begin{tabular}{ll} \textit{Iterator}(T) & \triangleq & \textit{Regular}(T) \land \\ & \textit{DistanceType}: \textit{Iterator} \rightarrow \textit{Integer} \land \\ & \textit{successor}: T \rightarrow T \land \\ & \textit{successor} \ \textit{is not necessarily} - \textit{regular} \\ \end{tabular}
```

 $ForwardIterator(T) \triangleq Iterator(T) \land regular_unary_function(successor)$

```
BidirectionalIterator(T) \triangleq ForwardIterator(T) \land
                                  predecessor: T \rightarrow T \land
                                  predecessor takes constant type \wedge
                                  (\forall i \in T)successor(i)isdefined \Longrightarrow
                                     predecessor(successor(i)) is defined
                                     and equals to i \wedge i
                                  (\forall i \in T) predecessor(i) is defined \implies
                                     successor(predecessor(i)) is defined
                                     and equals to i
```

```
IndexedIterator(T) \land
TotallyOrdered(T) \land
(\forall i, j \in T)i < j \iff i \prec j \land
DifferenceType:
   RandomAccessIterator \rightarrow Integer \land
+: T \times DifferenceType(T) \rightarrow T \wedge
-: T \times DifferenceType(T) \rightarrow T \wedge
-: T \times T \rightarrow DifferenceType(T) \wedge
< takes constant time \land

    between and iterator and an integer

   takes constant time
```

 $RandomAccessIterator(T) \triangleq BidirectionalIterator(T) \land$

FunctionalProcedure(F) \triangleq F is a regular procedure defined on regular types : replacing its inputs with equal objects results in equal output objects.

```
\label{eq:UnaryFunction} UnaryFunction(F) \triangleq FunctionalProcedure(F) \land Arity(F) = 1 \\ \land Domain: UnaryFunction \rightarrow Regular \\ F \mapsto InputType(F,0)
```

$$\label{eq:homogeneousFunction} \begin{split} \text{HomogeneousFunction}(F) &\triangleq \text{FunctionalProcedure}(F) \land \text{Arity}(F) > 0 \\ \land (\forall i,j \in \mathbb{N})(i,j < \text{Arity}(F)) &\Longrightarrow (\text{InputType}(F,i) = \text{InputType}(F,j)) \\ \land \text{Domain}: \text{HomogeneousFunction} \rightarrow \text{Regular} \\ F &\Longrightarrow \text{InputType}(F,0) \end{split}$$

 $Predicate(P) \triangleq Functional Procedure(F) \land Codomain(P) = bool$

 $Homogeneous Predicate(P) \triangleq Predicate(P) \land Homogeneous Function(P)$

 $Relation(R) \triangleq HomogeneousPredicate(R) \land Arity(R) = 2$

 $TotallyOrdered(T) \triangleq Regular(T) \land <: T \times T \rightarrow bool \land total_ordering(<)$

property(R:Relation) $total_ordering:R$ $r\mapsto transitive(r) \ \land (\forall a,b \in Domain(R)) \ exactly \ one \ of \ following \ holds:$ $r(a,b), \ r(b,a), \ or \ a=b$

Journey 1

- 1. min
- 2. max

```
int min(int x, int y) {
    if (y < x) {
        return y;
    }
    return x;
}</pre>
```

```
int min(int x, int y) {
    if (y < x) {
        return y;
    return x;
double min(double x, double y) {
    if (y < x) {
        return y;
    return x;
```

```
template < typename T>
T min(T x, T y) {
    if (y < x) {
        return y;
    }
    return x;
}</pre>
```

Journey #1

```
Dealing with large objects
template < typename T >
const T& min(const T& x, const T& y) {
    if (y < x) {
        return y;
    }
    return x;
}</pre>
```

```
template < typename T, typename P>
const T& min(const T& x, const T& y, P pred) {
    if (pred(y, x)) {
        return y;
    }
    return x;
}
```

```
struct employee {
    std::string full_name;
    int64 t salary:
};
void usage() {
  employee e0{"Bjarne_Stroustrup", 999999911};
  employee e1{"Alex_Stepanov", 999999911};
  min(e0, e1, [](const auto& x, const auto& y) {
    return x.salary < y.salary;</pre>
  }).salarv += 10000ll;
```

```
template < typename T, typename P>
T& min(T& x, T& y, P pred) {
    if (pred(y, x)) {
        return y;
    }
    return x;
}
```

```
template < Regular T, Relation r>
const T& min(const T& x, const T& y, Relation r) {
    if (r(v, x)) { return v: }
    return x;
template < Regular T, Relation r>
T\& min(T\& x, T\& y, Relation r) {
    if (r(y, x)) { return y; }
    return x;
```

```
template < TotallyOrdered T >
const T& min(const T& x, const T& y) {
    return min(x, y, std::less < T > ());
}

template < TotallyOrdered T >
T& min(T& x, T& y) {
    return min(x, y, std::less < T > ());
}
```

```
namespace cppcon {
template < totally_ordered T>
const T& min(const T& x, const T& y) {
    if (y < x) {
        return y;
    return x;
template < totally_ordered T>
T\& min(T\& x, T\& y) {
    if (y < x)
        return y;
    return x;
```

- 1. unique
- 2. unique_count

```
template < forward_iterator It , relation < ValueType (It) > R>
It unique(It first, It last, R r) {
    if (first == last) { return last; }
    It result = first; ++ first;
    while (first != last) {
        if (*result == *first) {
           ++ first:
        } else {
            ++result:
            *result = *first:
            ++ first:
    ++result:
    return result;
```

- 1. frequencies
- 2. transform_subgroups
- 3. squash_subgroups

- 1. split
- 2. transform_splits

- 1. remove_if
- 2. partition_semistable

More examples of concepts

- 1. Regular Type
- 2. Semiegular Type
- 3. Functional Procedure
- 4. Homogeneous Function
- 5. Homogeneous Predicate
- 6. Semiring
- 7. Sequence
- 8. Totally Ordered
- 9. Input Iterator
- 10. Forfward Iterator
- 11. Bidirectional Iterator

Properties

- 1. Associative
- 2. Distributive
- 3. Transitive
- 4. Semiegular Type
- 5. Functional Procedure

Techniques

- 1. Transformation-action duality
- 2. Operation-accumulation procedure duality
- 3. Memory adaptivity
- 4. Reduction to constrained subproblem

Conclusion

Conclusion

- 1. Concreteness costs
- 2. Abstracting algorithms to their most general setting without losing efficiency
- 3. Know your algorithms
- 4. If y