The Fast File System

M.K. McKusick, W. Joy, S.L. Leffler, R.S. Fabry, "A Fast File System for UNIX", ACM Transactions on Computer Systems, Vol. 2, No. 3, Aug. 1984

- Recall: Limitations of Original UNIX File System
- Fast File System (FFS)
 - Increase Block Size
 - Make File System "disk-aware"
 - Some Functional Improvements

Recall: Problems with Original Unix File System

- Original file system uses 512/1024-byte blocks.
- Inodes kept separately from data, causing long seeks to access data.
- Inodes of files in a common directory not kept together, causing low performance when searching directories.
- Data blocks of a file are not stored together.
- Free list quickly scrambles, increasing overhead of finding free blocks.

Original file system treats disk as Random Access Device!

"Fast FS" (FFS, ca. 1984)

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Three-pronged approach:

- Increase block size
- 2. Make file system disk-aware
- 3. Functional improvements

FFS: Increase Block Size

Increase block size from 512 byte to at least 4096 byte.

- Block size power-of-two multiple of 4096 byte
- Defined during file system creation
- Stored as parameter in superblock

FFS: Block Size and Fragmentation

Observation: Typical UNIX files are small compared to large blocks.

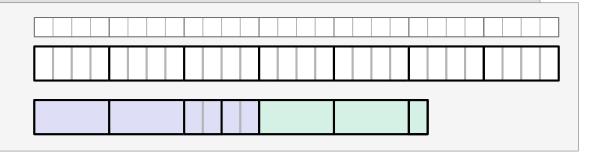
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	Table I.	Amount of Wasted Space as a Function of Block Size		
Space used (Mbytes)	i	% Waste	Organization	
775.2		0.0	Data only, no separation between files	
807.8		4.2	Data only, each file starts on 512-byte boundary	
828.7		6.9	Data + inodes, 512-byte block UNIX file system	
866.5		11.8	Data + inodes, 1024-byte block UNIX file system	
948.5		22.4	Data + inodes, 2048-byte block UNIX file system	
1128.3		45.6	Data + inodes, 4096-byte block UNIX file system	

Blocks vs. Fragments

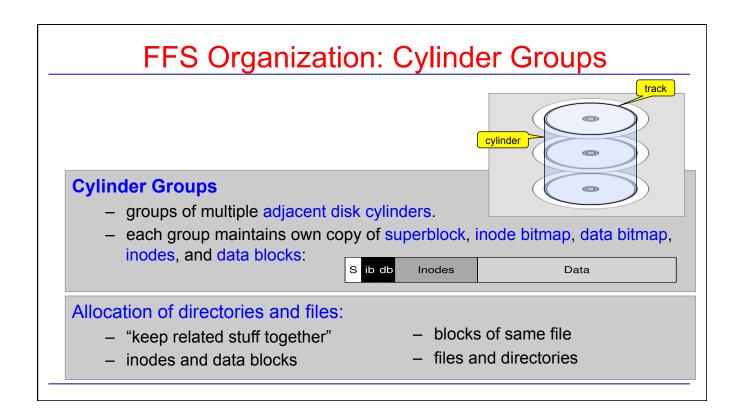
Solution: Split blocks into Fragments

- Fragment size defined at file system creation time
- Example: 4096/1024 has 4kB blocks and 1kB fragments.



FFS Organization: Making FS Disk-Aware

- 1. Cylinder Groups
- 2. File System Parameterization



FFS: Minimizing Rotational Latency

File System Parameterization

Goal: Parameterize **processor capabilities** and **disk characteristics** so that blocks can be **allocated** in an optimal, configuration-dependent way.

Allocate new block rotationally well-positioned.

Disk Parameters:

number of blocks per **track** disk **spin rate**.





CPU Parameters:

expected time to service interrupt and schedule new disk transfer

FFS: Performance

Table IIa. Reading Rates of the Old and New UNIX File Systems

Type of file system	Processor and bus measured	Speed (Kbytes/s)	Read bandwidth %	% CPU
Old 1024	750/UNIBUS	29	29/983 3	11
New 4096/1024	750/UNIBUS	221	221/983 22	43
New 8192/1024	750/UNIBUS	233	233/983 24	29
New 4096/1024	750/MASSBUS	466	466/983 47	73
New 8192/1024	750/MASSBUS	466	466/983 47	54

Table IIb. Writing Rates of the Old and New UNIX File Systems

Type of file system	Processor and bus measured	Speed (Kbytes/s)	Write bandwidth %	% CPU
Old 1024	750/UNIBUS	48	48/983 5	29
New 4096/1024	750/UNIBUS	142	142/983 14	43
New 8192/1024	750/UNIBUS	215	215/983 22	46
New 4096/1024	750/MASSBUS	323	323/983 33	94
New 8192/1024	750/MASSBUS	466	466/983 47	95

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FFS: Functional Enhancements

- 1. Arbitrary-length File Names
- 2. File Locking
- 3. Symbolic Links
- 4. "Rename" system call added
- 5. Quotas

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