# GEH1036/GEK1505 Tutorial 5(week 7)

Tsien Lilong February 27, 2017

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    Go to the website, and scroll down, see the **Work** entry, in the "Tutor" project, click the link for more information. I will provide this slide and the source .tex file on that site.
- Course information
  - EVERY WEEK MONDAY 10:00-11:00 \$16-0431

#### I. Trees with I vertices

#### Question

Draw all the trees with 7 unlabelled vertices.

We draw them according to the maximal degree (denote by  $\Delta$ ) among all the vertices. Note that the number of the edges is 6, then the sum of the degrees is  $2\times 6=12$ .

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 $\Delta=6$ : the distribution of degrees of the 7 vertices is 6+1+1+1+1+1.

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 $\Delta=5$ : the distribution of degrees of the 7 vertices is 5+2+1+1+1+1+1.

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 $\Delta=4$ : the distribution of degrees of the 7 vertices is 4+2+2+1+1+1+1 or 4+3+1+1+1.



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 $\Delta=3$ : the distribution of degrees of the 7 vertices is 3+3+2+1+1+1+1 or 3+2+2+1+1+1.



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We draw them according to the maximal degree (denote by  $\Delta$ ) among all the vertices. Note that the number of the edges is 6, then the sum of the degrees is  $2\times 6=12$ .

 $\Delta=2$ : the distribution of degrees of the 7 vertices is 2+2+2+2+1+1.

$$\Delta = 2$$
:

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Let n be the number of vertices. Since a tree must have n-1 edges, by the Degree Theorem,

$$5 + 4 \times 2 + (n - 5) = 2(n - 1) \Rightarrow n = 10.$$

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An example of this graph is,



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Draw two different trees (with unlabeled vertices) that satisfy the properties given in the previous Question.

 First draw one vertex 5 degree, other 5 vertices are connected to it.



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- Choose 4 vertices to be of degree 2 from the five vertices.



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- Or choose 3.



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- Using the **Degree Theorem**, the sum of all degrees is

$$2\times 4 = \sum_{i\in V(G)} d(v) \geqslant 2\times 3 + 3 = 9$$

(every vertex in a tree must has at least 1 degree), which is impossible.



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# 5. Spanning Tree

#### Question

Let T be a spanning tree in a graph G. Suppose G contains a cycle C and that ab is an edge of C that is also in T. Show that there is always at least one edge uv in C that can replace ab in T so that the resulting graph is still a spanning tree of G.

• First consider a example.



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- Delete ab from T, the resulting graph, denoted by T-ab, consists of two disjoint trees  $T_1$  and  $T_2$ .



## 5. Spanning Tree

#### Question

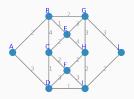
Let T be a spanning tree in a graph G. Suppose G contains a cycle C and that ab is an edge of C that is also in T. Show that there is always at least one edge uv in C that can replace ab in T so that the resulting graph is still a spanning tree of G.

- First consider a example.
- Delete ab from T, the resulting graph, denoted by T-ab, consists of two disjoint trees  $T_1$  and  $T_2$ .
- Since C is a cycle containing the edge ab, there will be an edge uv of C different from ab that joins one vertex of  $T_1$  and one vertex of  $T_2$ .

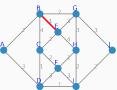


# Question

Find a minimal spanning tree and its weight in the following graph with Prim's Algorithm.

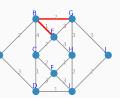


 $lack \$  Using the Prim's algorithm, chose a minimal weighted edge, BE for example,  $S=\{B,E\}.$ 

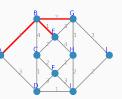


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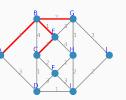
• Consider all edges connected to set S, choose one mimimal wighted edge, BG,  $S = \{B, E, G\}$ .



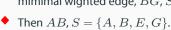
- Using the Prim's algorithm, chose a minimal weighted edge, BE for example,  $S=\{B,E\}$ .
- Consider all edges connected to set S, choose one mimimal wighted edge, BG,  $S = \{B, E, G\}$ .
- Then AB,  $S = \{A, B, E, G\}$ .



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- Consider all edges connected to set S, choose one mimimal wighted edge, BG,  $S = \{B, E, G\}$ .
- Then  $AB, S = \{A, B, E, G\}$ .
- Then EC,  $S = \{A, B, C, E, G\}$ .

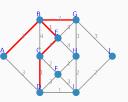


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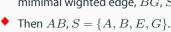


• Then 
$$EC$$
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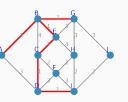
• Then 
$$CD$$
,  $S = \{A, B, C, D, E, G\}$ .



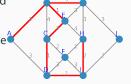
- Using the Prim's algorithm, chose a minimal weighted edge, BE for example,  $S=\{B,E\}$ .
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- Then EC,  $S = \{A, B, C, E, G\}$ .
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- Then DI,  $S = \{A, B, C, D, E, G, I\}$ .

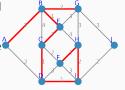


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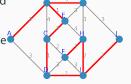
- Then AB,  $S = \{A, B, E, G\}$ .
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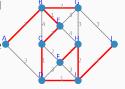
- Then  $AB, S = \{A, B, E, G\}$ .
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- Consider all edges connected to set S, choose one mimimal wighted edge, BG,  $S = \{B, E, G\}$ .



- Then  $AB, S = \{A, B, E, G\}.$
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- ♦ The final tree is AB, BE, BG, EC, CD, DI, IH, HF, IJ. The weight is therefore 2+1+2+2+1+1+2+1+2=14.

## 7 Vulnerable tree

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#### Question

Find a minimal set of edges that can be added to each of the following trees so that in the resulting graph, every edge belongs to some cycle.

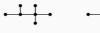


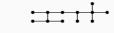
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 Since every vertex in a cycle cannot have degree 1.





Note the number of them as n, then we have at least  $\left[\frac{n}{2}\right]$  edges to connect them.

### 7. Vulnerable tree

### Question

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- Since every vertex in a cycle cannot have degree 1. Note the number of them as n, then we have at least  $\left[\frac{n}{2}\right]$  edges to connect them.
- The case with minimal edges is to link those one-degree vertices pairly. And can be checked that every edge is contained in a cycle.

#### 8. End

Thank you for listening!