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A Sparse Quasi-Newton Method Based on Automatic Differentiation for Solving Unconstrained Optimization Problems

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Abstract: In our paper, we introduce a sparse and symmetric matrix completion quasi-Newton model using automatic differentiation, for solving unconstrained optimization problems where the sparse structure of the Hessian is available. The proposed method is a kind of matrix completion quasi-Newton method and has some nice properties. Moreover, the presented method keeps the sparsity of the Hessian exactly and satisfies the quasi-Newton equation approximately. Under the usual assumptions, local and superlinear convergence are established. We tested the performance of the method, showing that the new method is effective and superior to matrix completion quasi-Newton updating with the Broyden–Fletcher–Goldfarb–Shanno (BFGS) method and the limited-memory BFGS method.

Keywords: symmetric quasi-Newton method; unconstrained optimization problems; matrix completion; automatic differentiation; superlinear convergence; Broyden–Fletcher–Goldfarb–Shanno method

MSC: 65K05; 90C06; 90C53



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1. Introduction

We concentrated on the unconstrained optimization problem

$$\min f(x), \quad x \in \mathbb{R}^n, \tag{1}$$

where $f: \mathbb{R}^n \to \mathbb{R}$ is a twice continuously differentiable function; and $\nabla f(x)$ and $\nabla^2 f(x)$ denote the gradient and Hessian of f at x, respectively. The first order necessary condition of (1) is

$$\nabla f(x) = 0$$
,

which can be written as the symmetric nonlinear equations

$$F(x) = 0, (2)$$

where $F: \mathbb{R}^n \to \mathbb{R}^n$ is a continuously differentiable mapping and the symmetry implies that the Jacobian F'(x) satisfies $F'(x) = F'(x)^T$. That symmetric nonlinear system has close relationships with many practical problems, such as the gradient mapping of unconstrained optimization problems, the Karush–Kuhn–Tuckrt (KKT) system of equality constrained optimization problem, the discretized two-point boundary value problem, and the saddle point problem (2) [1–5].

For small or medium-scale problems, classical quasi-Newton methods enjoy superlinear convergence without the calculation of the Hessian [6,7]. Let x_k be the current Symmetry **2021**, 13, 2093 2 of 21

iterative point and B_k be the symmetric approximation of the Hessian; then the iteration $\{x_k\}$ generated by quasi-Newton methods is

$$x_{k+1} = x_k + \alpha_k d_k \nabla f(x_k),$$

where $\alpha_k > 0$ is a step length obtained by some line search or other strategies. The search direction d_k can be gotten by solving the equations

$$B_k d_k + \nabla f(x_k) = 0$$
,

where the quasi-Newton matrix B_k is an approximation of $\nabla^2 f(x_k)$ and satisfies the secant condition:

$$B_{k+1}s_k = y_k$$

where $s_k = x_{k+1} - x_k$, $y_k = \nabla f(x_{k+1}) - \nabla f(x_k)$. The matrix B_k can be updated by different update formulae. The Davidon–Fletcher–Powell (DFP) update,

$$B_{k+1} = \left(I - \frac{y_k s_k^T}{y_k^T s_k}\right) B_k \left(I - \frac{y_k s_k^T}{y_k^T s_k}\right) + \frac{y_k y_k^T}{y_k^T s_k}$$

$$= B_k + \frac{(y_k - B_k s_k) s_k^T + s_k (y_k - B_k s_k)^T}{y_k^T s_k}$$

$$- \frac{(y_k - B_k s_k)^T s_k}{(y_k^T s_k)^2} y_k y_k^T$$

$$= \left(I - \frac{y_k s_k^T}{y_k^T s_k}\right) B_k \left(I - \frac{y_k s_k^T}{y_k^T s_k}\right) + \frac{y_k y_k^T}{y_k^T s_k},$$

was first proposed by Davidon [8] and developed by Fletcher and Powell [9]. The Broyden–Fletcher–Goldfard–Shanno (BFGS) update,

$$B_{k+1} = B_k - \frac{B_k s_k s_k^T B_k}{s_k^T B_k s_k} + \frac{y_k y_k^T}{y_k^T s_k},$$

was proposed independently by Broyden [10], Fletcher [11], Goldfarb [12], and Shanno [13]. One can find more on the topic in references [14–17].

If we assume that $H_k = B_k^{-1}$, then using Sherman–Morrison formula, we have the Broyden's family update:

$$H_{k+1} = H_k - \frac{H_k y_k y_k^T H_k}{y_k^T H_k y_k} + \frac{s_k s_k^T}{s_k^T y_k} + \phi_k v_k v_k^T,$$

where $\phi_k \in [0,1]$ is

$$\phi_k = \sqrt{y_k^T H_k y_k} \left(\frac{s_k}{s_k^T y_k} - \frac{H_k y_k}{y_k^T H_k y_k} \right),$$

When $\phi_k \equiv 1$, we have a BFGS update. When $\phi_k \equiv 0$, we have a DFP update.

However, a quasi-Newton method is not desirable when applied to solve large-scale problems, because we need to store the full matrx B_k . To overcome such drawback, the so-called sparse quasi-Newton methods [14] have received much attention. Early in 1970, Schubert [18] has proposed a sparse Broyden's rank one method. Then Powell and Toint [19], Toint [20] studied the sparse quasi-Newton method.

Existing sparse quasi-Newton methods usually use a sparse symmetric matrix as an approximation of the Hessian so that both matrices take the same form or have similar structures. If the limited memory technique [21,22] is adopted, which only stores several pairs (s_k, y_k) to construct a matrix H_k by updating the initial matrix H_0 m times, the method

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can be widely used in practical optimization problems. On the other hand, there are many large-scale problems in scientific fields take the partially separable form

$$f(x) = \sum_{i=1}^{m} f_i(x),$$

where function f_i , $i=1,\ldots,m$ is related to a few variables. For the partially separable unconstrained optimization problems, the partitioned BFGS method [23,24] was proposed and has better performance in practice. The partitioned BFGS method updates each matrix B_k^i of each element function $f_i(x)$ separately via BFGS updating and sums these matrices to construct the next quasi-Newton matrix B_{k+1} . Since the size of x in $f_i(x)$ is smaller than that of n, the matrix B_{k+1}^i will be a small matrix, and then the matrix B_{k+1} will be sparse. The quasi-Newton direction is the solution of the linear equations:

$$\left(\sum_{k=1}^m B_k^i\right) d_k = -\nabla f(x_k).$$

However, the partitioned BFGS method cannot always preserve the positive definiteness of the matrix B_k , only if that each element function $f_i(x)$ is convex, so the partitioned BFGS method is implemented with the trust region strategy [25]. Recently, for the partially separable nonlinear equations, Cao and Li [26] have introduced two kinds of partitioned quasi-Newton methods and given their global and superlinear convergence.

Another efficient sparse quasi-Newton method is designed to exploit the sparsity structures of the Hessian. We assume that for all $x \in \mathbb{R}^n$,

$$(\nabla^2 f(x))_{i,j} = 0, (i,j) \in F,$$

where $F \subseteq \{1, ..., n\} \times \{1, ..., n\}$. References [27,28] have proposed sparse quasi-Newton methods, where H_{k+1} satisfies the secant equation

$$H_{k+1}y_k = s_k$$

and sparse condition

$$(H_{k+1})_{ij}=0, \quad (i,j)\in F$$

simultaneously, where H_{k+1} is an approximate inverse Hessian. Recently, Yamashita [29] proposed another type of matrix completion quasi-Newton (MCQN) update for solving problem (1) with a sparse Hessian and proved the local and superlinear convergence for MCQN updates with the DFP method. Reference [30] established the convergence of MCQN updates with all of Broyden's convex family. However, global convergence analysis [31] was presented for two-dimensional functions with uniformly positive definite Hessians.

Another kind of quasi-Newton method for solving large scale unconstrained optimization problems is the diagonal quasi-Newton method, where the Hessian of an objective function is approximated by a diagonal matrix with positive elements. The first version was developed by Nazareth [32], where the quasi-Newton matrix satisfies the least change and weak secant condition [33]:

min
$$||H_{k+1} - H_k||$$

s.t. $y_k^T H_{k+1} y_k = y_k^T s_k$, (3)

where $\|\cdot\|_F$ is the standard Frobenius norm. Recently, Andrei N. [34] developed a diagonal quasi-Newton method, where the diagonal elements satisfy the least change weak secant condition (3) and minimize the trace of the update. Besides, lots of other techniques, such as forward and central finite differences, the variational principle with a weighted norm, and the generalized Frobenius norm, can be used to derive different kinds of diagonal quasi-

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Newton method [35–37]. Under usual assumptions, the diagonal quasi-Newton method is linearly convergent. The authors of [38] adopted a similar technique to derivation with the DFP method and got a low memory diagonal quasi-Newton method. Using the Armijo line search, they established the global convergence and gave the sufficient conditions for the method to be superlinearly convergent.

The main contribution of our paper is to propose a sparse quasi-Newton algorithm based on automatic differentiation for solving (1). Firstly, similarly to the derivation of BFGS update, we can perform a symmetric rank-two quasi-Newton update:

$$B_{k+1} = B_k - \frac{B_k \sigma_k \sigma_k^T B_k}{\sigma_k^T B_k \sigma_k} + \frac{\nabla^2 f(x_{k+1}) \sigma_k \sigma_k^T \nabla^2 f(x_{k+1})}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k},\tag{4}$$

where $\sigma_k \in \mathbb{R}^n$ and B_{k+1} satisfying the adjoint tangent condition [39]

$$\sigma_k^T B_{k+1} = \sigma_k^T \nabla^2 f(x_{k+1}).$$

For an $n \times n$ matrix, we denote $A \succeq 0$, as A is positive definite. Then, when $B_k \succeq 0$, $B_{k+1} \succeq 0$ if and only if $\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k > 0$, which means that the proposed update (4) keeps the positive definiteness, as in BFGS updating. Moreover, when B_0 is positive definite, the matrices $\{B_k\}$ updated by the proposed update (4) are positive definite for solving (1) with uniformly positive definite Hessians. In our work, we pay attention to $\sigma_k = s_k$; then the proposed rank-two quasi-Newton update (4) method satisfies

$$B_{k+1}s_k = \nabla^2 f(x_{k+1})s_k,$$

which means that B_{k+1} equals $\nabla^2 f(x_{k+1})$ in the direction s_k exactly. Several lemmas have been given to present the properties of the proposed rank-two quasi-Newton update formula. Secondly, combined with the idea of MCQN method [29], we propose a sparse and symmetric quasi-Newton algorithm for solving (1). Under appropriate conditions, local and superlinear convergence are established. Finally, our numerical results illustrate that the proposed algorithm has satisfying performance.

The paper is organized as follows. In Section 2, we introduce a symmetric rank-two quasi-Newton update based on automatic differentiation and prove several nice properties. In Section 3, by using the idea of matrix completion, we present a sparse quasi-Newton algorithm and show some nice properties. In Section 4, we prove the local and superlinear convergence of the algorithm proposed in Section 3. Numerical results are listed in Section 5, which verify that the proposed algorithm is very encouraging. Finally, we give the conclusion.

2. A New Symmetric Rank-Two Quasi-Newton Update

Similarly to the derivation of BFGS update, we will derive a new symmetric rank-two quasi-Newton update and show several lemmas. Let

$$B_{k+1} = B_k + \Delta_k$$

where Δ_k is a rank-two matrix and B_{k+1} satisfies the condition

$$\sigma_k^T B_{k+1} = \sigma_k^T \nabla^2 f(x_{k+1}), \tag{5}$$

where $\sigma_k \in \mathbb{R}^n$ and $\sigma_k \neq 0$. Similarly to the derivation of BFGS, we have the following symmetric rank-two update:

$$B_{k+1} = B_k - \frac{B_k \sigma_k \sigma_k^T B_k}{\sigma_k^T B_k \sigma_k} + \frac{\nabla^2 f(x_{k+1}) \sigma_k \sigma_k^T \nabla^2 f(x_{k+1})}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k}. \tag{6}$$

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If we denote $H_k = B_k^{-1}$ and $H_{k+1} = B_{k+1}^{-1}$, then (6) can be expressed as

$$H_{k+1} = H_{k} - \frac{H_{k} \nabla^{2} f(x_{k+1}) \sigma_{k} \sigma_{k}^{T} + \sigma_{k} \sigma_{k}^{T} \nabla^{2} f(x_{k+1}) H_{k}}{\sigma_{k}^{T} \nabla^{2} f(x_{k+1}) \sigma_{k}} + \left(1 + \frac{\sigma_{k}^{T} \nabla^{2} f(x_{k+1}) H_{k} \nabla^{2} f(x_{k+1}) \sigma_{k}}{\sigma_{k}^{T} \nabla^{2} f(x_{k+1}) \sigma_{k}}\right) \cdot \frac{\sigma_{k} \sigma_{k}^{T}}{\sigma_{k}^{T} \nabla^{2} f(x_{k+1}) \sigma_{k}}.$$
 (7)

It can be seen that the update (6) involves the Hessian $\nabla^2 f(x)$, but we do not need to compute them in practice. For given vectors x, s, and σ , we can get $\nabla^2 f(x)s$ and $\sigma^T \nabla^2 f(x)$ exactly by the forward and reverse mode of automatic differentiation.

Next, several lemmas are presented.

Lemma 1. We suppose that $B_k \succeq 0$ and B_{k+1} is updated by (6); then $B_{k+1} \succeq 0$ if and only if $\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k > 0.$

Proof. According to the condition (5), one has

$$\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k = \sigma_k^T B_{k+1} \sigma_k.$$

If B_{k+1} is positive definite, one has $\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k > 0$. Let $\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k > 0$ and $B_k \succeq 0$. Then for $\forall d_k \in R^n$, $d_k \neq 0$, it can be derived from (6) that

$$d_k^T B_{k+1} d_k = d_k^T B_k d_k - \frac{(d_k^T B_k \sigma_k)^2}{\sigma_k^T B_k \sigma_k} + \frac{(d_k^T \nabla^2 f(x_{k+1}) \sigma_k)^2}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k}.$$

According to that $B_k \succeq 0$, there is a symmetric matrix $B_k^{1/2} \succeq 0$, such that $B_k = B_k^{1/2} B_k^{1/2}$. Then we have from Cauchy-Schwarz inequality that

$$(d_k^T B_k \sigma_k)^2 = \left((B_k^{1/2} d_k)^T (B_k^{1/2} \sigma_k) \right)^2$$

$$\leq \|B_k^{1/2} d_k\|^2 \cdot \|B_k^{1/2} \sigma_k\|^2$$

$$= (d_k^T B_k d_k) (\sigma_k^T B_k \sigma_k),$$
(8)

where the equality holds if and only if $d_k = \lambda_k \sigma_k$, $\lambda_k \neq 0$.

If the inequality (8) holds strictly, one has

$$d_k^T B_{k+1} d_k > d_k^T B_k d_k - d_k^T B_k d_k + \frac{(d_k^T \nabla^2 f(x_{k+1}) \sigma_k)^2}{\sigma_k^T \nabla^2 f(x_{k+1} \sigma_k)} \geq 0.$$

If the equality (8) holds; i.e., there exists a $\lambda_k \neq 0$ such that $d_k = \lambda_k \sigma_k$, then it can be deduced from (8) that

$$d_k^T B_{k+1} d_k \ge \frac{(d_k^T \nabla^2 f(x_{k+1}) \sigma_k)^2}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k} = \lambda_k^2 \sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k > 0.$$

In conclusion, $d_k^T B_{k+1} d_k > 0$ for $\forall d_k \in \mathbb{R}^n$ and $d_k \neq 0$. \square

Lemma 2. If we rewrite update Formula (7) as $H_{k+1} = H_k + E$, where H_k is symmetric and satisfies $\sigma_k^T = \sigma_k^T \nabla^2 f(x_{k+1}) H_k$, then E is the solution of the following minimization problem:

$$\min_{E} \qquad ||E||_{W}$$
s.t.
$$E^{T} = E,$$

$$\sigma_{k}^{T} \nabla^{2} f(x_{k+1}) E = \eta^{T},$$

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where $\eta = \sigma_k^T - \sigma_k^T \nabla^2 f(x_{k+1}) H_k$ and W satisfies $\sigma_k^T W = \sigma_k^T \nabla^2 f(x_{k+1})$.

Proof. A suitable Lagrangian function of the convex programming problem is

$$\varphi = \frac{1}{4} \operatorname{trace}(WE^TWE) + \operatorname{trace}(\Lambda^T(E^T - E)) - \lambda^T W(E\nabla^2 f(x_{k+1})\sigma_k - \eta),$$

where Λ and λ are Lagrange multipliers. Moreover,

$$\begin{split} \frac{\partial \varphi}{\partial E_{ij}} &= \frac{1}{4}(\operatorname{trace}(We_je_i^TWE) + \operatorname{trace}(WE^TWe_ie_j^T)) \\ &+ \operatorname{trace}(\Lambda(e_je_i^T - e_ie_j^T)) - \lambda^TWe_ie_i^TF'(x_{k+1})\sigma_k = 0, \end{split}$$

or according to the symmetry and cyclic permutations, one has

$$\frac{1}{2}[WEW]_{ij} + \Lambda_{ij} - \Lambda_{ji} = [W\lambda\sigma_k^T\nabla^2 f(x_{k+1})]_{ij}.$$

Taking the transpose and accumulating eliminates Λ to yield

$$WEW = W\lambda \sigma_k^T \nabla^2 f(x_{k+1}) + \nabla^2 f(x_{k+1}) \sigma_k \lambda^T W,$$

and by $\sigma_k^T W = \sigma_k^T \nabla^2 f(x_{k+1})$ and the nonsingularity of W we have that

$$E = \lambda \sigma_k^T + \sigma_k \lambda^T. \tag{9}$$

Substituting (9) into $\sigma_k^T \nabla^2 f(x_{k+1}) E = \eta^T$ and rewriting gives

$$\lambda = \frac{\eta - \sigma_k \lambda^T \nabla^2 f(x_{k+1}) \sigma_k}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k}.$$

Postmultiplying by $\sigma_k^T \nabla^2 f(x_{k+1})$ gives

$$\lambda^T \nabla^2 f(x_{k+1}) \sigma_k = \frac{1}{2} \frac{\sigma_k^T \nabla^2 f(x_{k+1}) \eta}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k},$$

so we have

$$\lambda = \frac{\eta - \frac{1}{2} \frac{\sigma_k^T \sigma_k \nabla^2 f(x_{k+1}) \eta}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k}}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k} = \frac{H_k \nabla^2 f(x_{k+1}) \sigma_k - \frac{1}{2} \frac{\sigma_k^T \sigma_k \nabla^2 f(x_{k+1}) H \nabla^2 f(x_{k+1}) \sigma_k}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k}}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k}.$$

Substituting this into (9) gives the result (7). \Box

Lemma 3. If $H_k = B_k^{-1} > 0$ and $\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k > 0$. Then B_{k+1} given by (6) solves the variational problem

$$\min_{B>0} \qquad \psi(H_k^{1/2}BH_k^{1/2})$$
s.t.
$$B^T = B,$$

$$\sigma_k^T B = \sigma_k^T \nabla^2 f(x_{k+1}).$$

Proof. According to the definition of ψ , where $\psi : \mathbb{R}^{n \times n} \to \mathbb{R}$ [40] is given by

$$\psi(A) = \operatorname{tr}(A) - \ln \det(A),\tag{10}$$

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so we have

$$\psi(H_k^{1/2}BH_k^{1/2}) = \text{trace}(H_k B) - \ln(\text{det}H_k \text{det}B) = \psi(H_k B) = \psi(BH_k). \tag{11}$$

We have the Lagrangian function

$$L(B,\Lambda,\lambda) = \frac{1}{2}\psi(H_k^{1/2}BH_k^{1/2}) + \operatorname{trace}(\Lambda^T(B^T - B) + (\sigma_k^T B - \sigma_k^T \nabla^2 f(x_{k+1}))\lambda_k$$

$$= \frac{1}{2}(\psi(H_k B) - \ln(\det H_k) - \ln(\det B))$$

$$+ \operatorname{trace}(\Lambda^T(B^T - B)) + (\sigma_k^T B - \sigma_k^T \nabla^2 f(x_{k+1}))\lambda_k,$$

where Λ and λ are the Lagrange multipliers. Moreover, one has

$$\frac{\partial L}{\partial B_{ij}} = \frac{1}{2} \operatorname{trace} \left(H_k e_i e_j^T - (B^{-1})_{ji} \right) + \operatorname{trace} \left(\Lambda^T (e_k e_i^T - e_i e_j^T) \right) + \sigma_k^T e_i e_j^T \lambda$$

$$= \frac{1}{2} (H_k)_{ji} - (B^{-1})_{ji} + \Lambda_{ji} - \Lambda_{ij} + (\sigma_k^T \lambda)_{ij} = 0. \tag{12}$$

Transposing and adding in (12) that

$$H_k - B^{-1} + \sigma_k^T \lambda + \lambda^T \sigma_k = 0,$$

$$B^{-1} = H_k + \sigma_k^T \lambda + \lambda^T \sigma_k.$$
 (13)

Combined with the tangent condition, we have that

$$\sigma_k^T = \sigma_k^T \nabla^2 f(x_{k+1}) H_k + \sigma_k^T \nabla^2 f(x_{k+1}) \lambda \sigma_k^T + \sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k \lambda^T,$$

and hence

$$\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k = \frac{1}{2} \left(1 - \frac{\sigma_k^T \nabla^2 f(x_{k+1}) H_k \nabla^2 f(x_{k+1}) \sigma_k}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k} \right),$$

and so

$$\lambda = \frac{\sigma_k - H_k \nabla^2 f(x_{k+1}) - \frac{1}{2} \left(1 - \frac{\sigma_k^T \nabla^2 f(x_{k+1}) H_k \nabla^2 f(x_{k+1}) \sigma_k}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k} \right)}{\sigma_k^T \nabla^2 f(x_{k+1}) \sigma_k}.$$

Combined with (6), one has the Formula (7).

According to the Sherman–Morrison formula, (7) is equivalent to (6). Since the function $\psi(H_k^{1/2}BH_k^{1/2})$ is strictly convex on $B\succeq 0$, the update formula (6) is the unique solution of the variational problem. \square

In this paper, we set $\sigma_k = s_k$, so one has

$$B_{k+1}s_k = \nabla^2 f(x_{k+1})s_k,$$

which means that B_{k+1} is an exact approximation to $\nabla^2 f(x_{k+1})$ in direction s_k . Then we have the symmetric rank-two update formula

$$B_{k+1} = B_k - \frac{B_k s_k s_k^T B_k}{s_k^T B_k s_k} + \frac{\nabla^2 f(x_{k+1}) s_k s_k^T \nabla^2 f(x_{k+1})}{s_k^T \nabla^2 f(x_{k+1}) s_k}.$$
 (14)

It can be seen that B_{k+1} can preserve the symmetry when B_k is symmetric. If we denote $w_k = \nabla^2 f(x_{k+1}) s_k$, then we can obtain a similar Broyden convex family update formula:

$$H_{k+1} = H_k - \frac{H_k w_k w_k^T H_k}{w_k^T H_k w_k} + \frac{s_k s_k^T}{s_k^T w_k} + \phi_k v_k v_k^T, \tag{15}$$

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where the parameter $\phi_k \in [0,1]$ is defined as

$$v_k = \sqrt{w_k^T H_k w_k} \left(\frac{s_k}{s_k^T w_k} - \frac{H_k w_k}{w_k^T H_k w_k} \right). \tag{16}$$

The choice $\phi_k \equiv 1$ corresponds to the BFGS update

$$H_{k+1} = H_k + \left(1 + \frac{w_k^T H_k w_k}{s_k^T w_k}\right) \frac{s_k^T s_k}{s_k^T w_k} - \frac{s_k^T w_k H_k + H_k w_k s_k^T}{s_k^T w_k}$$

$$= H_k + \frac{(s_k - H_k w_k) s_k^T + s_k (s_k - H_k w_k)^T}{s_k^T w_k}.$$
(17)

3. Algorithm and Related Properties

For the update formula (15), we adopt the idea of matrix completion. The next quasi-Newton matrix H_{k+1} is the solution of the following minimization problem:

min
$$\psi(H_k^{-1/2}HH_k^{-1/2})$$

s.t. $H_{ij} = H_{i,j}^{AD}, (i,j) \in F,$
 $(H^{-1})_{i,j} = 0, (i,j) \notin F,$
 $H^T = H, H \succeq 0.$ (18)

When $G(V, \bar{F})$ is chordal, the minimization problem (18) can be solved by solving the problem

max
$$\det(H)$$

s.t. $H_{i,j} = H_{i,j}^{AD}, (i,j) \in F,$
 $H^T = H, H \succeq 0.$ (19)

Then H_{k+1} can be expressed as the sparse clique-factorization formula [29]. Then Algorithm 1 is stated as follows.

Algorithm 1 (Sparse Quasi-Newton Algorithm based on Automatic Differentiation)

Step 0. Compute \bar{F} according to F such that $G(V, \bar{F})$ is a chordal graph, where $V = \{1, 2, \dots, n\}$. Choose $x_0 \in R^n$, $\epsilon > 0$ and a matrix $H_0 \in R^{n \times n}$, $H_0 \succeq 0$ with $(H_0^{-1})_{ij} = 0$, $\forall (i, j) \notin F$. Let k := 0.

Step 1 If $\|\nabla f(x_k)\| \le \epsilon$, stop.

Step 2 $x_{k+1} = x_k - H_k \nabla f(x_k)$.

Step 3 Update H_k to get H_{ij}^{AD} , $\phi_k \in [0,1]$, $(i,j) \in F$ by update Formula (15).

Step 4 Get H_{k+1} by the minimization problem (18). When $G(V, \bar{F})$ is a chordal graph, the problem (18) can be solved by solving the problem (19).

Step 5 Let k := k + 1, go to Step 1.

When the H_k in step 3 is updated by Broyden's class method, the method corresponds to the method in [29]. In the present paper, we focus on the MCQN update with $H^{AD} = H_{k+1}$, where H_{k+1} is given by (15).

In what follows, we give some notation for the convenience of analysis. For a nonsingular matrix *P* satisfying

$$(P^{-1})_{ij} = 0, \ \forall (i,j) \in F,$$
 (20)

we let

$$\bar{s}_k = P^{-1/2} s_k$$
, $\bar{w}_k = P^{1/2} w_k$, $\bar{H}_k = P^{-1/2} H_k P^{-1/2}$, $\bar{H}^{AD} = P^{-1/2} H^{AD} P^{-1/2}$

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where $H^{AD} = H_{k+1}$ is given by (15). Then we can get from (15) that

$$\bar{H}^{AD} = \bar{H}_k - \frac{\bar{H}_k \bar{w}_k \bar{w}_k^T \bar{H}_k}{\bar{w}_k^T \bar{H}_k \bar{w}_k} + \frac{\bar{s}_k \bar{s}_k^T}{\bar{s}_k^T \bar{w}_k} + \phi_k \bar{v}_k \bar{v}_k^T, \tag{21}$$

where

$$\bar{v}_k = \sqrt{\bar{w}_k^T \bar{H}_k \bar{w}_k} \left(\frac{\bar{s}_k}{\bar{s}_k^T \bar{w}_k} - \frac{\bar{H}_k \bar{w}_k}{\bar{w}_k^T \bar{H}_k \bar{w}_k} \right). \tag{22}$$

Similarly to that in [30], we can assume that

$$\tau_{k} = \frac{\bar{w}_{k}^{T} \bar{H}_{k} \bar{w}_{k}^{T}}{\|\bar{w}_{k}\| \cdot \|\bar{H}_{k} \bar{w}_{k}\|}, \ q_{k} = \frac{\bar{w}_{k}^{T} \bar{H}_{k} \bar{w}_{k}}{\|\bar{w}_{k}\|^{2}}, \ \eta_{k} = \frac{\bar{s}_{k}^{T} \bar{H}_{k} \bar{w}_{k}}{\bar{s}_{k}^{T} \bar{w}_{k}}, \ m_{k} = \frac{\bar{s}_{k}^{T} \bar{w}_{k}}{\bar{w}_{k}^{T} \bar{w}_{k}},$$

$$M_k = \frac{\|\bar{s}_k\|^2}{\bar{s}_k^T \bar{w}_k}, \; \beta_k = \frac{\bar{s}_k^T \bar{H}_k^{-1} \bar{s}_k^T}{\bar{s}_k^T \bar{w}_k}, \; \gamma_k = \frac{\bar{w}_k^T \bar{H}_k \bar{w}_k}{\bar{s}_k^T \bar{w}_k}.$$

According to [41] and (21), we have

$$tr(\bar{H}^{AD}) = tr(\bar{H}_k) - (1 - \phi_k) \frac{q_k}{\tau_k^2} - 2\phi_k \eta_k + \left(1 + \phi_k \frac{q_k}{m_k}\right) M_k$$
 (23)

and

$$\det(\bar{H}^{AD}) = \det(\bar{H}_k) \left(1 + \phi_k (\beta_k \gamma_k - 1) \right) / \gamma_k. \tag{24}$$

Next, we establish a relation between \bar{H}_{k+1} and \bar{H}^{AD} , which is very important in the establishment of the local and superlinear convergence of Algorithm 1.

Proposition 1. For the Algorithm 1, we have the following relation:

$$tr(\bar{H}_{k+1}) = tr(\bar{H}^{AD}), det(\bar{H}_{k+1}) \geqslant det(\bar{H}^{AD}).$$
 (25)

Proof. We can obtain from (18) that

$$(H_{k+1})_{i,j}=(H^{AD})_{i,j},\ \forall (i,j)\in F.$$

Combined with (20), one has that for any $(i,j) \in F$, there at least exists one of the $(H_{k+1} - H^{AD})_{ij}$ and $P_{i,j}^{-1}$ equals to zero. Then we can get that

$$\operatorname{tr}(\bar{H}_{k+1} - \bar{H}^{AD}) = \operatorname{tr}(P^{-1}(H_{k+1} - H^{AD}))$$

$$= \sum_{i=1}^{n} \sum_{j=1}^{n} (P^{-1})_{i,j} (H_{k+1} - H^{AD})_{i,j} = 0.$$
(26)

Moreover, since H^{AD} satisfies (19), we must have

$$\det(H_{k+1}) \geqslant \det(H^{AD})$$

Consequently, one has

$$\det(\bar{H}_{k+1}) = \bar{H}^{AD}. \tag{27}$$

Remark 1. According to the definition of ψ (10) and the relation between \bar{H}_{k+1} and \bar{H}^{AD} (25), one has that

$$\psi(\bar{H}_{k+1}) \leqslant \psi(\bar{H}^{AD}). \tag{28}$$

When we substitute (23) and (24) into (28), the $\psi(\bar{H}_{k+1})$ and $\psi(\bar{H}_k)$ has the relation

$$\psi(\bar{H}_{k+1}) \leqslant \psi(\bar{H}_k) - (1 - \phi_k) \frac{q_k}{\tau_k^2} - 2\phi_k \eta_k + \left(1 + \phi_k \frac{q_k}{m_k}\right) M_k
- \ln\left(1 + \phi_k(\beta_k \gamma_k - 1)\right) + \ln\gamma_k.$$
(29)

4. The Local and Superlinear Convergence

Based on the discussion in Section 3, we prove the local and superlinear convergence of Algorithm 1. First, we list the assumptions.

Assumption 1. Assume that x^* is a solution of (1) and

$$]\Omega = \{x \in R^n | ||x - x^*|| \le b\},$$

where b > 0.

- (1) The function $f \in \mathbb{R}^n \to \mathbb{R}$ is twice continuously differentiable on Ω .
- (2) There exist two constants, m > 0 and M > 0, satisfying

$$m||u||^2 \le u^T (\nabla^2 f(x))^{-1} u \le M||u||^2, \ \forall u \in \mathbb{R}^n, x \in \Omega.$$
 (30)

According to Assumption 1, we have constants $\bar{L} > 0$ and L > 0 such that

$$\|\nabla f(x) - \nabla f(y)\| \leqslant \bar{L} \|x - y\|, \ \forall x, y \in \Omega, \tag{31}$$

$$\|\nabla^2 f(x) - \nabla^2 f(y)\| \le L\|x - y\|, \ \forall x, y \in \Omega.$$
 (32)

We define

$$\epsilon_k = \max\{\|x_k - x^*\|, \|x_{k+1} - x^*\|\},$$
 (33)

and get from (32) that

$$||w_{k} - \nabla^{2} f(x^{*}) s_{k}|| = ||\nabla^{2} f(x_{k+1}) s_{k} - \nabla^{2} f(x^{*}) s_{k}||$$

$$\leq ||\nabla^{2} f(x_{k+1}) - \nabla^{2} f(x^{*})|| \cdot ||s_{k}||$$

$$\leq L ||x_{k+1} - x^{*}|| \cdot ||s_{k}||$$

$$= L\epsilon_{k} ||s_{k}||.$$
(34)

If we take $P = H^*$, then one has from (34) that

$$\|\bar{w}_k - \bar{s}_k\| = \|P^{1/2}w_k - P^{-1/2}s_k\|$$

$$= \|H^{*1/2}\| \cdot \|w_k - \nabla^2 f(x^*)s_k\| \le L\|H^*\|^{1/2}\epsilon_k\|s_k\|, \tag{35}$$

Furthermore, it is easy to deduce that

$$M_k - 1, \mu_k = \frac{2 - M_k - m_k}{m_k}, \hat{\mu}_k = \frac{(\bar{w}_k - \bar{s}_k)^T \bar{H}_k \bar{w}_k}{\operatorname{tr}(\bar{H}_k) \bar{s}_k^T \bar{w}_k}, \ln m_k \leqslant \frac{1}{2} c_1 \epsilon_k, \tag{36}$$

where $c_1 > 0$, $c_2 \in (0, b)$, and $\epsilon_k < c_2$. We define

$$\rho_{k} = q_{k} - 1 - \ln q_{k},
\zeta_{k} = (1 - \phi_{k})q_{k}(\tau_{k}^{-2} - 1),
\xi_{k} = \ln(1 + \phi_{k}(\beta_{k}\gamma_{k} - 1)),$$
(37)

and rewrite (29) as

$$\psi(\bar{H}_{k+1}) \leq \psi(\bar{H}_k) - \rho_k - \zeta_k - \xi_k + (M_k - 1)
+ \phi_k q_k \mu_k + \phi_k \operatorname{tr}(\bar{H}_k) \hat{\mu}_k + \ln m_k.$$
(38)

As $\gamma_k = q_k/m_k$ and $0 \le q_k \le \operatorname{tr}(\bar{H}_k)$, we can obtain from the above inequality and (36) that

$$\psi(\bar{H}_{k+1}) \leqslant \psi(\bar{H}_k) - \rho_k - \zeta_k - \xi_k + c_1 \Big(1 + \operatorname{tr}(\bar{H}_k) \Big) \epsilon_k. \tag{39}$$

Considering

$$\lambda - \ln \lambda \geqslant \max\left(\left(1 - \frac{1}{e}\right)\lambda, 1\right), \ \forall \lambda > 0,$$
 (40)

one has

$$\psi(A) \geqslant \max\left(\left(1 - \frac{1}{e}\right) \operatorname{tr}(A), n\right),$$

where $A^T = A$ and A > 0. Moreover, it follows from (40) that

$$\psi(\bar{H}_{k+1}) \leqslant (1 + c_3 \epsilon_k) \psi(\bar{H}_k) - \rho_k - \zeta_k - \xi_k, \tag{41}$$

where $c_3 = c_1 \left(\frac{1}{n} + \frac{e}{e-1} \right)$. Since $\tau_k^2 \leqslant 1$ and $\beta_k \gamma_k \geqslant 1$, it is obvious that $\rho_k, \zeta_k, \xi_k > 0$, and

$$\psi(\bar{H}_{k+1}) \leqslant (1 + c_3 \epsilon_k) \psi(\bar{H}_k). \tag{42}$$

The theorem given bellow shows that Algorithm 1 converges locally and linearly, where the relation (42) plays an essential role.

Theorem 1. Let Assumption 1 hold and sequence $\{x_k\}$ be generated by Algorithm 1 with $\alpha_k \equiv 1$, where H_k is updated by (15). Then for any $\rho \in (0,1)$, there is a constant $\tau \|x_0 - x^*\| \leq \tau$, $\|H_0 - H^*\| \leq \tau$, such that

$$||x_{k+1} - x^*|| \le \rho ||x_k - x^*||. \tag{43}$$

Proof. According to the Lemma 4 [29], there are constants $\bar{\tau} \in (0, b)$ and $\delta > 0$ such that when $||x_0 - x^*|| \leq \bar{\tau}$, one has

$$\psi(\bar{H}_0) - n \leqslant \delta/2,\tag{44}$$

and

$$||H - H^*|| \leqslant \rho/(2\bar{L}),\tag{45}$$

where $H \succeq 0$ and $\bar{H} = H^{*-1/2}H\bar{H} = H^{*-1/2}$. Define

$$\tau = \min \left\{ \bar{\tau}, c_2, \frac{\rho}{\bar{L}}, \frac{\rho}{LM}, \frac{1-\rho}{c_3} \ln \left(\frac{2(n+\delta)}{2n+\delta} \right) \right\}. \tag{46}$$

We will prove the inequalities (43) and

$$||H_k - H^*|| \leqslant \frac{\rho}{2\bar{L}} \tag{47}$$

hold for any $k \geqslant 0$ by induction. By the Lipstchitz continuity of $\nabla^2 f(x)$, we have for $\forall x \in \Omega$,

$$||x - x^* - H^* \nabla f(x)|| \leq ||H^*|| \cdot \int_0^1 ||\nabla^2 f(x + t(x - x^*)) - \nabla^2 f(x^*)|| \cdot ||x - x^*|| dt$$

$$\leq \frac{1}{2} LM ||x - x^*||^2. \tag{48}$$

Then, when k = 0, it is easy to deduce (43) by (44) and (45). Moreover, when we take $\alpha_k \equiv 1$ and substitute x_0 into (48), we can obtain

$$||x_{1} - x^{*}|| = ||x_{0} - H_{0}\nabla f(x_{0}) - x^{*}||$$

$$\leq ||x_{0} - x^{*} - H^{*}\nabla f(x_{0})|| + ||(H_{0} - H^{*})(\nabla f(x_{0})) - \nabla f(x^{*}))||$$

$$\leq \frac{1}{2}LM||x_{0} - x^{*}||^{2} + ||H_{0} - H^{*}|| \cdot ||\nabla f(x_{0})) - \nabla f(x^{*})||$$

$$\leq \left(\frac{1}{2}LM||x_{0} - x^{*}|| + \frac{\rho}{2}\right)||x_{0} - x^{*}||$$

$$\leq \left(\frac{1}{2}LM\tau + \frac{\rho}{2}\right)||x_{0} - x^{*}||$$

$$\leq \rho||x_{0} - x^{*}||. \tag{49}$$

So we have that (43) and (47) hold for k = 1. Assume that (43) and (47) hold for k = 0, 1, ..., l; then one has

$$\epsilon_k = ||x_k - x^*||, \ \epsilon_k \leqslant \rho^k \epsilon_0 \leqslant \rho^k \tau, \ k = 0, 1, \dots, l,$$

and

$$||x_{l+1} - x^*|| = ||x_l - H_l \nabla f(x_l) - x^*||$$

$$\leq ||x_l - x^* - H^* \nabla f(x_l)|| + ||(H_l - H^*)(\nabla f(x_l)) - \nabla f(x^*))||$$

$$\leq \frac{1}{2} LM ||x_l - x^*||^2 + ||H_l - H^*|| \cdot ||\nabla f(x_l)) - \nabla f(x^*)||$$

$$\leq \left(\frac{1}{2} LM ||x_l - x^*|| + \frac{\rho}{2}\right) ||x_l - x^*||$$

$$\leq \left(\frac{1}{2} LM \rho^l \tau + \frac{\rho}{2}\right) ||x_l - x^*||$$

$$\leq \rho ||x_l - x^*||.$$
(50)

Then by the definition of τ (46), one has

$$c_3 \sum_{k=0}^{l} \epsilon_k \leqslant c_3 \tau \sum_{k=0}^{l} \rho^k = c_3 \tau \frac{1 - \rho^{l+1}}{1 - \rho} \leqslant \frac{c_3 \tau}{1 - \rho} \leqslant \ln \frac{2(n+\delta)}{2n + \delta}.$$
 (51)

Combine (42) and (44). It can seen that

$$\psi(\bar{H}_{l+1}) - n \leqslant (\psi(\bar{H}_0) - n) + \left(\prod_{k=0}^{l} (1 + c_3 \epsilon_k) - 1\right) \psi(\bar{H}_0)$$

$$\leqslant \frac{\delta}{2} + \left(n + \frac{\delta}{2}\right) \left(\prod_{k=0}^{l} e^{c_3 \epsilon_k} - 1\right)$$

$$\leqslant \frac{\delta}{2} + \left(n + \frac{\delta}{2}\right) \left(e^{c_3 \sum_{k=0}^{l} \epsilon_k} - 1\right)$$

$$\leqslant \frac{\delta}{2} + \left(n + \frac{\delta}{2}\right) \left(\frac{2(n+\delta)}{2n+\delta} - 1\right) = \delta.$$
(52)

Thus, we can get that (47) holds for all k = l + 1. This completes the proof. \Box

Based on the above discussion and the relation (42), we can show the superlinear convergence of the Algorithm 1.

Theorem 2. Let Assumption 1 hold and sequence $\{x_k\}$ be generated by Algorithm 1 with $\alpha_k \equiv 1$, where H_k is updated by (15). Then there is a constant $\tau > 0$ such that when $\|x_0 - x^*\| \leq \tau$, $\|H_0 - H^*\| \leq \tau$, one has

$$\lim_{k \to \infty} \frac{\|(H_k - H^*)w_k\|}{\|w_k\|} = 0.$$
 (53)

Then the sequence $\{x_k\}$ is superlinearly convergent

Proof. Let τ be defined as in (1), and for all k one has

$$\psi(\bar{H}_k) - n \leqslant \delta. \tag{54}$$

It follows from (41) that

$$\rho_k + \zeta_k + \xi_k \leqslant \left(\psi(\bar{H}_{k+1}) - \psi(\bar{H}_k)\right) + c_3 \epsilon_k \psi(\bar{H}_k).$$

Summing the above inequality and combining (51) and (54), we can deduce

$$\sum_{k \ge 1} (\rho_k + \zeta_k + \xi_k) \le c_3 \sum_{k \ge 1} \epsilon_k \psi(\bar{H}_k) \le c_3 (n + \delta) \ln \frac{2(n + \delta)}{2n + \delta} < \infty, \tag{55}$$

which means that the nonnegative constants ρ_k , ζ_k and ξ_k all tend to zero when $k \to +\infty$. Furthermore, according to the definition of (37), we have that

(1)
$$q_k \to 1$$
; (2) if $\phi_k \leqslant \frac{1}{2}$, $\tau \to 1$; (3) if $\phi_k > \frac{1}{2}$, $\beta_k \gamma_k \to 1$.

First, we have

$$\frac{\|H^{*-1/2}(H_k - H^*)w_k\|^2}{\|H^{*1/2}w_k\|^2} = \frac{\|\bar{H}_k\bar{w}_k\|^2 - 2\bar{w}_k^T\bar{H}_k\bar{w}_k + \|\bar{w}_k\|^2}{\|\bar{H}_k\|^2} \\
= \frac{q_k}{\tau_k^2} - 2q_k + 1.$$
(56)

For the case $\{k_i: \phi_{k_i} \leqslant \frac{1}{2}\}$, one has $q_k \to 1$, and $\tau_{k_i} \to 1$; and then (53) is true. Moreover, it is easy to deduce that

$$\frac{\|\bar{H}_{k}\bar{w}_{k} - \bar{s}_{k}\|^{2}}{\|\bar{w}_{k}\|^{2}} \leq \frac{\|\bar{H}_{k}^{1/2}\|^{2} \cdot \|\bar{H}_{k}^{1/2}\bar{w}_{k} - (\bar{H}_{k})^{-1/2}\bar{s}_{k}\|^{2}}{\|\bar{w}_{k}\|^{2}}$$

$$= \frac{\|\bar{H}_{k}^{1/2}\|^{2}(\bar{w}_{k}^{T}\bar{H}_{k}\bar{w}_{k} - 2\bar{s}_{k}^{T}\bar{w}_{k} + \bar{s}_{k}^{T}(\bar{H}_{k})^{-1}\bar{s}_{k})}{\|\bar{w}_{k}\|^{2}}$$

$$= \|\bar{H}_{k}^{1/2}\|^{2}\left(q_{k} - 2m_{k} + \frac{\beta_{k}\gamma_{k}}{q_{k}}\right). \tag{57}$$

We also have

$$\left| \frac{\|\bar{H}_k \bar{w}_k - \bar{w}_k\|}{\|\bar{w}_k\|} - \frac{\|\bar{H}_k \bar{w}_k - \bar{s}_k\|}{\|\bar{w}_k\|} \right| \leqslant \frac{\|\bar{w}_k - \bar{s}_k\|}{\|\bar{w}_k\|} \to 0.$$
 (58)

For the case $\{k_i: \phi_{k_i} > \frac{1}{2}\}$, one has $q_k \to 1$, $\beta_k \gamma_k \to 1$, $m_k \to 1$; then (53) is true by (56)–(58). Thus, the relation (53) holds for all k.

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Next, we will show that (53) indicates that the sufficient condition [6]

$$\lim_{k \to \infty} \frac{\|(B_k - \nabla^2 f(x^*))s_k\|}{\|s_k\|} = 0$$
 (59)

holds. According to (47), one has that there is a constant $\lambda_{\min} > 0$ such that $(\lambda_k)_i \geqslant \lambda_{\min}$, where $(\lambda_k)_i$ denotes the eigenvalues of H_k , $i = 1, 2, \cdots, n$. When we let $w_k = \nabla^2 f(x_{k+1}) s_k$, one has

$$\begin{aligned} & \| (H_k - H^*) w_k \| \\ &= \| (H_k - H^*) \nabla^2 f(x^*) s_k + (H_k - H^*) (\nabla^2 f(x_{k+1}) - \nabla^2 f(x^*)) s_k \| \\ &\geqslant \| H_k (\nabla^2 f(x^*) - B_k) s_k \| - \| (H_k - H^*) (\nabla^2 f(x_{k+1}) - \nabla^2 f(x^*)) s_k \| \\ &\geqslant \lambda_{\min} \| (\nabla^2 f(x^*) - B_k) s_k \| - \| (H_k - H^*) (\nabla^2 f(x_{k+1}) - \nabla^2 f(x^*)) s_k \|_{\ell} \end{aligned}$$

and

$$\begin{split} &\frac{\|(H_k-H^*)w_k\|}{\|w_k\|} \\ &= \frac{\lambda_{\min}\|(\nabla^2 f(x^*) - B_k)s_k\|}{\|\nabla^2 f(x_{k+1})s_k\|} - \frac{\|(H_k-H^*)(\nabla^2 f(x_{k+1}) - \nabla^2 f(x^*))s_k\|}{\|\nabla^2 f(x_{k+1})s_k\|} \\ &\geqslant \frac{\lambda_{\min}\|(\nabla^2 f(x^*) - B_k)s_k\|}{\frac{1}{\lambda_{\min}}\|s_k\|} - \frac{\|(H_k-H^*)(\nabla^2 f(x_{k+1}) - \nabla^2 f(x^*))s_k\|}{\frac{1}{\lambda_{\min}}\|s_k\|} \\ &= \lambda_{\min}^2\|(\nabla^2 f(x^*) - B_k)\| - \lambda_{\min}\|(H_k-H^*)(\nabla^2 f(x_{k+1}) - \nabla^2 f(x^*))\|. \end{split}$$

When $k \to \infty$, since $x_k \to x^*$, then one has from (53) that

$$\lim_{k \to \infty} \frac{\|(B_k - \nabla^2 f(x^*)) s_k\|}{\|s_k\|} = 0,$$

which is the well-known Dennis–Moré condition. Thus, we get the superlinear convergence. $\ \ \Box$

5. Numerical Experiments

The performance in [29] shows that the MCQN update with the BFGS method has better numerical performance than the MCQN update with DFP method. Hence, we compare the numerical performance of Algorithm 1 with the MCQN update with BFGS method and the limited-memory BFGS method.

The 24 test problems with initial points are given in Table 1, which are from [29,42–44]. It can be seen that all the test problems have special Hessian structures such as band matrices, so the chordal extension of the sparsity could be obtained easily. Then H_{k+1} in Algorithm 1 can be written as the sparse clique-factorization formula.

All the methods were coded in MATLAB R2016a on a Core (TM) i5 PC. The automatic differentiation was computed by ADMAT 2.0, which is available on the cayuga research GitHUB page. In Tables 1–4 and Figures 1 and 2, we report the numerical performances of the three methods. For the convenience of statement, we use the following notation in our numerical results.

Pro: the problems;

Dim: the dimensions of the test problem;

Init: the initial points;

Method: the algorithm used to solve the problem; MCQN-BFGS: MCQN update with the BFGS method; L-BFGS: limited-memory with the BFGS method. Symmetry **2021**, 13, 2093 15 of 21

We adopted the termination criterion as follows:

$$\frac{\|\nabla f(x)\|}{n} \leqslant 10^{-5} \text{ or ite } \geqslant 5000.$$

Table 1. The test problems.

Pro	the Test Functions	Init
1	TRIDIA [29]	$x_0 = (1, 1, \cdots, 1)^T$
2	the chained Rosenbrock problem [29]	$x_0 = (-1.2, -1, \dots, -1.2, -1)^T$
3	the boundary value problem [29]	$x_0 = (\frac{1}{n+1}, \frac{2}{n+1}, \cdots, \frac{n}{n+1})^T$
4	Broyden tridiagonal function [42]	$x_0 = (-1, -1, \cdots, -1)^T$
5	DQRTIC [43]	$x_0=(2,2,\cdots,2)^T$
6	EDENSCH [43]	$x_0 = (0,0,\cdots,0)^T$
7	ENGVAL1 [43]	$x_0 = (2, 2, \cdots, 2)^T$
8	COSINE [43]	$x_0=(1,1,\cdots,1)^T$
9	ERRINROS-modified [43]	$x_0 = (-1, -1, \cdots, -1)^T$
10	FREUROTH [43]	$x_0 = (0.5, -2, 0, \cdots, 0)^T$
11	MOREBV- different start point [43]	$x_0 = (0.5, 0.5, \cdots, 0.5)^T$
12	TOINTGSS [43]	$x_0 = (3, 3, \cdots, 3)^T$
13	SCHMVETT [43]	$x_0 = (3,3,\cdots,3)^T$
14	Extended Freudenstein and Roth function [44]	$x_0 = (0.5, -2, \cdots, 0.5, -2)^T$
15	Raydan 1 function [44]	$x_0 = (1, 1, \cdots, 1)^T$
16	Generalized Tridiagonal function [44]	$x_0=(2,2,\cdots,2)^T$
17	Extended Himmelblau function [44]	$x_0=(1,1,\cdots,1)^T$
18	Generalized PSCI function [44]	$x_0 = (3, 0.1, \cdots, 3, 0.1)^T$
19	Extended Tridiagonal 2 function [44]	$x_0=(1,1,\cdots,1)^T$
20	Raydan 2 function [44]	$x_0=(1,1,\cdots,1)^T$
21	Extended Freudenstein and Roth function [44]	$x_0=(1,1,\cdots,1)^T$
22	DQDRTIC function [44]	$x_0=(3,3,\cdots,3)^T$
23	Generalized Quartic function [44]	$x_0=(1,1,\cdots,1)^T$
24	HIMMELBG function [44]	$x_0 = (1.5, 1.5, \cdots, 1.5)^T$

Firstly, we tested all three methods on the above 24 problems, whose dimensions are 10, 20, 50, 100, 200, 5000, 1000, 2000, 5000, and 1000. We set m=15 in the limited-memory BFGS method. Tables 2 and 3 contain the numbers of iterations of the three methods for the test problems. Taking account of the total number of iterations, Algorithm 1 outperformed the MCQN update with BFGS method on 11 problems (2, 4, 5, 7, 9, 10, 12, 14, 18, 23, 24). Additionally, Algorithm 1 outperformed the limited memory BFGS method on 13 problems (1, 2, 3, 7, 9, 12, 15, 16, 18, 19, 20, 21, 23).

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Table 2. Numbers of iterations for problems 1-12.

(1) Algorithm 1 30 38 51 78 96 146 217 301 424 527 (1) MCQN-BFGS 29 38 51 72 95 146 192 298 424 528 (1) L-BFGS 26 39 96 158 360 864 1042 1759 3153 3152 (2) Algorithm 1 49 90 166 308 595 1345 2699 5437 3218 2725 (2) MCQN-BFGS 60 95 200 384 683 1668 3249 6486 4562 3207 (2) L-BFGS 59 113 260 504 999 2481 4947 9887 24,732 49,391 (3) Algorithm 1 16 26 42 58 59 51 49 60 102 399 (3) MCQN-BFGS 15 26 42 50 59 71 54 69 101 402 (3) L-BFGS 39 114 279 700 1503 1659 2695 3370 8867 27,471 (4) Algorithm 1 31 25 34 49 43 44 43 49 52 53 (4) MCQN-BFGS 30 29 43 45 49 58 61 62 63 56 (4) L-BFGS 21 27 40 54 41 38 38 56 52 50 (5) Algorithm 1 30 48 60 92 109 99 92 89 81 81 81 (5) MCQN-BFGS 35 49 67 94 111 108 112 98 84 84 (5) L-BFGS 28 27 34 31 33 39 41 43 54 81 (6) Algorithm 1 23 26 38 44 54 55 54 75 1 61 51 (6) MCQN-BFGS 17 27 36 53 60 55 54 57 51 50 54 (6) L-BFGS 17 19 19 22 21 23 24 26 24 25 (7) Algorithm 1 16 21 21 19 17 17 16 15 16 15 16 15 (7) MCQN-BFGS 20 22 23 22 15 15 15 15 17 17 16 (7) MCQN-BFGS 20 22 26 21 27 28 20 22 26 21 27 28 30 (20 20 20 20 20 20 20 20 20 20 20 20 20 2	Dim	10	20	50	100	200	500	1000	2000	5000	10,000
(1) L-BFGS 26 39 96 158 360 864 1042 1759 3153 3152 (2) Algorithm 1 49 90 166 308 595 1345 2699 5437 3218 2725 (2) MCQN-BFGS 60 95 200 384 683 1668 3249 6486 4562 3207 (2) L-BFGS 59 113 260 504 999 2481 4947 9887 24,732 49,391 (3) Algorithm 1 16 26 42 58 59 51 49 60 102 399 (3) MCQN-BFGS 15 26 42 50 59 71 54 69 101 402 (3) L-BFGS 39 114 279 700 1503 1659 2695 3370 8867 27,471 (4) Algorithm 1 31 25 34 49 43 44 43 49 52 53 (4) MCQN-BFGS 30 29 43 45 49 58 61 62 63 56 (4) L-BFGS 21 27 40 54 41 38 38 56 52 50 (5) Algorithm 1 30 48 60 92 109 99 92 89 81 81 (5) MCQN-BFGS 35 49 67 94 111 108 112 98 84 84 (5) L-BFGS 28 27 34 31 33 39 41 43 54 81 (6) Algorithm 1 23 26 38 44 54 55 54 51 50 54 (6) L-BFGS 17 27 36 53 60 55 54 51 50 54 (6) L-BFGS 17 19 19 22 21 23 24 26 24 25 (7) Algorithm 1 16 21 21 19 17 17 16 15 16 15 (7) MCQN-BFGS 20 22 23 22 15 15 15 15 17 17 17 16 (7) L-BFGS 20 22 26 21 22 21 25 27 28 30											
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(2) L-BFGS 59 113 260 504 999 2481 4947 9887 24,732 49,391 (3) Algorithm 1 16 26 42 58 59 51 49 60 102 399 (3) MCQN-BFGS 15 26 42 50 59 71 54 69 101 402 (3) L-BFGS 39 114 279 700 1503 1659 2695 3370 8867 27,471 (4) Algorithm 1 31 25 34 49 43 44 43 49 52 53 (4) MCQN-BFGS 30 29 43 45 49 58 61 62 63 56 (4) L-BFGS 21 27 40 54 41 38 38 56 52 50 (5) Algorithm 1 30 48 60 92 109 99 92 89 81 81 <tr< td=""><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td></tr<>											
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(3) L-BFGS 39 114 279 700 1503 1659 2695 3370 8867 27,471 (4) Algorithm 1 31 25 34 49 43 44 43 49 52 53 (4) MCQN-BFGS 30 29 43 45 49 58 61 62 63 56 (4) L-BFGS 21 27 40 54 41 38 38 56 52 50 (5) Algorithm 1 30 48 60 92 109 99 92 89 81 81 (5) MCQN-BFGS 35 49 67 94 111 108 112 98 84 84 (5) L-BFGS 28 27 34 31 33 39 41 43 54 81 (6) Algorithm 1 23 26 38 44 54 55 47 51 61 51 (6) L-BFGS 17 27 36 53 60 55 54 51 50 54 (6) L-BFGS 17 19 19 22 21 23 24 26 24 25 <											
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(4) L-BFGS 21 27 40 54 41 38 38 56 52 50 (5) Algorithm 1 30 48 60 92 109 99 92 89 81 81 (5) MCQN-BFGS 35 49 67 94 111 108 112 98 84 84 (5) L-BFGS 28 27 34 31 33 39 41 43 54 81 (6) Algorithm 1 23 26 38 44 54 55 47 51 61 51 (6) MCQN-BFGS 17 27 36 53 60 55 54 51 50 54 (6) L-BFGS 17 19 19 22 21 23 24 26 24 25 (7) Algorithm 1 16 21 21 19 17 17 16 15 16 15 (7) MCQN-BFGS 20 22 23 22 15 15 15 17 17 16 (7) L-BFGS 20 22 26 21 22 21 25 27 28 30											
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(5) L-BFGS 28 27 34 31 33 39 41 43 54 81 (6) Algorithm 1 23 26 38 44 54 55 47 51 61 51 (6) MCQN-BFGS 17 27 36 53 60 55 54 51 50 54 (6) L-BFGS 17 19 19 22 21 23 24 26 24 25 (7) Algorithm 1 16 21 21 19 17 17 16 15 16 15 (7) MCQN-BFGS 20 22 23 22 15 15 15 17 17 16 (7) L-BFGS 20 22 26 21 22 21 25 27 28 30	(5) Algorithm 1			60		109					
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(6) L-BFGS 17 19 19 22 21 23 24 26 24 25 (7) Algorithm 1 16 21 21 19 17 17 16 15 16 15 (7) MCQN-BFGS 20 22 23 22 15 15 15 17 17 16 (7) L-BFGS 20 22 26 21 22 21 25 27 28 30	(6) Algorithm 1	23	26	38				47			
(7) Algorithm 1 16 21 21 19 17 17 16 15 16 15 (7) MCQN-BFGS 20 22 23 22 15 15 15 17 17 16 (7) L-BFGS 20 22 26 21 22 21 25 27 28 30	(6) MCQN-BFGS	17	27	36		60	55	54	51	50	
(7) MCQN-BFGS 20 22 23 22 15 15 15 17 17 16 (7) L-BFGS 20 22 26 21 22 21 25 27 28 30	(6) L-BFGS	17	19	19	22	21	23	24	26	24	25
(7) L-BFGS 20 22 26 21 22 21 25 27 28 30	(7) Algorithm 1	16		21	19	17	17	16	15	16	15
	(7) MCQN-BFGS	20		23	22	15	15			17	16
	(7) L-BFGS	20	22	26	21	22	21	25	27	28	30
	(8) Algorithm 1	22	23	24	21	23	27	27	28	29	30
(8) MCQN-BFGS 23 25 26 26 26 27 28 28 29 30	(8) MCQN-BFGS	23	25	26	26	26	27	28	28	29	30
(8) L-BFGS 9 9 9 10 10 10 10 10 10	(8) L-BFGS	9	9	9	9	10	10	10	10	10	10
(9) Algorithm 1 74 122 134 149 137 180 148 153 172 170	(9) Algorithm 1	74	122	134	149	137	180	148	153	172	170
(9) MCQN-BFGS 106 125 145 171 199 181 168 171 174 179	(9) MCQN-BFGS	106	125	145	171	199	181	168	171	174	179
(9) L-BFGS 163 245 216 196 189 190 163 169 171 192	(9) L-BFGS	163	245	216	196	189	190	163	169	171	192
(10) Algorithm 1 45 45 48 39 45 43 45 145 161 145	(10) Algorithm 1	45	45	48	39	45	43	45	145	161	145
(10) MCQN-BFGS 47 48 49 43 47 48 41 244 204 279	(10) MCQN-BFGS	47	48	49	43	47	48	41	244	204	279
(10) L-BFGS 24 25 24 24 22 22 22 20 22	(10) L-BFGS	24	25	24	24	24	22	22	22	20	22
(11) Algorithm 1 24 45 97 121 82 67 35 34 21 11	(11) Algorithm 1	24	45	97	121	82	67	35	34	21	11
(11) MCQN-BFGS 24 45 98 121 82 67 35 34 21 11		24	45	98	121	82	67	35	34	21	11
(11) L-BFGS 33 103 136 127 85 49 32 21 10 9	(11) L-BFGS	33	103	136	127	85	49	32	21	10	9
(12) Algorithm 1 13 11 11 12 9 5 6 2 3 2	(12) Algorithm 1	13	11	11	12	9	5	6	2	3	2
(12) MCQN-BFGS 15 11 13 12 12 7 6 2 3 2		15	11	13	12	12	7	6		3	2
(12) L-BFGS 6 8 9 10 13 11 10 10 9 10	(12) L-BFGS	6	8	9	10	13	11	10	10	9	10

Table 3. Numbers of iterations for problems 13–24.

Dim	10	20	50	100	200	500	1000	2000	5000	10,000
(13) Algorithm 1	19	20	21	22	18	17	15	14	13	12
(13) MCQN-BFGS	19	20	21	22	18	17	15	14	13	12
(13) L-BFGS	16	18	17	18	18	17	18	18	17	18
(14) Algorithm 1	36	52	95	111	169	262	612	567	1062	1114
(14) MCQN-BFGS	37	54	86	120	190	300	594	679	1062	1126
(14) L-BFGS	10	10	10	10	10	10	10	10	10	11
(15) Algorithm 1	11	13	23	30	39	45	60	97	196	295
(15) MCQN-BFGS	11	12	21	28	37	45	64	95	196	295
(15) L-BFGS	13	21	32	50	79	122	207	338	402	770

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Table 3. Cont.

Dim	10	20	50	100	200	500	1000	2000	5000	10,000
(16) Algorithm 1	29	31	47	69	110	149	165	174	170	163
(16) MCQN-BFGS	29	35	51	70	100	146	164	173	171	164
(16) L-BFGS	25	65	80	162	160	156	151	150	144	140
(17) Algorithm 1	16	15	14	11	13	12	12	12	10	9
(17) MCQN-BFGS	14	11	18	12	13	12	12	12	10	9
(17) L-BFGS	8	8	8	8	8	8	8	8	8	8
(18) Algorithm 1	49	21	23	21	19	14	22	21	13	11
(18) MCQN-BFGS	43	28	22	29	20	15	20	25	23	26
(18) L-BFGS	36	34	37	36	39	36	40	34	41	37
(19) Algorithm 1	13	12	11	13	12	12	12	10	7	5
(19) MCQN-BFGS	13	12	11	13	12	12	12	10	7	5
(19) L-BFGS	12	14	15	16	16	17	16	15	16	17
(20) Algorithm 1	5	5	4	4	4	4	4	3	3	3
(20) MCQN-BFGS	5	5	4	4	4	4	4	3	3	3
(20) L-BFGS	7	7	7	7	7	7	7	7	7	7
(21) Algorithm 1	14	11	11	21	18	18	35	28	26	43
(21) MCQN-BFGS	13	10	11	22	18	18	35	28	26	43
(21) L-BFGS	10	13	15	19	22	30	36	51	56	64
(22) Algorithm 1	36	36	26	28	26	27	27	30	28	29
(22) MCQN-BFGS	36	36	26	28	26	27	27	30	28	29
(22) L-BFGS	13	13	16	16	17	16	16	15	19	19
(23) Algorithm 1	13	17	12	21	12	14	8	11	10	9
(23) MCQN-BFGS	16	18	22	25	15	13	12	13	12	10
(23) L-BFGS	14	14	16	15	17	24	27	27	27	31
(24) Algorithm 1	11	13	18	24	31	37	28	21	13	7
(24) MCQN-BFGS	12	15	19	25	32	37	28	21	13	7
(24) L-BFGS	10	10	10	10	10	10	10	10	10	10

For the sake of precise comparison, we adopted the performance profiles from [45], which are distribution functions of a performance metric. We denote P and S as the test set and the set of solvers; and N_p and N_s as the umber of problems and number of solvers, respectively. For solver $s \in S$ and problem $p \in P$, we define $t_{p,s}$ as the number of iterations or number of function evaluations required for solve problem p using solver s. Then, using the performance ration

$$r_{p,s} = \frac{t_{p,s}}{\min\{t_{p,q} : q \in S\}},$$

we define

$$\rho_s(t) = \frac{1}{N_p} \operatorname{size} \{ p \in P : r_{p,s} \le t \},\,$$

where $r_{p,s} \le r_M$ for some constant for all p and s. The equality holds if and only if solver s cannot solve problem p. Therefore, $\rho_s : R \to [0,1]$ was the probability for $s \in S$ satisfying $r_{p,s} \le t$, $t \in R$ among the best possible ratios.

Figure 1 evaluates the number of iterations of and the MCQN update with BFGS method by using performance profiles. It can be seen that the top curve corresponds to Algorithm 1, which shows that Algorithm 1 had better performance than the MCQN update with BFGS method. Additionally, Figure 2 demonstrates that Algorithm 1 had better performance than the limited-memory BFGS method.

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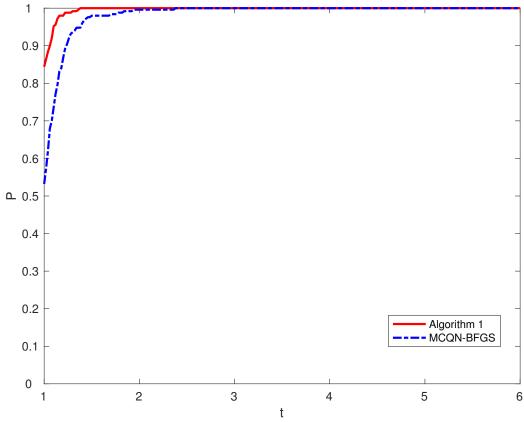


Figure 1. Performance profiles based on the numbers of iterations.

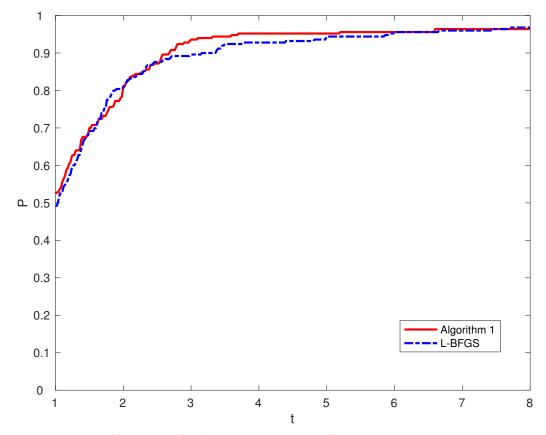


Figure 2. Performance profiles based on the numbers of iterations.

Secondly, for a further comparison of Algorithm 1 and the MCQN update with BFGS method, we tested five different initial points, x_0 , $2x_0$, $4x_0$, $7x_0$, and $10x_0$, where x_0 is specified in Table 1. The dimensions of the test problems was 1000. Table 4 reports the number of iterations required of the two methods for 24 test problems, which also demonstrates that Algorithm 1 was effective and superior to the MCQN update with BFGS method.

Table 4. Res	sults of Dim =	= 1000 with	different initial	points.
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Pro	Algorithm 1						MCQN-BFGS					
Init	x_0	$2x_0$	$4x_0$	$7x_0$	$10x_0$	x_0	$2x_0$	$4x_0$	$7x_0$	$10x_0$		
(1)	217	189	192	194	196	192	210	213	220	213		
(2)	2699	2684	2641	1192	2694	3249	4850	5056	2157	4961		
(3)	49	47	55	68	65	54	213	210	228	294		
(4)	43	47	36	81	80	60	213	210	228	294		
(5)	92	83	86	91	89	112	106	85	94	94		
(6)	47	47	47	47	47	54	54	54	54	54		
(7)	16	30	53	19	21	15	19	27	21	22		
(8)	27	31	16	35	54	28	31	16	33	56		
(9)	148	159	154	184	157	168	151	147	191	154		
(10)	45	45	164	170	175	41	263	203	192	194		
(11)	35	70	96	112	127	35	70	96	112	127		
(12)	6	6	2	2	2	6	6	2	2	2		
(13)	15	16	17	14	14	15	16	17	14	14		
(14)	612	312	504	511	318	594	523	503	481	563		
(15)	60	57	208	452	1024	64	58	203	532	941		
(16)	165	180	173	153	129	164	183	191	197	215		
(17)	12	11	10	7	21	12	11	8	7	30		
(18)	22	23	25	60	41	20	23	30	65	75		
(19)	12	12	14	16	24	12	12	14	22	20		
(20)	4	6	5	6	4	4	6	5	6	4		
(21)	35	30	35	112	719	35	30	35	112	719		
(22)	27	28	28	29	29	27	28	28	29	29		
(23)	8	22	23	14	35	12	13	24	19	70		
(24)	28	20	21	21	21	28	20	21	21	21		

6. Conclusions

In this paper, we presented a symmetric rank-two quasi-Newton update method based on an adjoint tangent condition for solving unconstrained optimization problems. Combined with the idea of matrix completion, we proposed a sparse quasi-Newton algorithm and established its local and superlinear convergence. Extensive numerical results demonstrated that the proposed algorithm outperformed other methods and can be used to solve large-scale unconstrained optimization problems.

Author Contributions: Conceptualization, H.C.; methodology, H.C. and X.A.; software, H.C. and X.A.; formal analysis, H.C.; writing—original draft preparation, H.C. and X.A.; writing—review and editing H.C. and X.A. All authors have read and agreed to the published version of the manuscript.

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Institutional Review Board Statement: Not applicable.

Informed Consent Statement: Not applicable.

Data Availability Statement: The date used to support the research plan and all the code used in this study are available from the corresponding author upon request.

Conflicts of Interest: The authors declare no conflict of interest.

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