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HW4

CS231

1. a) $E = \{ (a,b), (a,d), (b,c), (b,f), (c,d), (d,e), (e,f), (f,a), (e,b) \}$ 9 edges

b)

vertex	indegree	outdegree
a	1	2
b	2	2
c	1	1
d	2	1
e	1	2
f	2	1

c)

	a	b	c	d	e	f
a	0	1	0	1	0	0
b	0	0	1	0	0	1
c	0	0	0	1	0	0
d	0	0	0	0	1	0
e	0	1	0	0	0	1
f	1	0	0	0	0	0

d)

$a \rightarrow b \rightarrow d$
 $b \rightarrow c \rightarrow f$
 $c \rightarrow d$
 $d \rightarrow e$
 $e \rightarrow b \rightarrow f$
 $f \rightarrow a$

e) Distinct paths from e to d

1. e, b, c, d
2. e, f, a, d
3. e, f, a, b, c, d
4. e, b, f, a, d

f)

a) 3 length for shortest

b) e, b, c, d
e, f, a, d

g) end/start with d

a) 4

b) d, e, f, a, b, c, d

d, e, b, c, d

d, e, f, a, d

d, e, b, f, a, d

2.



initial graph

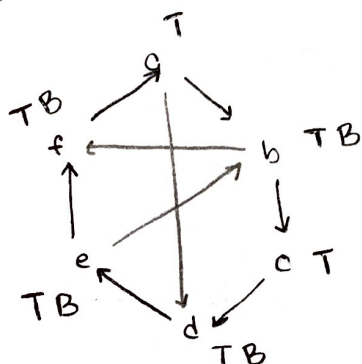
A. Discovered, in order:



e, b, c, d, e, f, a, b, d, f
e, b, c, d, f, a

assuming we don't show vertices when we call them back recursively

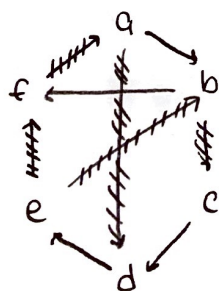
B.



3.

A. Vertex Labels

	shortest-path length	predecessor
a	2	f
b	1	e
c	2	b
d	3	a/a (2+1)
e	0	none
f	1	e



5. If a given non-empty directed graph has at least 1 sink, 1 source, it must be a DAG.

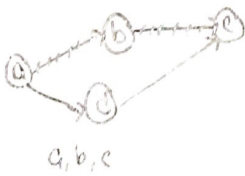
↑
FALSE



Consider the above graph, where we have a source a and sink d. We have an intermediate vertex c that cycles back to vertex b. Because of this, it is not acyclic and as a result is not a DAG, thus proving the claim to be false.

4.

A.



B.



Adjacency List 1

$a \rightarrow b \rightarrow d$
 $b \rightarrow c$
 c
 $d \rightarrow c$

source vertex a:
 destination: c

$\langle a \rangle$
 $\langle b, d \rangle$
 $\langle c \rangle$
 $\langle c \rangle$

Adjacency List 2

$a \rightarrow d \rightarrow b$
 $b \rightarrow c$
 c
 $d \rightarrow c$

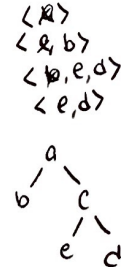
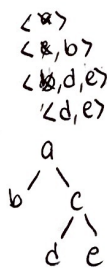
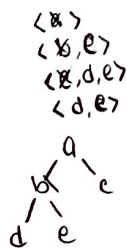
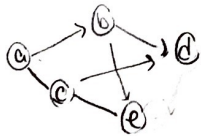
source vertex a:
 destination: c

$\langle a \rangle$
 $\langle b \rangle$
 $\langle b, c \rangle$
 $\langle c \rangle$

C. $G = (V, E)$ source $s \in V$

C.1. To show the above statement is false, I will construct a graph $G = (V, E)$ and specifying a source $s \in V$ such that we will have multiple nodes to the destinations d and e and both destinations cannot show their shortest paths for a certain tree, thus no matter the adjacency list, we cannot use BFS to represent it.

C.2.



Consider the possible shortest path tree:



This cannot be represented by a BFS approach, no matter how we set the adjacency list.