

Large Graph Mining: Power Tools and a Practitioner's guide

Task 1: Node importance

Faloutsos, Miller, Tsourakakis

CMU



Outline

• Introduction – Motivation

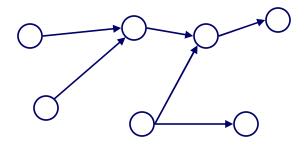


- Task 1: Node importance
- Task 2: Community detection
- Task 3: Recommendations
- Task 4: Connection sub-graphs
- Task 5: Mining graphs over time
- Task 6: Virus/influence propagation
- Task 7: Spectral graph theory
- Task 8: Tera/peta graph mining: hadoop
- Observations patterns of real graphs
- Conclusions



Node importance - Motivation:

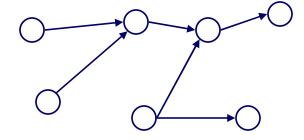
- Given a graph (eg., web pages containing the desirable query word)
- Q: Which node is the most important?





Node importance - Motivation:

- Given a graph (eg., web pages containing the desirable query word)
- Q: Which node is the most important?
- A1: HITS (SVD = Singular Value Decomposition)
- A2: eigenvector (PageRank)





Node importance - motivation

- SVD and eigenvector analysis: very closely related
- See 'theory Task', later



SVD - Detailed outline



- Motivation
- Definition properties
- Interpretation
- Complexity
- Case studies



- problem #1: text LSI: find 'concepts'
- problem #2: compression / dim. reduction



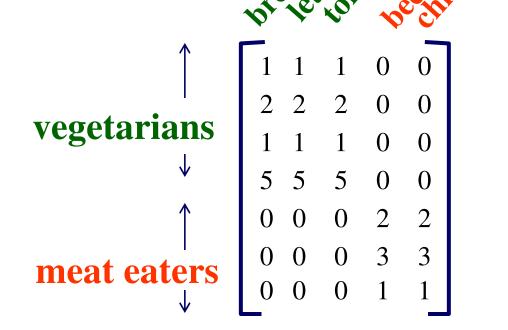
• problem #1: text - LSI: find 'concepts'

\mathbf{term}	data	information	retrieval	brain	lung
document					
CS-TR1	1	1	1	0	0
CS-TR2	2	2	2	0	0
CS-TR3	1	1	1	0	0
CS-TR4	5	5	5	0	0
MED-TR1	0	0	0	2	2
MED-TR2	0	0	0	3	3
MED-TR3	0	0	0	1	1



Customer-product, for recommendation

system:







• problem #2: compress / reduce dimensionality





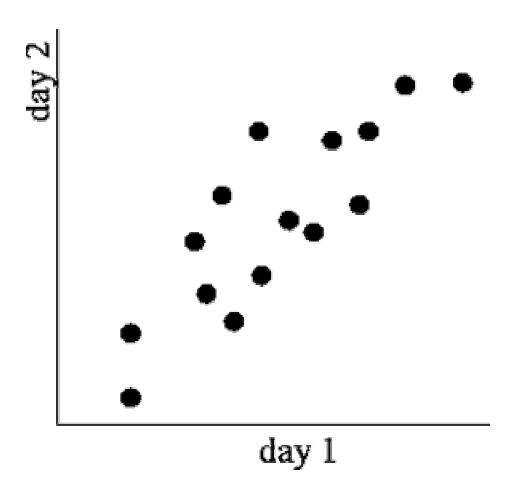
Problem - specs

- $\sim 10**6$ rows; $\sim 10**3$ columns; no updates;
- random access to any cell(s); small error: OK

day	Wc	\mathbf{Th}	$\mathbf{F}\mathbf{r}$	\mathbf{Sa}	Su
customer	7/10/96	7/11/96	7/12/96	7/13/96	7/14/96
ABC Inc.	1	1	1	0	0
DEF Ltd.	2	2	2	0	0
GHI Inc.	1	1	1	0	0
KLM Co.	5	5	5	0	0
Smith	0	0	0	2	2
Johnson	0	0	0	3	3
Thompson	0	0	0	1	1

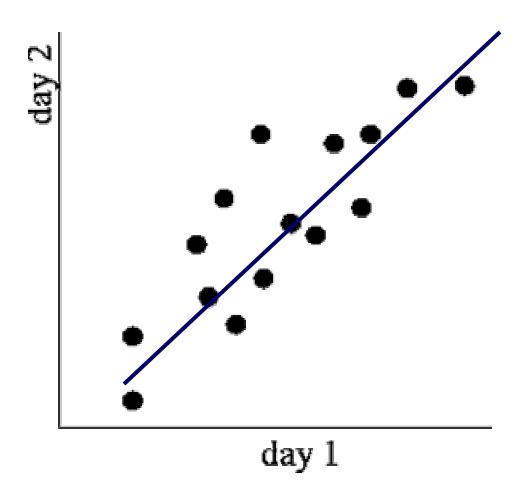














SVD - Detailed outline

Motivation



- Definition properties
 - Interpretation
 - Complexity
 - Case studies
 - Additional properties



(reminder: matrix multiplication

$$\begin{bmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{bmatrix} \quad x \quad \begin{bmatrix} 1 \\ -1 \end{bmatrix} = \begin{bmatrix} 1 \\ -1 \end{bmatrix}$$

 3×2 2×1

KDD'09



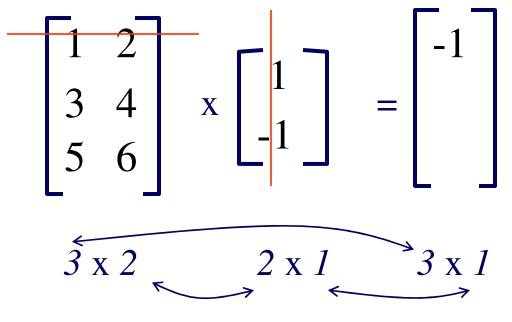
(reminder: matrix multiplication

$$\begin{bmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{bmatrix} \times \begin{bmatrix} 1 \\ -1 \end{bmatrix} = \begin{bmatrix} \\ \\ \end{bmatrix}$$

$$3 \times 2 \times 1 \quad 3 \times 1$$



(reminder: matrix multiplication

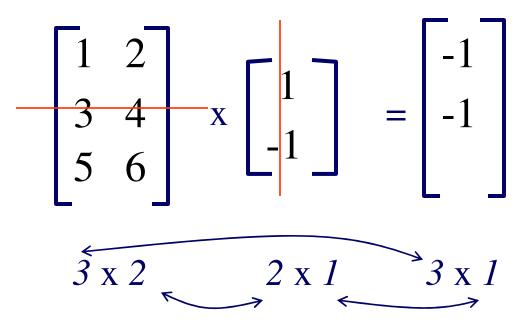


KDD'09

Faloutsos, Miller, Tsourakakis



(reminder: matrix multiplication





(reminder: matrix multiplication

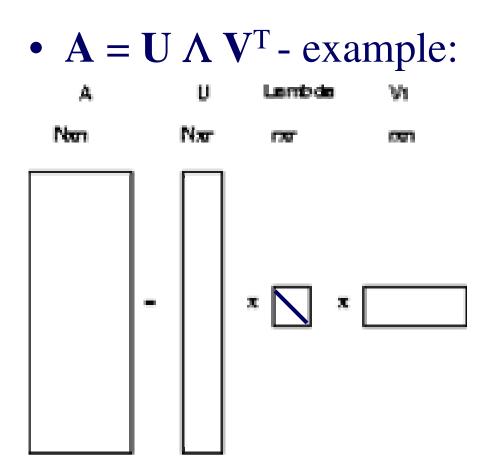
$$\begin{bmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{bmatrix} \times \begin{bmatrix} 1 \\ -1 \\ -1 \end{bmatrix} = \begin{bmatrix} -1 \\ -1 \\ -1 \end{bmatrix}$$



$\mathbf{A}_{[\mathbf{n} \times \mathbf{m}]} = \mathbf{U}_{[\mathbf{n} \times \mathbf{r}]} \Lambda_{[\mathbf{r} \times \mathbf{r}]} (\mathbf{V}_{[\mathbf{m} \times \mathbf{r}]})^{\mathrm{T}}$

- A: n x m matrix (eg., n documents, m terms)
- U: n x r matrix (n documents, r concepts)
- Λ : r x r diagonal matrix (strength of each 'concept') (r : rank of the matrix)
- V: m x r matrix (m terms, r concepts)







SVD - Properties

THEOREM [Press+92]: always possible to decompose matrix \mathbf{A} into $\mathbf{A} = \mathbf{U} \Lambda \mathbf{V}^{\mathrm{T}}$, where

- **U**, Λ, **V**: unique (*)
- U, V: column orthonormal (ie., columns are unit vectors, orthogonal to each other)
 - $\mathbf{U}^{\mathrm{T}} \mathbf{U} = \mathbf{I}$; $\mathbf{V}^{\mathrm{T}} \mathbf{V} = \mathbf{I}$ (I: identity matrix)
- A: singular are positive, and sorted in decreasing order

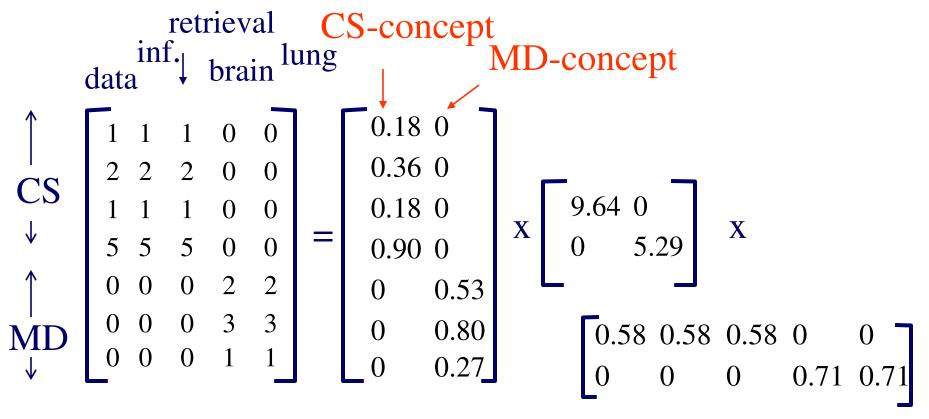


• $\mathbf{A} = \mathbf{U} \mathbf{\Lambda} \mathbf{V}^{\mathrm{T}}$ - example:

retrieval

$$= \begin{bmatrix} 0.18 & 0 \\ 0.36 & 0 \\ 0.18 & 0 \\ 0.90 & 0 \\ 0 & 0.53 \\ 0 & 0.80 \\ 0 & 0.27 \end{bmatrix}$$

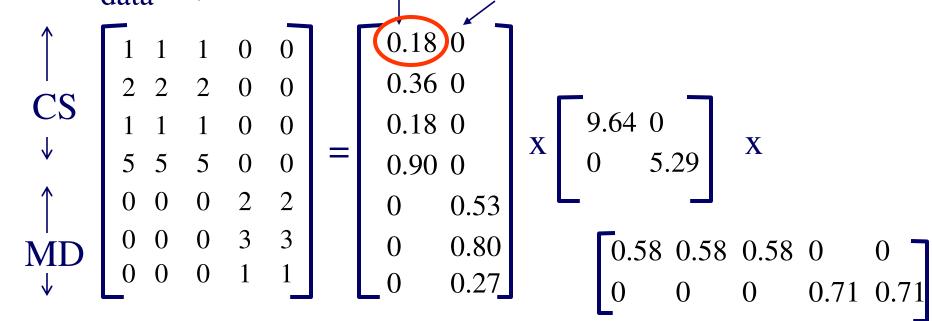




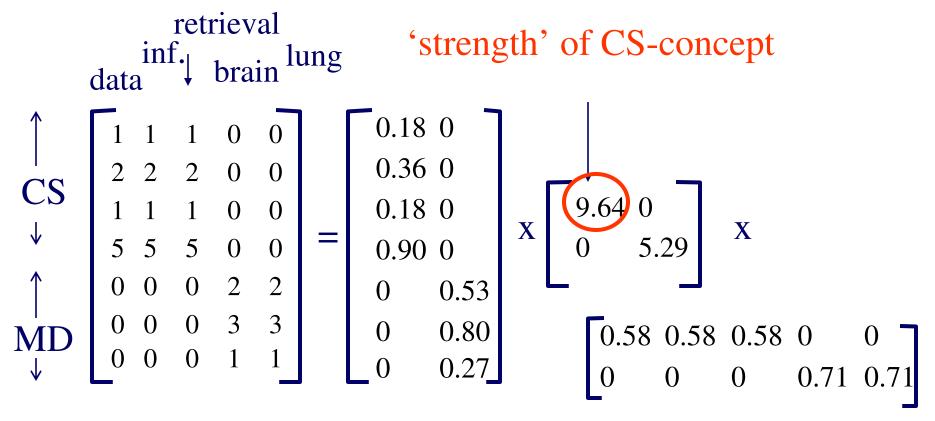


• $A = U \wedge V^T$ - example: doc-to-concept similarity matrix

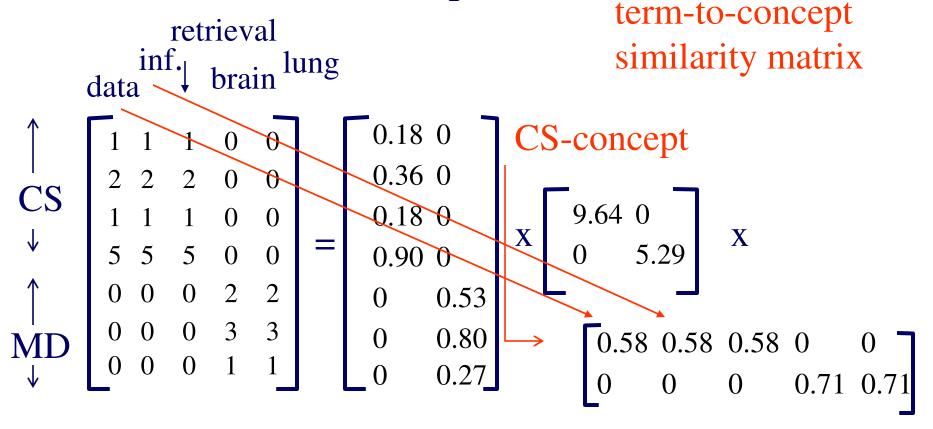
retrieval CS-concept inf. brain lung MD-concept



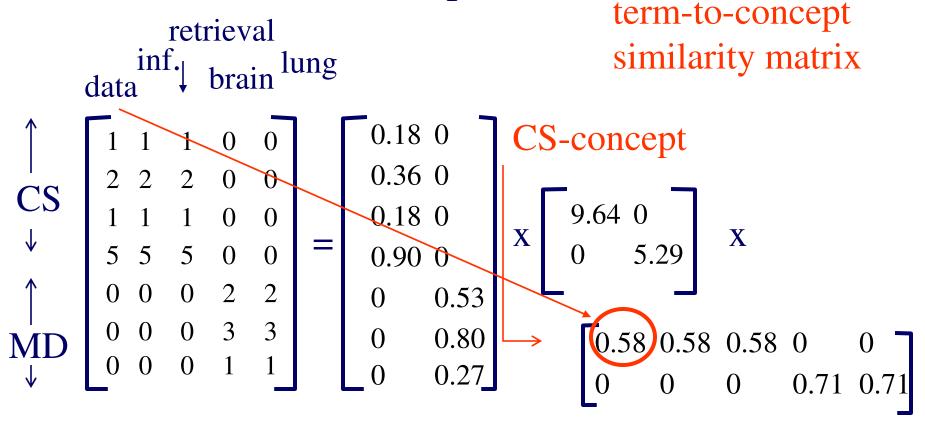














SVD - Detailed outline

- Motivation
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SVD - Interpretation #1

'documents', 'terms' and 'concepts':

- U: document-to-concept similarity matrix
- V: term-to-concept sim. matrix
- Λ: its diagonal elements: 'strength' of each concept



SVD – Interpretation #1

'documents', 'terms' and 'concepts':

Q: if A is the document-to-term matrix, what is A^T A?

A:

 $Q: A A^T$?

A:



SVD – Interpretation #1

'documents', 'terms' and 'concepts':

Q: if A is the document-to-term matrix, what is A^T A?

A: term-to-term ([m x m]) similarity matrix

 $Q: A A^T$?

A: document-to-document ([n x n]) similarity matrix



SVD properties

• V are the eigenvectors of the *covariance* $matrix A^{T}A$

• U are the eigenvectors of the Gram (inner-product) matrix $\mathbf{A}\mathbf{A}^{T}$

Further reading:

- 1. Ian T. Jolliffe, *Principal Component Analysis* (2nd ed), Springer, 2002.
- 2. Gilbert Strang, Linear Algebra and Its Applications (4th ed), Brooks Cole, 2005.



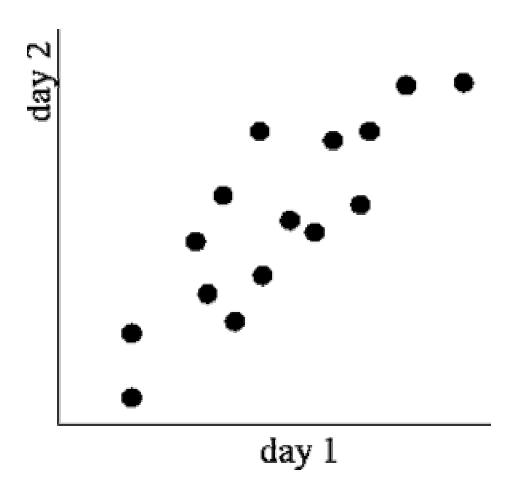


SVD - Interpretation #2

• best axis to project on: ('best' = min sum of squares of projection errors)





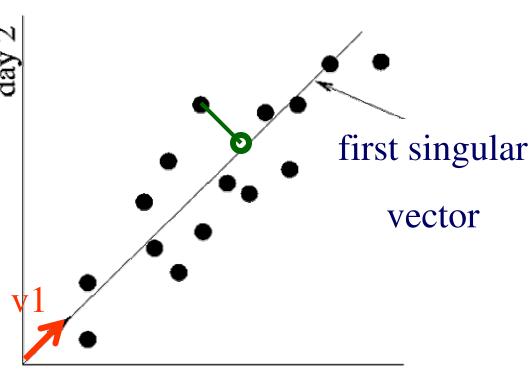






SVD - interpretation #2

SVD: gives best axis to project



minimum RMS error

day 1





day	We	\mathbf{Th}	\mathbf{Fr}	\mathbf{Sa}	Su
customer	7/10/96	7/11/96	7/12/96	7/13/96	7/14/96
ABC Inc.	1	1	1	0	0
DEF Ltd.	2	2	2	0	0
GHI Inc.	1	1	1	0	0
KLM Co.	5	5	5	0	0
Smith	0	0	0	2	2
Johnson	0	0	0	3	3
Thompson	0	0	0	1	1





• $\mathbf{A} = \mathbf{U} \boldsymbol{\Lambda} \mathbf{V}^{\mathrm{T}}$ - example:

$$\begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 2 & 2 & 2 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 5 & 5 & 5 & 0 & 0 \\ 0 & 0 & 0 & 2 & 2 \\ 0 & 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 0.18 & 0 \\ 0.36 & 0 \\ 0.18 & 0 \\ 0.90 & 0 \\ 0 & 0.53 \\ 0 & 0.80 \\ 0 & 0.27 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 5.29 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 5.29 \end{bmatrix} \times \begin{bmatrix} 0.58 & 0.58 & 0.58 & 0 \\ 0 & 0 & 0.71 & 0.71 \end{bmatrix}$$





• $\mathbf{A} = \mathbf{U} \boldsymbol{\Lambda} \mathbf{V}^{\mathrm{T}}$ - example:

variance ('spread') on the v1 axis

$$\begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 2 & 2 & 2 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 5 & 5 & 5 & 0 & 0 \\ 0 & 0 & 0 & 2 & 2 \\ 0 & 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 0.18 & 0 \\ 0.36 & 0 \\ 0.18 & 0 \\ 0.90 & 0 \\ 0 & 0.53 \\ 0 & 0.80 \\ 0 & 0.27 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 5.29 \end{bmatrix} \times \begin{bmatrix} 0.58 & 0.58 & 0.58 & 0 & 0 \\ 0 & 0 & 0.71 & 0.71 \end{bmatrix}$$





- $\mathbf{A} = \mathbf{U} \boldsymbol{\Lambda} \mathbf{V}^{\mathrm{T}}$ example:
 - U Λ gives the coordinates of the points in the projection axis

$$\begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 2 & 2 & 2 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 5 & 5 & 5 & 0 & 0 \\ 0 & 0 & 0 & 2 & 2 \\ 0 & 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 0.18 & 0 \\ 0.36 & 0 \\ 0.18 & 0 \\ 0.90 & 0 \\ 0 & 0.53 \\ 0 & 0.80 \\ 0 & 0.27 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 5.29 \end{bmatrix} \times \begin{bmatrix} 0.58 & 0.58 & 0.58 & 0 & 0 \\ 0 & 0 & 0.71 & 0.71 \end{bmatrix}$$





- More details
- Q: how exactly is dim. reduction done?

$$\begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 2 & 2 & 2 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 5 & 5 & 5 & 0 & 0 \\ 0 & 0 & 0 & 2 & 2 \\ 0 & 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 0.18 & 0 \\ 0.36 & 0 \\ 0.18 & 0 \\ 0.90 & 0 \\ 0 & 0.53 \\ 0 & 0.80 \\ 0 & 0.27 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 5.29 \end{bmatrix} \times \begin{bmatrix} 0.58 & 0.58 & 0.58 & 0 \\ 0 & 0 & 0.71 & 0.71 \end{bmatrix}$$





- More details
- Q: how exactly is dim. reduction done?
- A: set the smallest singular values to zero:

$$\begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 2 & 2 & 2 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 5 & 5 & 5 & 0 & 0 \\ 0 & 0 & 0 & 2 & 2 \\ 0 & 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 0.18 & 0 \\ 0.36 & 0 \\ 0.18 & 0 \\ 0.90 & 0 \\ 0 & 0.53 \\ 0 & 0.80 \\ 0 & 0.27 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 5.29 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 5.29 \end{bmatrix} \times \begin{bmatrix} 0.58 & 0.58 & 0.58 & 0 & 0 \\ 0 & 0 & 0.71 & 0.71 \end{bmatrix}$$

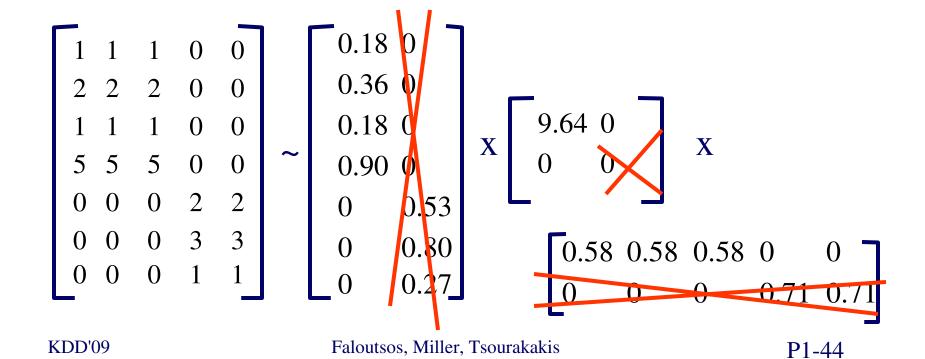




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\begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 2 & 2 & 2 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 5 & 5 & 5 & 0 & 0 \\ 0 & 0 & 0 & 2 & 2 \\ 0 & 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix} \sim \begin{bmatrix} 0.18 & 0 \\ 0.36 & 0 \\ 0.18 & 0 \\ 0.90 & 0 \\ 0 & 0.53 \\ 0 & 0.80 \\ 0 & 0.27 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 0 \end{bmatrix} \times \begin{bmatrix}
```















```
      1
      1
      1
      0
      0

      2
      2
      2
      0
      0

      1
      1
      1
      0
      0

      5
      5
      5
      0
      0

      0
      0
      0
      2
      2

      0
      0
      0
      3
      3

      0
      0
      0
      1
      1
```

```
      1
      1
      1
      0
      0

      2
      2
      2
      0
      0

      1
      1
      1
      0
      0

      5
      5
      5
      0
      0

      0
      0
      0
      0
      0

      0
      0
      0
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      0

      0
      0
      0
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      0

      0
      0
      0
      0
      0
```





Exactly equivalent:

$$\begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 2 & 2 & 2 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 5 & 5 & 5 & 0 & 0 \\ 0 & 0 & 0 & 2 & 2 \\ 0 & 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 0.18 & 0 \\ 0.36 & 0 \\ 0.18 & 0 \\ 0.90 & 0 \\ 0 & 0.53 \\ 0 & 0.80 \\ 0 & 0.27 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 5.29 \end{bmatrix} \times \begin{bmatrix} 0.58 & 0.58 & 0.58 & 0 \\ 0 & 0 & 0.71 & 0.71 \end{bmatrix}$$



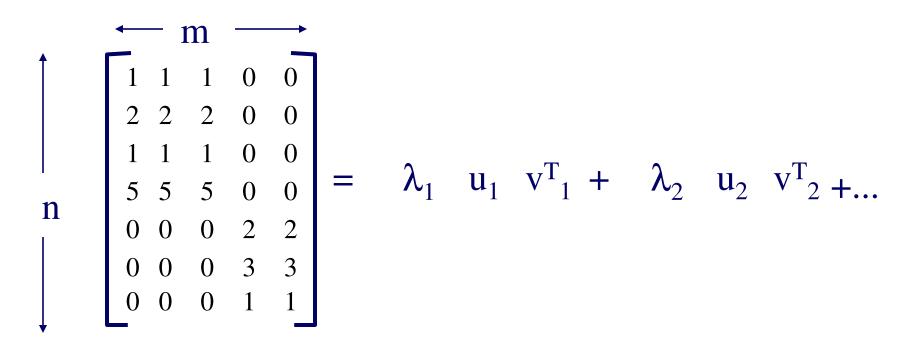


Exactly equivalent:





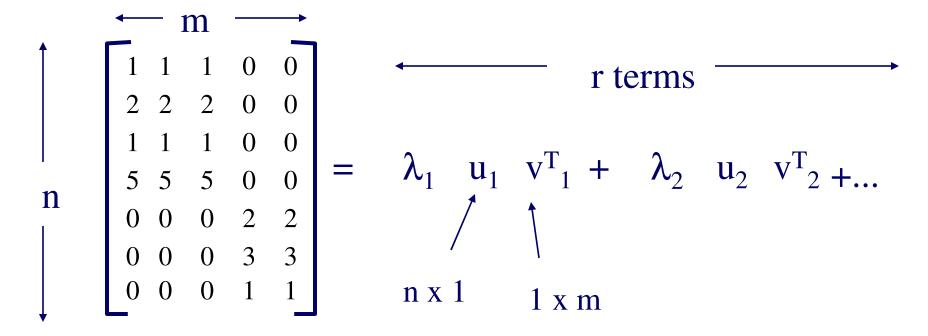
Exactly equivalent:







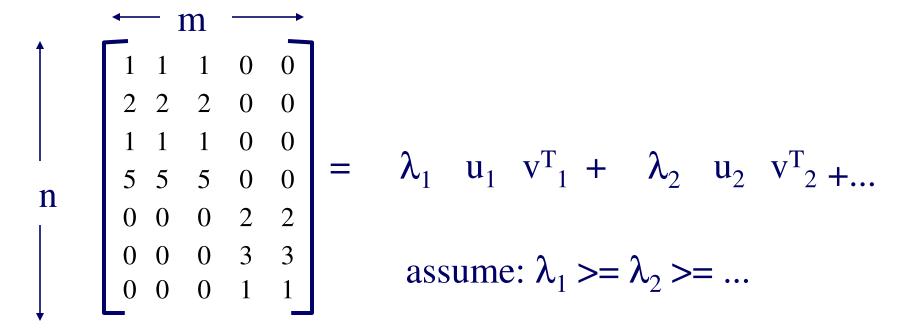
Exactly equivalent:







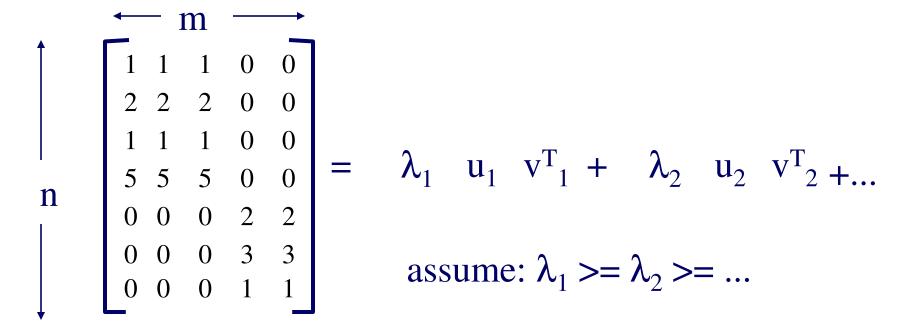
approximation / dim. reduction: by keeping the first few terms (Q: how many?)







A (heuristic - [Fukunaga]): keep 80-90% of 'energy' (= sum of squares of λ_i 's)





SVD - Detailed outline

- Motivation
- Definition properties
- Interpretation
 - #1: documents/terms/concepts
 - #2: dim. reduction



- #3: picking non-zero, rectangular 'blobs'
- Complexity
- Case studies
- Additional properties



• finds non-zero 'blobs' in a data matrix

$$\begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 2 & 2 & 2 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 5 & 5 & 5 & 0 & 0 \\ 0 & 0 & 0 & 2 & 2 \\ 0 & 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 0.18 & 0 \\ 0.36 & 0 \\ 0.18 & 0 \\ 0.90 & 0 \\ 0 & 0.53 \\ 0 & 0.80 \\ 0 & 0.27 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 5.29 \end{bmatrix} \times \begin{bmatrix} 0.58 & 0.58 & 0.58 & 0 \\ 0 & 0 & 0.71 & 0.71 \end{bmatrix}$$



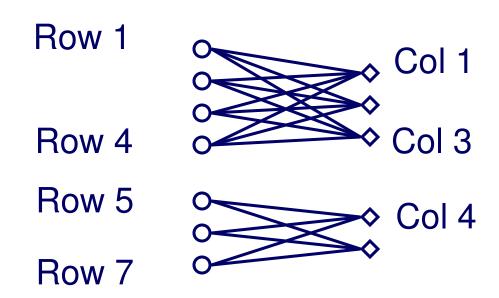
• finds non-zero 'blobs' in a data matrix

$$\begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 2 & 2 & 2 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 \\ 5 & 5 & 5 & 0 & 0 \\ \hline 0 & 0 & 0 & 2 & 2 \\ 0 & 0 & 0 & 3 & 3 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 0.18 & 0 \\ 0.36 & 0 \\ 0.18 & 0 \\ 0.90 & 0 \\ 0 & 0.53 \\ 0 & 0.80 \\ 0 & 0.27 \end{bmatrix} \times \begin{bmatrix} 9.64 & 0 \\ 0 & 5.29 \end{bmatrix} \times \begin{bmatrix} 0.58 & 0.58 & 0.58 & 0 & 0 \\ 0 & 0 & 0.71 & 0.71 \end{bmatrix}$$



- finds non-zero 'blobs' in a data matrix =
- 'communities' (bi-partite cores, here)

0	0 0
0	0
0	0
2	2
3	3
4	1





SVD - Detailed outline

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SVD - Complexity

- O(n * m * m) or O(n * n * m) (whichever is less)
- less work, if we just want singular values
- or if we want first *k* singular vectors
- or if the matrix is sparse [Berry]
- Implemented: in any linear algebra package (LINPACK, matlab, Splus, mathematica ...)



SVD - conclusions so far

- SVD: $A = U \wedge V^T$: unique (*)
- U: document-to-concept similarities
- V: term-to-concept similarities
- A: strength of each concept
- dim. reduction: keep the first few strongest singular values (80-90% of 'energy')
 - SVD: picks up linear correlations
- SVD: picks up non-zero 'blobs'





SVD - Detailed outline

- Motivation
- Definition properties
- Interpretation
- Complexity



- Case studies
- Conclusions



SVD - Other properties - summary

- can produce orthogonal basis (obvious) (who cares?)
- can solve over- and under-determined linear problems (see C(1) property)
- can compute 'fixed points' (= 'steady state prob. in Markov chains') (see C(4) property)





SVD -outline of properties

- (A): obvious
- (B): less obvious
- (C): least obvious (and most powerful!)





Properties - by defn.:

$$\mathbf{A}(0): \mathbf{A}_{[n \times m]} = \mathbf{U}_{[n \times r]} \mathbf{\Lambda}_{[r \times r]} \mathbf{V}^{\mathbf{T}}_{[r \times m]}$$

A(1):
$$\mathbf{U}^{\mathrm{T}}_{[r \times n]} \mathbf{U}_{[n \times r]} = \mathbf{I}_{[r \times r]}$$
 (identity matrix)

A(2):
$$V^{T}_{[r \times n]} V_{[n \times r]} = I_{[r \times r]}$$

A(3):
$$\Lambda^k = \text{diag}(\lambda_1^k, \lambda_2^k, ... \lambda_r^k)$$
 (k: ANY real number)

$$A(4)$$
: $A^T = V \wedge U^T$





$$\mathbf{A}(0): \mathbf{A}_{[n \times m]} = \mathbf{U}_{[n \times r]} \mathbf{\Lambda}_{[r \times r]} \mathbf{V}^{\mathbf{T}}_{[r \times m]}$$

B(1):
$$\mathbf{A}_{[n \times m]} (\mathbf{A}^T)_{[m \times n]} = ??$$





A(0):
$$\mathbf{A}_{[n \times m]} = \mathbf{U}_{[n \times r]} \Lambda_{[r \times r]} \mathbf{V}^{T}_{[r \times m]}$$

B(1): $\mathbf{A}_{[n \times m]} (\mathbf{A}^{T})_{[m \times n]} = \mathbf{U} \Lambda^{2} \mathbf{U}^{T}$
symmetric; Intuition?





$$\mathbf{A}(0): \mathbf{A}_{[n \times m]} = \mathbf{U}_{[n \times r]} \mathbf{\Lambda}_{[r \times r]} \mathbf{V}^{\mathbf{T}}_{[r \times m]}$$

B(1):
$$\mathbf{A}_{[n \times m]}(\mathbf{A}^T)_{[m \times n]} = \mathbf{U} \Lambda^2 \mathbf{U}^T$$

symmetric; Intuition?

'document-to-document' similarity matrix

B(2): symmetrically, for 'V'

(AT)
$$[m \times n] A [n \times m] = V L2 VT$$

Intuition?





A: term-to-term similarity matrix

B(3): $((A^T)_{[m \times n]} A_{[n \times m]})^k = V \Lambda^{2k} V^T$ and

B(4): $(A^T A)^k \sim v_1 \lambda_1^{2k} v_1^T$ for k >> 1 where

 \mathbf{v}_1 : [m x 1] first column (singular-vector) of \mathbf{V}

 λ_1 : strongest singular value





B(4): $(A^TA)^k \sim v_1 \lambda_1^{2k} v_1^T \text{ for } k >> 1$

B(5): $(\mathbf{A}^T \mathbf{A})^k \mathbf{v}' \sim \text{(constant)} \mathbf{v}_1$

ie., for (almost) any \mathbf{v} , it converges to a vector parallel to \mathbf{v}_1

Thus, useful to compute first singular vector/value (as well as the next ones, too...)



Less obvious properties - repeated:



A(0):
$$\mathbf{A}_{[n \times m]} = \mathbf{U}_{[n \times r]} \mathbf{\Lambda}_{[r \times r]} \mathbf{V}^{T}_{[r \times m]}$$

B(1):
$$\mathbf{A}_{[n \times m]} (\mathbf{A}^T)_{[m \times n]} = \mathbf{U} \Lambda^2 \mathbf{U}^T$$

B(2):
$$(\mathbf{A}^{\mathrm{T}})_{[m \times n]} \mathbf{A}_{[n \times m]} = \mathbf{V} \Lambda^{2} \mathbf{V}^{\mathrm{T}}$$

B(3):
$$((A^T)_{[m \times n]} A_{[n \times m]})^k = V \Lambda^{2k} V^T$$

B(4):
$$(A^T A)^k \sim v_1 \lambda_1^{2k} v_1^T$$

B(5):
$$(\mathbf{A}^{\mathrm{T}}\mathbf{A})^{\mathrm{k}}\mathbf{v}' \sim \text{(constant)}\mathbf{v}_1$$



Least obvious properties - cont'd

A(0):
$$\mathbf{A}_{[n \times m]} = \mathbf{U}_{[n \times r]} \mathbf{\Lambda}_{[r \times r]} \mathbf{V}^{\mathbf{T}}_{[r \times m]}$$

C(2): $\mathbf{A}_{[n \times m]} \mathbf{v}_{1[m \times 1]} = \lambda_1 \mathbf{u}_{1[n \times 1]}$ where \mathbf{v}_1 , \mathbf{u}_1 the first (column) vectors of \mathbf{V} , \mathbf{U} . (\mathbf{v}_1 == right-singular-vector)

C(3): symmetrically: $\mathbf{u}_1^T \mathbf{A} = \lambda_1 \mathbf{v}_1^T$ $\mathbf{u}_1 == \text{left-singular-vector}$

Therefore:



Least obvious properties - cont'd

A(0):
$$\mathbf{A}_{[n \times m]} = \mathbf{U}_{[n \times r]} \mathbf{\Lambda}_{[r \times r]} \mathbf{V}^{\mathbf{T}}_{[r \times m]}$$

C(4):
$$A^T A v_1 = \lambda_1^2 v_1$$

(**fixed point** - the dfn of eigenvector for a symmetric matrix)







A(0):
$$\mathbf{A}_{[n \times m]} = \mathbf{U}_{[n \times r]} \mathbf{\Lambda}_{[r \times r]} \mathbf{V}^{\mathbf{T}}_{[r \times m]}$$

C(1): $\mathbf{A}_{[n \times m]} \mathbf{x}_{[m \times 1]} = \mathbf{b}_{[n \times 1]}$ then, $\mathbf{x}_0 = \mathbf{V} \Lambda^{(-1)} \mathbf{U}^T \mathbf{b}$: shortest, actual or least-squares solution

C(2):
$$\mathbf{A}_{[n \times m]} \mathbf{v}_{1[m \times 1]} = \lambda_1 \mathbf{u}_{1[n \times 1]}$$

$$C(3): \mathbf{u_1}^{\mathrm{T}} \mathbf{A} = \lambda_1 \mathbf{v_1}^{\mathrm{T}}$$

C(4):
$$A^T A v_1 = \lambda_1^2 v_1$$





Properties - conclusions

A(0):
$$\mathbf{A}_{[n \times m]} = \mathbf{U}_{[n \times r]} \mathbf{\Lambda}_{[r \times r]} \mathbf{V}^{\mathbf{T}}_{[r \times m]}$$

B(5):
$$(\mathbf{A}^{\mathrm{T}}\mathbf{A})^{\mathrm{k}}\mathbf{v}$$
 ~ (constant) \mathbf{v}_{1}

C(1):
$$\mathbf{A}_{[n \times m]} \mathbf{x}_{[m \times 1]} = \mathbf{b}_{[n \times 1]}$$

then, $\mathbf{x}_0 = \mathbf{V} \Lambda^{(-1)} \mathbf{U}^T \mathbf{b}$: shortest, actual or least-squares solution

C(4):
$$A^T A v_1 = \lambda_1^2 v_1$$



SVD - detailed outline

- •
- SVD properties
- case studies



- Kleinberg's algorithm
- Google's algorithm
- Conclusions



Kleinberg's algo (HITS)



Kleinberg, Jon (1998).

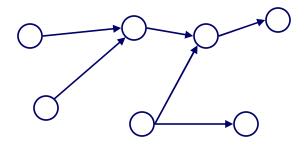
Authoritative sources in a hyperlinked environment.

Proc. 9th ACM-SIAM Symposium on Discrete Algorithms.



Recall: problem dfn

- Given a graph (eg., web pages containing the desirable query word)
- Q: Which node is the most important?





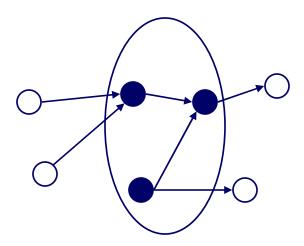
- Problem dfn: given the web and a query
- find the most 'authoritative' web pages for this query

Step 0: find all pages containing the query terms

Step 1: expand by one move forward and backward

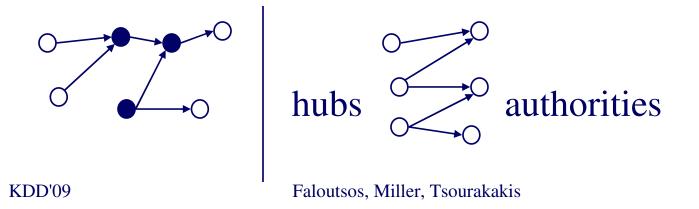


• Step 1: expand by one move forward and backward





- on the resulting graph, give high score (= 'authorities') to nodes that many important nodes point to
- give high importance score ('hubs') to nodes that point to good 'authorities')



P1-79



observations

- recursive definition!
- each node (say, 'i'-th node) has both an authoritativeness score a_i and a hubness score h_i



Let *E* be the set of edges and **A** be the adjacency matrix:

the (i,j) is 1 if the edge from i to j exists

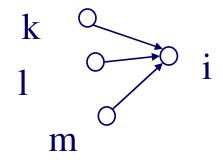
Let *h* and *a* be [n x 1] vectors with the 'hubness' and 'authoritativiness' scores.

Then:



Then:

$$a_i = h_k + h_l + h_m$$



that is

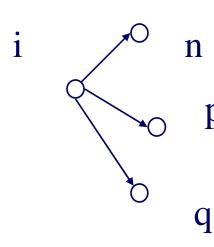
$$a_i = \text{Sum } (h_j)$$
 over all j that (j,i) edge exists

or

$$\mathbf{a} = \mathbf{A}^{\mathrm{T}} \mathbf{h}$$



symmetrically, for the 'hubness':



$$h_i = a_n + a_p + a_q$$

that is

 $h_i = \text{Sum}(q_j)$ over all j that (i,j) edge exists

or

$$h = A a$$



In conclusion, we want vectors **h** and **a** such that:

$$\mathbf{h} = \mathbf{A} \mathbf{a}$$

$$\mathbf{a} = \mathbf{A}^{\mathrm{T}} \mathbf{h}$$

Recall properties:

C(2):
$$\mathbf{A}_{[n \times m]} \mathbf{v}_{1[m \times 1]} = \lambda_1 \mathbf{u}_{1[n \times 1]}$$

C(3): $\mathbf{u}_1^T \mathbf{A} = \lambda_1 \mathbf{v}_1^T$



In short, the solutions to

$$h = A a$$

$$\mathbf{a} = \mathbf{A}^{\mathrm{T}} \mathbf{h}$$

are the <u>left- and right- singular-vectors</u> of the adjacency matrix **A**.

Starting from random a' and iterating, we'll eventually converge

(Q: to which of all the singular-vectors? why?)



(Q: to which of all the singular-vectors? why?)

A: to the ones of the strongest singular-value, because of property B(5):

B(5): $(\mathbf{A}^T \mathbf{A})^k \mathbf{v}' \sim \text{(constant)} \mathbf{v}_1$



Kleinberg's algorithm - results

Eg., for the query 'java':

- 0.328 www.gamelan.com
- 0.251 java.sun.com
- 0.190 www.digitalfocus.com ("the java developer")



Kleinberg's algorithm - discussion

• 'authority' score can be used to find 'similar pages' (how?)



SVD - detailed outline

- Complexity
- SVD properties
- Case studies
 - Kleinberg's algorithm (HITS)



- Google's algorithmConclusions



PageRank (google)



•Brin, Sergey and Lawrence Page (1998). *Anatomy of a Large-Scale Hypertextual Web Search Engine*. 7th Intl World Wide Web Conf.

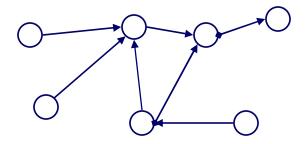
Larry Page

Sergey Brin



Problem: PageRank

Given a directed graph, find its most interesting/central node



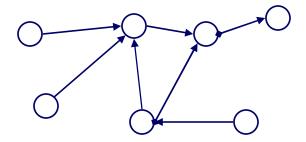
A node is important, if it is connected with important nodes (recursive, but OK!)



Problem: PageRank - solution

Given a directed graph, find its most interesting/central node

Proposed solution: Random walk; spot most 'popular' node (-> steady state prob. (ssp))

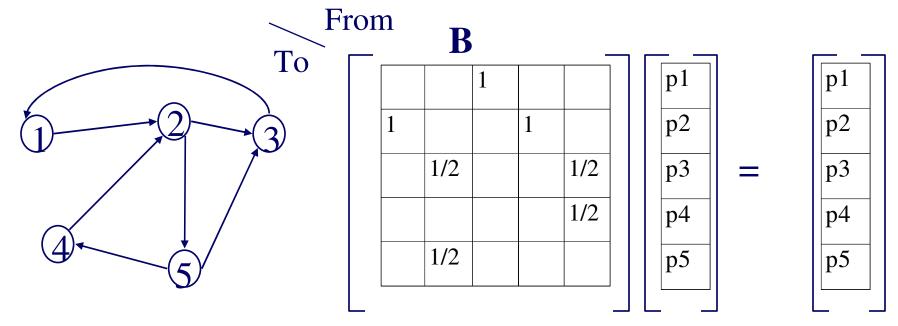


A node has high ssp, if it is connected with high ssp nodes (recursive, but OK!)



(Simplified) PageRank algorithm

- Let A be the adjacency matrix;
- let **B** be the transition matrix: transpose, column-normalized then



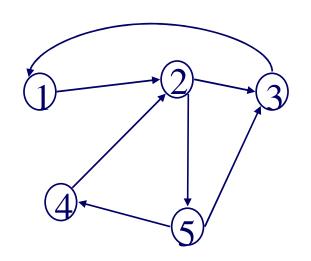


(Simplified) PageRank algorithm

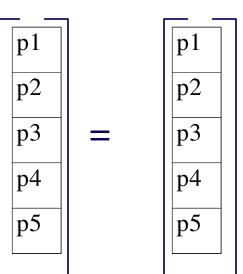
• $\mathbf{B} \mathbf{p} = \mathbf{p}$

B

 $\mathbf{p} = \mathbf{p}$



		1		
1			1	
	1/2			1/2
				1/2
	1/2			





Definitions

- **A** Adjacency matrix (from-to)
- **D** Degree matrix = (diag (d1, d2, ..., dn))
- **B** Transition matrix: to-from, column normalized

$$\mathbf{B} = \mathbf{A}^{\mathrm{T}} \mathbf{D}^{-1}$$



(Simplified) PageRank algorithm

- B p = 1 * p
- thus, **p** is the **eigenvector** that corresponds to the highest eigenvalue (=1, since the matrix is column-normalized)
- Why does such a **p** exist?
 - p exists if B is nxn, nonnegative, irreducible[Perron–Frobenius theorem]



(Simplified) PageRank algorithm

- In short: imagine a particle randomly moving along the edges
- compute its steady-state probabilities (ssp)

Full version of algo: with occasional random jumps

Why? To make the matrix irreducible

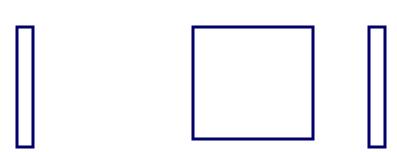


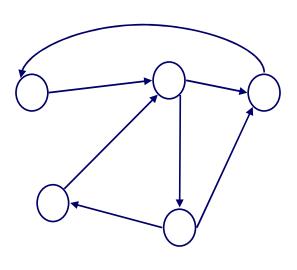
Full Algorithm

- With probability 1-c, fly-out to a random node
- Then, we have

$$p = c B p + (1-c)/n 1 =>$$

$$p = (1-c)/n [I - c B]^{-1} 1$$







Alternative notation

M Modified transition matrix

$$M = c B + (1-c)/n 1 1^T$$

Then

$$p = M p$$

That is: the steady state probabilities =

PageRank scores form the *first eigenvector* of the 'modified transition matrix'



Parenthesis: intuition behind eigenvectors



Formal definition

If A is a (n x n) square matrix (λ, \mathbf{x}) is an **eigenvalue/eigenvector** pair of A if

$$\mathbf{A} \mathbf{x} = \lambda \mathbf{x}$$

CLOSELY related to singular values:



Property #1: Eigen- vs singular-values

if

$$\mathbf{B}_{[\mathbf{n} \times \mathbf{m}]} = \mathbf{U}_{[\mathbf{n} \times \mathbf{r}]} \Lambda_{[\mathbf{r} \times \mathbf{r}]} (\mathbf{V}_{[\mathbf{m} \times \mathbf{r}]})^{\mathrm{T}}$$
then $\mathbf{A} = (\mathbf{B}^{\mathrm{T}}\mathbf{B})$ is symmetric and

C(4):
$$\mathbf{B}^{\mathbf{T}} \mathbf{B} \mathbf{v_i} = \lambda_i^2 \mathbf{v_i}$$

ie, v_1 , v_2 , ...: eigenvectors of $A = (B^TB)$



Property #2

- If $A_{[nxn]}$ is a real, symmetric matrix
- Then it has *n* real eigenvalues

(if **A** is not symmetric, some eigenvalues may be complex)



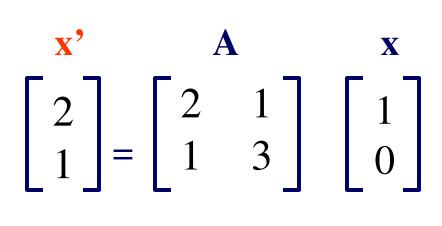
Property #3

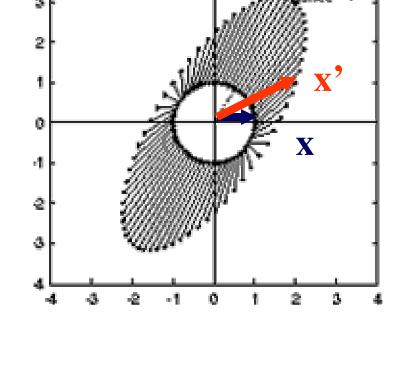
- If $A_{[nxn]}$ is a real, symmetric matrix
- Then it has *n* real eigenvalues
- And they agree with its *n* singular values, except possibly for the sign

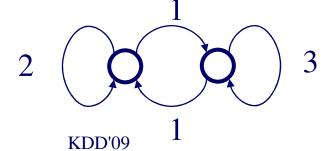


Intuition

• A as vector transformation





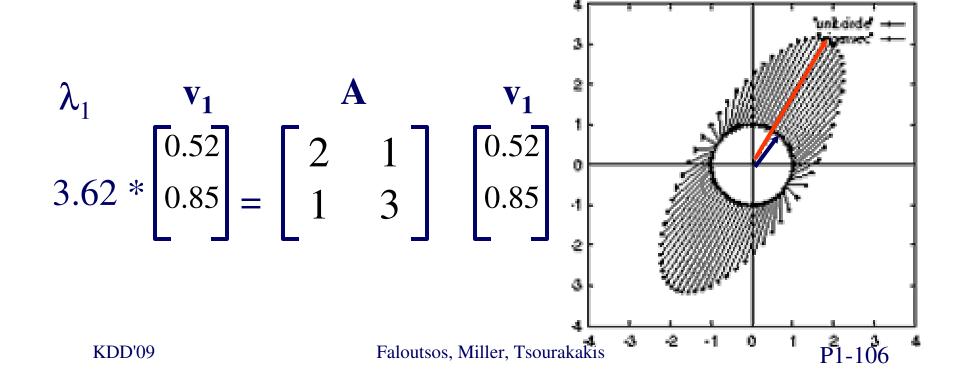


Faloutsos, Miller, Tsourakakis



Intuition

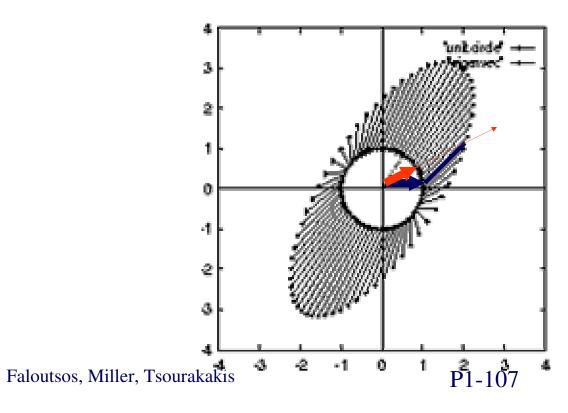
• By defn., eigenvectors remain parallel to themselves ('fixed points')





Convergence

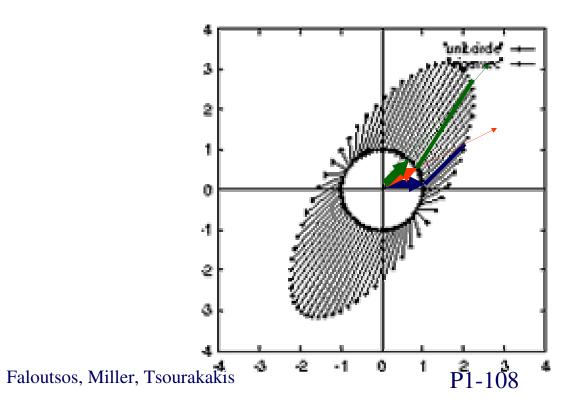
• Usually, fast:





Convergence

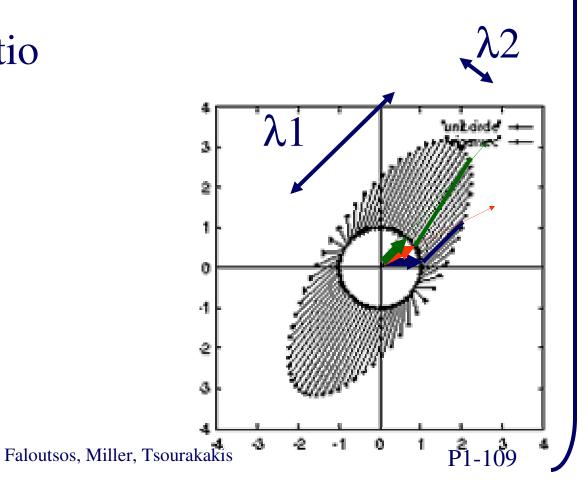
• Usually, fast:





Convergence

- Usually, fast:
- depends on ratio $\lambda 1 : \lambda 2$





Kleinberg/google - conclusions

SVD helps in graph analysis:

hub/authority scores: strongest left- and rightsingular-vectors of the adjacency matrix

random walk on a graph: steady state probabilities are given by the strongest eigenvector of the (modified) transition matrix



Conclusions

- SVD: a **valuable** tool
- given a document-term matrix, it finds 'concepts' (LSI)
- ... and can find fixed-points or steady-state probabilities (google/ Kleinberg/ Markov Chains)



Conclusions cont'd

(We didn't discuss/elaborate, but, SVD

- ... can reduce dimensionality (KL)
- ... and can find rules (PCA; RatioRules)
- ... and can solve optimally over- and underconstraint linear systems (least squares / query feedbacks)



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