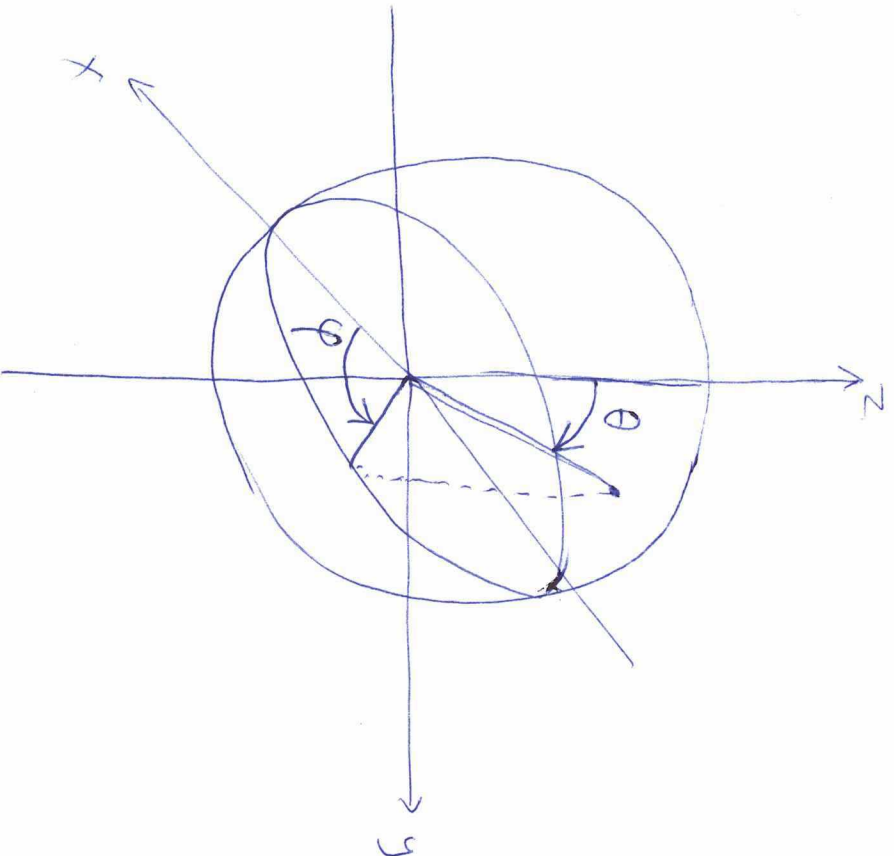


Bloch sphere, two qubits and Deutsch algorithmBloch sphere

- state of single qubit can be graphically represented on a so called Bloch sphere
- Bloch sphere is different than qsphere in ~~q~~ Quantum Composer



- the angles,  $\theta, \varphi$  fully parameterize the state of qubit
- $\theta \in [0, \pi]$ ,  $\varphi \in [0, 2\pi]$
- Positions of  $|0\rangle, |1\rangle, |+\rangle, |-\rangle$
- How gates work:  $X, Z, Y, H, R_z(\varphi) \rightarrow Z, S, T$

We can quickly check, that:

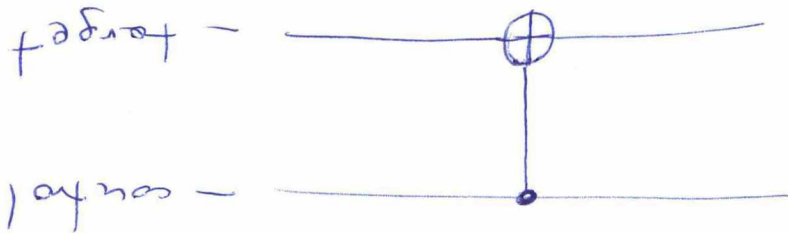
$$XX = \mathbb{I}$$

$$HH = \mathbb{I}$$

$$|\psi\rangle = \cos \frac{\theta}{2} |0\rangle + e^{i\varphi} \sin \frac{\theta}{2} |1\rangle$$

Two qubits

- Two qubits can give 4 possible outputs after measurements: 00, 01, 10, 11
- Basis:  $|00\rangle, |01\rangle, |10\rangle, |11\rangle$
- We represent state of such system as 4-element vectors
- the only really gate acting on 2 qubits which we will use is CNOT



$$|x, y\rangle \rightarrow |x, x \oplus y\rangle$$

$$\begin{array}{l} x \rightarrow x \\ y \rightarrow y \\ x, y \rightarrow x, y \\ 0, 0 \rightarrow 0, 0 \\ 0, 1 \rightarrow 0, 1 \\ 1, 0 \rightarrow 1, 1 \\ 1, 1 \rightarrow 1, 0 \end{array}$$

- Kladka operacie na n qubitoch  
more byc zaprisosa jako  
~~sekvencie~~ operacie: 4-qubitovych  
! branek CNOT

$$CNOT = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

$$|a\rangle \otimes |b\rangle = \begin{pmatrix} a_1 \\ a_2 \end{pmatrix} \otimes \begin{pmatrix} b_1 \\ b_2 \end{pmatrix} = \begin{pmatrix} a_1 b_1 \\ a_1 b_2 \\ a_2 b_1 \\ a_2 b_2 \end{pmatrix}$$

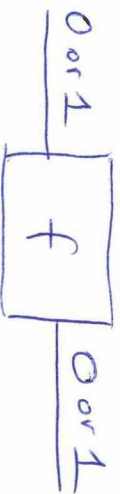
i.e.,

$$|1\rangle \otimes |0\rangle = |10\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 1 \\ 0 \end{pmatrix}$$

## Deutsch algorithm

Let's consider binary function:

$$f: \{0, 1\} \rightarrow \{0, 1\}$$



	$f_1$	$f_2$	$f_3$	$f_4$
0	0	0	1	1
1	0	1	0	1

$\Rightarrow$  4 such functions

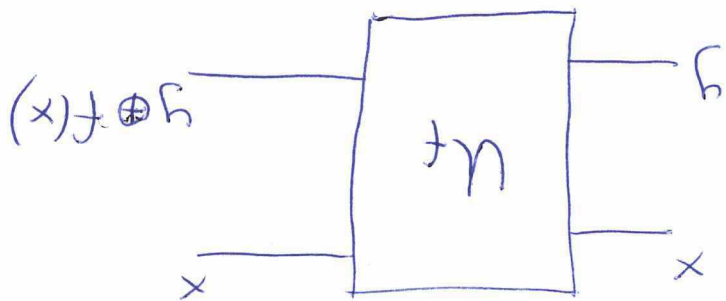
- 2 are constant (same values for all inputs  $\Rightarrow f_1$  and  $f_4$ )
- 2 are balanced (same number of 0s and 1s in outputs)  $\Rightarrow f_2, f_3$

• we have block box implementing one of these functions, now how to determine which of them it is?

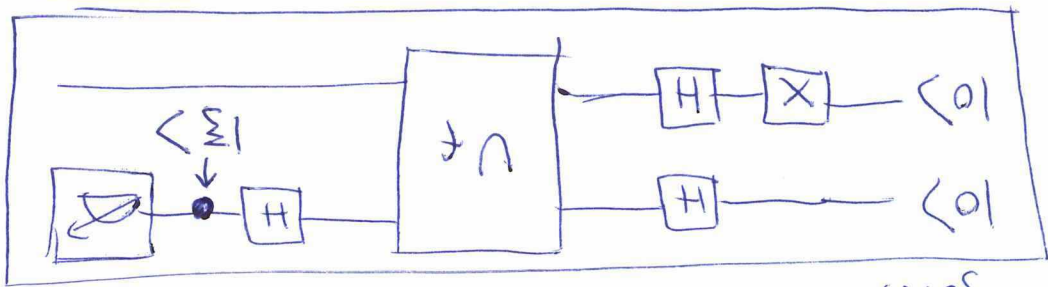
• We have to call the box twice, for inputs 0 and 1

• Deutsch algorithm requires quantum computer, but it can do this in just one call to the box implementing the function

• the box:



• the algorithm



$$|z\rangle = \text{const.} \left[ \left( (-1)^{f(0)} + (-1)^{f(1)} \right) |0\rangle + \left( (-1)^{f(0)} - (-1)^{f(1)} \right) |1\rangle \right]$$



- if we measure first qubit in state  $|0\rangle$  (3) then it means that  $f(0)=f(1) \Rightarrow f$  is constant
  - if we measure first qubit in state  $|1\rangle$  then it means, that  $f(0)=-f(1)$
- $f$  is balanced  $\Downarrow$

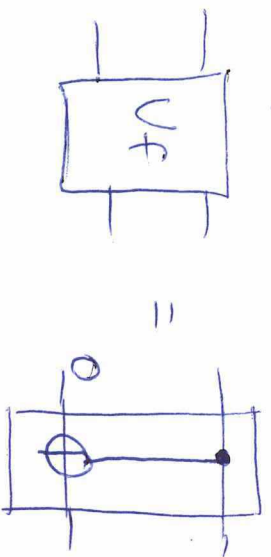
Classical complexity

2

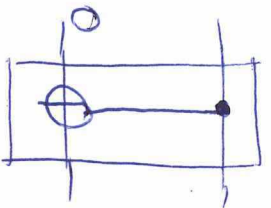
Quantum complexity

1

### Examples



=



$\Rightarrow$

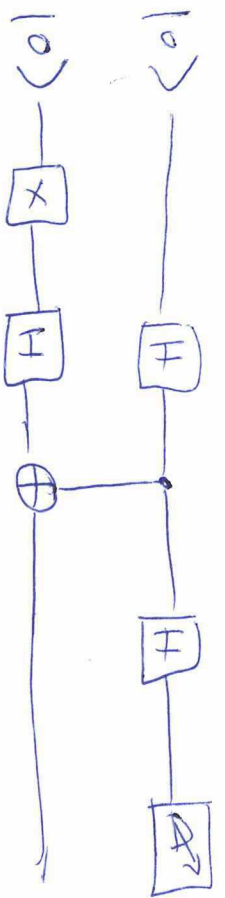
$\begin{matrix} 00 & \rightarrow & 00 \\ 10 & \rightarrow & 11 \end{matrix}$

$\Downarrow$

$f: \begin{matrix} 0 \rightarrow 0 \\ 1 \rightarrow 1 \end{matrix}$

$(f_2)$

balanced

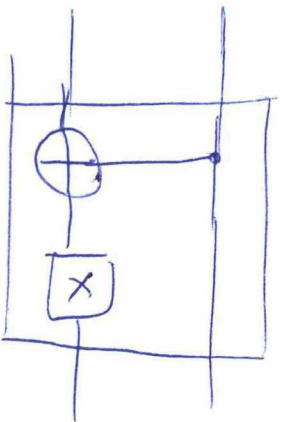


$\Downarrow$

do this in quantum composer

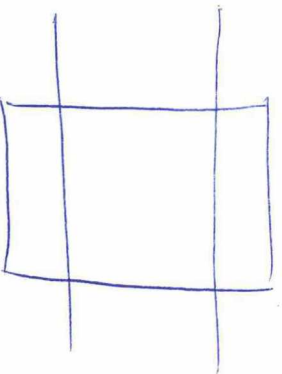
$\downarrow$

## Other functions



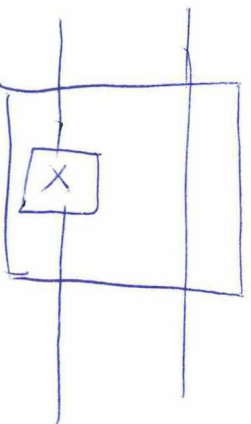
$$\begin{array}{l} 0 \rightarrow 1 \\ 1 \rightarrow 0 \end{array}$$

$\Rightarrow f_3$ , balanced



$$\begin{array}{l} 0 \rightarrow 0 \\ 1 \rightarrow 0 \end{array}$$

$\Rightarrow f_1$ , constant



$$\begin{array}{l} 0 \rightarrow 1 \\ 1 \rightarrow 1 \end{array}$$

$\Rightarrow f_4$ , constant

Generalization of the Deutsch algorithm is Deutsch-Jozsa algorithm, which is still not very practical, but we can have ~~it~~ (theoretically) huge gain in complexity reduction



However for volunteers to explain this