

DAA ASSIGNMENT-5

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PROBLEM STATEMENT

Given an array `arr[]` of n numbers and a number K , find the number of subsets of `arr[]` having XOR of elements as K .

ALGORITHM

- 1) We define a number m such that $m = m \mid a[i]$ for all i from 1 to n . This number is actually the maximum value any XOR subset will acquire.
- 2) We create a 2D array $dp[n+1][m+1]$, such that $dp[i][j]$ equals the number of subsets having XOR value j from subsets of $arr[0 \dots i-1]$.
- 3) We fill the dp array as following: We initialize all values of $dp[i][j]$ as 0. Set value of $dp[0][0] = 1$ since XOR of an empty set is 0. Iterate over all the values of $arr[i]$ from left to right and for each $arr[i]$, iterate over all the possible values of XOR i.e from 0 to m (both inclusive) and fill the dp array as following:

for $i = 1$ to n :

for $j = 0$ to m :

$dp[i][j] = dp[i-1][j] + dp[i-1][j \oplus arr[i-1]]$ Counting the number of subsets with XOR value k : Since $dp[i][j]$ is the number of subsets having j as XOR value from the subsets of $arr[0 \dots i-1]$, then the number of subsets from set $arr[0 \dots n]$ having XOR value as K will be $dp[n][K]$

PSEUDO CODE

BEGIN:

Take n and k input

a[n+1]

orr=0;

FOR i 1 to n :

INPUT(a[i])

orr=orr|a[i]

IF orr<k:

PRINT 0

dp[n+1][orr+1]

INITIALIZE(dp to 0)

dp[0][0]=1;

FOR i 1 to n:

FOR j 0 to orr:

dp[i][j]=dp[i-1][j]+dp[i-1][orr^a[i]]

PRINT(dp[n][k])

END:

CODE

```
30  int main()
31  {
32      lli t;
33      cout<<"ENTER TOTAL NUMBER OF TEST CASES:"<<endl;
34      cin>>t;
35      cout<<"ENTER THE TEST CASES:"<<endl;
36      while(t-->0)
37      {
38          lli n,k;
39          cout<<"ENTER LENGTH OF ARRAY AND THE VALUE OF K:"<<endl;
40          cin>>n>>k;
41          lli K=k;
42          lli a[n+1];
43          lli orr=0;
44          cout<<"ENTER ARRAY ELEMENTS:"<<endl;
45          For(i,1,n+1){cin>>a[i];orr=orr|a[i];}
46          if(orr<k)
47          {
48              cout<<"0"<<endl;continue;
49          }
50          k=orr;
51          lli** dp=new lli*[n+1];
52          For(i,0,n+1)
53          {
54              dp[i]=new lli[k+1]();
55          }
56          dp[0][0]=1;
57          For(i,1,n+1)
58          {
59              For(j,0,k+1)
60              {
61                  dp[i][j]=dp[i-1][j]+dp[i-1][j^a[i]];
62              }
63          }
64          cout<<"Answer = "<<dp[n][K]<<endl<<endl;
65      }
66  }
```

COMPLEXITY ANALYSIS

Here, we are going through a matrix of order n times r , where

n = number of elements in array

r = OR of all elements in array

Hence time complexity comes out to be **$O(nr)$**

And for space complexity, since we are making a matrix of order n times r ,

Space complexity is **$O(nr)$**

OUTPUT

```
ENTER TOTAL NUMBER OF TEST CASES:
2
ENTER THE TEST CASES:
ENTER LENGTH OF ARRAY AND THE VALUE OF K:
5
4
ENTER ARRAY ELEMENTS:
1 2 3 4 5
Answer = 4

ENTER LENGTH OF ARRAY AND THE VALUE OF K:
5 7
ENTER ARRAY ELEMENTS:
83 23 11 23 2
Answer = 0

...Program finished with exit code 0
Press ENTER to exit console.
```

THANK YOU