## **Deploying a Constraint Satisfaction Framework**

- Constraints are defined using a Constraint class.
- Each Constraint consists of the variables it constrains and a method that checks whether it's satisfied ().

The determination of whether a constraint is satisfied is the main logic that goes into defining a specific constraint-satisfaction problem.

- The default implementation must be overridden because we're defining our Constraint class as an abstract base class. Abstract base classes aren't meant to be instantiated; only the subclasses that override and implement their @abstractmethods are used.
- Please refer:

Information on typing module at:

https://docs.python.org/3/library/typing.html

```
from typing import Generic, TypeVar, Dict, List, Optional
from abc import ABC, abstractmethod

V = TypeVar('V') # variable type

D = TypeVar('D') # domain type

# Base class for all constraints

class Constraint(Generic[V, D], ABC): # The variables that

the constraint is between

def __init__(self, variables: List[V]) -> None:
    self.variables = variables

# Must be overridden by subclasses
    @abstractmethod

def satisfied(self, assignment: Dict[V, D]) -> bool:
```

- The center piece of our constraint-satisfaction framework is a class called CSP.
- CSP is the gathering point for variables, domains, and constraints.
- In terms of its type hints, it uses generics to make itself flexible enough to work with any kind of variables and domain values (v keys and domain values).
- Within CSP, the definitions of the collections variables, domains, and constraints are of types that you'd expect.
- The variables collection is a list of variables,

- domains is a dict mapping variables to lists of possible values (the domains of those variables), and
- Constraints is a dict that maps each variable to a list of the constraints imposed on it.

```
# A constraint satisfaction problem consists of variables of type V #
that have ranges of values known as domains of type D and
constraints # that determine whether a particular variable's domain
selection is valid
class CSP(Generic[V, D]):
 def init (self, variables: List[V], domains: Dict[V, List[D]]) -
> None:
    self.variables: List[V] = variables # variables to be
self.domains: Dict[V, List[D]] = domains # domain of each
variable
self.constraints: Dict[V, List[Constraint[V, D]]] = {}
 for variable in self.variables:
 self.constraints[variable] = []
  if variable not in self.domains:
      raise LookupError ("Every variable should have a domain assigned
to it.")
def add constraint(self, constraint: Constraint[V, D]) -> None:
  for variable in constraint.variables:
  if variable not in self.variables:
 raise LookupError("Variable in constraint not in CSP")
 else:
  self.constraints[variable].append(constraint)
```

- The init () initializer creates the constraints dict.
- The add constraint () method goes through all of the variables touched by a given constraint and adds itself to the constraints mapping for each of them.
- Both methods have basic error-checking in place, and raise an exception when a variable is missing a domain or a constraint is on a non-existent variable.

- A given configuration of variables and selected domain values satisfy the constraints The configuration is called assignment.
- We need a function that checks every constraint for a given variable against an assignment to see if the variable's value in the assignment works for the constraints. Here we implement a consistent () function as a method on CSP.

```
Check if the value assignment is consistent by checking all
constraints # for the given variable against it
  def consistent(self, variable: V, assignment: Dict[V, D]) ->
bool:
  for constraint in self.constraints[variable]:
    if not constraint.satisfied(assignment):
        return False
  return True
```

- consistent() goes through every constraint for a given variable (it's always a variable
  that was newly added to the assignment) and checks if the constraint is satisfied, given
  the new assignment.
- If the assignment satisfies every constraint, True is returned. If any constraint imposed on the variable isn't satisfied, False is returned.
- This constraint-satisfaction framework uses a simple backtracking search to find solutions to problems.
- *Backtracking* is the idea that once you hit a wall in your search, you go back to the last known point where you made a decision before the wall, and choose a different path.
- The backtracking search implemented in the following backtracking\_search () function is a kind of recursive depth-first search.
- This function's added as a method to the CSP class.

```
def backtracking search(self, assignment: Dict[V, D] = {}) ->
Optional[Dict[V, D]]:
# assignment is complete if every variable is assigned (our
if len(assignment) == len(self.variables):
 return assignment
# get all variables in the CSP but not in the
assignment unassigned: List[V] = [v for v in self.variables
if v not in assignment]
 # get the every possible domain value of the first unassigned
variable
first: V = unassigned[0]
for value in self.domains[first]:
local assignment=assignment.copy()
local assignment[first] = value
# if we're still consistent, we recurse (continue)
  if self.consistent(first, local assignment):
                      result: Optional[Dict[V,
                                                    D]]
self.backtracking search(local assignment)
                                                     # if we
didn't find the result, we will end up backtracking
       if result is not None:
```

```
return result
return None
```

• Let's check backtrackingSearch(), line by line

```
if len(assignment) == len(self.variables):
    return assignment
```

The base case for the recursive search is finding a valid assignment for every variable.
 Once we have, we return the first instance of a solution that was valid (we stop searching).

```
unassigned: List[V] = [v for v in self.variables if v not in
assignment] first: V = unassigned[0]
```

- To select a new variable whose domain we can explore, we go through all of the variables and find the first that doesn't have an assignment.
- To do this, we create a list of variables in self.variables but not in assignment through a list comprehension, and call it unassigned. Then we pull out the first value in unassigned.

- We try assigning every possible domain value for that variable, one at a time. The new assignment for each is stored in a local dictionary called <code>local\_assignment</code>.
- If the new assignment in <code>local\_assignment</code> is consistent with all of the constraints (which is what <code>consistent()</code> checks for), we continue recursively searching with the new assignment in place. If the new assignment turns out to be complete (the base case), we return the new assignment up the recursion chain.

```
return None # no solution
```

- Finally, if we've gone through every possible domain value for a particular variable, and there's no solution utilizing the existing set of assignments, we return None, indicating no solution.
- This leads to backtracking up the recursion chain to the point where a different prior assignment could have been made.