In-class Exercises: Properties of Decompositions

1. A lossy join decomposition. Suppose we have a relation with attributes cdf, name, grade. Here is an instance of that relation:

cdf	name	grade
g3tout	Amy	91
g4foobar	David	78
c0zhang	David	85

(a) Suppose we were to decompose this into two new relations: R1(cdf, name) and R2(name, grade). Project the data onto those two new relations.

	cdf	name
R1:	93-tout 94-foban cozhang	Army David David

	name	grade
R2:	Amy David David	9/ 78 85

(b) Now compute $R1 \bowtie R2$ to rebuild the original table.

	$^{\circ}$ cdf	name	grade		
	93tout	Arry	91		
	94 Golan	David	78		,
p p a series	54 foobar	David	85) uh oh	
	coshony	David	78	and the second s	
	coshony	David	85		
	3 1				

We didn't lose tyles. Every original type is there. We lost information. We no loyer know which David got which mark.

- 2. A decomposition that fails to preserve dependencies [Example 3.25 from the text.] Suppose we have a relation with attributes movie, theatre, city and FDs { theatre → city; movie, city → theatre }. The FD theatre → city violates BCNF, and applying the BCNF decomposition algorithm, we get two new relations:
 - R1(theatre, city) with one FD: theatre \rightarrow city
 - R2(theatre, movie) with no FDs
 - (a) Create small instances of R1 and R2 that satisfy their own FDs, but when natural-joined together, violate one of the original FDs.

	theatre	city
R1:	Varsity Royal	Toronto

[theatre	movie
R2:	Vavsity Royal	Boyhood

	theatre	city	movie
R1 ⋈ R2:	Vavsity Royal	Tovento	Boyhood

Zviolales Imovie, city - theatve

(b) In the original relation, with attributes movie, that re, city, does the functional dependency theatre \rightarrow city violate 3 NF?

(c) In the original relation, does the functional dependency theatre \rightarrow city violate BCNF?

Yes-because theatre is not a superkey. The BCNF algorithm would decoupose as shown, and lose a dependency.