

# 1 Solving Nonlinear Equations [by Root Finding $y = 0$ ]

Root Multiplicity,  $m$ :  $0 = f(\bar{x}) = f'(\bar{x}) = \dots = f^{(m-1)}(\bar{x})$  (Simple Root:  $m = 1$ )

$k$ -th Iteration Error:  $e_k = x_k - \bar{x}$       Convergence Rate,  $r$ :  $\lim_{k \rightarrow \infty} \frac{\|e_{k+1}\|}{\|e_k\|^r} = C$  ( $0 < C < 1$  if  $r = 1$ )

## 1.1 One Dimension/Equation skipped a lot

Interval Bisection (Finding  $y = 0$ ):  $[f(a) < 0], [f(b) > 0], [f \text{ is cont.}] \Rightarrow \exists m \text{ s.t. } f(m) = 0$

Fixed-Point Iteration (Finding  $y = x$ ):  $\text{cont. } f(x) = 0 \Rightarrow \text{Find } g(x) = x \rightarrow x_{k+1} = g(x_k)$

$\sim$  Banach-Fixed Point Theorem (there are many FP theorems)

- $g$  is Contractive (over a domain):  $\text{dist}(g(x), g(y)) \leq q \cdot \text{dist}(x, y) \quad q \in [0, 1)$
- $e_{k+1} = [x_{k+1} - \bar{x}] = [g(x_k) - g(\bar{x})] = g'(\xi_k)(x_k - \bar{x}) = g'(\xi_k)e_k$
- $\forall |g'(\xi_k)| < G < 1 \Rightarrow (|e_{k+1}| \leq G|e_k| \leq \dots \leq G^k|e_0|) \Rightarrow \lim_{k \rightarrow \infty} e_k = 0 \quad (G = \max g' \text{ over domain})$
- $\lim_{k \rightarrow \infty} |g'(\xi_k)| = \boxed{\begin{matrix} (0 < |g'(\bar{x})| < 1) \\ \text{(one contractive condition)} \end{matrix}} = C \quad (r = 1)$
- $\boxed{g'(\bar{x}) = 0} \Rightarrow [g(x_k) - g(\bar{x})] = \frac{g''(\xi_k)}{2}(x_k - \bar{x})^2 \Rightarrow \boxed{\left| \frac{g''(\bar{x})}{2} \right|} = C \quad (r = 2 \text{ if } \bar{x} \text{ is an } m = 2 \text{ root of } g)$

Newton's Method (Finding  $y = 0$ ):

$$f(\bar{x}) = 0 = f(x_k + h_k) \approx f(x_k) + f'(x_k)h_k \Rightarrow \boxed{x_{k+1} = x_k + h_k = x_k - \frac{f(x_k)}{f'(x_k)}}$$

- $\boxed{g(x) \equiv x - \frac{f(x)}{f'(x)}} \Rightarrow g(\bar{x}) = \bar{x}, \boxed{g'(\bar{x}) = \frac{f(\bar{x})f''(\bar{x})}{f'(\bar{x})^2} = 0}, \boxed{r = 2} \quad (\text{if } \bar{x} \text{ is a simple root of } f)$
- $\bar{x} \text{ is an } m > 1 \text{ root of } f \Rightarrow \boxed{r = 1, C = 1 - 1/m} \quad (\text{proof not given})$

Secant Method/Linear Interpolation (Finding  $y = 0$ ):

$$f'(x_k) \approx \frac{f(x_k) - f(x_{k-1})}{x_k - x_{k-1}} \quad \text{Approx. } f'(x_k) \text{ with a secant line's slope} \Rightarrow \boxed{x_{k+1} = x_k + h_k = x_k - \frac{x_k - x_{k-1}}{f(x_k) - f(x_{k-1})} f(x_k)}$$

- $\boxed{r = r_+ \approx 1.618} : r_+^2 - r_+ - 1 = 0 \quad (\text{proof hard})$
- Lower cost of iter. offsets the larger number of iter. compared to Newton's Method with derivatives

Inverse Parabolic Interpolation: Use 3 pts to approx. an inverse [sideways] parabola

## 1.2 $m$ Dimensions/System of Equations stuff skipped

Newton's Method (Solving  $\vec{y} = 0$ ):

$$\{J_f(\vec{x})\}_{ij} = \frac{\partial f_i(\vec{x})}{\partial x_j} : \quad J_f(\vec{x}_k)\vec{h}_k = -\vec{f}(\vec{x}_k) \Rightarrow \vec{x}_{k+1} = \vec{x}_k + \vec{h}_k = \vec{x}_k - J_f(\vec{x}_k)^{-1}\vec{f}(\vec{x}_k)$$

- $$\vec{g}(\vec{x}) \equiv \vec{x} - J_f(\vec{x})^{-1}\vec{f}(\vec{x}) \Rightarrow \begin{aligned} J_g(\vec{x}) &= \cancel{I - J_f(\vec{x})^{-1}J_f(\vec{x})} + \sum_{i=1}^n H_i(\vec{x})f_i(\vec{x}) \\ &\quad \text{(if } J_f(\vec{x}) \text{ is nonsingular)} \end{aligned} \quad \begin{array}{l} H_i = \text{component} \\ \text{matrix of the} \\ \text{tensor, } D_x J_f(\vec{x}) \end{array}$$

$$= \mathcal{O} \Rightarrow \boxed{r=2} \quad (\text{uh... idk})$$
- $LU$  fact. of the Jacobian costs  $\mathcal{O}(n^3)$

Broyden's [Secant Updating] Method (Solving  $\vec{y} = 0$ ):

$$B_k\vec{h}_k = -\vec{f}(\vec{x}_k) \Rightarrow \vec{x}_{k+1} = \vec{x}_k + \vec{h}_k, \quad B_{k+1} = B_k + \frac{f(\vec{x}_{k+1})h_k^T}{h_k^T h_k} \quad (\text{cost is } \mathcal{O}(n^3))$$

- $B_{k+1}(\vec{x}_{k+1} - \vec{x}_k) = B_{k+1}\vec{h}_k = f(\vec{x}_{k+1}) - f(\vec{x}_k)$
- $B_k$  factorization is updated to factorization of  $B_{k+1}$  at cost  $\mathcal{O}(n^2)$  instead of directly from the above eq.
- Lower cost of iter. offsets the larger number of iter. compared to Newton's Method with derivatives

## 2 Optimizing [By Finding $\min f(\vec{x}) = f(\vec{x})$ ]

### 2.1 Function Shape and Convexity

Coercive:  $\lim_{x \rightarrow \pm\infty} f(x) = \infty$       Unimodal:  $\begin{array}{l} a \leq \bar{x} \leq b \\ x_1 < x_2 \end{array} : \begin{array}{l} x_2 < \bar{x} \rightarrow f(x_1) > f(x_2) \\ \bar{x} < x_1 \rightarrow f(x_1) < f(x_2) \end{array}$

$\exists$  global min  $f$  if

- cont.  $f$  on a closed and bounded set
- cont.  $f$  is coercive on a closed, unbounded set
- cont.  $f$  on a set and has a nonempty, closed, and bounded sublevel set
- domain set is unbounded: cont.  $f$  is coercive  $\Leftrightarrow$  all sublevel sets are bounded

$f$  is convex [on a convex set] :

- any sublevel set is convex
- any local min. is a global min

$f$  is strictly convex [on a convex set] :

- any local min. is a unique global min.
- if set is unbounded:  $f$  is coercive  $\Leftrightarrow f$  has a unique global min.

## 2.2 Derivative Tests (Gradient, Jacobian, Hessian) and Lagrangians

Req. :  $\boxed{\text{cont. } f(\bar{x}) = \min f, \text{ cont. } \vec{\nabla} f(\bar{x}), \text{ cont. } H_f(\bar{x})}$

Taylor's Theorem: 
$$\begin{aligned} f(\bar{x} + \vec{s}) - f(\bar{x}) &= \vec{\nabla} f(\bar{x} + \alpha_1 \vec{s}) \cdot \vec{s} = \vec{\nabla} f(\bar{x}) \cdot \vec{s} + \frac{1}{2} \langle \vec{s} | H_f(\bar{x} + \alpha_2 \vec{s}) | \vec{s} \rangle \geq 0 \\ f(\bar{x} + s\hat{u}) - f(\bar{x}) &= \vec{\nabla} f(\bar{x} + \alpha_1 s\hat{u}) \cdot s\hat{u} = \vec{\nabla} f(\bar{x}) \cdot \vec{s} + \frac{s^2}{2} \langle \hat{u} | H_f(\bar{x} + \alpha_2 \vec{s}) | \hat{u} \rangle \end{aligned}$$

- $\lim_{s \rightarrow 0} \left( \frac{f(\bar{x} + \vec{s}) - f(\bar{x})}{s} = \vec{\nabla} f(\bar{x} + \alpha_1 s\hat{u}) \cdot \hat{u} \right) \Rightarrow \left( \vec{\nabla} f(\bar{x}) \cdot \hat{u} \geq 0 \rightarrow \boxed{\vec{\nabla} f(\bar{x}) \cdot \vec{s} \geq 0} \right)$  ,  $\boxed{\text{Cauchy-Schwarz} \rightarrow \max \vec{\nabla} f(\bar{x}) \cdot \hat{u} \text{ if } \vec{u} = \vec{\nabla} f(\bar{x})}$
- $\boxed{\vec{u} = \mp \vec{\nabla} f(\bar{x})} \Rightarrow \lim_{s \rightarrow 0} \left( \frac{f(\bar{x} + \vec{s}) - f(\bar{x})}{s} = \mp \frac{\vec{\nabla} f(\bar{x} + \alpha_1 s\hat{u}) \cdot \vec{\nabla} f(\bar{x})}{\|\vec{\nabla} f(\bar{x})\|} \right) = \mp \|\vec{\nabla} f(\bar{x})\| \leq 0$   $\boxed{\text{if } \pm \vec{\nabla} f(\bar{x}) \neq 0, \text{ its dir. is an ascent/descent.}}$
- $\lim_{s \rightarrow 0} \left( \frac{f(\bar{x} + \vec{s}) - f(\bar{x}) + f(\bar{x} - \vec{s}) - f(\bar{x})}{s^2} = \frac{\langle \hat{u} | H_f(\bar{x} + \alpha_2 \vec{s}) + H_f(\bar{x} - \alpha_3 \vec{s}) | \hat{u} \rangle}{2} \right) = \langle \hat{u} | H_f(\bar{x}) | \hat{u} \rangle \Rightarrow \boxed{\langle \vec{s} | H_f(\bar{x}) | \vec{s} \rangle \geq 0}$

### 2.2.1 Unconstrained Optimization Conditions

- $\boxed{f(\bar{x}) = \min f} \Leftrightarrow \left( \begin{array}{l} \vec{\nabla} f(\bar{x}) \cdot \vec{s} \geq 0, \vec{\nabla} f(\bar{x}) \cdot -\vec{s} \geq 0 \\ \Rightarrow \boxed{\vec{\nabla} f(\bar{x}) = 0} \end{array} \right), \quad \vec{u} = -\vec{\nabla} f(\bar{x}) \Rightarrow \boxed{\vec{\nabla} f(\bar{x}) = 0}, \quad \boxed{\begin{array}{l} \text{(for strict convexity)} \\ \langle \vec{s} | H_f(\bar{x}) | \vec{s} \rangle > 0 \end{array}} \right)$

Optimization  $f: \mathbb{R}^n \rightarrow \mathbb{R} \quad \boxed{\min f(\vec{x}) = y}$

$$\boxed{\mathcal{L}(\vec{x}) = f(\vec{x})}, \quad \boxed{\nabla \mathcal{L}(\bar{x}) = 0}, \quad \boxed{H_{\mathcal{L}} = \nabla_{xx} \mathcal{L}: \langle s | H_{\mathcal{L}}(\bar{x}) | s \rangle > 0} \Rightarrow \boxed{y = f(\bar{x})}$$

### 2.2.2 Constrained Optimization Conditions

- $\boxed{\vec{s} = \text{feasible direction}} \Leftrightarrow \left( \boxed{\vec{\nabla} f(\bar{x}) \cdot \vec{s} \geq 0}, \boxed{\langle \vec{s} | H_f(\bar{x}) | \vec{s} \rangle \geq 0} \right)$

Optimization  $\begin{array}{l} f: \mathbb{R}^n \rightarrow \mathbb{R} \\ g: \mathbb{R}^n \rightarrow \mathbb{R}^m \\ h: \mathbb{R}^n \rightarrow \mathbb{R}^p \end{array} \quad \min f(\vec{x}) = y \quad \text{w/} \quad \left( \begin{array}{l} \vec{g}(\vec{x}) = 0 \\ \vec{h}(\vec{x}) \leq 0 \end{array} \right) \quad \begin{array}{l} \text{active: } h_i(\bar{x}) = 0 \\ \text{inactive: } h_i(\bar{x}) < 0 \rightarrow \bar{\mu}_i = 0 \end{array}$  (see KKT)

$$\begin{aligned} \mathcal{L}(\bar{x}, \bar{\lambda}, \bar{\mu}) &= f(\bar{x}) + \bar{\lambda} \cdot \vec{g}(\bar{x}) + \bar{\mu} \cdot \vec{h}(\bar{x}) \\ &= f + \sum_i^m \lambda_i g_i + \sum_i^p \cancel{\mu_i h_i} \quad \text{(KKT) if } \bar{x} = \bar{x} \end{aligned}, \quad \nabla \mathcal{L}(\bar{x}, \bar{\lambda}, \bar{\mu}) = \begin{pmatrix} \nabla_x \mathcal{L} = 0 \\ \nabla_{\lambda} \mathcal{L} = 0 \\ \nabla_{\mu} \mathcal{L} \leq 0 \end{pmatrix} = \begin{pmatrix} \nabla f(\bar{x}) + J_g^T(\bar{x}) \bar{\lambda} + J_h^T(\bar{x}) \bar{\mu} \\ \vec{g}(\bar{x}) \\ \vec{h}(\bar{x}) \end{pmatrix}$$

$$H_{\mathcal{L}}(\bar{x}, \bar{\lambda}, \bar{\mu}) = \begin{pmatrix} \nabla_{xx} \mathcal{L} & \nabla_{x\lambda} \mathcal{L} & \nabla_{x\mu} \mathcal{L} \\ \nabla_{\lambda x} \mathcal{L} & \nabla_{\lambda\lambda} \mathcal{L} & \nabla_{\lambda\mu} \mathcal{L} \\ \nabla_{\mu x} \mathcal{L} & \nabla_{\mu\lambda} \mathcal{L} & \nabla_{\mu\mu} \mathcal{L} \end{pmatrix} = \begin{pmatrix} \nabla_{xx} \mathcal{L} & J_g^T & J_h^T \\ J_g & 0 & 0 \\ J_h & 0 & 0 \end{pmatrix}, \quad \nabla_{xx} \mathcal{L}(\bar{x}, \bar{\lambda}, \bar{\mu}) = H_f + \sum_i^m \bar{\lambda}_i H_{g_i} + \sum_i^{\text{act} \leq p} \bar{\mu}_i H_{h_i}$$

(can't be pos. def.)

- Assume  $m \leq n$  (not overdetermined)

- $y = f(\bar{x}) : \nabla \mathcal{L}(\bar{x}, \bar{\lambda}, \bar{\mu}) \dots$  ,  $p = 0 : Z^T (\nabla_{xx} \mathcal{L}) Z > 0$  col. of  $Z$  = basis of  $\text{null}(J_g)$

- Assume  $h_i$  don't contradict each other? Assume full rank( $J_{h_{\text{act}}}$ )

- $y = f(\bar{x}) : \nabla \mathcal{L}(\bar{x}, \bar{\lambda}, \bar{\mu}) \dots$  ,  $p > 0$ , Karush-Kuhn-Tucker (KKT) :  $\bar{\mu}_i \geq 0$ ,  $\bar{\mu}_i h_i(\bar{x}) = 0$  (2nd deriv. cond. not given)

## 2.3 Unconstrained One Dimension/Independent Variable

[Interval] Golden-Section Search (if Unimodal):  $\tau^2 = 1 - \tau = .382$  ,  $r = 1$  ,  $C = \tau$

$$[a < x_1 < x_2 < b] : \begin{cases} f(x_1) > f(x_2) \rightarrow [x_1 < x_2 < x_1 + \tau(b - x_1) < b] \\ f(x_1) \leq f(x_2) \rightarrow [a < a + (1 - \tau)(x_2 - a) < x_1 < x_2] \end{cases}$$

Newton's Method:  $f(\bar{x}) = f(x + h) \approx f(x) + f'(x)h + \frac{1}{2}f''(x)h^2 = g(h)$

$$g\left(\frac{-b}{2a}\right) = \min g \text{ (or max)} \Rightarrow x_{k+1} = x_k + h_k = x_k - \frac{b}{2a} = x_k - \frac{f'(x)}{f''(x)} , r = 2$$

Successive Linear Interpolation [Secant Method]: Not useful, since lines have no unique minimum

Successive Parabolic Interpolation: Use 3 pts to approx. a parabola w/  $r = 1.324$  (not guaranteed)

## 2.4 Unconstrained $m$ -Dimensions/Independent Variables

Steepest [Gradient] Descent/Line Search (go down  $-\nabla f(\vec{x}_k)$ ):

$$\phi(\alpha) = f(\vec{x} - \alpha \vec{\nabla} f(\vec{x})) , \phi(\alpha_k) = \min \phi \Rightarrow \vec{x}_{k+1} = \vec{x}_k - \alpha_k \vec{\nabla} f(\vec{x}_k) , r = 1 , C_{\text{varies}}$$

- $\vec{\nabla} f(\vec{x}_k) \cdot \vec{\nabla} f(\vec{x}_{k+1}) = 0 \Rightarrow$  Path will zig-zag to the min. (not too efficient)

Newton's Method:  $f(\bar{x}) = f(\vec{x} + \vec{h}) \approx f(\vec{x}) + \vec{\nabla} f(\vec{x}) \cdot \vec{h} + \frac{1}{2} \langle \vec{h} | H_f(\vec{x}) | \vec{h} \rangle$

$$H_f(\vec{x}_k) \vec{h}_k = -\vec{\nabla} f(\vec{x}_k) \Rightarrow \vec{x}_{k+1} = \vec{x}_k + \vec{h}_k , r = 2$$

BFGS [Secant Updating] Method:  $B_k \vec{h}_k = -\vec{\nabla} f(\vec{x}_k) , \vec{y}_k = \vec{\nabla} f(x_{k+1}) - \vec{\nabla} f(x_k)$

$$\Rightarrow \vec{x}_{k+1} = \vec{x}_k + \vec{h}_k , B_{k+1} = B_k + \frac{|y_k\rangle\langle y_k|}{\langle y_k | h_k \rangle} - \frac{B_k |h_k\rangle\langle h_k| B_k}{\langle h_k | B_k | H_k \rangle} \quad (\text{cost is } \mathcal{O}(n^3))$$

- Preserves symmetry and pos. def.
- $B_k$  factorization is updated to factorization of  $B_{k+1}$  at cost  $\mathcal{O}(n^2)$  instead of directly from the above eq.
- Lower cost of iter. offsets the larger number of iter. compared to Newton's Method with derivatives

Conjugate Gradient [Line Search] :

$$\boxed{\vec{h}_{k+1} = \vec{\nabla} f(\vec{x}_{k+1}) - \frac{\vec{\nabla} f(\vec{x}_{k+1}) \cdot \vec{\nabla} f(\vec{x}_{k+1})}{\vec{\nabla} f(\vec{x}_k) \cdot \vec{\nabla} f(\vec{x}_k)} \vec{h}_k} \quad (\text{Fletcher and Reeves}) \Rightarrow \boxed{\vec{x}_{k+1} = \vec{x}_k - \alpha_k \vec{h}_k}$$

- Seq. of conj. (where  $(a, b) = \langle a | H_f | b \rangle$ ) search directions implicitly accumulates info. about  $H_f$ .
- Better for nonlin. to use 
$$\boxed{\vec{h}_{k+1} = \vec{\nabla} f(\vec{x}_{k+1}) - \frac{\vec{\nabla} f(\vec{x}_{k+1}) \cdot \vec{\nabla} f(\vec{x}_{k+1}) - \vec{\nabla} f(\vec{x}_k) \cdot \vec{\nabla} f(\vec{x}_{k+1})}{\vec{\nabla} f(\vec{x}_k) \cdot \vec{\nabla} f(\vec{x}_k)} \vec{h}_k} \quad (\text{Polak and Ribiere})$$
- Restart algorithm after  $n$  iter. using last point as the new initial; a quadratic func. finishes after at most  $n$  iter.

### 2.4.1 Nonlinear Least Squares, $\{ \min \|\vec{r}(\vec{x})\|^2 : \vec{f}(\vec{a}, \vec{x}) + \vec{r}(\vec{x}) = \vec{b} \}$

Linear Least Squares	Nonlinear Least Squares
$\begin{pmatrix} \vdots \\ -\vec{a}_i \\ \vdots \end{pmatrix} \begin{pmatrix}   \\ \vec{x} \\   \end{pmatrix} + \begin{pmatrix}   \\ \vec{r} \\   \end{pmatrix} = \begin{pmatrix}   \\ \vec{b} \\   \end{pmatrix} \Rightarrow \begin{pmatrix}   \\ \vec{f}(\vec{a}, \vec{x})_i \\   \end{pmatrix} + \begin{pmatrix}   \\ \vec{r} \\   \end{pmatrix} = \begin{pmatrix}   \\ \vec{b} \\   \end{pmatrix}$	

$\boxed{\phi(\vec{x}) \equiv \frac{1}{2} \vec{r} \cdot \vec{r}}, \quad \boxed{-\vec{\nabla} \phi(\vec{x}) = -J_r^T \vec{r}}$	<p>Newton's Method</p> $\boxed{H_\phi(\vec{x}_k) \vec{h}_k = -\vec{\nabla} \phi(\vec{x}_k)} \Rightarrow \boxed{\vec{x}_{k+1} = \vec{x}_k + \vec{h}_k}$ <p>(usually expensive to compute)</p>	
$\boxed{H_\phi(\vec{x}) = J_r^T J_r + \sum_i H_{r_i} \vec{r}_i}$	:	

Gauss-Newton Method: If  $\vec{r}$  is small  $\Rightarrow H_\phi \approx J_r^T J_r \Rightarrow \boxed{J_r^T (J_r \vec{h}_k) = -J_r^T \vec{r}(\vec{x}_k)}$  System of Normal Equations

Levenberg-Marquardt Method (Gauss-Newton + Line Search):

$\boxed{(J_r^T J_r + \mu_k I) \vec{h}_k = -J_r^T \vec{r}(\vec{x}_k) \Rightarrow \vec{x}_{k+1} = \vec{x} + \vec{h}_k}$ $\Rightarrow \boxed{\begin{pmatrix} J_r^T(\vec{x}) & \sqrt{\mu_k} I \end{pmatrix} \begin{pmatrix} J_r(\vec{x}) \\ \sqrt{\mu_k} I \end{pmatrix} \vec{h}_k = \begin{pmatrix} J_r^T(\vec{x}) & \sqrt{\mu_k} I \end{pmatrix} \begin{pmatrix} -\vec{r}(\vec{x}_k) \\ 0 \end{pmatrix}}$	<p><u>Regularization</u></p> <ul style="list-style-type: none"> <li>• Replacing <math>H_{r_i} \vec{r}_i</math> terms with a scalar mult. of <math>I</math>.</li> <li>• Shifting the Gauss-Newton Hessian to make it pos. def (or boosting its rank).</li> </ul>
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## 2.5 Constrained $m$ -Dimensions/Independent Variables

<p>Newton's Method</p> $\boxed{H_{\mathcal{L}} \vec{h}_k = -\vec{\nabla} \mathcal{L}}$	<p>KKT Matrix (Eq. Constr)</p> $\begin{pmatrix} \nabla_{xx} \mathcal{L} & J_g^T \\ J_g & 0 \end{pmatrix} \begin{pmatrix} \vec{s}_k \\ \vec{\delta}_k \end{pmatrix} = - \begin{pmatrix} \nabla f(\vec{x}) + J_g^T(\vec{x}) \bar{\lambda} \\ \vec{g}(\vec{x}) \end{pmatrix} \Rightarrow$ $\boxed{\begin{pmatrix} B & J^T \\ J & 0 \end{pmatrix} \begin{pmatrix} \vec{s} \\ \vec{\delta} \end{pmatrix} = - \begin{pmatrix} w \\ g \end{pmatrix}}$	<p>[Sequential] Quadratic Programming (SQP) Problem</p> $\min_s \left( \vec{s}_k \cdot \vec{\nabla}_x \mathcal{L} + \frac{1}{2} \langle \vec{s}_k   \vec{\nabla}_{xx} \mathcal{L}   \vec{s}_k \rangle \right)$ <p>s.t. <math>J_g(\vec{x}_k) \vec{s}_k + \vec{g}(\vec{x}_k) = 0</math></p>
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Direct Solution: KKT Matrix is sym. and sparse  $\rightarrow$  solve for  $\vec{h}_k$  using sym. indef. factorization w/ some pivoting

(Column-Space)

Range-Space Method:  $Bs = -w - J^T \delta$  ,  $Js = -g \rightarrow JB^{-1}(-w - J^T \delta) = -g$   
 $\rightarrow (JB^{-1}J^T)\delta = g - JB^{-1}w$

- Solve for  $\delta$ , then for  $s$ .
- Forming  $(JB^{-1}J^T)_{m \times m}$  leads to issues similar to forming  $A^T A$  (loss of info. and degrades conditioning).
- $B$  must be nonsingular and  $J$  full rank.
- Useful if  $m$  is small.

Null-Space Method:  $J^T = (Q_{\parallel} \ Q_{\perp}) \begin{pmatrix} R \\ 0 \end{pmatrix} \quad (Q_{\parallel} \in \mathbb{R}^{n \times m}) \Rightarrow \begin{cases} JQ_{\parallel} = R^T \\ JQ_{\perp} = 0 \end{cases}$

Find  $u_{\parallel}$  :  $Js \equiv (JQ_{\parallel}u_{\parallel} + \cancel{JQ_{\perp}u_{\perp}}) = \boxed{R^T u_{\parallel} = -g}$

Find  $u_{\perp}$  :  $Q_{\perp}^T(Bs + J^T \delta = -w) \rightarrow (Q_{\perp}^T BQ_{\parallel})u_{\parallel} + (Q_{\perp}^T BQ_{\perp})u_{\perp} = -Q_{\perp}^T w - \cancel{(JQ_{\perp})^T} \delta$   
 $\boxed{(Q_{\perp}^T BQ_{\perp})u_{\perp} = -Q_{\perp}^T w - (Q_{\perp}^T BQ_{\parallel})u_{\parallel}}$

Find  $\delta$  :  $Q_{\parallel}^T(J^T \delta = -w - Bs) \rightarrow \boxed{R\delta = -Q_{\parallel}^T w - Q_{\parallel}^T B(Q_{\parallel}u_{\parallel} - Q_{\perp}u_{\perp})}$

- Near a min.,  $(Q_{\perp}^T BQ_{\perp})$  can be Cholesky factored.
- Avoids issues with loss of info. and degraded conditioning.
- $J$  must be full rank and  $R$  nonsingular.
- Useful if  $m$  is large, so  $n - m$  is small.

Decent Initial  $\vec{\lambda}_0$  Guess Given an  $\vec{x}_0$ :  $J_g^T(\vec{x}_0)\vec{\lambda}_0 + \vec{r} = -\vec{\nabla} f(\vec{x}_0) \quad (\text{Linear Least Sq.})$

Penalty Func. Method

$\lim_{\rho \rightarrow \infty} \vec{x}_{\rho} = \vec{x}$  (not explained)

(“Under approp. conds.”)

One Simple Function  
(Ill-conditioned  $\rho \gg 1$ ) :  $\min_{\vec{x}} \phi_{\rho}(\vec{x}) = f(\vec{x}) + \frac{1}{2}\rho \|g(\vec{x})\|^2$

Augmented Lagrangian  
(Less Ill-conditioned) :  $\min_{\vec{x}} \mathcal{L}_{\rho}(\vec{x}) = f(\vec{x}) + \vec{\lambda}_0 \cdot \vec{g}(\vec{x}) + \frac{1}{2}\rho \|g(\vec{x})\|^2$

Barrier Func. Method

$\lim_{\rho \rightarrow 0} \vec{x}_{\rho} = \vec{x}$

(“Under approp. conds.”)

Inverse :  $\min_{\vec{x}} \phi_{\rho}(\vec{x}) = f(\vec{x}) - \rho \sum_i^p \frac{1}{h_i(\vec{x})}$

Logarithmic :  $\min_{\vec{x}} \phi_{\rho}(\vec{x}) = f(\vec{x}) - \rho \sum_i^p \log(-h_i(\vec{x}))$

(For Ineq. Constr.)

- Along with line search and trust region (not explained), a merit func. - using perhaps a penalty func. - can be used to make an algorithm more robust.
- An active set strategy (not explained) can be used with an SQP method for ineq.-constr. problems.
- A penalty method penalizes points that violates constraints, but doesn't avoid them. Barrier methods do.

### 3 [Polynomial] Interpolation, $f(t_i) = \hat{f}(t_i) = \sum_j x_j \phi_j(t_i)$

$$\hat{f}(t_i) = \sum_j x_j \phi_j(t_i) \quad \left| \begin{array}{l} \det(A) \neq 0 \\ \text{Given } \vec{\phi}, \\ \text{solve for } \vec{x} \end{array} \right| \quad A\vec{x} = \begin{pmatrix} \vdots \\ -\vec{\phi}(t_i) \\ \vdots \end{pmatrix} \begin{pmatrix} | \\ \vec{x} \\ | \end{pmatrix} = \vec{y} = \begin{pmatrix} \vdots \\ f(t_i) \\ \vdots \end{pmatrix}$$

- Runge Phenom.: As  $n$  increases, evenly-spaced  $t_i$  could produce a high-dimensional polynomial  $\hat{f}(t)$  that tends to be extremely wavy near the endpoints (like Gibbs phenom.). Choosing  $t_i$  to be Chebyshev nodes between the two endpoints mitigates this.
- Interpolation w/ other func. like rationals are possible.
- Error:  $\max_{t \in [t_1, t_n]} \left| \hat{f} - f = \frac{f^{(n)}(\xi)}{n!} \prod_i (t - t_i) \right| \leq \left| \max_{t \in [t_1, t_n]} \frac{f^{(n)}(t)}{n!} \right| \left| \frac{(n-1)! h^n}{4} \right| = \boxed{\max_{t \in [t_1, t_n]} \left| f^{(n)}(t) \frac{h^n}{4n} \right|} \rightarrow \text{error decreases if } f^{(n)} \text{ is well behaved}$

#### 3.1 Taylor Series Polynomial Interpolation

$$\begin{aligned} \hat{f}_n(t) &= f(t_0) + f'(t_0)(t - t_0) + \frac{f''(t_0)}{2}(t - t_0)^2 + \dots + \frac{f^{(n-1)}(t_0)}{(n-1)!}(t - t_0)^{n-1} \\ \hat{f}_n(t + h) &= f(t) + f'(t)h + \frac{f''(t)}{2}h^2 + \dots + \frac{f^{(n-1)}(t)}{(n-1)!}h^{n-1} \end{aligned}$$

- Can interpolate an  $n$ -polynomial from  $n + 1$  points/derivatives/info.

#### 3.2 Monomial Basis Functions $\rightarrow$ Vandermonde Matrix

$$\vec{\phi}(t) = (1, t, t^2, \dots, t^{n-1})^T$$

$$\hat{f}(t) = x_1 + x_2 t + \dots + x_n t^{n-1}$$

(Full, Dense  
Vandermonde Matrix)

$$\begin{pmatrix} 1 & t_1 & \dots & t_1^{n-1} \\ \vdots & \vdots & & \vdots \\ 1 & t_n & \dots & t_n^{n-1} \end{pmatrix} \begin{pmatrix} \vdots \\ x_i \\ \vdots \end{pmatrix} = \vec{y}$$

- Solved with  $\mathcal{O}(n^3)$  work using Gauss. Elim. ( $\mathcal{O}(n^2)$  is possible with other tech.).
- Ill-conditioned since successive  $t^j$  look the same at higher  $j$ .

#### 3.3 Lagrange Basis Functions (Fund. Polynomials) $\rightarrow$ Identity Matrix

$$l(t) = (t - t_1)(t - t_2) \dots (t - t_n)$$

$$w_j = (t_j - t_j)/l(t_j) \quad (\text{barycentric weights})$$

(Diag. Iden. Matrix)

$$\begin{pmatrix} 1 & 0 & \dots \\ 0 & 1 & \ddots \\ \vdots & \ddots & \ddots \end{pmatrix} \vec{x} = \vec{y}$$

$$\phi_j(t) = \frac{l(t)/(t - t_j)}{l(t_j)/(t_j - t_j)} = l(t) \frac{w_j}{t - t_j}$$

$$\phi_j(t_i) = \delta_{ij} \Rightarrow \boxed{\vec{\phi}(t_i) = \vec{e}_i}$$

$$\hat{f}(t) = \vec{x} \cdot \vec{\phi}(t) = l(t) \left[ x_1 \frac{w_1}{t - t_1} + \dots + x_n \frac{w_n}{t - t_n} \right]$$

$$\hat{f}(t_j) = x_j = y_i$$

- Finding  $w_j$  is  $\mathcal{O}(n^2)$  work.
- Finding  $\hat{f}(t)$  from  $w_j$ 's is  $\mathcal{O}(n)$  work.
- Updating with an extra point  $(t_{n+1}, y_{n+1})$  is  $\mathcal{O}(n)$  work by changing  $w_j = w_j/(t_j - t_{n+1})$  and finding  $w_{n+1}$ .
- Basis func. are more varied  $\rightarrow$  better-conditioned.

$$\int_{t_1}^{t_n} \hat{f}(t) dt = \sum_{i=1}^n y_i \int_{t_1}^{t_n} \phi_i(t) dt$$

### 3.4 Newton Basis Functions → Low. Triang. Matrix

$$\left. \begin{array}{l} \phi_j(t) = (t - t_1)(t - t_2) \dots (t - t_{j-1}) \\ \vec{\phi}(t) = [1, (t - t_1), (t - t_1)(t - t_2), \dots]^T \\ \hat{f}(t) = x_1 + x_2(t - t_1) + \dots + x_n \phi_n(t) \end{array} \right| \begin{array}{l} \text{(Low. Triang. Matrix)} \\ \begin{pmatrix} 1 & 0 & 0 & \dots \\ 1 & t_1 - t_2 & 0 & \dots \\ 1 & t_3 - t_2 & (t_3 - t_1)(t_3 - t_2) & \ddots \\ \vdots & \vdots & \vdots & \ddots \end{pmatrix} \begin{pmatrix} \vdots \\ x_i \end{pmatrix} = \vec{y} \end{array}$$

- For. sub. is  $\mathcal{O}(n^2)$ .
- Cond. of  $A$  depends on ordering of points → best to order points from their dist. to their mean/other num.
- Basis func. are more varied → better-conditioned.

#### Incremental Updating Newton Interpolation:

$$\hat{f}_{n+1}(t) = \hat{f}_n(t) + x_{n+1} \phi_{n+1}(t)$$

$$y_{n+1} = \hat{f}_{n+1}(t_{n+1})$$

$$= \hat{f}_n(t_{n+1}) + x_{n+1} \phi_{n+1}(t_{n+1})$$

$$\Rightarrow \hat{f}_{j+1}(t) = \hat{f}_j(t) + \frac{y_{j+1} - \hat{f}_j(t_{j+1})}{\phi_{j+1}(t_{j+1})} \phi_{j+1}(t)$$

#### Divided Differences Newton Interpolation:

$$g[t_1, \dots, t_k] \equiv \frac{g[t_2, \dots, t_k] - g[t_1, \dots, t_{k-1}]}{t_k - t_1}$$

$$\vec{x} = \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \end{pmatrix} = \begin{pmatrix} g[t_1] \\ g[t_1, t_2] \\ g[t_1, t_2, t_3] \\ \vdots \end{pmatrix}$$

- Also costs  $\mathcal{O}(n^2)$ .
- Less prone to over/underflow.

### 3.5 Orthogonal Polynomial Basis (no method given)

Inner Product:  $\langle \vec{u} | \vec{v} \rangle_{ab}^w = \int_a^b [u(t)v(t)] w(t) dt$

Orthogonal Polynomials:  $\langle u_i | u_j \rangle = \delta_{ij}$

Three-Term Recurrence:  $\hat{f}_{k+1}(t) = [A(k)t + B(k)] \hat{f}_k(t) - C(k) \hat{f}_{k-1}(t) \quad (A(k) \neq 0)$

### 3.6 Piecewise [Hermite] Cubic Interpolation

#### Piecewise Cubic:

$n$  knots/pts.  $\Rightarrow n - 1$  cubics

$$\Rightarrow 4(n - 1) \text{ param./eq.}$$

#### Hermite Interpolation:

Using  $k$ -th derivatives as info.

Extra equations can be used for monotonicity/convexity.

#### Hermite Cubic Interpolation:

Continuous 0th and 1st derivatives;  $n - 1$  cubics

$$\Rightarrow [2(n - 1)]_{1\text{st deriv. eq}} + [n - 2]_{2\text{nd deriv. eq.}}$$

$$= 3n - 4 \text{ eq.} \Rightarrow n \text{ free/extra param./eq.}$$

### 3.7 Piecewise Cubic [Spline] Interpolation

#### Spline:

A piecewise func. of  $n$ -polynomials that is  $n$ -differentiable (of differentiability class  $C^{n-1}$ , or  $n - 1$  cont. differentiable).

#### Cubic Spline Interpolation:

Cont. 0th, 1st, and 2nd derivatives;  $n - 1$  cubics

$$\Rightarrow [2(n - 1)]_{1\text{st}} + [n - 2]_{2\text{nd}} + [n - 2]_{3\text{rd}}$$

$$= 4n - 6 \text{ eq.} \Rightarrow 2 \text{ free/extra param./eq.}$$

#### $B$ -splines (basis func.):

Orthog.  $\{\phi_j(t)\}$  are  $j$ -poly. splines w/ local compact support and look like bells. (not much detail here).



## 4 Numerical Integration/Quadrature, $I(f) \equiv \int_a^b f(x)dx$

### 4.1 $\infty$ -Norm and Condition Number

Function  $\infty$ -Norm:

[Abs.] Integration Condition Number if  $\hat{b}$ :

$$\|f(x)\|_{\infty} = \max_{x \in [a,b]} f(x)$$

$$\left| \int_a^{\hat{b}} f(x)dx - \int_a^b f(x)dx \right| = \left| \int_b^{\hat{b}} f(x)dx \right| \leq (\hat{b} - b) \|f(x)\|_{\infty}$$

[Abs.] Integration Condition Number if  $\hat{f}$ :

[Rel.] Integration Condition Number if  $\hat{f}$ :

$$\begin{aligned} \left| \int_a^b \hat{f}(x) - f(x) dx \right| &\leq \int_a^b |\hat{f}(x) - f(x)| dx \\ &\leq (b-a) \|\hat{f}(x) - f(x)\|_{\infty} \\ \left| \frac{\Delta I}{\Delta f} \right| &\leq \boxed{b-a} \end{aligned}$$

$$\begin{aligned} \left| \frac{\Delta I/I}{\Delta f/f} \right| &\leq \frac{(b-a)/\left| \int_a^b f(x)dx \right|}{1/\|f(x)\|_{\infty}} \\ &= \frac{(b-a)\|f(x)\|_{\infty}}{\left| \int_a^b f(x)dx \right|} \end{aligned}$$

### 4.2 1-D [Interpolary] Quadrature Rule for $f \approx \hat{f}$

$$\hat{f} \in P_{n-1} : \hat{f}(x) = \left( \begin{array}{c} \vec{y} \cdot \vec{\phi}(x) = \sum_{i=1}^n f(x_i) \phi_i(x) \\ \text{(Lagrange Basis Vectors)} \end{array} \right) = \left( \begin{array}{c} \sum_{j=0}^{n-1} c_j x^j \\ \text{(Monomial Basis Vectors)} \end{array} \right) \quad \begin{array}{l} \bullet x_1 < \dots < x_n \\ \bullet f(x_i) = \hat{f}(x_i) \end{array}$$

$$\Rightarrow Q_n(f) \equiv I(\hat{f}) = \int_a^b \hat{f}(x)dx = \sum_{i=1}^n f(x_i) \int_a^b \phi_i(x)dx = \sum_{i=1}^n f(x_i) w_i \quad \begin{array}{l} \bullet x_i, w_i \rightarrow 2n \text{ max param.} \\ \bullet a \leq x_1 < \dots < x_n \leq b \\ \bullet \text{closed if equality, open if not} \end{array}$$

Method of Undetermined Coefficients

$$\int_a^b \left( \sum_{j=0}^{n-1} c_j x^j \right) dx = \sum_{i=1}^n \left( \sum_{j=0}^{n-1} c_j x_i^j \right) w_i \quad \int_a^b x^j dx = \sum_{i=1}^n x_i^j w_i = \frac{b^{j+1} - a^{j+1}}{j+1} \equiv z_j \quad \rightarrow \boxed{z_0 = \sum w_i = b-a}$$

$$\sum_{j=0}^{n-1} c_j \left( \int_a^b x^j dx \right) = \sum_{j=0}^{n-1} c_j \left( \sum_{i=1}^n x_i^j w_i \right) \Rightarrow \begin{array}{l} \text{(Vandermode Matrix)} \\ \left( \begin{array}{cccc} 1 & 1 & 1 & \dots \\ x_1 & x_2 & x_3 & \dots \\ x_1^2 & x_2^2 & x_3^2 & \dots \\ \vdots & \vdots & \vdots & \end{array} \right) \vec{w} = \vec{z} \quad \begin{array}{l} \text{System of Moment} \\ \text{Equations} \end{array} \end{array}$$

(maybe some dot product to isolate terms)

Error  $I$  :  $|\Delta I| \leq (b-a) \|f - \hat{f}\|_{\infty} \leq \frac{b-a}{4n} h^n \|f^{(n)}\|_{\infty} \leq \frac{h^{n+1}}{4} \|f^{(n)}\|_{\infty} \rightarrow \text{error decreases if } f^{(n)} \text{ is well behaved}$

Error  $Q_n$  :  $g \approx f \rightarrow |Q_n(f) - Q_n(g)| \leq \boxed{\sum |w_i| \|f-g\|_{\infty}} \Rightarrow \boxed{\forall w_i \geq 0 \rightarrow \text{cond}(Q_n) = b-a}$   
 $= \left| \sum w_i [f(x_i) - g(x_i)] \right|$  (otherwise using  $Q_n$  might be unstable.)

[Rule] Degree,  $d$ :  $\forall p(x) \in P_d$ , rule  $Q(p) = I(p)$ , but not  $\forall p \in P_{d+1}$

Newton-Cotes Quadrature [Rule]:  $n$  evenly-spaced  $x_i \rightarrow n$  param. for  $w_i$

Midpoint Rule ( $Q_1$ ) :	$M(f) = \frac{b-a}{1} f(\frac{a+b}{2})$	$\vec{w} = (b-a)[1]^T$
Trapezoidal Rule ( $Q_2$ ) :	$T(f) = \frac{b-a}{2} [f(a) + f(b)]$	$\vec{w} = (b-a)[\frac{1}{2}, \frac{1}{2}]^T$
Simpsons's Rule ( $Q_3$ ) :	$S(f) = \frac{b-a}{6} [f(a) + 4f(\frac{a+b}{2}) + f(b)]$	$\vec{w} = (b-a)[\frac{1}{6}, \frac{4}{6}, \frac{1}{6}]^T$

• Taylor Expansion and Error

$f(x) = \sum_{m=0} \frac{f^{(m)}(\frac{a+b}{2})}{m!} (x - \frac{a+b}{2})^m$	$I(f) = \sum_{m=0} \frac{f^{(m)}(\frac{a+b}{2})}{(m+1)!} \frac{(2x-a-b)^{m+1}}{2^{m+1}} \Big _a^b$	
$T(f) = \frac{b-a}{2} \sum_{m=0} \frac{f^{(m)}(\frac{a+b}{2})}{m!} \frac{(b-a)^m}{2^m} [(-1)^m + 1]$	$= \sum_{m=0}^{\text{even}} \frac{f^{(m)}(\frac{a+b}{2})}{2^m(m+1)!} (b-a)^{m+1}$	
$= \sum_{m=0}^{\text{even}} \left[ \frac{f^{(m)}(\frac{a+b}{2})}{2^m m!} \right] (b-a)^{m+1}$	$= M(f) + \sum_{m=2}^{\text{even}} \frac{E_m(f)}{m+1} h^{m+1}$	$Q_1$ error is $f^{(2)}$ derivative, not $f^{(1)}$ !
$= M(f) + \sum_{m=2}^{\text{even}} \left[ E_m(f) \right] h^{m+1}$	$= T(f) - \sum_{m=2}^{\text{even}} m \frac{E_m(f)}{m+1} h^{m+1}$	$Q_2$ error is $f^{(2)}$ & twice as large as $Q_1$
$S(f) = \left[ \frac{2}{3}M(f) + \frac{1}{3}T(f) \right]$	$= S(f) - \sum_{m=4}^{\text{even}} \frac{m-2}{3} \frac{E_m(f)}{m+1} h^{m+1}$	$Q_3$ error is $f^{(4)}$ derivative, not $f^{(3)}$ !

- $n$  is even :  $Q_n$  error is expected  $f^{(n)}$  derivative  $Q(p_{n-1}) = I(p_{n-1}) \rightarrow \boxed{d = n - 1}$
- $n$  is odd :  $Q_n$  error is  $f^{(n+1)}$  derivative  $Q(p_n) = I(p_n) \rightarrow \boxed{d = n}$
- 2 Rule Error: Est. diff. between  $T(f)$  and  $M(f)$  can be used to est.  $I(f)$  error in using either.
- Can use subinterval, so can be progressive.
- Evenly-spaced  $x_i$  exhibit the Runge Phenom.  $\rightarrow \boxed{Q_\infty(f) \text{ isn't always } I(f)}$
- $\boxed{\text{Ill-conditioned and unstable}}$  :  $(n \geq 11 \Rightarrow \exists w_i < 0), (\sum_i^\infty |w_i| \rightarrow \infty)$

Curtis-Clenshaw Quadrature [Rule]:  $n$  Chebyshev Nodes,  $x_i \rightarrow n$  param. for  $w_i$

- |  |   |
|--|---|
| <ul style="list-style-type: none"> <li>• <math>\forall n : \forall w_i &gt; 0 \Rightarrow \text{cond}(Q) = b - a</math></li> <li>• <math>\lim_{n \rightarrow \infty} C_n(f) = I(f)</math></li> <li>• <math>\boxed{d_n = n - 1}</math></li> </ul> | <ul style="list-style-type: none"> <li>• <math>\exists</math> an algorithm w/ Chebyshev polynomials to find integrand w/o solving for <math>w_i</math>.</li> <li>• Using Chebyshev polynomial zeroes is the classical CCQ.</li> <li>• Using Chebyshev extrema leads to a progressive rule [practical CCQ].</li> </ul> |
|--|---|

Gaussian Quadrature [Rule]:  $\boxed{2n \text{ free param. for } x_i, w_i} \Rightarrow \boxed{d_n = 2n - 1}$

- $x_i, w_i : x_{n < i \leq 2n} = w_{n < i \leq 2n} = 0 \rightarrow \boxed{\begin{pmatrix} 1 & \dots & 1 & 0 & \dots \\ x_1 & \dots & x_n & 0 & \dots \\ x_1^2 & \dots & x_n^2 & 0 & \dots \\ \vdots & & \vdots & \vdots & \end{pmatrix} \begin{pmatrix} \vdots \\ w_n \\ 0 \\ \vdots \end{pmatrix} = \vec{z}(a, b)}$  usually  $x_i \notin \mathbb{Q}$

- Ortho. Poly. :  $\langle p_n(x) | x^k \rangle_{ab} = 0 \Rightarrow \boxed{x_i : \begin{matrix} p_n(x_i) = 0, & x_i \in \mathbb{R}, \\ x_i \neq x_{j \neq i}, & x_i \in (a, b) \end{matrix}}$  [−1,1] – Legendre  
(−∞,∞) – Hermite  
[0,∞) – Laguerre

- Interval Transform :  $\boxed{\int_a^b f(t) dt = \frac{b-a}{\beta-\alpha} \int_\alpha^\beta f(t) dx \quad t = \frac{(b-a)x + a\beta - b\alpha}{\beta - \alpha}}$

- $\forall n : \forall w_i > 0 \Rightarrow \text{cond}(Q) = b - a$       •  $\lim_{n \rightarrow \infty} G_n(f) = I(f)$

- $n = 2m + 1 \rightarrow \frac{a+b}{2} \in \{x_i\}_n$ ; otherwise usually  $\{x_i\}_n \cup \{x_i\}_{\neq n} = 0 \rightarrow \boxed{\text{Not progressive}}$

- Progressive Gauss-Kronrod,  $K_{2n+1}$  :  $n \text{ from } G_n \rightarrow \frac{n+1}{2n+1} \text{ param for } x_{i>n} \Rightarrow \boxed{d_{2n+1} = 3n + 1 < 4n + 1}$   
GK 2-Rule Error :  $\boxed{\Delta I(f) \approx (200|G_n - K_{2n+1}|)^{1.5}}$

Progressive Gauss-Patterson,  $P_{4n+3}$  :  $2n + 1 \text{ from } K_{2n+1} \rightarrow \frac{2n+2}{4n+3} \text{ param for } x_{i>n} \Rightarrow \boxed{d_{4n+3} = 6n + 4 < 8n + 5}$

- Closed Gauss-Randau :  $x_i \in [a, b) \text{ or } (a, b] \rightarrow \boxed{d = 2n - 2}$

Closed Gauss-Lobatto :  $x_i \in [a, b] \rightarrow \boxed{d = 2n - 3}$

Composite [ $k$ -Subintervals] Quadrature for Rule  $Q_n$ :  $Q_n \rightarrow Q_{kn} \text{ or } Q_{kn-(k-1)},$

- $\lim_{k \rightarrow \infty} C_{k,n} = \sum_{j=1}^{k \rightarrow \infty} \left[ \sum_{i=1}^n w_i f(x_{ji}) \right] = \sum_{i=1}^n \frac{w_i}{h_k} \left[ \sum_{j=1}^{k \rightarrow \infty} h_k f(x_{ji}) \right] = I(f) \sum_{i=1}^n \frac{w_i}{h_k} = I(f)$   $h_k = (b-a)/k$   
 $\geq (x_{jn} - x_{j1})$   
 $\boxed{d \geq 0} \Rightarrow \sum w_i = h_k$
- Error :  $\mathcal{O}(h^{m+1}) \rightarrow \mathcal{O}(kh_k^{m+1}) = \boxed{\mathcal{O}(h_k^m)}$  ( $k > 1$ )

Adaptive Quadrature for Rule  $Q_n$ : Divide subinterval until a tolerance is met.

### 4.3 $n$ -D Integration

Double Integral: Use a pair of 1-D routines for the inner/outer integral.

( $n > 2$ )-Dimension Integral: Monte Carlo is best (error  $1/\sqrt{n} \rightarrow 0$ ).

### 4.4 Other Integrals

Tabular Data: Integrate a piecewise interpolant.

Improper Integral: Separate the integral, do a variable change,  
or add/subtract a term to remove singularities.

(Fredholm) Integral Equations: skipped

## 4.5 Richardson Extrapolation [for Integration]

$$\begin{aligned} F(h) &= I(f) + a_1 h^p + \mathcal{O}(h^{q>p}) \\ F(\frac{h}{k}) &= I(f) + a_1 (\frac{h}{k})^p + \mathcal{O}(h^{r\geq q}) \end{aligned} \Rightarrow \boxed{I(f) = \frac{k^p F(\frac{h}{k}) - F(h)}{k^p - 1} + \mathcal{O}(h^{q>p})}$$

- Romberg Integration [Quadratic Extrapolation for Comp. Trapezoidal Rule] :

$$\begin{aligned} T(f, \frac{h}{2^k}) &= I(f) + 2^k \left[ a_1 (\frac{h}{2^k})^3 + \mathcal{O}(\frac{h}{2^k}^5) \right] \\ T_{k,j=0} &= I(f) + h a_1 [\frac{h}{2^k}]^2 + h \mathcal{O}([\frac{h}{2^k}]^4) \end{aligned} \Rightarrow T_{k+1,j+1} \equiv \frac{4^{j+1} T_{k+1,j} - T_{k,j}}{4^{j+1} - 1} \quad (1 \leq j \leq k)$$

$$4T_{k+1,0} = 4I(f) + h a_1 [\frac{h}{2^k}]^2 + \frac{h}{4} \mathcal{O}([\frac{h}{2^k}]^4)$$

$$\boxed{I(f) = T_{k,j} + \mathcal{O}(h^{2j+2})}$$

## 5 Numerical Differentiation

Conditioning: Inverse of Integration - which smooths noisy data - so derivatives are inherently sensitive to small changes.

### 5.1 Finite-Difference Approx

$$\begin{aligned} f'(x) &= \frac{f(x+h)-f(x)}{h} - \sum_{n=2}^{\infty} \frac{f^{(n)}(x)}{n!} h^{n-1} \\ &= \frac{f(x)-f(x-h)}{h} - \sum_{n=2}^{\infty} \frac{f^{(n)}(x)}{n!} (-h)^{n-1} \\ &= \frac{f(x+h)-f(x-h)}{2h} - \sum_{n=3}^{\text{odd}} \frac{f^{(n)}(x)}{n!} h^{n-1} \end{aligned}$$

- Use more points  $n$  for higher order approx.

### 5.2 Deriving Interpolant

$$\begin{aligned} f(x) &\approx \hat{f}_n(x) = p_{n-1}(x) \in P_{n-1} \\ f^{(m)}(x) &\approx \hat{f}_n^{(m)}(x) \end{aligned}$$

- Equivalent but easier than finite-diff. approach.
- Using more points  $n$  leads to better accuracy.
- Polynomials, or other interpolants like trig. func. can be used.

### 5.3 Richardson Extrapolation [for Differentiation]

$$\begin{aligned} F(h) &= D(f) + a_1 h^p + \mathcal{O}(h^{q>p}) \\ F(\frac{h}{k}) &= D(f) + a_1 (\frac{h}{k})^p + \mathcal{O}(h^{r\geq q}) \end{aligned} \Rightarrow \boxed{D(f) = \frac{k^p F(\frac{h}{k}) - F(h)}{k^p - 1} + \mathcal{O}(h^{q>p})}$$

- E.g.  $D(f) = \frac{f(x+h)-f(x)}{h} + \mathcal{O}(h)$

$$\begin{aligned} F(h) &= \frac{f(x+h)-f(x)}{h} \\ F(\frac{h}{2}) &= \frac{f(x+\frac{h}{2})-f(x)}{h/2} \end{aligned} \Rightarrow \boxed{D(f) = \frac{2 \cdot \frac{f(x+h/2)-f(x)}{h/2} - \frac{f(x+h)-f(x)}{h}}{2-1} + \mathcal{O}(h^2)}$$