

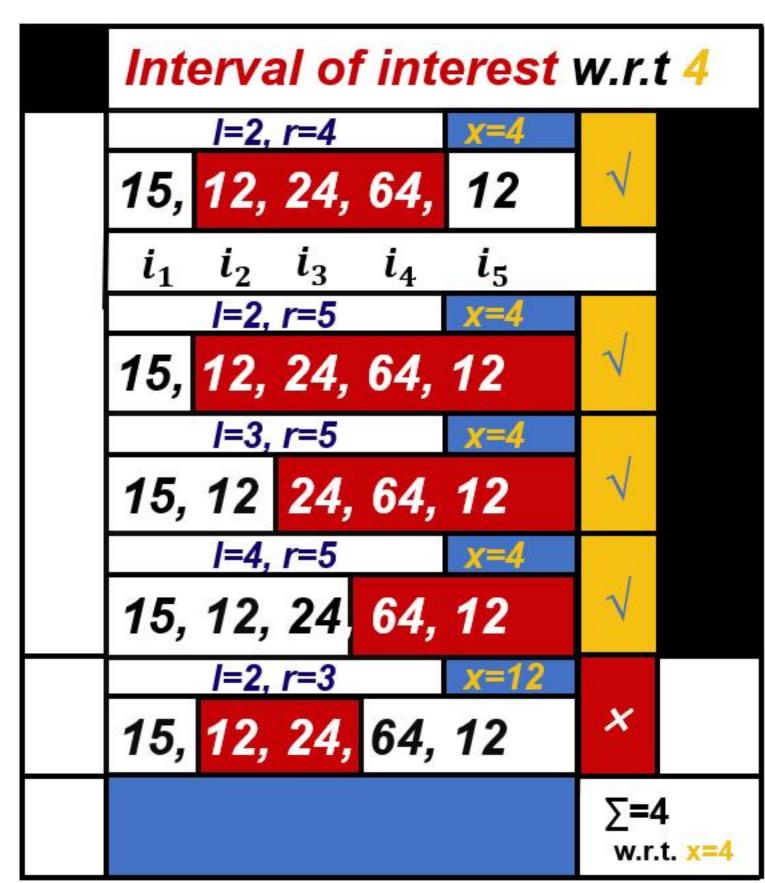
# Motivation

A sequence of intergers are given: i\_1, i\_2, ..., i\_n. Define the interval of interest(IOI) with respect to x, parameterized by[I, r] subjected to: gcd(i\_I, i\_I+1, ..., i\_r)=x, where 1<=I<=r<=n.

Gcd here means the greatest common divisor.

## Goal

Count all possible intervals with respect to x=x\_1, x\_2, ..., x\_q for the given sequence i\_1, i\_2, ..., i\_n.



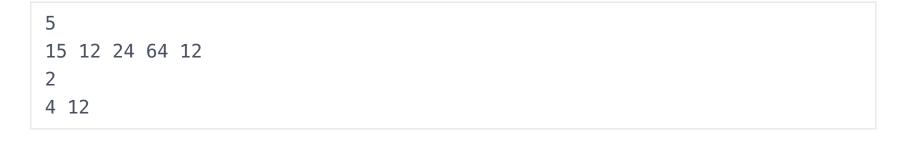
### 输入

- 1. The first line: an integer n,  $(1 \le n \le 10^5)$ , indicating the length of the given sequence.
- 2. Next line: n integers separated by space: i\_1, i\_2, ..., i\_n (1<=i\_{}<=10^9).
- 3. The third line: an integer q,  $(1 \le q \le 3*10^5)$ , indicating the number of xs
- 4. The forth line: q integers separted by space:  $x_1$ ,  $x_2$ , ...,  $x_q$ ,  $(1 <= x_{\{\}} <= 10^9)$ .

输出

For each x, output the number of possible IOIs with respect to x\_1, x\_2, ...

### 输入样例 1 🖺



### 输出样例 1



## 提示

Observation1: gcd(a, b, c) = gcd(gcd(a, b), c)

Observation2: for the sequence i\_1, ..., i\_n, define the sequence k\_1=i\_1, k\_j=gcd(k\_{j-1},i\_{j}) for 2<=j<=n. The number of distinct values in k sequence is no more than 1+log\_2^{i\_1}.

**Hint1**: Divide & conquer may be useful:

- 1. count IOIs at left half(by recursion)
- 2. count IOIs at right half(by recursion)
- 3. merge and conut intervals that start from the left half and end at right half(how to count efficiently?, hint: observation 2)

Hint2: try to set a AVL tree / hash, count gcd for all x in one time.

Hint3: the solution of nlogn is possible