CS101 Algorithms and Data Structures

Divide and Conquer, Merge/Quick Sort Textbook Ch 4, 7



Outline

- Insertion sort
- Bubble sort
- Heap sort
- Merge sort
- Quicksort

Outline

This topic covers merge sort

- A recursive divide-and-conquer algorithm
- Merging two lists
- The merge sort algorithm
- A run-time analysis

Merge Sort

The merge sort algorithm is defined recursively:

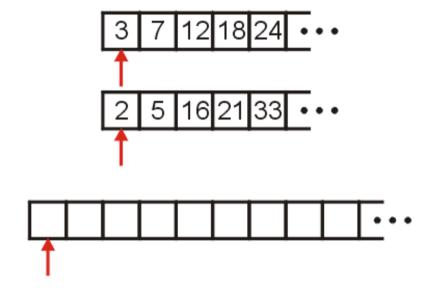
- If the list is of size 1, it is sorted—we are done;
- Otherwise:
 - Divide an unsorted list into two sub-lists,
 - Sort each sub-list recursively using merge sort, and
 - Merge the two sorted sub-lists into a single sorted list

This strategy is called *divide-and-conquer*

Question: How can we merge two sorted sub-lists into a single sorted list?

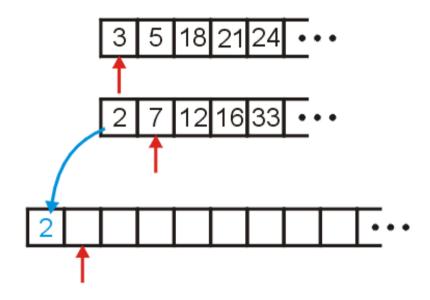
Consider the two sorted arrays and an empty array

Define three indices at the start of each array

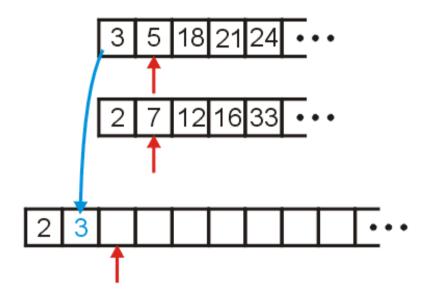


We compare 2 and 3: 2 < 3

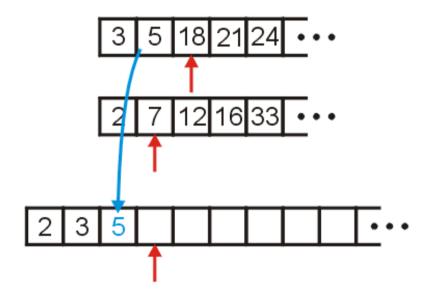
- Copy 2 down
- Increment the corresponding indices



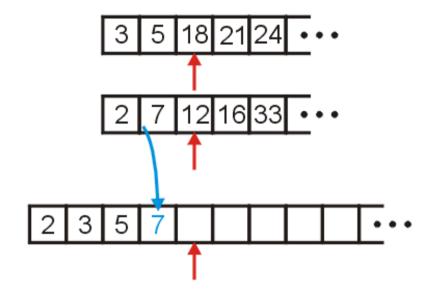
- Copy 3 down
- Increment the corresponding indices



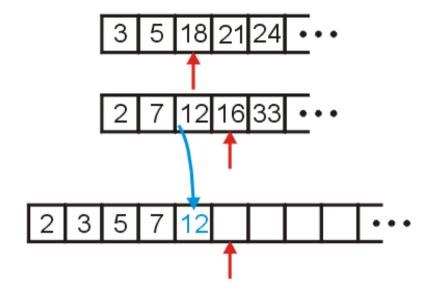
- Copy 5 down
- Increment the appropriate indices



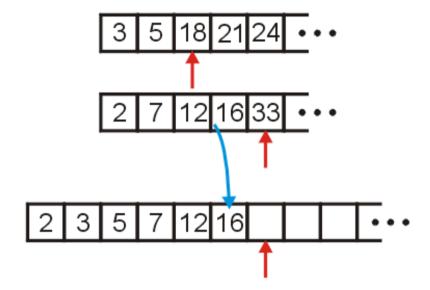
- Copy 7 down
- Increment...



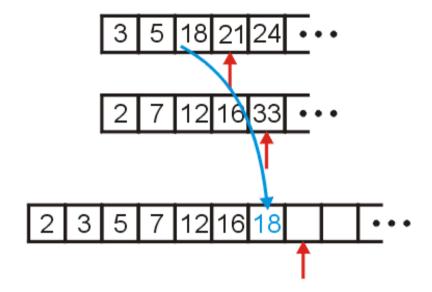
- Copy 12 down
- Increment...



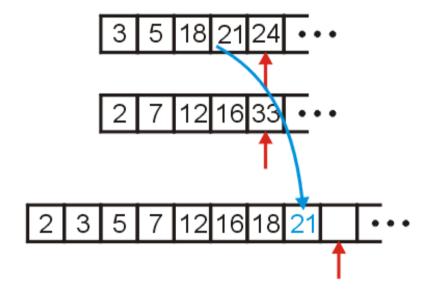
- Copy 16 down
- Increment...



- Copy 18 down
- Increment...

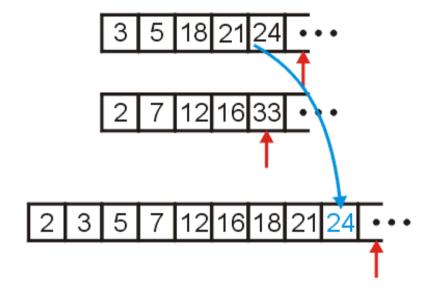


- Copy 21 down
- Increment...

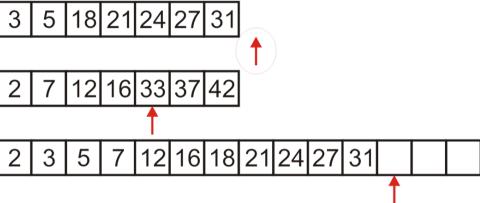


We compare 24 and 33

- Copy 24 down
- Increment...

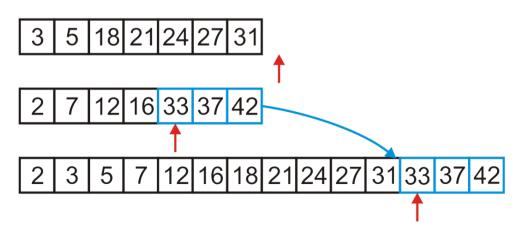


We would continue until we have passed beyond the limit of one of the two arrays



After this, we simply copy over all remaining entries in the non-

empty array



Merging Two Lists

Programming a merge is straight-forward:

- the sorted arrays, array1 and array2, are of size n1 and n2, respectively, and
- we have an empty array, arrayout, of size n1 + n2

Define three variables

int
$$i1 = 0$$
, $i2 = 0$, $k = 0$;

which index into these three arrays

Merging Two Lists

We can then run the following loop:

```
#include <cassert>
//...
int i1 = 0, i2 = 0, k = 0;
while ( i1 < n1 \&\& i2 < n2 ) {
    if ( array1[i1] < array2[i2] ) {</pre>
        arrayout[k] = array1[i1];
        ++i1;
    } else {
        assert( array1[i1] >= array2[i2] );
        arrayout[k] = array2[i2];
        ++i2;
    ++k;
```

Merging Two Lists

We're not finished yet, we have to empty out the remaining array

```
for ( ; i1 < n1; ++i1, ++k ) {
    arrayout[k] = array1[i1];
}

for ( ; i2 < n2; ++i2, ++k ) {
    arrayout[k] = array2[i2];
}</pre>
```

Analysis of merging

Time: we have to copy $n_1 + n_2$ elements

- Hence, merging may be performed in $\Theta(n_1 + n_2)$ time
- If the arrays are approximately the same size, $n = n_1 \approx n_2$, we can say that the run time is Θ(n)

Space: we cannot merge two arrays in-place

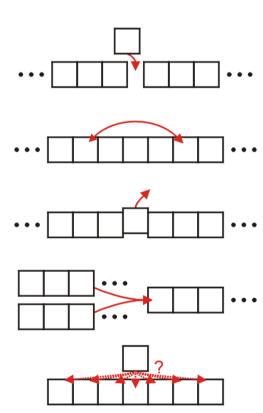
- This algorithm always required the allocation of a new array
- Therefore, the memory requirements are also $\Theta(n)$

The Algorithm

Recall the five sorting techniques:

- Insertion
- Exchange
- Selection
- Merging
- Distribution

Clearly merge sort falls into the fourth category



The Algorithm

The merge sort algorithm is defined recursively:

- If the list is of size 1, it is sorted—we are done;
- Otherwise:
 - Divide an unsorted list into two sub-lists,
 - Sort each sub-list recursively using merge sort, and
 - Merge the two sorted sub-lists into a single sorted list

In practice:

- If the list size is less than a threshold, use an algorithm like insertion sort
- Otherwise: 与寄存器互换相关、是快
 - Divide...

```
Suppose we already have a function
   template <typename Type>
   void merge( Type *array, int a, int b, int c );
that assumes that the entries
   array[a] through array[b - 1], and
   array[b] through array[c - 1]
are sorted and merges these two sub-arrays into a single sorted
array from index a through index c - 1, inclusive
```

For example, given the array,

9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33
13	77	49	35	61	3	23	48	73	89	95	17	32	37	57	94	99	28	15	55	7	51	88	97	62

a call to

void merge(array, 14, 20, 26);

merges the two sub-lists

9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33
13						23																		62

forming

9	1	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32	33
1	3	77	49	35	61	3	17	23	32	37	48	57	73	89	94	95	99	28	15	55	7	51	88	97	62

```
We implement a function
    template <typename Type>
    void merge_sort( Type *array, int first, int last );
that will sort the entries in the positions first <= i and i < last</pre>
```

- If the number of entries is less than N, call insertion sort
- Otherwise:
 - Find the mid-point,
 - · Call merge sort recursively on each of the halves, and
 - Merge the results

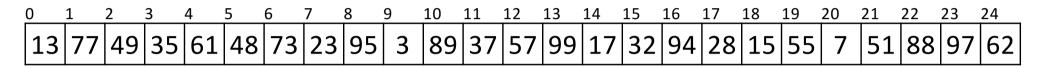
```
template <typename Type>
void merge_sort( Type *array, int first, int last ) {
    if ( last - first <= N ) {
        insertion_sort( array, first, last );
    } else {
        int midpoint = (first + last)/2;

        merge_sort( array, first, midpoint );
        merge_sort( array, midpoint, last );
        merge( array, first, midpoint, last );
    }
}</pre>
```

Like merge sort, insertion sort will sort a sub-range of the array:

```
template <typename Type>
void insertion sort( Type *array, int first, int last ) {
    for ( int k = first + 1; k < last; ++k ) {
        Type tmp = array[k];
        for ( int j = k; k > first; --j ) {
            if ( array[j - 1] > tmp ) {
                array[j] = array[j - 1];
            } else {
                array[j] = tmp;
                goto finished;
        array[first] = tmp;
        finished: ;
}
```

Consider the following is of unsorted array of 25 entries



We will call insertion sort if the list being sorted of size N=6 or less

We call merge_sort(array, 0, 25)

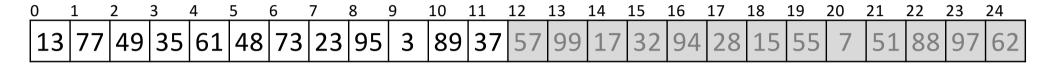
```
merge_sort( array, 0, 25 )
```

We are calling merge_sort(array, 0, 25)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
13	77	49	35	61	48	73	23	95	3	89	37	57	99	17	32	94	28	15	55	7	51	88	97	62

```
First, 25-0>6, so find the midpoint and call merge_sort recursively midpoint = (0 + 25)/2; // == 12 merge_sort( array, 0, 12 );
```

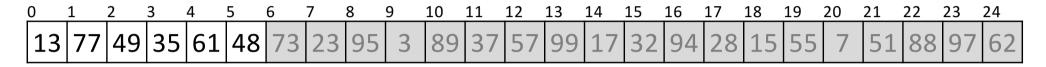
We are now executing merge_sort(array, 0, 12)



First, 12-0>6, so find the midpoint and call merge_sort recursively midpoint = (0 + 12)/2; // == 6 merge_sort(array, 0, 6);

```
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

We are now executing merge_sort(array, 0, 6)



Now, $6-0 \le 6$, so find we call insertion sort

```
merge_sort( array, 0, 6 )
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

Insertion sort just sorts the entries from 0 to 5

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
13	77	49	35	61	48	73	23	95	3	89	37	57	99	17	32	94	28	15	55	7	51	88	97	62

```
insertion_sort( array, 0, 6 )
merge_sort( array, 0, 6 )
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

Insertion sort just sorts the entries from 0 to 5



This function call completes and so we exit

```
insertion_sort( array, 0, 6 )
merge_sort( array, 0, 6 )
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

This call to merge_sort is now also finished, so it, too, exits

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
13	35	48	49	61	77	73	23	95	3	89	37	57	99	17	32	94	28	15	55	7	51	88	97	62

```
merge_sort( array, 0, 6 )
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

We return to continue executing merge_sort(array, 0, 12)

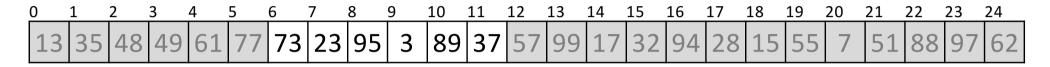
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
13	35	48	49	61	77	73	23	95	3	89	37	57	99	17	32	94	28	15	55	7	51	88	97	62

We continue calling

```
midpoint = (0 + 12)/2; // == 6
merge_sort( array, 0, 6 );
merge_sort( array, 6, 12 );
```

```
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

We are now executing merge_sort(array, 6, 12)



Now, $12-6 \le 6$, so find we call insertion sort

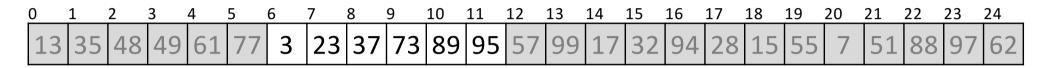
```
merge_sort( array, 6, 12 )
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

Insertion sort just sorts the entries from 6 to 11

0	1	2	3	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
1	3 3	5 4	18	49	61	77	73	23	95	3	89	37	57	99	17	32	94	28	15	55	7	51	88	97	62

```
insertion_sort( array, 6, 12 )
merge_sort( array, 6, 12 )
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

Insertion sort just sorts the entries from 6 to 11



This function call completes and so we exit

```
insertion_sort( array, 6, 12 )
merge_sort( array, 6, 12 )
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

This call to merge_sort is now also finished, so it, too, exits

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
13	35	48	49	61	77	3	23	37	73	89	95	57	99	17	32	94	28	15	55	7	51	88	97	62

```
merge_sort( array, 6, 12 )
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

We return to continue executing merge_sort(array, 0, 12)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
13	35	48	49	61	77	3	23	37	73	89	95	57	99	17	32	94	28	15	55	7	51	88	97	62

We continue calling

```
midpoint = (0 + 12)/2; // == 6
merge_sort( array, 0, 6 );
merge_sort( array, 6, 12 );
merge( array, 0, 6, 12 );
```

```
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

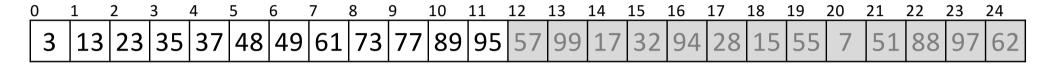
We are executing merge (array, 0, 6, 12)



These two sub-arrays are merged together

```
merge( array, 0, 6, 12 )
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

We are executing merge (array, 0, 6, 12)

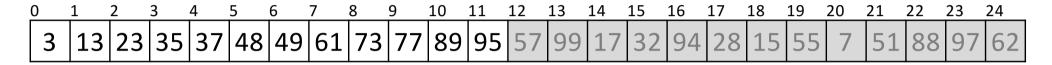


These two sub-arrays are merged together

This function call exists

```
merge( array, 0, 6, 12 )
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

We return to executing merge_sort(array, 0, 12)



We are finished calling this function as well

```
midpoint = (0 + 12)/2; // == 6
merge_sort( array, 0, 6 );
merge_sort( array, 6, 12 );
merge( array, 0, 6, 12 );
```

Consequently, we exit

```
merge_sort( array, 0, 12 )
merge_sort( array, 0, 25 )
```

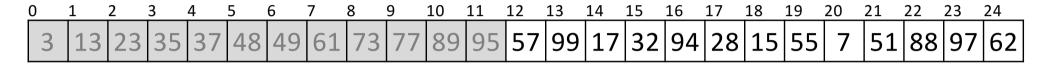
We return to executing merge_sort(array, 0, 25)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	13	23	35	37	48	49	61	73	77	89	95	57	99	17	32	94	28	15	55	7	51	88	97	62

We continue calling

```
midpoint = (0 + 25)/2; // == 12
merge_sort( array, 0, 12 );
merge_sort( array, 12, 25 );
```

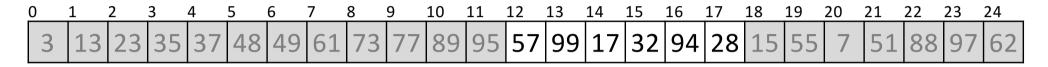
We are now executing merge_sort(array, 12, 25)



```
First, 25 – 12 > 6, so find the midpoint and call merge_sort recursively midpoint = (12 + 25)/2; // == 18 merge_sort( array, 12, 18 );
```

```
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We are now executing merge_sort(array, 12, 18)



Now, $18 - 12 \le 6$, so find we call insertion sort

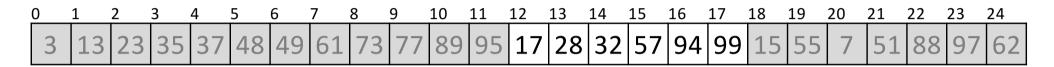
```
merge_sort( array, 12, 18 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

Insertion sort just sorts the entries from 12 to 17

0		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	3	13	23	35	37	48	49	61	73	77	89	95	57	99	17	32	94	28	15	55	7	51	88	97	62

```
insertion_sort( array, 12, 18 )
merge_sort( array, 12, 18 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

Insertion sort just sorts the entries from 12 to 17



This function call completes and so we exit

```
insertion_sort( array, 12, 18 )
merge_sort( array, 12, 18 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

This call to merge_sort is now also finished, so it, too, exits

0	1		2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	1	13	23	35	37	48	49	61	73	77	89	95	17	28	32	57	94	99	15	55	7	51	88	97	62

```
merge_sort( array, 12, 18 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We return to continue executing merge_sort(array, 12, 25)

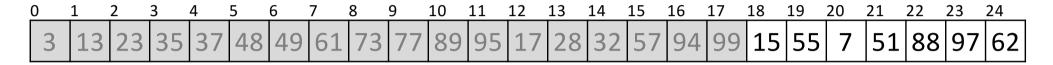
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	13	23	35	37	48	49	61	73	77	89	95	17	28	32	57	94	99	15	55	7	51	88	97	62

We continue calling

```
midpoint = (12 + 25)/2; // == 18
merge_sort( array, 12, 18 );
merge_sort( array, 18, 25 );
```

```
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

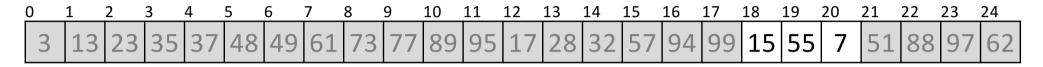
We are now executing merge_sort(array, 18, 25)



```
First, 25-18>6, so find the midpoint and call merge_sort recursively midpoint = (18 + 25)/2; // == 21 merge_sort( array, 18, 21 );
```

```
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We are now executing merge_sort(array, 18, 21)



Now, $21 - 18 \le 6$, so find we call insertion sort

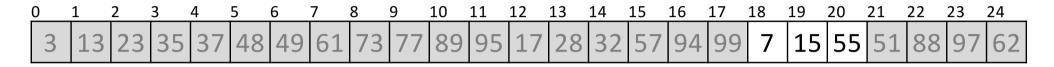
```
merge_sort( array, 18, 21 )
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

Insertion sort just sorts the entries from 18 to 20

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	13	23	35	37	48	49	61	73	77	89	95	17	28	32	57	94	99	15	55	7	51	88	97	62

```
insertion_sort( array, 18, 21 )
merge_sort( array, 18, 21 )
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

Insertion sort just sorts the entries from 18 to 20



This function call completes and so we exit

```
insertion_sort( array, 18, 21 )
merge_sort( array, 18, 21 )
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

This call to merge_sort is now also finished, so it, too, exits

0		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
	3	13	23	35	37	48	49	61	73	77	89	95	17	28	32	57	94	99	7	15	55	51	88	97	62

```
merge_sort( array, 18, 21 )
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We return to executing merge_sort(array, 18, 25)

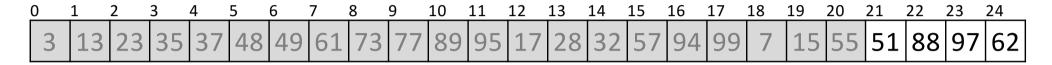
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	13	23	35	37	48	49	61	73	77	89	95	17	28	32	57	94	99	7	15	55	51	88	97	62

We continue calling

```
midpoint = (18 + 25)/2; // == 21
merge_sort( array, 18, 21 );
merge_sort( array, 21, 25 );
```

```
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We are now executing merge_sort(array, 21, 25)



Now, $25 - 21 \le 6$, so find we call insertion sort

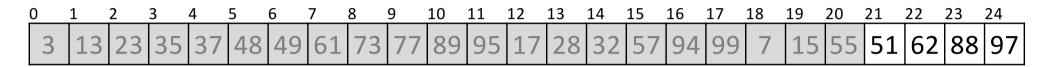
```
merge_sort( array, 21, 25 )
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

Insertion sort just sorts the entries from 21 to 24

0		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
	3	13	23	35	37	48	49	61	73	77	89	95	17	28	32	57	94	99	7	15	55	51	88	97	62

```
insertion_sort( array, 21, 25 )
merge_sort( array, 21, 25 )
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

Insertion sort just sorts the entries from 21 to 24



This function call completes and so we exit

```
insertion_sort( array, 21, 25 )
merge_sort( array, 21, 25 )
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

This call to merge_sort is now also finished, so it, too, exits

0		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
	3	13	23	35	37	48	49	61	73	77	89	95	17	28	32	57	94	99	7	15	55	51	62	88	97

```
merge_sort( array, 21, 25 )
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We return to continue executing merge_sort(array, 18, 25)

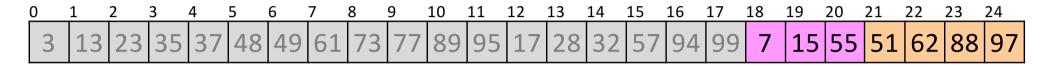
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	13	23	35	37	48	49	61	73	77	89	95	17	28	32	57	94	99	7	15	55	51	62	88	97

We continue calling

```
midpoint = (18 + 25)/2; // == 21
merge_sort( array, 18, 21 );
merge_sort( array, 21, 25 );
merge( array, 18, 21, 25 );
```

```
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

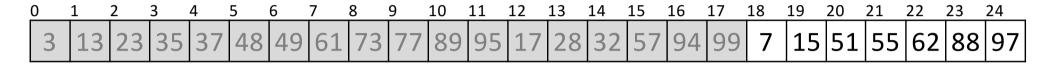
We are executing merge (array, 18, 21, 25)



These two sub-arrays are merged together

```
merge( array, 18, 21, 25 )
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We are executing merge (array, 18, 21, 25)

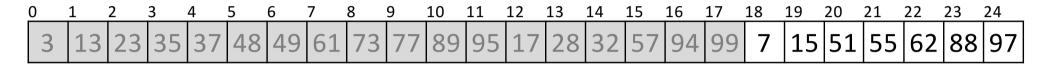


These two sub-arrays are merged together

This function call exists

```
merge( array, 18, 21, 25 )
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We return to executing merge_sort(array, 18, 25)



We are finished calling this function as well

```
midpoint = (18 + 25)/2; // == 21
merge_sort( array, 18, 21 );
merge_sort( array, 21, 25 );
merge( array, 18, 21, 25 );
```

Consequently, we exit

```
merge_sort( array, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We return to continue executing merge_sort(array, 12, 25)

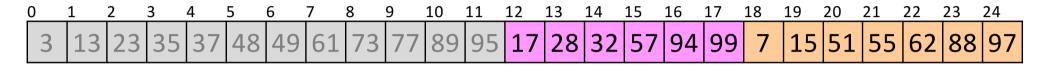
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	13	23	35	37	48	49	61	73	77	89	95	17	28	32	57	94	99	7	15	51	55	62	88	97

We continue calling

```
midpoint = (12 + 25)/2; // == 18
merge_sort( array, 12, 18 );
merge_sort( array, 18, 25 );
merge( array, 12, 18, 25 );
```

```
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

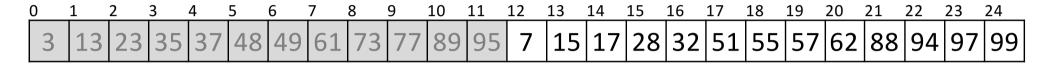
We are executing merge (array, 12, 18, 25)



These two sub-arrays are merged together

```
merge( array, 12, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We are executing merge (array, 12, 18, 25)

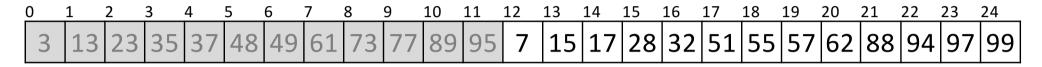


These two sub-arrays are merged together

This function call exists

```
merge( array, 12, 18, 25 )
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

We return to executing merge_sort(array, 12, 25)



We are finished calling this function as well

```
midpoint = (12 + 25)/2; // == 18
merge_sort( array, 12, 18 );
merge_sort( array, 18, 25 );
merge( array, 12, 18, 25 );
```

Consequently, we exit

```
merge_sort( array, 12, 25 )
merge_sort( array, 0, 25 )
```

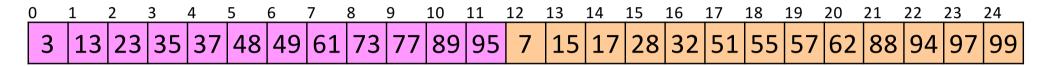
We return to continue executing merge_sort(array, 0, 25)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	13	23	35	37	48	49	61	73	77	89	95	7	15	17	28	32	51	55	57	62	88	94	97	99

We continue calling

```
midpoint = (0 + 25)/2; // == 12
merge_sort( array, 0, 12 );
merge_sort( array, 12, 25 );
merge( array, 0, 12, 25 );
```

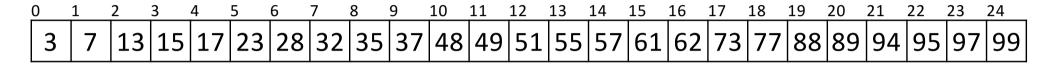
We are executing merge (array, 0, 12, 25)



These two sub-arrays are merged together

```
merge( array, 0, 12, 25 )
merge_sort( array, 0, 25 )
```

We are executing merge (array, 0, 12, 25)



These two sub-arrays are merged together

This function call exists

```
merge( array, 0, 12, 25 )
merge_sort( array, 0, 25 )
```

We return to executing merge_sort(array, 0, 25)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	73	77	88	89	94	95	97	99

We are finished calling this function as well

```
midpoint = (0 + 25)/2; // == 12
merge_sort( array, 0, 12 );
merge_sort( array, 12, 25 );
merge( array, 0, 12, 25 );
```

Consequently, we exit

Example

The array is now sorted

0		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	73	77	88	89	94	95	97	99

Run-time Analysis of Merge Sort

The time required to sort an array of size n > 1 is:

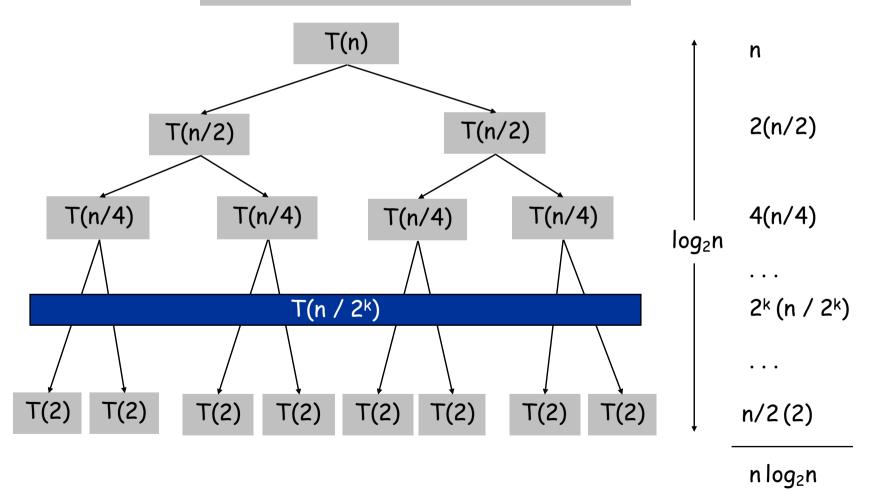
- the time required to sort the first half,
- the time required to sort the second half, and
- the time required to merge the two lists

That is:
$$T(n) = \begin{cases} \mathbf{\Theta}(1) & n = 1 \\ 2T(\frac{n}{2}) + \mathbf{\Theta}(n) & n > 1 \end{cases}$$

Solution: $T(n) = \Theta(n \ln(n))$

Proof by Recursion Tree

$$T(n) = \begin{cases} 0 & \text{if } n = 1 \\ 2T(n/2) + n & \text{otherwise} \\ \text{sorting both halves merging} \end{cases}$$



Run-time Summary

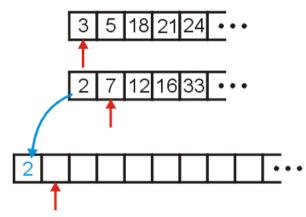
The following table summarizes the run-times of merge sort

Case	Run Time	Comments
Worst	$\Theta(n \ln(n))$	No worst case
Average	$\Theta(n \ln(n))$	
Best	$\Theta(n \ln(n))$	No best case

Why is it not $O(n^2)$

When we are merging, we are comparing values

- What operation prevents us from performing $O(n^2)$ comparisons?
- During the merging process, if 2 came from the second half, it was only compared to 3 and it was not compared to any other of the other n-1 entries in the first array



— In this case, we remove n inversions with one comparison

Comments

In practice, merge sort is faster than heap sort, though they both have the same asymptotic run times

Merge sort requires an additional array

Heap sort does not require

Next we see quick sort

- Faster, on average, than either heap or quick sort
- Requires o(n) additional memory

Merge Sort

The (likely) first proposal of merge sort was by John von Neumann

in 1945

 The creator of the von Neumann architecture used by all modern computers:



http://en.wikipedia.org/wiki/Von_Neumann

Summary

This topic covered merge sort:

- Divide an unsorted list into two equal or nearly equal sub lists,
- Sorts each of the sub lists by calling itself recursively, and then
- Merges the two sub lists together to form a sorted list

Outline

- Insertion sort
- Bubble sort
- Heap sort
- Merge sort
- Quicksort

Outline

In this topic we will look at quicksort:

- The idea behind the algorithm
- The run time and worst-case scenario
- Strategy for avoiding the worst-case: median-of-three
- Implementing quicksort in place
- Examples

Strategy

We have seen two $\Theta(n \ln(n))$ sorting algorithms:

- Heap sort which allows in-place sorting, and
- Merge sort which is faster but requires more memory

We will now look at a recursive algorithm which may be done *almost* in place but which is faster than heap sort

- Use an object in the array (a pivot) to divide the two
- Average case: $\Theta(n \ln(n))$ time and $\Theta(\ln(n))$ memory
- Worst case: $\Theta(n^2)$ time and $\Theta(n)$ memory

We will look at strategies for avoiding the worst case

Quicksort

Merge sort splits the array into two sub-lists and sorts them

 It splits the larger problem into two sub-problems based on location in the array

Consider the following alternative:

 Chose an object in the array and partition the remaining objects into two groups relative to the chosen entry

Quicksort

For example, given

|--|

we can select the middle entry, 44, and sort the remaining entries into two groups, those less than 44 and those greater than 44:

	38	10	26	12	43	3	44	80	95	84	66	79	87	96	81
- 1			l	l											

Notice that 44 is now in the correct location if the list was sorted

 Proceed by recursively applying the algorithm to the first six and last eight entries

Run-time analysis

Like merge sort, we can either:

- Sort the sub-lists using quicksort
- If the size of the sub-list is sufficiently small, apply insertion sort

In the best case, the list will be split into two approximately equal sub-lists, and thus, the run time could be very similar to that of merge sort: $\Theta(n \ln(n))$

What happens if we don't get that lucky?

Worst-case scenario

Suppose we choose the middle element as our pivot and we try ordering a sorted list:

80	38	95	84	66	10	79	2	26	87	96	12	43	81	3
			1		1									ı

Using 2, we partition into

We still have to sort a list of size n-1

The run time is
$$T(n) = T(n-1) + \Theta(n) = \Theta(n^2)$$

- Thus, the run time drops from $n \ln(n)$ to n^2

Worst-case scenario

Our goal is to choose the median element in the list as our pivot:

80	38	95	84	66	10	79	2	26	87	96	12	43	81	3
	1				1									

Using the median element 66, we can get two equal-size sub-lists

3	38	43	12	2	10	26	66	79	87	96	84	95	81	80
---	----	----	----	---	----	----	----	----	----	----	----	----	----	----

Unfortunately, median is difficult to find

Consider another strategy:

Choose the median of the first, middle, and last entries in the list

This will usually give a better approximation of the actual median



Sorting the elements based on 44 results in two sub-lists, each of which must be sorted (again, using quicksort)

Select the 26 to partition the first sub-list:



Select 81 to partition the second sub-list:



If we choose a random pivot, this will, on average, divide a set of n items into two sets of size 1/4 n and 3/4 n, why?

Choosing the median-of-three, this will, on average, divide the n items into two sets of size 5/16 n and 11/16 n

- Median-of-three helps speed the algorithm
- This requires order statistics:

$$2\int_{0}^{\frac{1}{2}} x \cdot (6x(1-x)) dx = \frac{5}{16} = 0.3125$$

Recall that merge sort always divides a list into two equal halves:

The median-of-three will require $\frac{\ln\left(\frac{1}{2}\right)}{\ln\left(\frac{11}{16}\right)} \approx 1.8499$ or 85 % more recursive steps, how to get this?

A single random pivot will require recursive steps

$$\frac{\ln\left(\frac{1}{2}\right)}{\ln\left(\frac{3}{4}\right)} \approx 2.4094 \text{ or } 141 \% \text{ more}$$

Question: what is the affect on run time?

- Surprisingly, not so much
- Here we see the ratios of the recurrence relations for large values of n

$$T_{\text{random pivot}}(n) = T\left(\left\lfloor \frac{1}{4}n \right\rfloor\right) + T\left(\left\lceil \frac{3}{4}n \right\rceil\right) + n$$

$$T_{\text{median of 3}}(n) = T\left(\left\lfloor \frac{5}{16}n \right\rfloor\right) + T\left(\left\lceil \frac{11}{16}n \right\rceil\right) + n$$

20 billion

$$T_{\text{median}}(n) = T\left(\left[\frac{1}{2}n\right]\right) + T\left(\left[\frac{1}{2}n\right]\right) + n$$

$$T_{\text{median}}(n)$$

$$T_{\text{median}}(n)$$

$$T_{\text{median}}(n)$$

$$T_{\text{median}}(n)$$

$$T_{\text{median}}(n)$$

$$T_{\text{median}}(n)$$

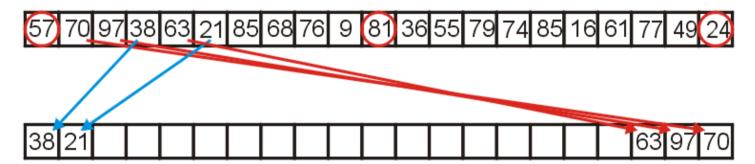
10 billion

If we choose to allocate memory for an additional array, we can implement the partitioning by

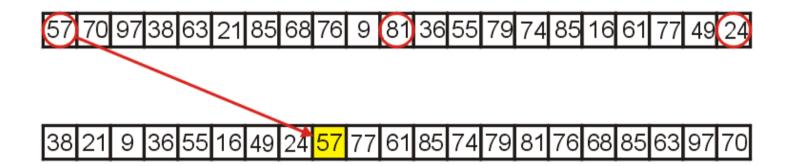
- copying elements either to the front or the back of the additional array
- placing the pivot into the resulting hole

For example, consider the following:

- 57 is the median-of-three
- we go through the remaining elements, assigning them either to the front or the back of the second array



Once we are finished, we copy the median-of-three, 57, into the resulting hole



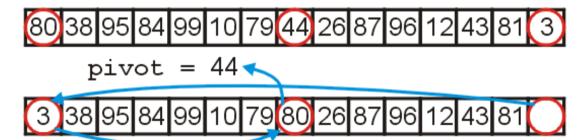
Can we implement quicksort in place?

Yes!

- Swap the pivot to the last slot of the list
- We repeatedly try to find two entries:
 - Staring from the front: an entry larger than the pivot
 - Starting from the back: an entry smaller than the pivot
- Such two entries are out of order, so we swap them
- Repeat until all the entries are in order
- Move the leftmost entry larger than the pivot into the last slot of the list and fill the hole with the pivot

First, we have already examined the first, middle, and last entries and chosen the median of these to be the pivot In addition, we can:

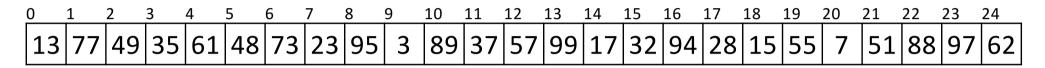
- move the smallest entry to the first entry
- move the largest entry to the middle entry



The implementation is straight-forward

```
template <typename Type>
void quicksort( Type *array, int first, int last ) {
    if ( last - first <= N ) {</pre>
        insertion sort( array, first, last );
    } else {
        Type pivot = find pivot( array, first, last );
        int low =
                       find next( pivot, array, first + 1 );
        int high = find previous( pivot, array, last - 2 );
        while ( low < high ) {</pre>
            std::swap( array[low], array[high] );
                    find next( pivot, array, low + 1 );
            low =
            high = find previous( pivot, array, high - 1 );
        }
        array[last - 1] = array[low];
        array[low] = pivot;
        quicksort( array, first, low );
        quicksort( array, high, last );
}
```

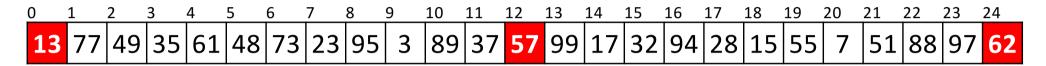
Consider the following unsorted array of 25 entries



We will call insertion sort if the list being sorted of size N=6 or less

We call quicksort(array, 0, 25)

We are calling quicksort (array, 0, 25)



```
First, 25 - 0 > 6, so find the midpoint and the pivot midpoint = (0 + 25)/2; // == 12
```

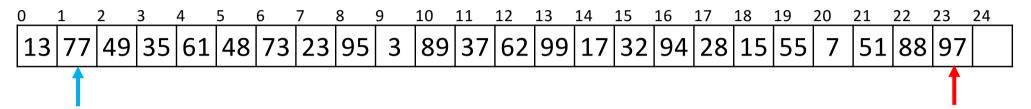
We are calling quicksort (array, 0, 25)

```
    0
    1
    2
    3
    4
    5
    6
    7
    8
    9
    10
    11
    12
    13
    14
    15
    16
    17
    18
    19
    20
    21
    22
    23
    24

    13
    77
    49
    35
    61
    48
    73
    23
    95
    3
    89
    37
    62
    99
    17
    32
    94
    28
    15
    55
    7
    51
    88
    97
```

```
First, 25-0>6, so find the midpoint and the pivot midpoint = (0 + 25)/2; // == 12 pivot = 57;
```

We are calling quicksort (array, 0, 25)



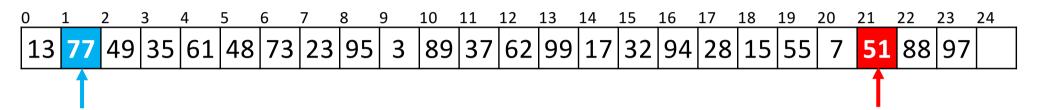
Starting from the front and back:

- Find the next element greater than the pivot
- The last element less than the pivot

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)



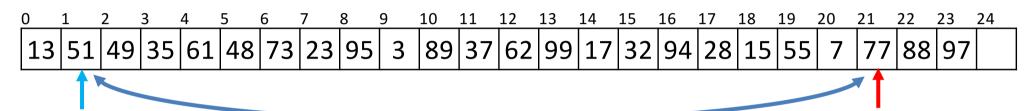
Searching forward and backward:

```
low = 1;
high = 21;
```

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)



Searching forward and backward:

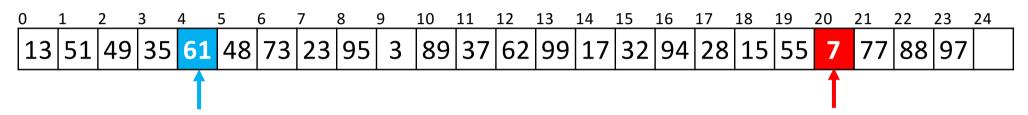
```
low = 1;
high = 21;
```

Swap them

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)

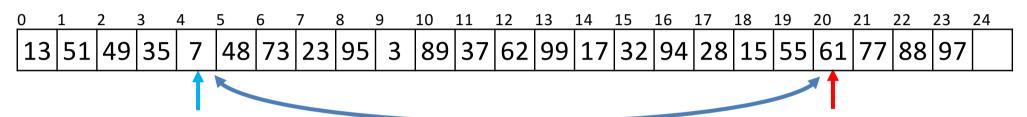


Continue searching

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)



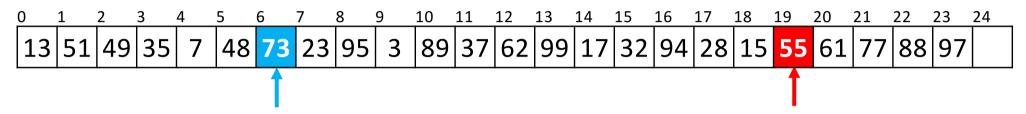
Continue searching

Swap them

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)

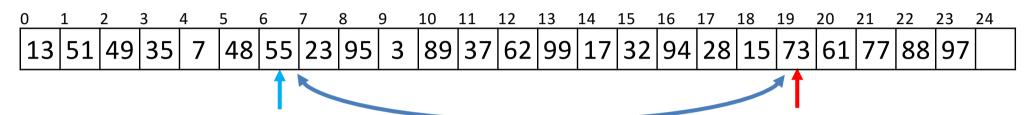


Continue searching

```
pivot = 57;
```

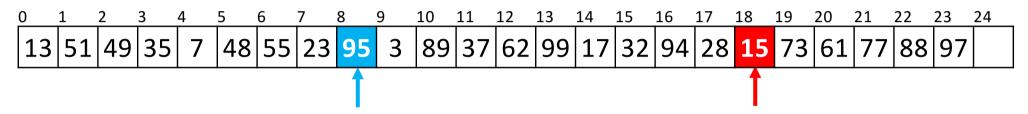
```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)



Continue searching

We are calling quicksort (array, 0, 25)

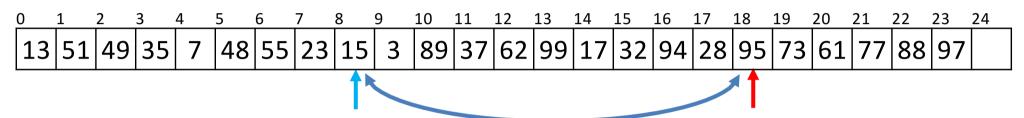


Continue searching

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)

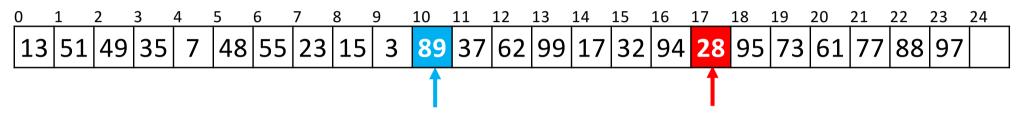


Continue searching

$$pivot = 57;$$

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)

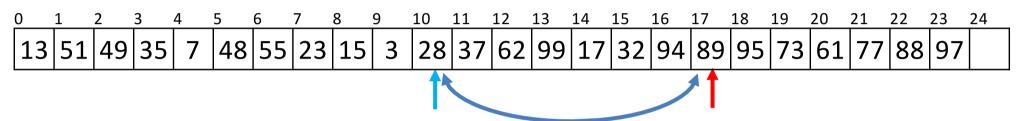


Continue searching

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)

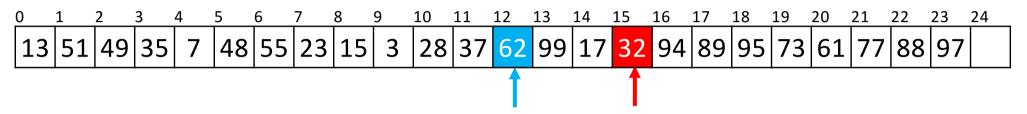


Continue searching

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)

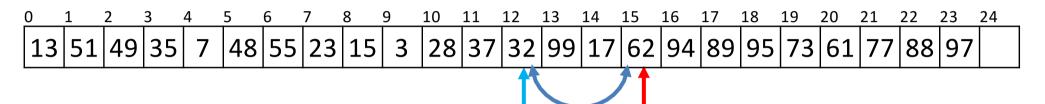


Continue searching

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)

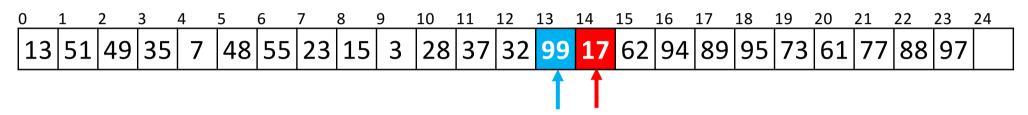


Continue searching

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)

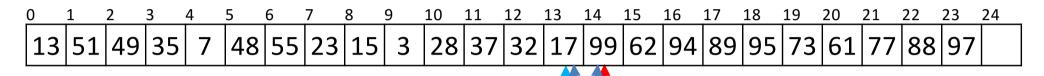


Continue searching

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort(array, 0, 25)

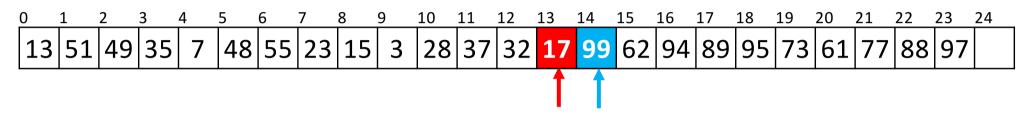


Continue searching

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)



Continue searching

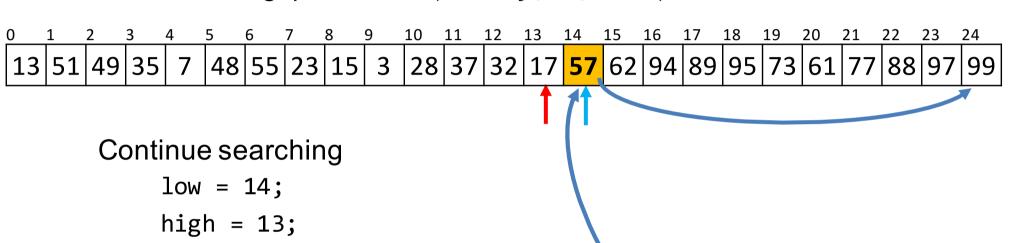
Now, low > high, so we stop

```
pivot = 57;
```

```
quicksort( array, 0, 25 )
```

We are calling quicksort (array, 0, 25)

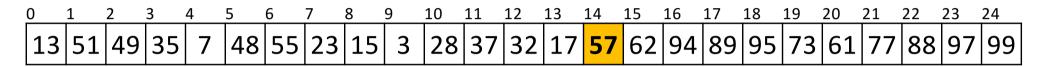
Now, low > high, so we stop



pivot = 57;

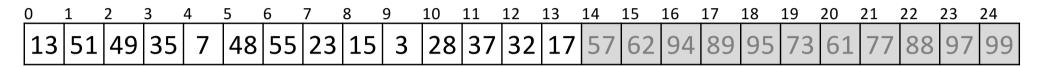
quicksort(array, 0, 25)

We are calling quicksort (array, 0, 25)



We now begin calling quicksort recursively on the first half quicksort(array, 0, 14);

We are executing quicksort(array, 0, 14)



```
First, 14-0 > 6, so find the midpoint and the pivot midpoint = (0 + 14)/2; // == 7
```

```
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 0, 14)

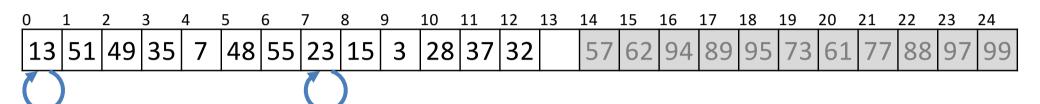
```
    0
    1
    2
    3
    4
    5
    6
    7
    8
    9
    10
    11
    12
    13
    14
    15
    16
    17
    18
    19
    20
    21
    22
    23
    24

    13
    51
    49
    35
    7
    48
    55
    23
    15
    3
    28
    37
    32
    17
    57
    62
    94
    89
    95
    73
    61
    77
    88
    97
    99
```

```
First, 14-0>6, so find the midpoint and the pivot midpoint = (0 + 14)/2; // == 7 pivot = 17
```

```
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

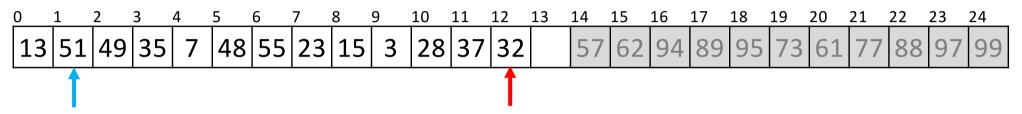
We are executing quicksort (array, 0, 14)



```
First, 14-0 > 6, so find the midpoint and the pivot midpoint = (0 + 14)/2; // == 7
```

```
pivot = 17;
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 0, 14)

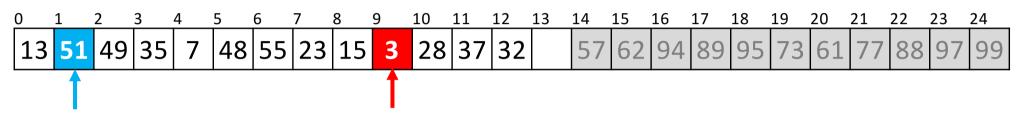


Starting from the front and back:

- Find the next element greater than the pivot
- The last element less than the pivot

```
pivot = 17;
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 0, 14)

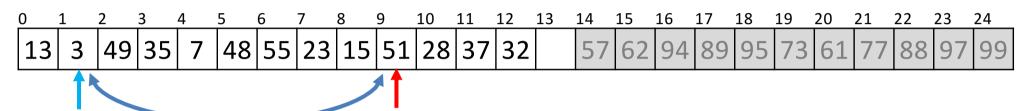


Searching forward and backward:

```
low = 1;
high = 9;
```

```
pivot = 17;
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 0, 14)

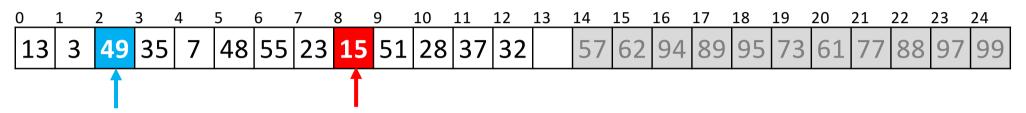


Searching forward and backward:

```
low = 1;
high = 9;
```

```
pivot = 17;
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

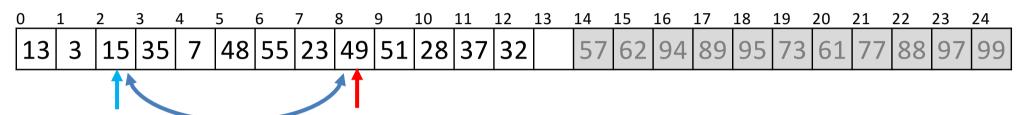
We are executing quicksort (array, 0, 14)



Searching forward and backward:

```
pivot = 17;
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 0, 14)

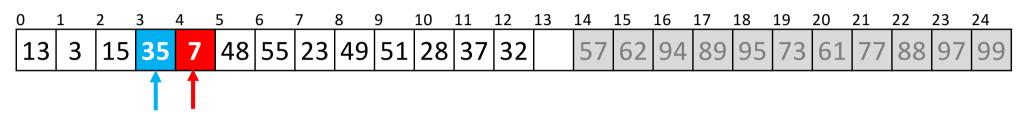


Searching forward and backward:

```
low = 2;
high = 8;
```

```
pivot = 17;
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

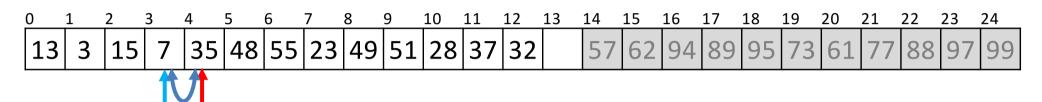
We are executing quicksort (array, 0, 14)



Searching forward and backward:

```
pivot = 17;
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 0, 14)

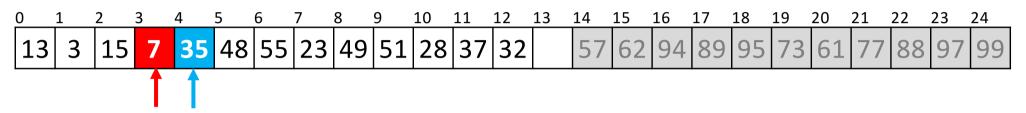


Searching forward and backward:

```
low = 3;
high = 4;
```

```
pivot = 17;
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 0, 14)



Searching forward and backward:

Now, low > high, so we stop

```
pivot = 17;
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 0, 14)

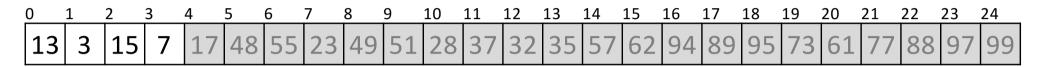
```
    0
    1
    2
    3
    4
    5
    6
    7
    8
    9
    10
    11
    12
    13
    14
    15
    16
    17
    18
    19
    20
    21
    22
    23
    24

    13
    3
    15
    7
    17
    48
    55
    23
    49
    51
    28
    37
    32
    35
    57
    62
    94
    89
    95
    73
    61
    77
    88
    97
    99
```

We continue calling quicksort recursively quicksort(array, 0, 4);

```
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 0, 4)



Now, $4-0 \le 6$, so find we call insertion sort

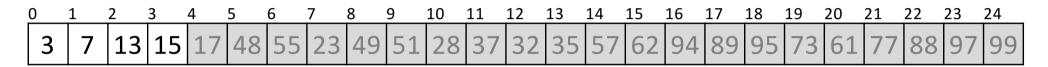
```
quicksort( array, 0, 4 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

Insertion sort just sorts the entries from 0 to 3

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
13	3	15	7	17	48	55	23	49	51	28	37	32	35	57	62	94	89	95	73	61	77	88	97	99

```
insertion_sort( array, 0, 4 )
quicksort( array, 0, 4 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

Insertion sort just sorts the entries from 0 to 3



This function call completes and so we exit

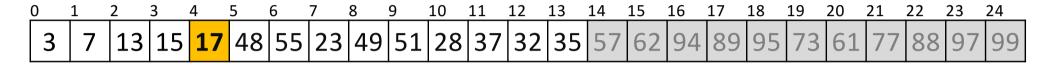
```
insertion_sort( array, 0, 4 )
quicksort( array, 0, 4 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

This call to quicksort is now also finished, so it, too, exits

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	48	55	23	49	51	28	37	32	35	57	62	94	89	95	73	61	77	88	97	99

```
quicksort( array, 0, 4 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are back to executing quicksort (array, 0, 14)



We continue calling quicksort recursively on the second half

```
quicksort( array, 0, 4 );
quicksort( array, 5, 14 );
```

```
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We now are calling quicksort(array, 5, 14)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	48	55	23	49	51	28	37	32	35	57	62	94	89	95	73	61	77	88	97	99

```
First, 14-5>6, so find the midpoint and the pivot midpoint = (5 + 14)/2; // == 9
```

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

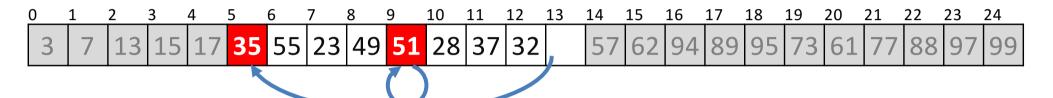
We now are calling quicksort(array, 5, 14)

```
    3
    7
    13
    15
    17
    48
    55
    23
    49
    51
    28
    37
    32
    35
    57
    62
    94
    89
    95
    73
    61
    77
    88
    97
    99
```

```
First, 14-5>6, so find the midpoint and the pivot midpoint = (5 + 14)/2; // == 9 pivot = 48
```

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

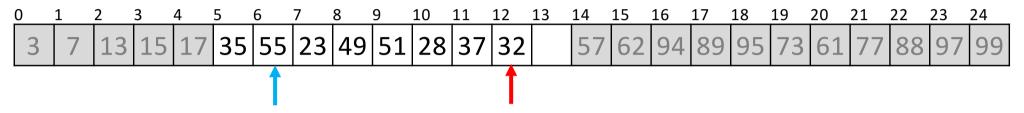
We now are calling quicksort(array, 5, 14)



```
First, 14-5>6, so find the midpoint and the pivot midpoint = (5 + 14)/2; // == 9 pivot = 48
```

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We now are calling quicksort(array, 5, 14)



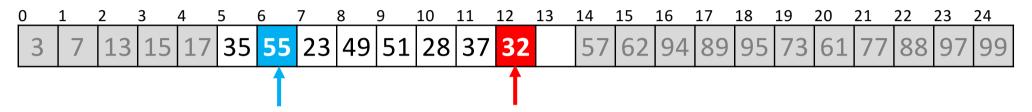
Starting from the front and back:

- Find the next element greater than the pivot
- The last element less than the pivot

```
pivot = 48;

quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We now are calling quicksort(array, 5, 14)

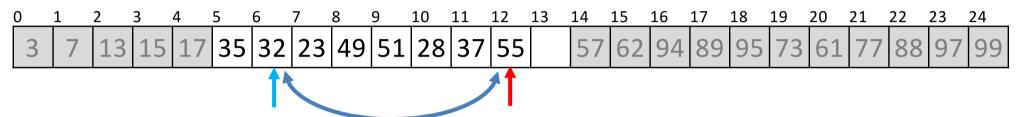


Searching forward and backward:

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

pivot = 48;

We now are calling quicksort(array, 5, 14)



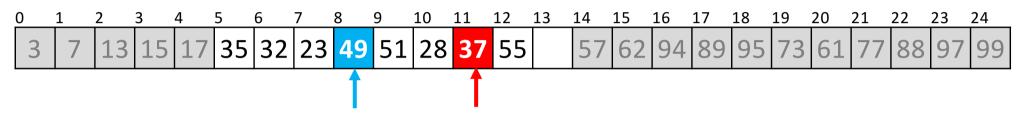
Searching forward and backward:

Swap them

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

pivot = 48;

We now are calling quicksort(array, 5, 14)

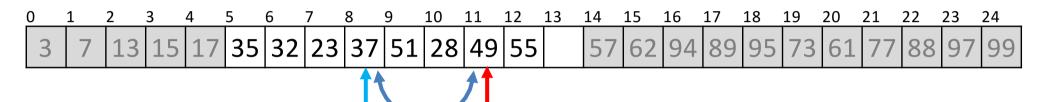


Continue searching

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

pivot = 48;

We now are calling quicksort(array, 5, 14)



Continue searching

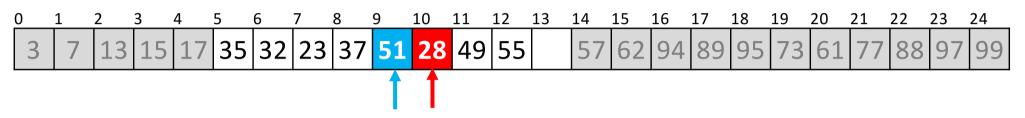
Swap them

```
pivot = 48;
```

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
```

quicksort(array, 0, 25)

We now are calling quicksort(array, 5, 14)

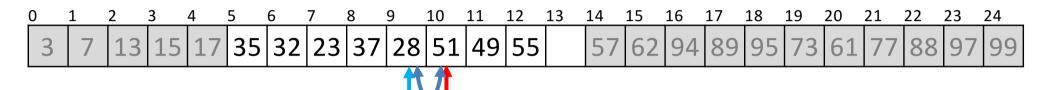


Continue searching

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

pivot = 48;

We now are calling quicksort(array, 5, 14)



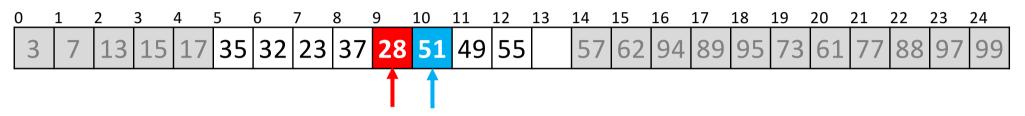
Continue searching

Swap them

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

pivot = 48;

We now are calling quicksort(array, 5, 14)



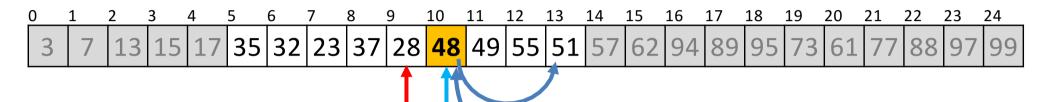
Continue searching

Now, low > high, so we stop

```
pivot = 48;
```

```
quicksort( array, 5, 14 )
```

We now are calling quicksort(array, 5, 14)



Continue searching

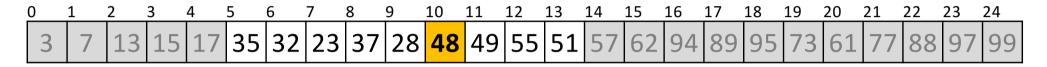
$$low = 8;$$

Now, low > high, so we stop

```
pivot = 48;
```

```
quicksort( array, 5, 14 )
```

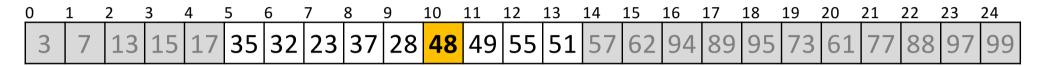
We now are calling quicksort(array, 5, 14)



We now begin calling quicksort recursively on the first half quicksort(array, 5, 10);

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We now are calling quicksort(array, 5, 14)



We now begin calling quicksort recursively quicksort(array, 5, 10);

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 5, 10)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	35	32	23	37	28	48	49	55	51	57	62	94	89	95	73	61	77	88	97	99

Now, $10-5 \le 6$, so find we call insertion sort

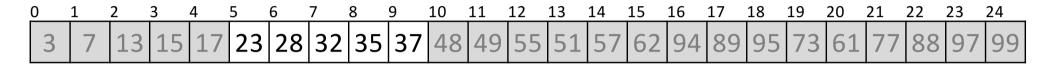
```
quicksort( array, 5, 10 )
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

Insertion sort just sorts the entries from 5 to 9

```
    3
    7
    13
    15
    17
    35
    32
    23
    37
    28
    48
    49
    55
    51
    57
    62
    94
    89
    95
    73
    61
    77
    88
    97
    99
```

```
insertion_sort( array, 5, 10 )
quicksort( array, 5, 10 )
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

Insertion sort just sorts the entries from 5 to 9



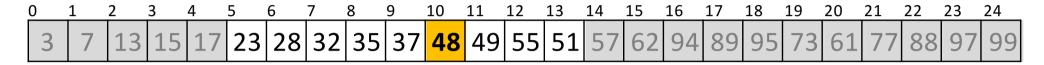
This function call completes and so we exit

```
insertion_sort( array, 5, 10 )
quicksort( array, 5, 10 )
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

0	1	L	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	3	7	13	15	17	23	28	32	35	37	48	49	55	51	57	62	94	89	95	73	61	77	88	97	99

```
quicksort( array, 5, 10 )
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are back to executing quicksort (array, 5, 14)

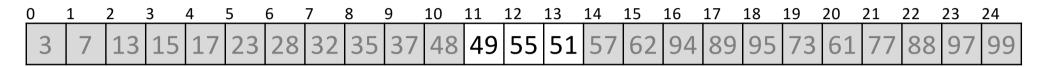


We continue calling quicksort recursively on the second half

```
quicksort( array, 5, 10 );
quicksort( array, 11, 14 );
```

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 11, 15)



Now, $15 - 11 \le 6$, so find we call insertion sort

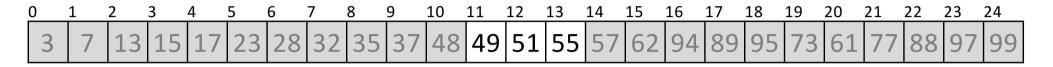
```
quicksort( array, 6, 14 )
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

Insertion sort just sorts the entries from 11 to 14

```
    3
    7
    13
    15
    17
    23
    28
    32
    35
    37
    48
    49
    55
    51
    57
    62
    94
    89
    95
    73
    61
    77
    88
    97
    99
```

```
insertion_sort( array, 11, 14 )
quicksort( array, 11, 14 )
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

Insertion sort just sorts the entries from 11 to 14



This function call completes and so we exit

```
insertion_sort( array, 11, 14 )
quicksort( array, 11, 14 )
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

0		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
	3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	62	94	89	95	73	61	77	88	97	99

```
quicksort( array, 11, 14 )
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

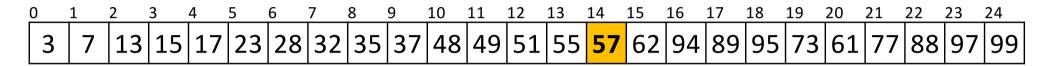
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	62	94	89	95	73	61	77	88	97	99

```
quicksort( array, 5, 14 )
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	62	94	89	95	73	61	77	88	97	99

```
quicksort( array, 0, 14 )
quicksort( array, 0, 25 )
```

We are back to executing quicksort (array, 0, 25)



We continue calling quicksort recursively on the second half

```
quicksort( array, 0, 14 );
quicksort( array, 15, 25 );
```

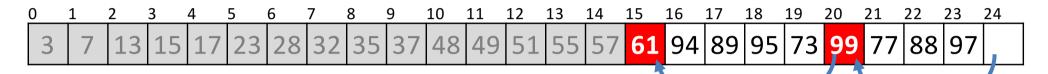
We are back to executing quicksort (array, 15, 25)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	62	94	89	95	73	61	77	88	97	99

```
First, 25-15>6, so find the midpoint and the pivot midpoint = (15 + 25)/2; // == 20
```

```
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

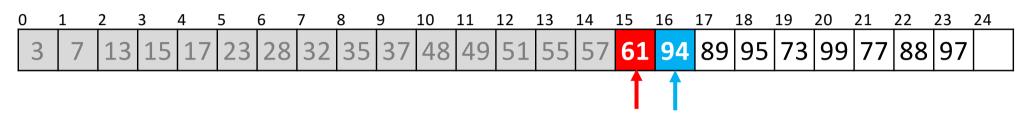
We are back to executing quicksort (array, 15, 25)



```
First, 25-15>6, so find the midpoint and the pivot midpoint = (15 + 25)/2; // == 20 pivot = 62;
```

```
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

We are back to executing quicksort (array, 15, 25)

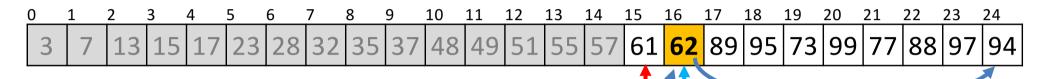


Searching forward and backward:

Now, low > high, so we stop

```
pivot = 62;
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

We are back to executing quicksort (array, 15, 25)



Searching forward and backward:

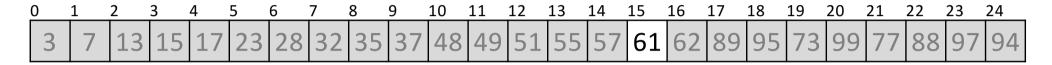
Now, low > high, so we stop

- Note, this is the worst-case scenario
- The pivot is the second smallest element

```
pivot = 62;
```

```
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

We are back to executing quicksort (array, 15, 25)



We continue calling quicksort recursively on the first half quicksort(array, 15, 16);

```
quicksort( array, 15, 16 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

We are executing quicksort (array, 15, 16)

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	89	95	73	99	77	88	97	94

Now, $16-15 \le 6$, so find we call insertion sort

```
quicksort( array, 15, 16 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

Insertion sort immediately returns

```
    0
    1
    2
    3
    4
    5
    6
    7
    8
    9
    10
    11
    12
    13
    14
    15
    16
    17
    18
    19
    20
    21
    22
    23
    24

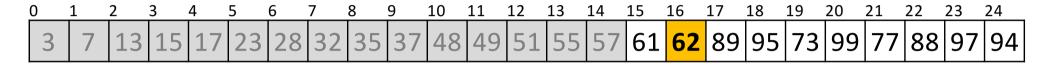
    3
    7
    13
    15
    17
    23
    28
    32
    35
    37
    48
    49
    51
    55
    57
    61
    62
    89
    95
    73
    99
    77
    88
    97
    94
```

```
insertion_sort( array, 15, 16 )
quicksort( array, 15, 16 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

0		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
	3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	89	95	73	99	77	88	97	94

```
quicksort( array, 15, 16 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

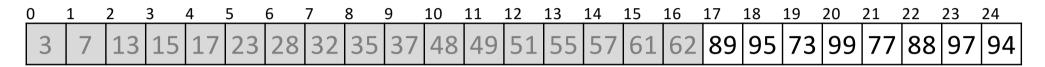
We are back to executing quicksort (array, 15, 25)



We continue calling quicksort recursively on the second half

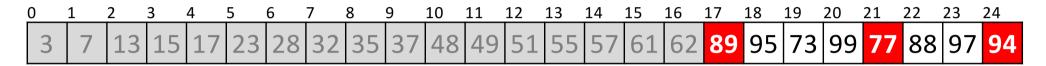
```
quicksort( array, 15, 16 );
quicksort( array, 17, 25 );
```

```
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```



```
First, 25-17>6, so find the midpoint and the pivot midpoint = (17 + 25)/2; // == 21
```

```
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```



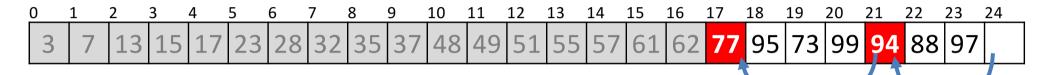
```
First, 25-17 > 6, so find the midpoint and the pivot midpoint = (17 + 25)/2; // == 21
```

```
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```



```
First, 25-17>6, so find the midpoint and the pivot midpoint = (17 + 25)/2; // == 21 pivot = 89
```

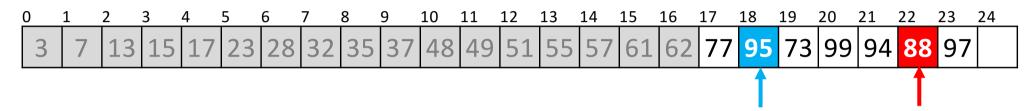
```
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```



```
First, 25-17>6, so find the midpoint and the pivot midpoint = (17 + 25)/2; // == 21 pivot = 89
```

```
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

We are now calling quicksort(array, 17, 25)

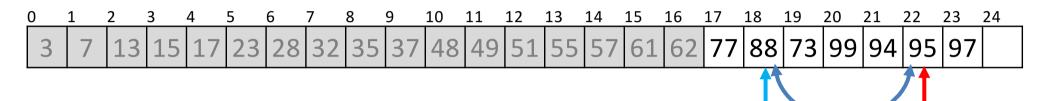


Searching forward and backward:

```
low = 18;
high = 22;
```

```
pivot = 89;
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

We are now calling quicksort(array, 17, 25)



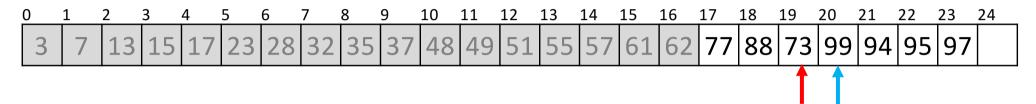
Searching forward and backward:

Swap them

```
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

pivot = 89;

We are now calling quicksort(array, 17, 25)



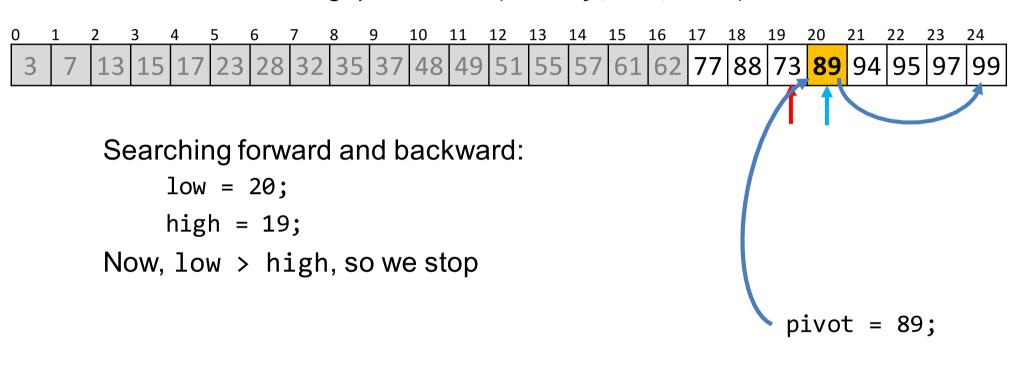
Searching forward and backward:

Now, low > high, so we stop

```
pivot = 89;
```

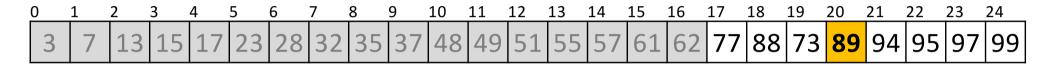
```
quicksort( array, 17, 25 )
```

We are now calling quicksort(array, 17, 25)



```
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

We are now calling quicksort(array, 17, 25)



We start by calling quicksort recursively on the first half quicksort(array, 17, 20);

```
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

We are now executing quicksort (array, 17, 20)

0	1	2	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	77	88	73	89	94	95	97	99

Now, $4-0 \le 6$, so find we call insertion sort

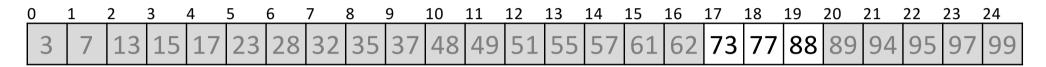
```
quicksort( array, 17, 20 )
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

Insertion sort just sorts the entries from 17 to 19

```
    3
    7
    13
    15
    17
    23
    28
    32
    35
    37
    48
    49
    51
    55
    57
    61
    62
    77
    88
    73
    89
    94
    95
    97
    99
```

```
insertion_sort( array, 17, 20 )
quicksort( array, 17, 20 )
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

Insertion sort just sorts the entries from 17 to 19



This function call completes and so we exit

```
insertion_sort( array, 17, 20 )
quicksort( array, 17, 20 )
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

0)	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
	3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	73	77	88	89	94	95	97	99

```
quicksort( array, 17, 20 )
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

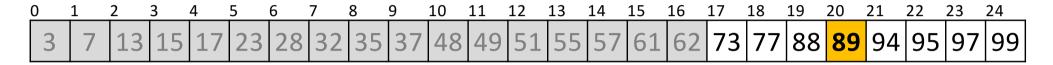
We are back to executing quicksort (array, 17, 25)

```
    0
    1
    2
    3
    4
    5
    6
    7
    8
    9
    10
    11
    12
    13
    14
    15
    16
    17
    18
    19
    20
    21
    22
    23
    24

    3
    7
    13
    15
    17
    23
    28
    32
    35
    37
    48
    49
    51
    55
    57
    61
    62
    73
    77
    88
    89
    94
    95
    97
    99
```

```
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

We are back to executing quicksort (array, 17, 25)

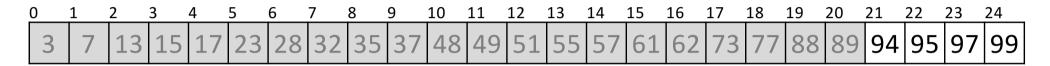


We continue by calling quicksort on the second half

```
quicksort( array, 17, 20 );
quicksort( array, 21, 25 );
```

```
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

We are now calling quicksort(array, 21, 25)



Now, $25 - 21 \le 6$, so find we call insertion sort

```
quicksort( array, 21, 25 )
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

Insertion sort just sorts the entries from 21 to 24

```
    3
    7
    13
    15
    17
    23
    28
    32
    35
    37
    48
    49
    51
    55
    57
    61
    62
    73
    77
    88
    89
    94
    95
    97
    99
```

```
insertion_sort( array, 21, 25 )
quicksort( array, 21, 25 )
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

Insertion sort just sorts the entries from 21 to 24

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	73	77	88	89	94	95	97	99

- In this case, the sub-array was already sorted
- This function call completes and so we exit

```
insertion_sort( array, 21, 25 )
quicksort( array, 21, 25 )
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	73	77	88	89	94	95	97	99

```
quicksort( array, 21, 25 )
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

0	1		2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3		7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	73	77	88	89	94	95	97	99

```
quicksort( array, 17, 25 )
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

0	1		2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3		7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	73	77	88	89	94	95	97	99

```
quicksort( array, 15, 25 )
quicksort( array, 0, 25 )
```

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	73	77	88	89	94	95	97	99

We have now used quicksort to sort this array of 25 entries

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
3	7	13	15	17	23	28	32	35	37	48	49	51	55	57	61	62	73	77	88	89	94	95	97	99

Black Board Example

Sort the following list using quicksort

Use insertion sort for any sub-list of size 4 or less

	_	_							9	
34	15	65	59	68	42	40	80	50	65	23

Memory Requirements

The additional memory?

- Function call stack
 - Each recursive function call places its local variables, parameters, etc., on a stack
- Average case: the depth of the recursion is $\Theta(\ln(n))$
- Worst case: the depth of the recursion is $\Theta(n)$

Run-time Summary

To summarize all three $\Theta(n \ln(n))$ algorithms

	Average Run Time	Worst-case Run Time	Average Memory	Worst-case Memory
Heap Sort	O(n	ln(n))	ϵ	$\Theta(1)$
Merge Sort	$\Theta(n)$	ln(n))	ϵ	$\Theta(n)$
Quicksort	$\Theta(n \ln(n))$	$\Theta(n^2)$	$\Theta(\ln(n))$	$\Theta(n)$

Further modifications

Our implementation is by no means optimal:

An excellent paper on quicksort was written by Jon L. Bentley and M. Douglas McIlroy:

Engineering a Sort Function

found in Software—Practice and Experience, Vol. 23(11), Nov 1993

Summary

This topic covered quicksort

- On average faster than heap sort or merge sort
- Uses a pivot to partition the objects
- Using the median of three pivots is a reasonably means of finding the pivot
- Average run time of $\Theta(n \ln(n))$ and $\Theta(\ln(n))$ memory
- Worst case run time of $\Theta(n^2)$ and $\Theta(n)$ memory