

1 (5') Stack, Queue and Complexity Analysis

Each question has one or more correct answer(s). Select all the correct answer(s). For each question, you get 0 point if you select one or more wrong answers, but you get 0.5 point if you select a non-empty subset of the correct answers.

Note that you should write you answers of section 1 in the table below.

Question 1	Question 2	Question 3	Question 4	Question 5
D	A D	ABCD	A	ABC

Question 1. *In the lectures of Week 2, suppose we implement a circular queue by using an array with the index range from 0 to $(n - 1)$, then what the size of this queue would be? We assume that the queue is non-empty.*

- (A) $rear - front + 1$
- (B) $(rear - front + 1) \% n$
- (C) $(rear - front + n) \% n$
- (D) $(rear - front + n) \% n + 1$

Question 2. *Which of the following is known to be correct?*

- (A) *Stack is a linear data structure and the operations on stacks are more restricted, the same is true for queue.*
- (B) *Lists store elements in sequential locations in memory.*
- (C) *Both stacks and queues allow us insert or delete an element at the front.*
- (D) *We can use two queues to implement stack.*

Question 3. *Which of the following is/are applications of queue and stack respectively?*

- (A) **Queue:** *A resource shared by multiple users/processes;* **Stack:** *Handling function calls*
- (B) **Queue:** *Loading Balancing;* **Stack:** *Reverse-Polish Notation*
- (C) **Queue:** *Handling of interrupts in real-time systems;* **Stack:** *Compilers/Word Processors*
- (D) **Queue:** *IO Buffers;* **Stack:** *Arithmetic expression evaluation*

Question 4. *Read the following code, what function does it realize?*

```
void Q4( Queue &Q)
{
    Stack S;
    int d;
    InitStack(S);
    while (!QueueEmpty(Q))
    {
```

```

        DeQueue(Q, d)
        Push(S, d);
    }
    while (!StackEmpty(S))
    {
        Pop(S, d);
        EnQueue(Q, d);
    }
}

```

- (A) Use stack to reverse the queue.
- (B) Use queue to reverse the stack.
- (C) Use stack to implement the queue.
- (D) Use queue to implement the stack.

Question 5. Which of the following comparison is correct?

- (A) $n^2 + n^3 = O(n^4)$
- (B) $\log_2 n = \Theta(\log n)$
- (C) $\log^2 n = \Omega(\log \log n)$
- (D) $n! = \omega(n^n)$

2 (10') Stack and Queue

Question 6. (2') The following post-fix expression (Reverse-Polish Notation) with single digit operands is evaluated using a stack and the final result:

$$8 \ 2 \ 3 \ ^ / \ 2 \ 3 \ * \ + \ 5 \ 1 \ *$$

Note that $^$ is the exponentiation operator. Please write down the corresponding in-fix notation:

$$\frac{8}{2^3} + 2 * 3 - 5 * 1$$

2

Question 7. (4') Describe how to implement a queue using a singly-linked list.

Check whether Queue is empty.

ENQUEUE: inserts an element at the end of the list. In this case we need to keep track of the last element of the list. We can do that with a sentinel.

DEQUEUE: removes an element from the beginning of the list.

Question 8. (1') If we use an array with size N to implement a normal queue, it gets full when $rear = N - 1$

Question 9. (1') By implementing the following operations on stack, the value of x is **a**
 $InitStack(st)$; $Push(st, a)$; $Push(st, b)$; $Pop(st, x)$; $Top(st, x)$;

Question 10. (2') When will "stack overflow" and "stack underflow" happen? (Give a short explanation.)

Overflow: Adding items to a full stack.

Underflow: Removing items from an empty stack.

3 (8') Complexity Analysis

Question 11. (3') Given a fraction of a code as the following, write down the time complexity for each **for** loop.

```

for ( i=1; i<n; i*=2) {
    for ( j=n; j>0; j/=2) {
        for ( k=j; k<n; k+=2) {
            sum += ( i + j*k )
        }
    }
}

```

----- $O(\log n)$ -----
----- $O(\log n)$ -----
----- $O(n)$ -----

Question 12. (5') Calculate the average processing time $T(n)$ of the following recursive algorithm. Suppose that it takes one unit time for **random(int n)** to return a random integer which is uniformly distributed in the range $[0, n]$. Also note that $T(0) = 0$.

*Hints: The equation $\frac{1}{1*2} + \frac{1}{2*3} + \dots + \frac{1}{n*(n+1)} = \frac{n}{n+1}$ might be needed.*

```

int hw( int n ) {
    if ( n <= 0 ) return 0;
    else {
        int i = random( n-1 );
        return hw( i ) + hw( n-1-i );
    }
}

```

The algorithm suggests that $T(n) = T(i) + T(n-1-i) + 1$. By summing this relationship for all the possible random values $i = 0, 1, \dots, n-1$, we obtain that in average

$nT(n) = 2(T(0) + T(1) + \dots + T(n-2) + T(n-1)) + n$. Because

$(n-1)T(n-1) = 2(T(0) + T(1) + \dots + T(n-2)) + (n-1)$, the basic recurrence is as follows:

$nT(n) - (n-1)T(n-1) = 2T(n-1) + 1$, or $nT(n) = (n+1)T(n-1) + 1$, or

$\frac{T(n)}{n+1} = \frac{T(n-1)}{n} + \frac{1}{n*(n+1)}$. The telescoping results in the following system of expressions:

$$\frac{T(n)}{n+1} = \frac{T(n-1)}{n} + \frac{1}{n*(n+1)}$$

$$\frac{T(1)}{2} = \frac{T(0)}{1} + \frac{1}{1*2}$$

Since $T(0) = 0$, the explicit expression for $T(n)$ is:

$$\frac{T(n)}{n+1} = \frac{1}{1*2} + \frac{1}{2*3} + \dots + \frac{1}{(n-1)*n} + \frac{1}{n*(n+1)}$$

So that $T(n) = n$