

# PHIL 3something - Logic II

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## Misc. Notation

- The set of positive integers  $\{x : x \text{ is a positive integer} \}$
- The set of positive integers less than 3  $\{x : x \text{ is a positive integer and } x \text{ is less than } 3 \} = \{1, 2\}$ .
- The empty set:  $\emptyset$  or  $\Delta$
- Member of:  $A \subseteq B$  iff  $\forall x(x \in A \implies x \in B)$
- Union of A and B:  $A \cup B$  iff  $\{x : x \in A \vee x \in B\}$
- Intersection of A and B:  $A \cap B$  iff  $\{x : x \in A \wedge x \in B\}$
- Difference of A and B:  $\{x : x \in A \wedge x \notin B\}$
- For any non-empty sets A, B: Cartesian product: A of B:  $A \times B: \{ \langle x, y \rangle : x \in A \wedge y \in B \}$  (ALL OF THE POSSIBILITIES)
- TOTAL FUNCTION: Every element in the domain is valid
- PARTIAL FUNCTION: Not every element in the domain is valid.
- for any set of sets A:
  - $\cup A = \{x : \exists y(y \in A \wedge x \in y)\}$
  - $\cap A = \{x : \forall y(y \in A \rightarrow x \in y)\}$
- Relations: R is
  - reflexive :  $\forall x Rxx$
  - symmetric :  $\forall x \forall y (Rxy \implies Ryx)$
  - transitive :  $\forall x \forall y \forall z ((Rxy \wedge Ryz) \implies Rxz)$
  - Euclidean :  $\forall x \forall y \forall z ((Rxy \wedge Rxz) \implies Ryz)$
  - a equivalence relation : it's symmetric, reflexive, transitive.
  - a equivalence relation (alt) : it's symmetric, and euclidean.
  - a (partial) function :  $\exists x$  and there is at most one y:  $Rxy$  : denoted  $f$
  - a (partial) function<sup>1</sup> :  $\exists x, \exists y | Rxy$  : denoted  $f$ .
  - a (total) function: assigns a value to each number of A : denoted  $f$
  - a (total) function<sup>2</sup>:  $\forall x, \exists y | Rxy$ : denoted  $f$ .
- Domain: The set of a functions arguments.
- Range: The set of its values. (Results)

- $f$  is a function from a set  $A$  iff the domain of  $f$  is included in  $A$
- $f$  is a function to a set  $B$  iff its range is included in  $B$ .
- $f^{-1}$  is the inverse of the function  $f$  from the set  $A$  to the set  $B$  iff: if for every member  $b \in B$ , there is exactly one member of  $a \in A$  such that  $f(a) = b$ , then  $f^{-1}(b) = a$ , otherwise  $f^{-1}(b)$  is undefined.
- $f$  is onto  $B$  iff  $B$  is the range of  $f$  (Surjective)  
Alt: (Wikipedia) :  $\forall y \in Y, \exists x \in X | y = f(x)$
- $f$  is one-to-one iff  $\forall x \forall y (f(x) = f(y) \implies x = y)$  (Injective)
- $f$  is a bijection iff  $f$  is onto and one-to-one.
- $f$  is a correspondence iff  $f$  is total, one-to-one and onto.
- Sets  $A$  and  $B$  are equinumerous iff there is a correspondence from  $A$  to  $B$ .

*Equinumerous is transitive.* Prove: if  $A$  is equinumerous with  $B$  and  $B$  is equinumerous with  $C$ , then  $A$  is equinumerous with  $C$ . Proof: Suppose  $A$  is equinumerous to  $B$ , and  $B$  is equinumerous to  $C$ . Then: There is a total, one-to-one function  $f$  from  $A$  onto  $B$ , and a total one-to-one function  $g$  from  $B$  to  $C$ . Prove equinumerous via  $h=g(f)$ , such that  $h(n)=g(f(n))$

- $h$  is total: Let  $a$  be a member of  $A$ .  $h(a) = g(f(a))$ . Since  $f$  is total there is a member of  $b$  of  $B$  such that  $f(a) = b$ . since  $g$  is total, there is a member of  $c \in C$  such that  $g(b) = c$ . Hence,  $h$  is total.
- $h$  is onto  $C$ . WLOG Let  $c$  be a member of  $C$ , as  $g$  is onto,  $\exists b \in B$  such that  $g(b) = c$ . As  $f$  is onto, then  $\exists a \in A$  such that  $f(a) = b$ . Hence, the composition of  $h = f(g)$  is onto  $C$ .
- $h$  is one-to-one: Suppose  $h$  is not one-to-one.  
Then there  $\exists a_1, a_2 \in A$  such that  $h(a_1) = h(a_2), a_1 \neq a_2$ .  
Giving  $g(f(a_1)) = g(f(a_2)), a_1 \neq a_2$   
Since  $g$  is one-to-one  $g(b_1) = g(b_2)$  iff  $b_1 = b_2$ .  
So the issue must lie in  $f$ . However  $f$  is one-to-one  $f(a_1) = f(a_2)$  iff  $f(a) = f(b)$ . Which is a contradiction, giving us that  $h$  is one-to-one.

□

$A^n$  : the  $n$ th Cartesian product of  $A$  with itself.

Suppose that the set of real Numbers  $r, r < 1$ , is enumerable. Then  $L_r : r_1, r_2, r_3, \dots$  written in a notation of  $0.n_1n_2n_3 \dots$  (being natural numbers)

The set of functions from the set of positive integers to positive integers is not enumerable.

*Proof.* Suppose  $S$  is enumerable.

Then there is a list  $L_s$  of the members of  $S$ .

$L_s = \{s_1, s_2, s_3, \dots\}$

Let  $\forall n \in \mathbb{N}, n \in k \iff n \notin S_n$

$k$  is a set of positive integers.

so There is a number  $j$  such that  $k = s_j$ . So  $j \in S_j \iff j \notin S_j$

Hence  $S$  is not enumerable.

□

The set of total recursive functions from the set of positive integers,  $F^1$ , is not enumerable.

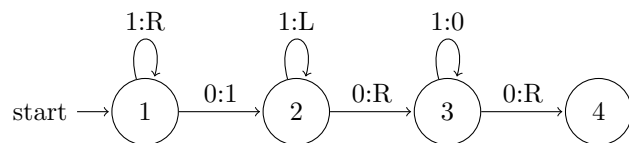
It's a Proof by contradiction.

Turing machines are in the following form:  $q_n, S_{1/0}, S_{1/0}/R/L, q_m$  where  $q_n$  is our current state, and you see  $S_{1/0}$ , perform function  $S_{1/0}/R/$  and move to state  $q_m$ . If there is no operation specified on the current state for a scan, then it halts. (Also Called the Turing Alphabet)

Example with notation:  $\text{start} \rightarrow \begin{array}{c} \circ \\ n \end{array} \quad \begin{array}{c} \circ \\ m \end{array}$

ex: (These are the same)

$Q_1S_1RQ_1, Q_1S_0S_1Q_2, Q_2S_1LQ_2, Q_2S_0RQ_3, Q_3S_1S_0Q_3, Q_3S_0RQ_4$



**Remark** (Turing Machines). • Each Turing machine is a finite set of Turing instructions.

- Each instruction is a 4 letter word of the Turing Alphabet.
- The set of Turing machine is enumerable. (Proof: exercise)

**Definition** (Standard initial configuration). A Turing machine is in a standard Initial configuration  $\iff$

- for some positive integer  $k$ , there are  $k$  blocks of 1's on the tape.
- separated by a blank,
- and the rest of the tape is blank.
- the machine is scanning the left-most 1 on the tape.
- the machine is in it's lowest numbered state.

ex:  $\dots 0010110111000\dots$  is a *SIC*. (if it's in lowest state) ex:  $\dots 00010000\dots$  is a *SIC*.

**Definition** (Standard final configuration). A Turing machine is in a standard final configuration  $\iff$

- there is a single block of  $k$  1's
- and the rest of the tape is blank.
- the machine is scanning the left-most 1 on the tape.

ex:  $\dots 00111111000\dots$  is a *SIC*. (if it's in lowest state) ex:  $\dots 00010000\dots$  is a *SIC*.

**Definition** (Computes a one-place function  $f^1$ ). A Turing  $M$  computes a one-place function  $f^1$ : if  $M$  is started in a *SIC* with a single block of  $k$  1's and

- if  $f^1$  is defined for the argument  $k$ , then  $M$  eventually halts in a *SFC*
- or if  $f^1$  is not defined for the argument  $k$ , then either  $M$  never halts or it halts in a non-standard final configuration.

**Remark.** Every Turing machine computes exactly one function of two arguments.

**Remark.** For any  $n$ , each Turing Machine computes exactly one function of  $n$  arguments.

The set of one-place Turing computable functions is enumerable

$\vdots$

The set of Turing computable functions is enumerable.

**Definition** (The halting problem). The problem of finding an effective method to determine whether a Turing machine will eventually halt or not after it is started with some input.

The halting problem is unsolvable. Ex:  $L_M : M_1, \dots$

$h(m, n) =$

1 if  $M_m$  eventually halts after starting with input  $n$

2 if  $M_m$  never halts after starting with input  $n$

The halting problem is solvable iff  $h$  is computable. Show:  $h$  is not Turing computable.

Let  $C$  be a copying machine.

Let  $F$  be  $\frac{1}{2}$  flipper.

Suppose  $h$  is Turing Computable.

Let  $H$  be a Turing machine that computes  $h$ .

If  $h$  is a Turing computable, then  $H$  exists.

If  $H$  exists, then  $D(C - H - F)$  exists.

Let  $D = M_k$ , for some  $k$ .  $M_k \in L_M$ .

Start  $D$  with input  $k$ . The  $C$ -part of  $D$  will produce a copy of  $k$ , Then the  $H$ -part will do its job:

- If  $M_k$  will eventually halt after starting with input  $k$ , then  $H$  will produce output 1.
- If  $M_k$  will never halt after starting with input  $k$ , then  $H$  will produce output 2.

Then the  $F$ -part will do its job.

- If output from  $H$  is 1,  $F$  will never halt.
- If output from  $H$  is 2,  $F$  will eventually halt.

Giving us:

- If  $M_k$  will eventually halt after starting with input  $k$ , then  $D$  will never halt after starting with input  $k$ .
- If  $M_k$  will never halt after starting with input  $k$ , then  $D$  will eventually halt after starting with input  $k$ .

So  $M_k$  will halt, after starting with input  $k$ ,  $\iff$   $D$  will never halt after starting with input  $k$ .

Then  $M_k$  isn't identical with  $D$ , which is a contradiction! Hence  $D$  doesn't exist. So  $H$  does not exist. So  $h$  is not Turing computable.  $\square$

Another halting problem..?  $L_M : M_1, \dots$   $L_F : F_1, \dots$

$g(n) = 1$ , if  $f_n(n) = 2$

$g(n) = 2$ , otherwise.

$g \neq f_k \forall k$

$h(m, n) = 1$  if  $M_m$  eventually halts after starting with input  $n$ .

$h(m, n) = 2$  if  $M_m$  never halts after starting with input  $n$ .

$s(m) = 1$ , if  $M_m$  eventually halts after starting with input  $m$ .

$s(m) = 2$ , if  $M_m$  never halts after starting with input  $m$ .

1. The halting problem is solvable iff  $h$  is computable.
2. If  $h$  is computable, then  $s$  is computable.
3. If  $s$  is computable, and  $TT$ 's is true, then  $g$  is computable.
4.  $g$  is not Turing computable.
5. Turing Thesis is true (Whatever is not Turing Computable is not computable)

6. The halting problem is not solvable.

3. Suppose  $S$  is computable and  $TT$  is true.

Then: There is a Turing machine  $S^*$  that computes  $s$ .

Suppose that we are to calculate  $g(n)$ , for some  $n$ .

Start  $S^*$  with input  $n$ .

- Case 1:  $S^*$  eventually halts with output 1

We know that  $M_n$  will eventually halt after it is started with input  $n$ . Start  $M_n$  with input  $n$ , when it halts, inspect the tape.

– Case 1.1: Halted in SFC  $f_n(n) = 2$   $g(n) = 1$

– Case 1.2: Halted in non-SFC:  $f_n$  is undefined.  $f_n(n) \neq 2$

And then Blake broke it:

As it's a halting problem to figure out if it's in  $SFC$ ?

$g(n) = 2$

- Case 2:  $S^*$  eventually halts with output 2 We know that  $M_n$  will never halt after it is started with input  $n$ .

So we know that  $f_n$  is undefined for the argument  $n$ .

So we know that  $g(n) = 2$

□

□