PHIL 3something - Logic II

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Misc. Notation

- The set of positive integers $\{x : x \text{ is a positive integer }\}$
- The set of positive integers less than3 $\{x : x$ is a positive integer and x is less than 3 $\}. = \{1, 2\}$.
- The empty set: $\emptyset$ or Δ
- Member of: $A \subseteq B$ iff $\forall X (x \in A \implies x \in B)$
- Union of A and B: $A \cup B$ iff $\{x : x \in A \vee x \in B\}$
- Intersection of A and B: $A \cap B$ iff $\{x : x \in A \wedge x \in B\}$
- Difference of A and B: $\{x : x \in A \wedge x \notin B\}$
- For any non-empty sets A, B: Cartesian product: A of B: $A \times B$: $\{< x, y >: x \in A \wedge y \in B\}$ (ALL OF THE POSSIBILITIES)
- TOTAL FUNCTION: Every element in the domain is valid
- PARTIAL FUNCTION: Not every element in the domain is valid.
- for any set of sets A:
  - $\Cup A = \{x : \exists y (y \in A \wedge x \in y)\}$
  - $\Cap A = \{x : \forall y (y \in A \rightarrow x \in y)\}$
- Relations: R is
  - reflexive : $\forall x Rxx$
  - symmetric : $\forall x \forall y (Rxy \implies Ryx)$
  - transitive : $\forall x \forall y \forall z ((Rxy \wedge Ryz) \implies Rxz)$
  - Euclidean : $\forall x \forall y \forall z ((Rxy \wedge Rxz) \implies Ryz)$
  - a equivalence relation : it's symmetric,reflexive,transitive.
  - a (partial) function : $\exists x$ and there is at most one y: Rxy : denoted f
  - a (partial) function[1] : $\exists x, \exists y | Rxy$: denoted f.
  - a (total) function: assigns a value to each number of A : denoted f
  - a (total) function[2]: $\forall x, \exists y | Rxy$: denoted f.
- Domain: The set of a functions arguments.
- Range: The set of its values. (Results)
- f is a function from a set A iff the domain of f is included in A

- f is a function to a set B iff its range is included in B.
- f^{-1} is the inverse of the function f from the set A to the set B iff:if for every member $b \in B$, there is exactly one member of $a \in A$ such that $f(a) = b$, then $f^{-1}(b) = a$, otherwise $f^{-1}(b)$ is undefined.
- f is onto B iff B is the range of f (Surjective)
  Alt: (Wikipedia) : $\forall y \in Y, \exists x \in X | y = f(x)$
- f is one-to-one iff $\forall x \forall y (f(x) = f(y) \implies x = y)$ (Injective)
- f is a bijection iff f is onto and one-to-one.
- f is a correspondence iff f is total, one-to-one and onto.
- Sets A and B are equinumerous iff there is a correspondence from A to B.

Equinumerous is transitive. Prove: if A is equinumerous with B and B is equinumerous wit C, then A is equinumerous with C. Proof: Suppose A is equinumerous to B, and B is equinumerous to C. Then: There is a total,one-to-one function f from A onto B, and a total one-to-one function g from B to C. Prove equinumerous via h=g(f), such that h(n)=g(f(n))

- h is total: Let a be a member of A. $h(a) = g(f(a))$. Since f is total there is a member of b of B such that $f(a) = b$). since g is total, there is a member of $c \in C$ such that $g(b) = c$. Hence, h is total.
- h is onto C. WLOG Let c be a member of C, as g is onto, $\exists b \in B$ such that $g(b) = c$. As f is onto, then $\exists a \in A$ such that $f(a) = b$. Hence, the composition of $h = f(g)$ is onto C.
- h is one-to-one: Suppose h is not one-to-one.
  Then there $\exists a_1, a_2 \in A$ such that $h(a_1) = h(a_2), a_1 \neq a_2$.
  Giving $g(f(a_1)) = g(f(a_2)), a_1/not = a_2$
  Since g is one-to-one $g(b_1) = g(b_2)$ iff $b_1 = b_2$.
  So the issue must lie in f. However f is one-to-one $f(a_1) = f(a_2)$ iff $f(a) = f(b)$. Which is a contradiction, giving us that h is one-to-one.

□

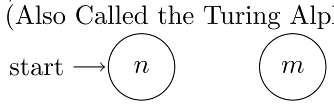
A^n : the nth Cartesian product of A with itself.
Suppose that the set of real Numbers $r, r < r < 1$, is enumerable. Then $L_r : r_1, r_2, r_3 ....$ written in a notation of $0.n_1n_2n_3.(nbeingnaturalnumbers)$
The set of functions form the set of positive integers to positive integers is not enumerable.

Proof. Suppose S is enumerable.
Then there is a list L_s of the members of S.
$L_s = \{s_1, s_2, s_3, \cdots\}$
Let $\forall n \in \mathbb{N}, n \in k \iff n \notin S_n$
k is a set of positive integers.
so There is a number $j$such that $k = s_j$. So $j \in S_j \iff j \not inS_j$
Hence S is not enumerable.

□

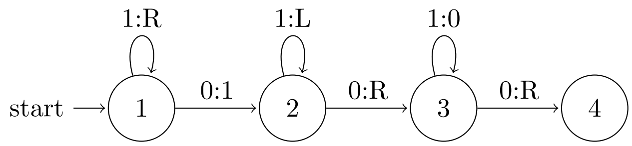
The set of total nomadic functions from the set of positive integers, F^1, is not enumerable.
It's a Proof by contradiction.
Turing machines are in the following form: $q_n, S_{1/0}, S_{1/0}/R/L, q_m$ where q_n is our current state, and you see $S_{1/0}$, perform function $S_{1/0}/R/$ and move to state q_m. If there is no operation specified on the current state for a scan, then it halts. (Also Called the Turing Alphabet)

Example with notation:



ex: (These are the same)

$$Q_1S_1RQ_1, Q_1S_0S_1Q_2, Q_2S_1LQ_2, Q_2S_0RQ_3, Q_3S_1S_0Q_3, Q_3S_0RQ_4$$



Remark (Turing Machines)**.**

- Each Turing machine is a finite set of Turing instructions.
- Each instruction is a 4 letter word of the Turing Alphabet.
- The set of Turing machine is enumerable. (Proof: exercise)

Definition (Standard inital configuration)**.** A Turing machine is in a standard Initial configuration $\iff$

- for some positive integer k, there are k blocks of 1's on the tape.
- separated by a blank,
- and the rest of the tape is blank.
- the machine is scanning the left-most 1 on the tape.
- the machine is in it's lowest numbered state.

ex: $\cdots 0010110111000 \cdots$ is a SIC. (if it's in lowest state) ex: $\cdots 00010000 \cdots$ is a SIC.

Definition (Standard final configuration)**.** A Turing machine is in a standard final configuration $\iff$

- there is a single block of k 1's
- and the rest of the tape is blank.
- the machine is scanning the left-most 1 on the tape.

ex: $\cdots 00111111000 \cdots$ is a SIC. (if it's in lowest state) ex: $\cdots 00010000 \cdots$ is a SIC.

Definition (Computes a one-place function f^1)**.** A Turing M computes a one-place function f^1:
if M is started in a SIC with a single block of k 1's and

- if f^1 is defined for the argument k, then M eventually halts in a SFC
- or if f^1 is not defined for the argument k, then either M never halts or it halts in a non-standard final configuration.

Remark. Every Turing machine computes exactly one function of two arguments.

Remark. For any n, each Turing Machine computes exactly one function of n arguments.
The set of one-place Turing computable functions is enumerable

$\vdots$

The set of Turing computable functions is enumerable.

Definition (The halting problem)**.** The problem of finding an effective method to determine whether a Turing macine will eventually ahlt or not after it is started with some input.

The halting problem is unsolvable. Ex: $L_M : M_1, \ldots$
$h(m, n) =$
1 if M_m eventually ahlts after starting with input n
2 if M_m never halts after starting with input n
The halting problem is solvable iff h is computable. Show: h is not Turing computable.
Let C be a copying machine.
Let F be $\frac{1}{2}$ flipper.
Suppose h is Turing Computable.
Let H be a Turing machine that computes h.
If h is a Turing computable, then H exists.
If H exists, then $D(C - H - F)$ exists.
Let $D = M_k$, for some k. $M_k \in L_M$.

Start D with input k. The C-part of D will produce a copy of k, Then the $H-$part will do its job:

- If M_k will eventually halt after starting with input k, then H will produce output 1.
- If M_k will never halt after starting with input k, then H will produce output 2.

Then the $F-$part will do it's job.

- If output from H is 1, F will never halt.
- If output from H is 2, F will eventually halt.

Giving us:

- If M_k will eventually halt after starting with input k, then D will never halt ater starting with input k.
- If M_k will never halt ater starting with input k, then D will eventually halt after starting with input k.

So M_k will halt, after starting with input k, $\iff$ D will nver halt after starting with input k.
Then M_k isn't identical with D, which is a contradiction! Hence$D$ doesn't exist. So H does not exists. So h is not turing computable. □

Another halting problem..? $L_M : M_1, \cdots$ $L_F : F_1, \cdots$
$g(n) = 1$, if $f_n(n) = 2$
$g(n) = 2$, otherwise.
$g \neq f_k \forall k$
$h(m, n) = 1$if M_m eventuall halys after starting with input n.
$h(m, n) = 2$ if M_m never halts after starting with input n.
$s(m) = 1$, if M_m eventually halts after starting with input m.
$s(m) = 2$, if M_m never halts after starting with input m.

1. The halting problem is solvable iff h is computable.
2. If h is computable, then s is computable.
3. If s is computable, and TT's is true, then g is computable.
4. g is not Turing computable.
5. Turing Thesis is true (Whatever is not Turing Computable is not computable)

6. The halting problem is not solvable.

3. Suppose S is computable and TT is true.
Then: There is a Turing machine $S*$ that computes s.
Suppose that we are to calculate $g(n)$, for some n.
Start $S*$ with input n.

- Case 1: S* eventually halts with output 1
  We know that M_n will eventually ahlt after it is started with input n Start M_n with input n, when it halts, inspect the tape.
  - Case 1.1: Halted in SFC $f_n(n) = 2$ $g(n) = 1$
  - Case 1.2: Halted in non-SFC: f_n is undefined. $f_n(n) \neq 2$

  $g(n) = 2$
- Case 2: S* eventually halts with output 2 We know that M_n will never halt after it is started with input n.
  So we know that f_n is undefined for the argument n.
  So we know that $g(n) = 2$

$\square$

$\square$