



Vietnam National University of HCMC  
International University  
School of Computer Science and Engineering



# Data Structures and Algorithms

## ★ Simple Sorting ★

Dr Vi Chi Thanh - [vcthanh@hcmiu.edu.vn](mailto:vcthanh@hcmiu.edu.vn)

<https://vichithanh.github.io>



SCAN ME

# Objectives

- Understand and know how to use basic sorting methods.
  - Bubble Sort
  - Selection Sort
  - Insertion Sort
- Compare their performance

# Major Topics

- Introductory Remarks
- Bubble Sort
- Selection Sort
- Insertion Sort
- Sorting Objects
- Comparing Sorts

# Introduction

- Why do we need to sort data?
  - To get the lowest price
  - To get the most crowded country
  - etc.
- So many lists are better dealt if ordered.
- Sorting data may be a preliminary step to searching

# Introduction

- Sorting is time-consuming task
  - ➔ many sorting algorithms are developed
- Will look at simple sorting first.
  - Note: there are books written on many advanced sorting techniques. E.g. shell sort; quicksort; heapsort; etc.
- Will start with 'simple' sorts.
  - Relatively slow,
  - Easy to understand, and
  - Excellent performance under circumstances.
- All these are  $O(N^2)$  sorts.

# How would you do it?



ARRANGE PLAYER IN ORDER OF  
INCREASING HEIGHT

# Basic idea of these simple sorts

- Compare two items
- Swap or Copy over
- Depending on the specific algorithm...
- Don't need additional space

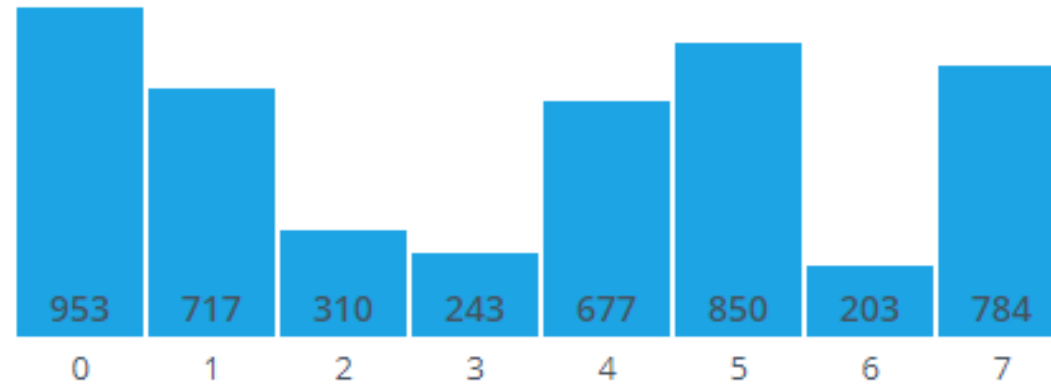
# Bubble sort



# Bubble Sort

- Very slow but simple
- Basic idea:
  - Compare the first item in the first two positions (e.g., the left most)
  - If the first one is larger, swap with the second
  - Move one position right.
- At end, the largest item is in the last position.
  - (e.g., the right most)
- That why is name Bubble sort

# Simulation



Steps:  
Starting Bubble Sort

<https://www.hackerearth.com/practice/algorithms/sorting/bubble-sort/visualize/>

<https://visualgo.net/en/sorting>

# Bubble Sort process

- After first pass, we made
  - $n - 1$  comparisons
  - 0 to  $n - 1$  swaps (depend on data)
- Continue this process.
- Next time we do not check the last entry ( $N - 1$ ), because we know it is in the right spot. We stop comparing at ( $N - 2$ ).
- See the simulation:  
<https://www.hackerearth.com/practice/algorithms/sorting/bubble-sort/visualize/>
- See Some code (p85-86).

# Bubble Sort process

```
for(out=nElems-1; out>1; out--) // outer loop (backward)
    for(in=0; in<out; in++)      // inner loop (forward)
        if( a[in] > a[in+1] )    // out of order?
            swap(in, in+1);      // swap them
```

# Hand-on

+ Try with some examples

<b>89</b>	<b>58</b>	<b>29</b>	<b>40</b>	<b>12</b>	<b>42</b>	<b>10</b>	<b>1</b>
<b>1</b>	<b>32</b>	<b>12</b>	<b>53</b>	<b>11</b>	<b>76</b>	<b>23</b>	<b>89</b>

# Efficiency of the Bubble Sort - **Comparisons**

- Can readily see that there are fewer comparisons each 'pass.'
- Thus, number of comparisons is computed as:  
 $(n-1) + (n-2) + \dots + 1 = n(n-1)/2;$
- For 10 elements, the number is  $10 \cdot 9 / 2 = 45$ .
- So, the algorithm makes about  $n^2/2$  comparisons
- (ignoring the -1 which is negligible especially if N is large)

# Efficiency of the Bubble Sort - **Swaps**

- Fewer SWAPS than COMPARISONS
  - since every comparison does not result in a swap.
- In general, a swap will occur half the time.
  - For  $\mathbf{n^2/2}$  comparisons,  
we have  $\mathbf{n^2/4}$  swaps.
- Worst case, every compare results in a swap (which case?)

# Overall – Bubble Sort

- Both swaps and compares are proportional to  $n^2$ .
  - Ignore the 2 and 4
- ➔ Complexity of Bubble Sort is  $O(n^2)$
- ➔ Rather slow.
- Hint to determine Big-O:
  - 2 nested-loop ➔  $O(n^2)$
  - Outer loop executes  $n$  times and inner loop executes in  $n$  times PER execution of the outer  $n$  times: hence  $n^2$



# Question

- The bubble sort algorithm alternates between:
  - A) Comparing and swapping
  - B) Moving and copying
  - C) Moving and comparing
  - D) Copying and comparing
- What can be improved in Bubble Sort algorithm?

# Selection sort

# Selection Sort

- Fewer swaps, same comparisons.

- Swaps  $O(n^2) \rightarrow O(n)$
- Comparisons  $O(n^2)$

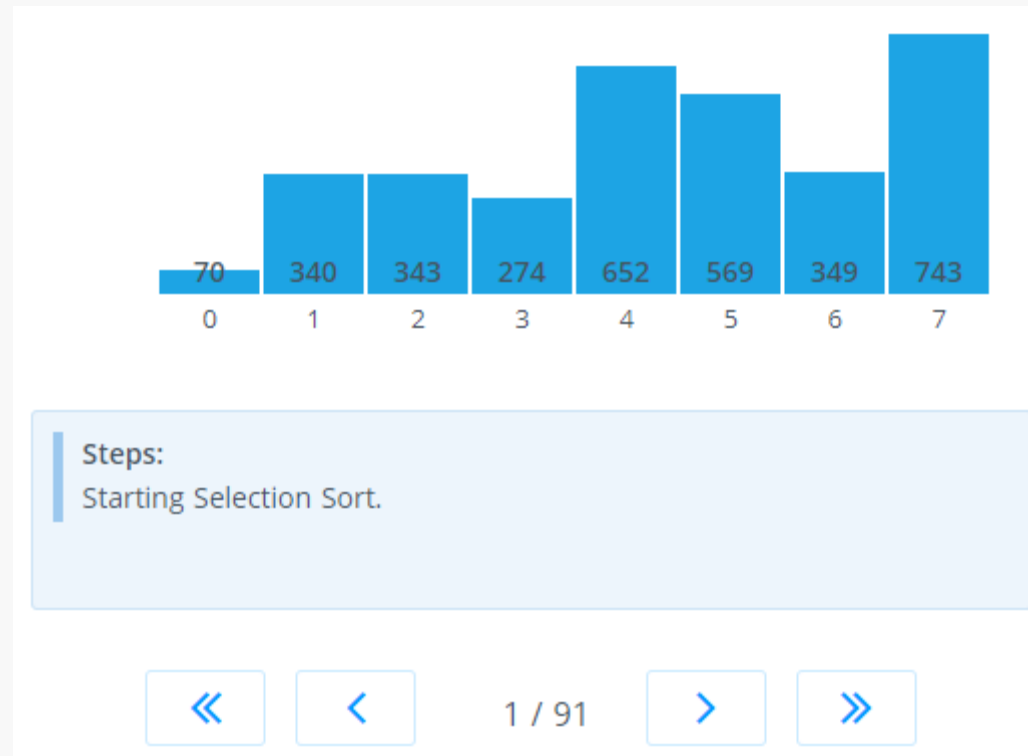
➔ Important while dealing with large records

➔ Reduction in swap time is more important than one in comparison time

# How does the Selection Sort work?

- Start from the first element (e.g., the left end)
- Scan all elements to selecting the smallest (largest) item.
- Swap with the first element
- Next pass, move one position right
- Repeat until all are sorted.

# Animation



<https://www.hackerearth.com/practice/algorithms/sorting/selection-sort/visualize/>

<https://visualgo.net/en/sorting>

# Selection Sort – in more detail

- So, in one pass, you have made  $n$  comparisons but possibly ONLY ONE Swap!
- With each succeeding pass,
  - one more item is sorted and in place;
  - one fewer item needs to be considered.
- Java code for the Selection Sort (p93-94).

# Selection Sort

```
for(out=0; out<nElems-1; out++)    // outer loop
{
    min = out;                      // minimum
    for(in=out+1; in<nElems; in++) // inner loop
        if(a[in] < a[min] )        // if min greater,
            min = in;              // we have a new min
    swap(out, min);                 // swap them
} // end for(out)
} // end selectionSort()
```

# Hand-on

+ Try with some examples

<b>89</b>	<b>58</b>	<b>29</b>	<b>40</b>	<b>12</b>	<b>42</b>	<b>10</b>	<b>1</b>
<b>1</b>	<b>32</b>	<b>12</b>	<b>53</b>	<b>11</b>	<b>76</b>	<b>23</b>	<b>89</b>



# Selection Sort itself - more

- + Algorithm implies it is an  $O(n^2)$  sort (and it is).

  - + How did we see this?

- + In comparison with Bubble Sort

  - + Same number of comparisons  **$n^2/2$**

  - + Fewer swap  **$n$**

- Faster than Bubble Sort

# Insertion sort

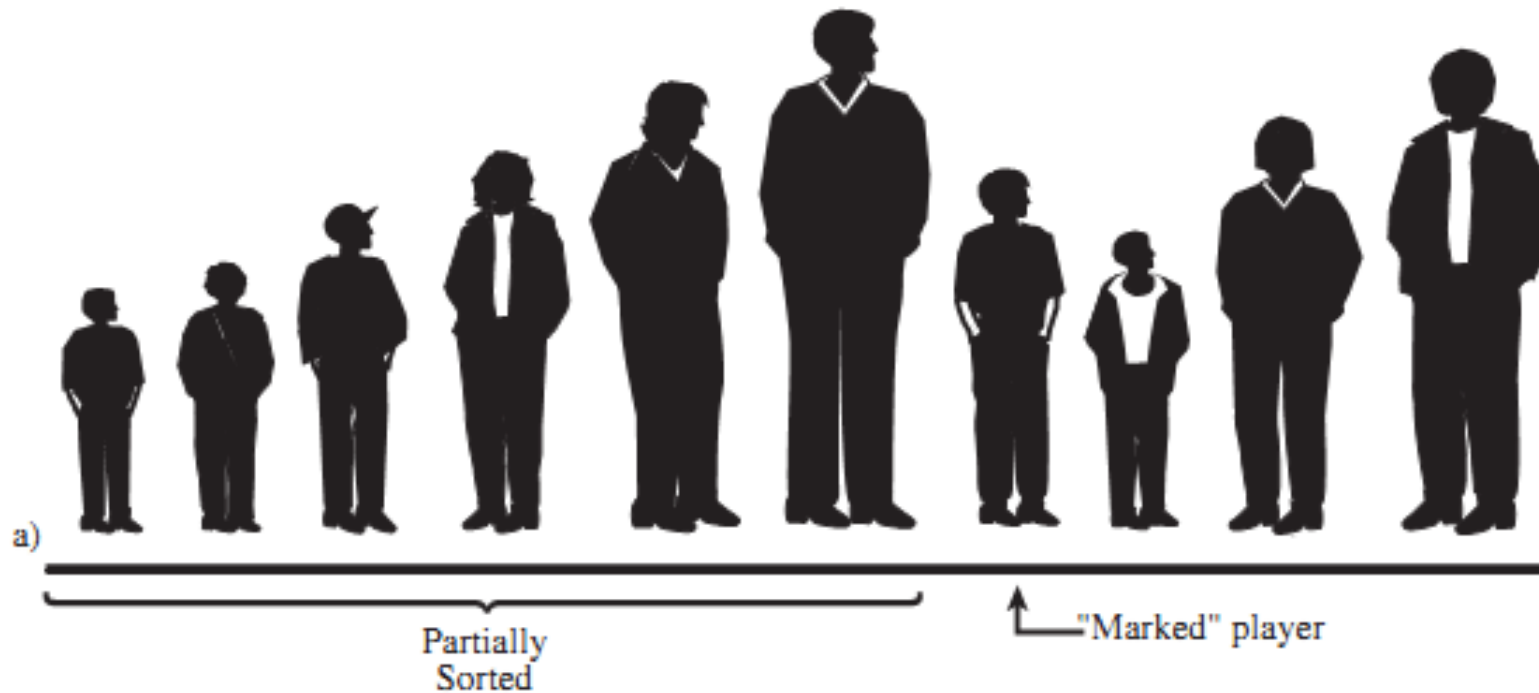
# Insertion Sort

- In many cases, this sort is considered the best of these elementary sorts.
- Still an  $O(n^2)$  but:
  - about twice as fast as bubble sort and
  - somewhat faster than selection sort in most situations.
- Easy, but a bit more complex than the others
- Sometimes used as **final stage** of some more sophisticated sorts, such as a **QuickSort** (coming).

# Insertion Sort – The idea

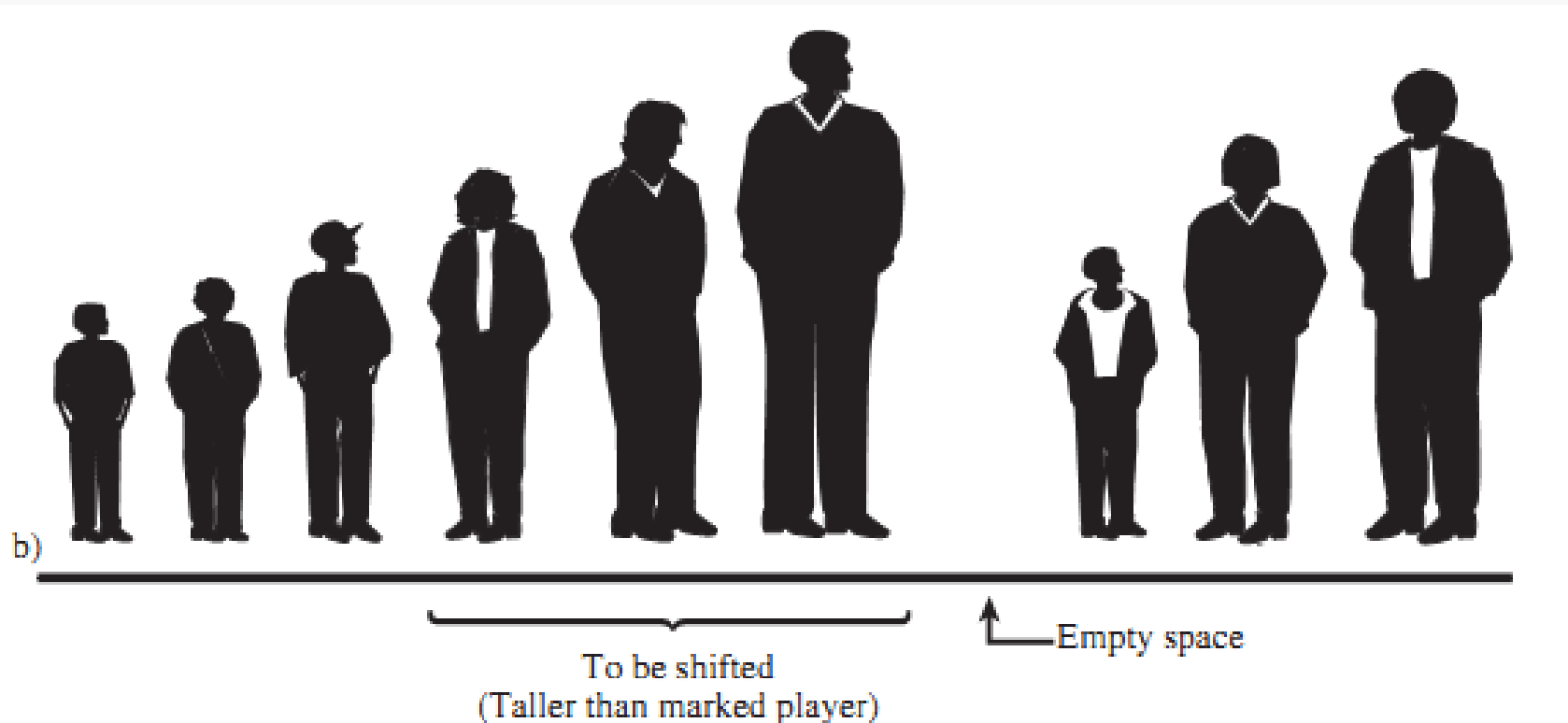
28

+ Thinking that 'half' of the list of items to be sorted. (*Partially sorted*, *Marked*, *Unsorted* items)

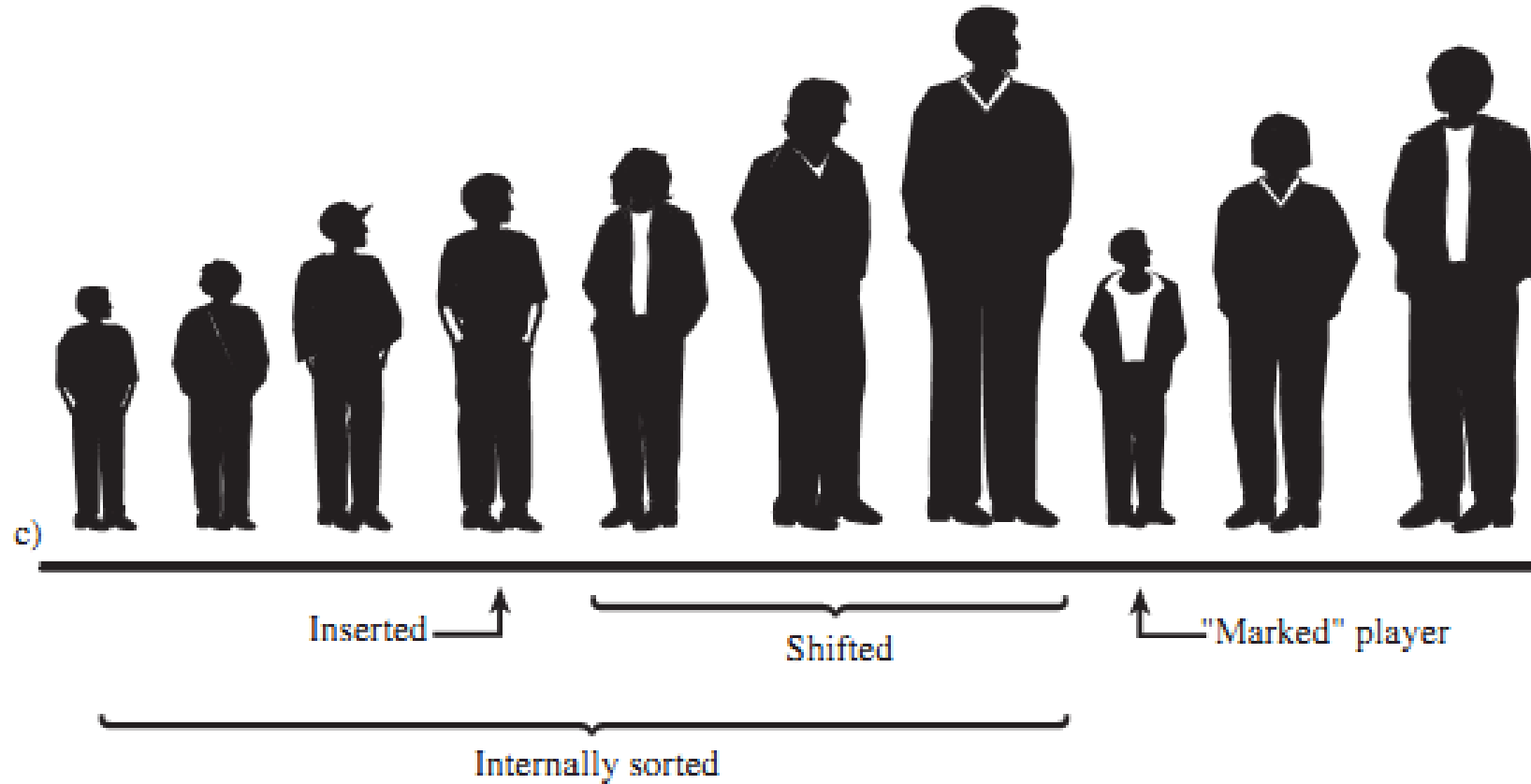


# Insert *marked* item to *partially sorted* list

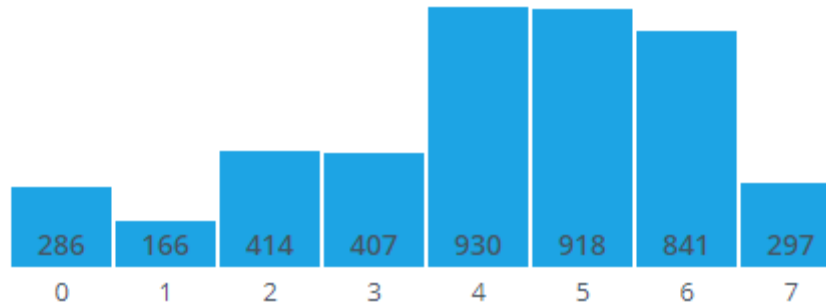
+ Take out of line → Shift right until appropriate place → Insert



# Insertion sort – Intermediate result



# Animation



Steps:  
Starting Insertion Sort.

<https://www.hackerearth.com/practice/algorithms/sorting/insertion-sort/visualize/>

<https://visualgo.net/en/sorting>

# Insertion Sort

- Result after each round:
  - Partially-ordered list is now one item larger and
  - The unsorted list is now one item smaller.
  - Marked item moves one slot to the right, so once more it is again in front of the leftmost unsorted item.
- Continue process until all unsorted items have been inserted.
- Hence the name 'insertion sort.'
- Code page 99 - 100



# Insertion sort

```
for(out=1; out<nElems; out++)    // out is dividing line
{
    long temp = a[out];           // remove marked item
    in = out;                     // start shifts at out
    while(in>0 && a[in-1] >= temp) // until one is smaller,
    {
        a[in] = a[in-1];         // shift item right,
        --in;                    // go left one position
    }
    a[in] = temp;                 // insert marked item
} // end for
} // end insertionSort()
```

# Discussion– How it really implemented!!

- Start with  $out = 1$ , which means there is only a single element to its 'left.'
  - We infer that this item to its left is sorted unto itself.
  - Hard to argue this is not true. (This is  $out = 0$ )
- $a[out]$  is the marked item, and it is moved into temp.
  - $a[out]$  is the leftmost unsorted item.

# Hand-on

+ Try with some examples

<b>89</b>	<b>58</b>	<b>29</b>	<b>40</b>	<b>12</b>	<b>42</b>	<b>10</b>	<b>1</b>
<b>1</b>	<b>32</b>	<b>12</b>	<b>53</b>	<b>11</b>	<b>76</b>	<b>23</b>	<b>89</b>

# Efficiency of the Insertion Sort

- Comparisons:

- On pass one, max of one;  
pass two, max of two, etc.  
Up to a max of  $n-1$  comparisons.

→  $1+2+3+\dots+n-1 = n*(n-1)/2$  comparisons.

- But, on average

$$n*(n-1)/4.$$

# Efficiency of the Insertion Sort (2 of 3)

- Copy:
  - Have lots of 'copies'
  - (same as number of comparisons – about)
  - ➔ But a copy is not nearly as time-consuming as a **swap**. **Think about this!!**
- For random data,
  - twice as fast as the bubble sort
  - faster than the selection sort.
- Still runs on  $O(n^2)$  time for random data.

# Efficiency of the Insertion Sort (3 of 3)

- If data is *nearly sorted* → quite well.
  - Condition in while loop is never true, so it almost runs in  $O(n)$  time; much better than  $O(n^2)$  time!
- If data is in *very unsorted* order (nearly backward)
  - → No faster than the bubble sort
  - as every possible comparison and shift is carried out.

# Sorting Objects

# Sorting Objects

- + Very important to be able to sort objects.
- + Must be careful, especially in noting that
  - + An array are **objects**, and
  - + The sort is based on **values of String attributes** inside of the object.
- + Use inherited String method, *compareTo*.



# Sorting Objects

```
while(in>0 &&                // until smaller one found,
      a[in-1].getLast().compareTo(temp.getLast())>0)
{
    a[in] = a[in-1];          // shift item to the right
    --in;                     // go left one position
}
```

**TABLE 3.1** Operation of the compareTo() Method

s2.compareTo(s1)	Return Value
s1 < s2	< 0
s1 equals s2	0
s1 > s2	> 0

# Secondary Sort Fields?

- Equal Keys Problems ? E.g., Last name
- Solutions
  - Using Secondary Key? e.g., ZIP code
  - Using 'Stable' Sort? What does this mean?
    - Only sort what needs to be sorted
    - Leave everything else in its original order
- All of the algorithms so far are stable
- Think Windows:
  - Arrange by Type; then by date modified...

# Comparing the Simple Sorts


- Bubble Sort, Selection Sort, and Insertion Sort all have a worst-case time complexity of  $O(n^2)$ .
- Bubble Sort, Selection Sort, and Insertion Sort all have a space complexity of  $O(1)$  as they are in-place sorting algorithms.

# Comparing the Simple Sorts

- Bubble Sort – simplest.
  - Use only if you don't have other algorithms available and 'n' is small.
- Selection Sort
  - Minimizes number of swaps, but number of comparisons still high.
  - Useful when amount of data is small, and swapping is very time consuming – like when sorting records in tables – internal sorts.
- Insertion Sort
  - The most versatile, and is usually best bet in most situations, if amount of data is small or data is almost sorted.
- For large n, other sorts are better. We will cover advanced sorts later...

# Comparing the Simple Sorts

- All require very little space, and they sort in place.
- Can't see the real efficiencies / differences unless you apply the different sources to large amounts of data.



Vietnam National University of HCMC  
International University  
School of Computer Science and Engineering



**THANK YOU**

Dr Vi Chi Thanh - [vcthanh@hcmiu.edu.vn](mailto:vcthanh@hcmiu.edu.vn)

<https://vichithanh.github.io>



SCAN ME