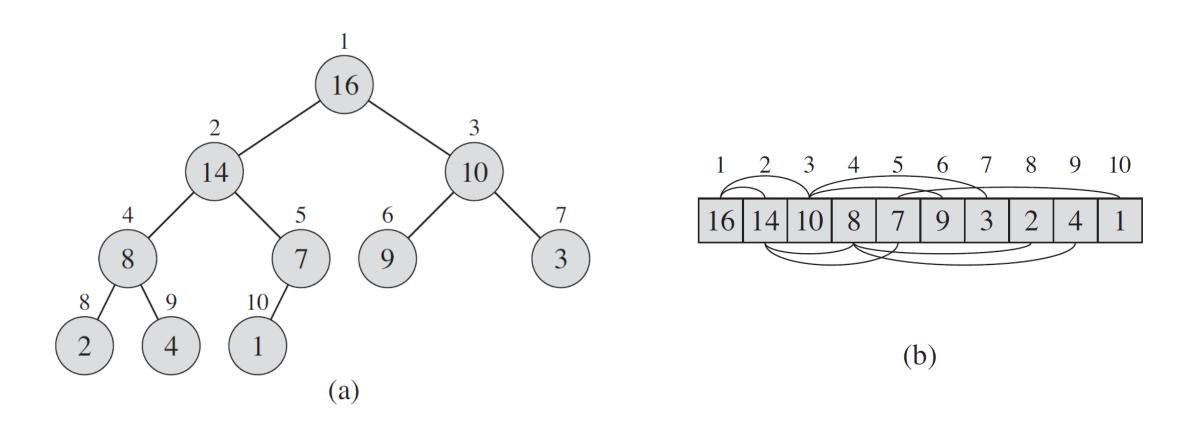
# GDI2 – 5 - Heapsort

# Heap and ist representation as array



PARENT(i)

1 return  $\lfloor i/2 \rfloor$ 

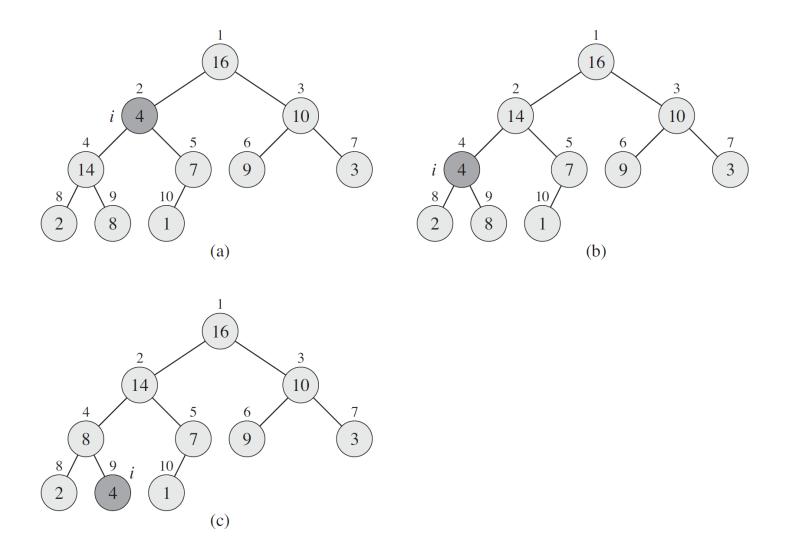
LEFT(i)

1 return 2i

RIGHT(i)

1 return 2i + 1

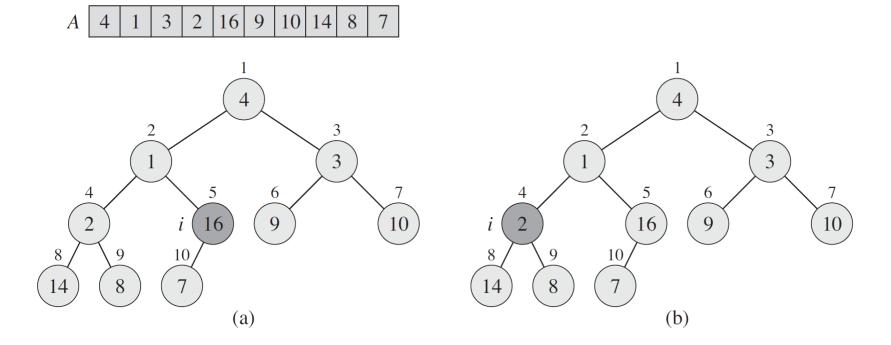
## MAX-HEAPIFY

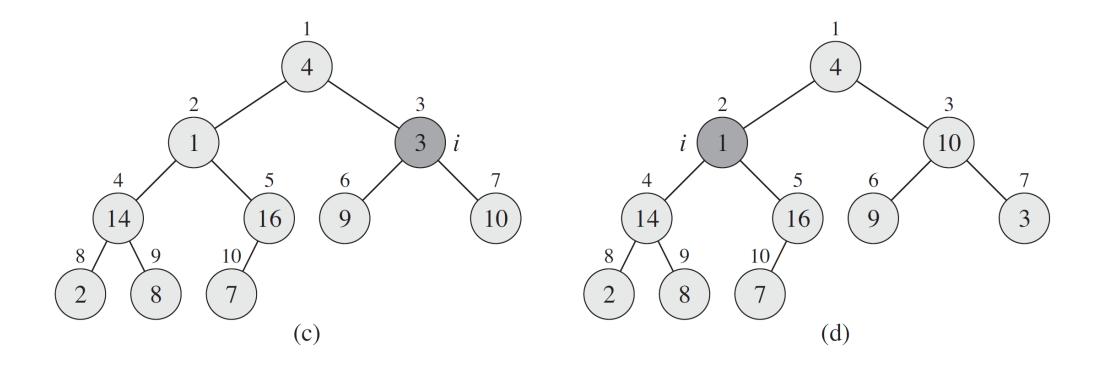


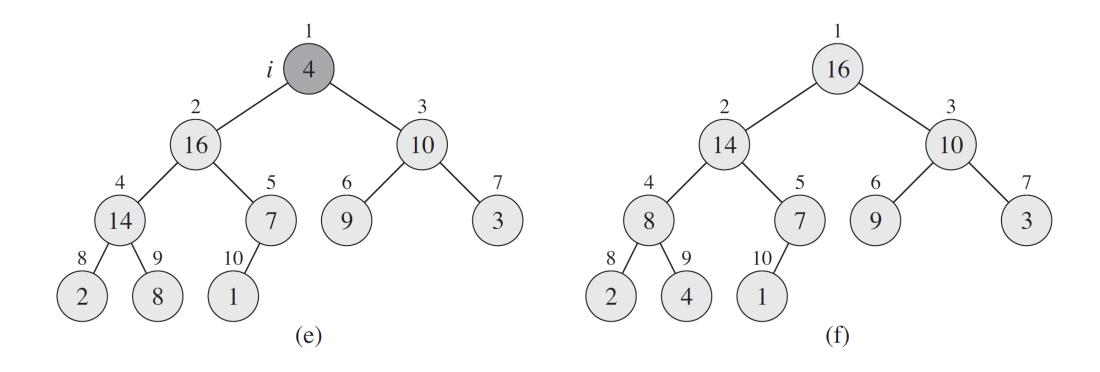
```
MAX-HEAPIFY (A, i)
 1 \quad l = \text{Left}(i)
 2 \quad r = RIGHT(i)
 3 if l \leq A. heap-size and A[l] > A[i]
         largest = l
   else largest = i
 6 if r \leq A. heap-size and A[r] > A[largest]
         largest = r
    if largest \neq i
         exchange A[i] with A[largest]
         MAX-HEAPIFY(A, largest)
10
```

### BUILD-MAX-HEAP (A)

- 1 A.heap-size = A.length
- 2 for  $i = \lfloor A.length/2 \rfloor$  downto 1
- 3 MAX-HEAPIFY(A, i)







At the start of each iteration of the **for** loop of lines 2–3, each node i + 1, i + 2, ..., n is the root of a max-heap.

**Initialization:** Prior to the first iteration of the loop,  $i = \lfloor n/2 \rfloor$ . Each node  $\lfloor n/2 \rfloor + 1, \lfloor n/2 \rfloor + 2, \ldots, n$  is a leaf and is thus the root of a trivial max-heap.

**Maintenance:** To see that each iteration maintains the loop invariant, observe that the children of node i are numbered higher than i. By the loop invariant, therefore, they are both roots of max-heaps. This is precisely the condition required for the call MAX-HEAPIFY (A, i) to make node i a max-heap root. Moreover, the MAX-HEAPIFY call preserves the property that nodes i + 1, i + 2, ..., n are all roots of max-heaps. Decrementing i in the **for** loop update reestablishes the loop invariant for the next iteration.

---- -- r --- --- --- --- ----

**Termination:** At termination, i = 0. By the loop invariant, each node 1, 2, ..., n is the root of a max-heap. In particular, node 1 is.

#### HEAPSORT(A)

```
1 BUILD-MAX-HEAP(A)

2 for i = A.length downto 2

3 exchange A[1] with A[i]

4 A.heap-size = A.heap-size -1

5 MAX-HEAPIFY(A, 1)
```

