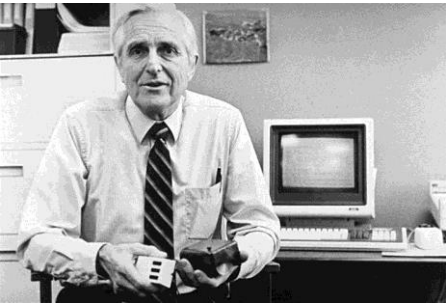
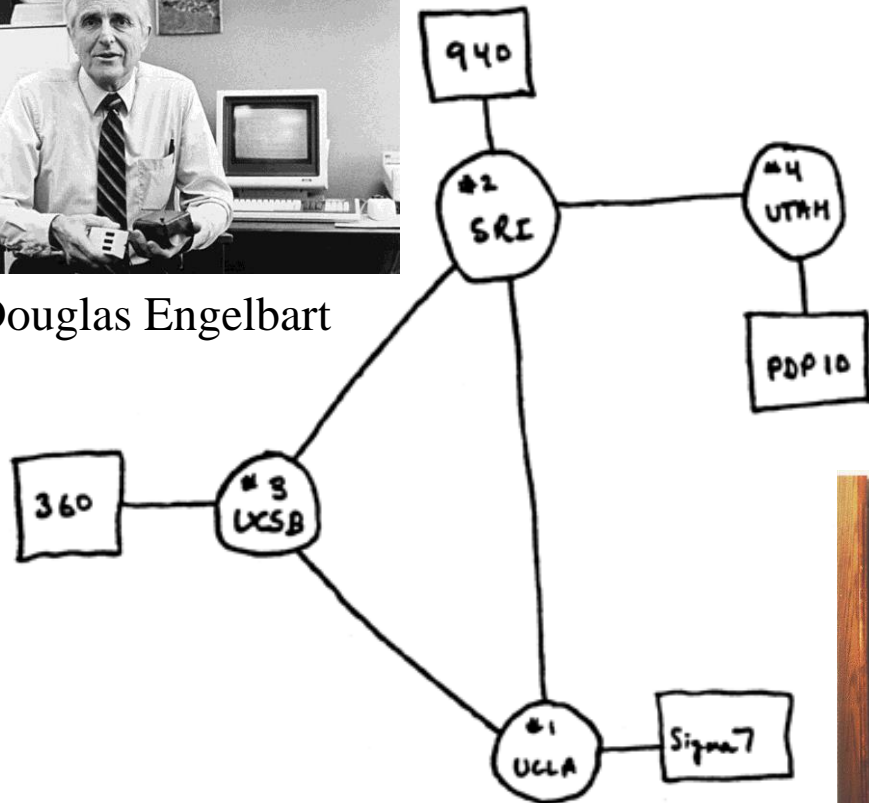


1. Introduction
2. Local Area Networking
3. Error Control and Retransmission Protocols
- 4. Architectural Principles of the Internet**
5. Flow Control in Internet: TCP
6. Routing Algorithms
7. Routing Protocols
8. The principles of ATM
9. Traffic Management in ATM
10. Scheduling
11. Quality of Service
12. Quality of Service routing
13. Peer-to-peer networks

Start of the Internet



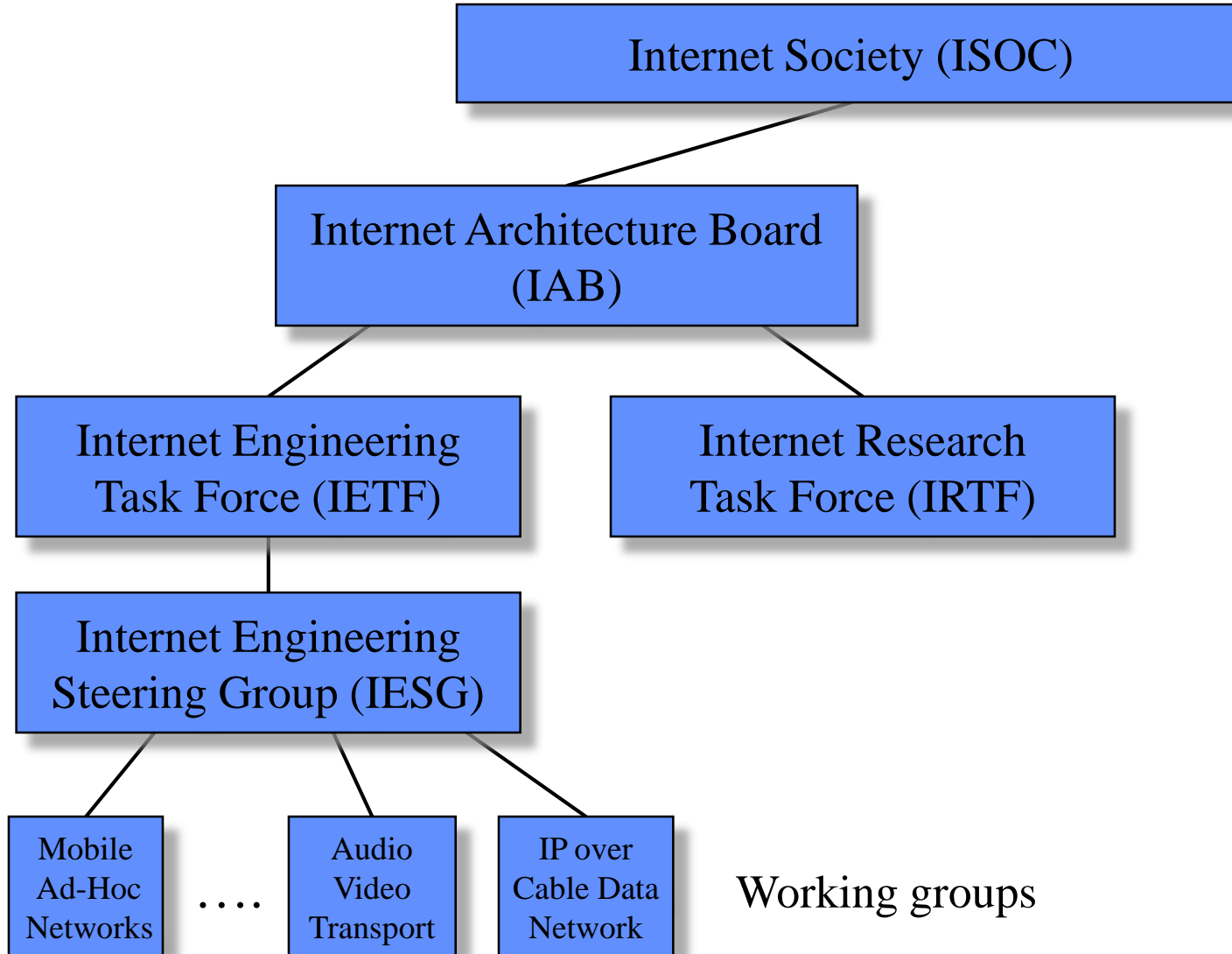
Douglas Engelbart



Leonard Kleinrock

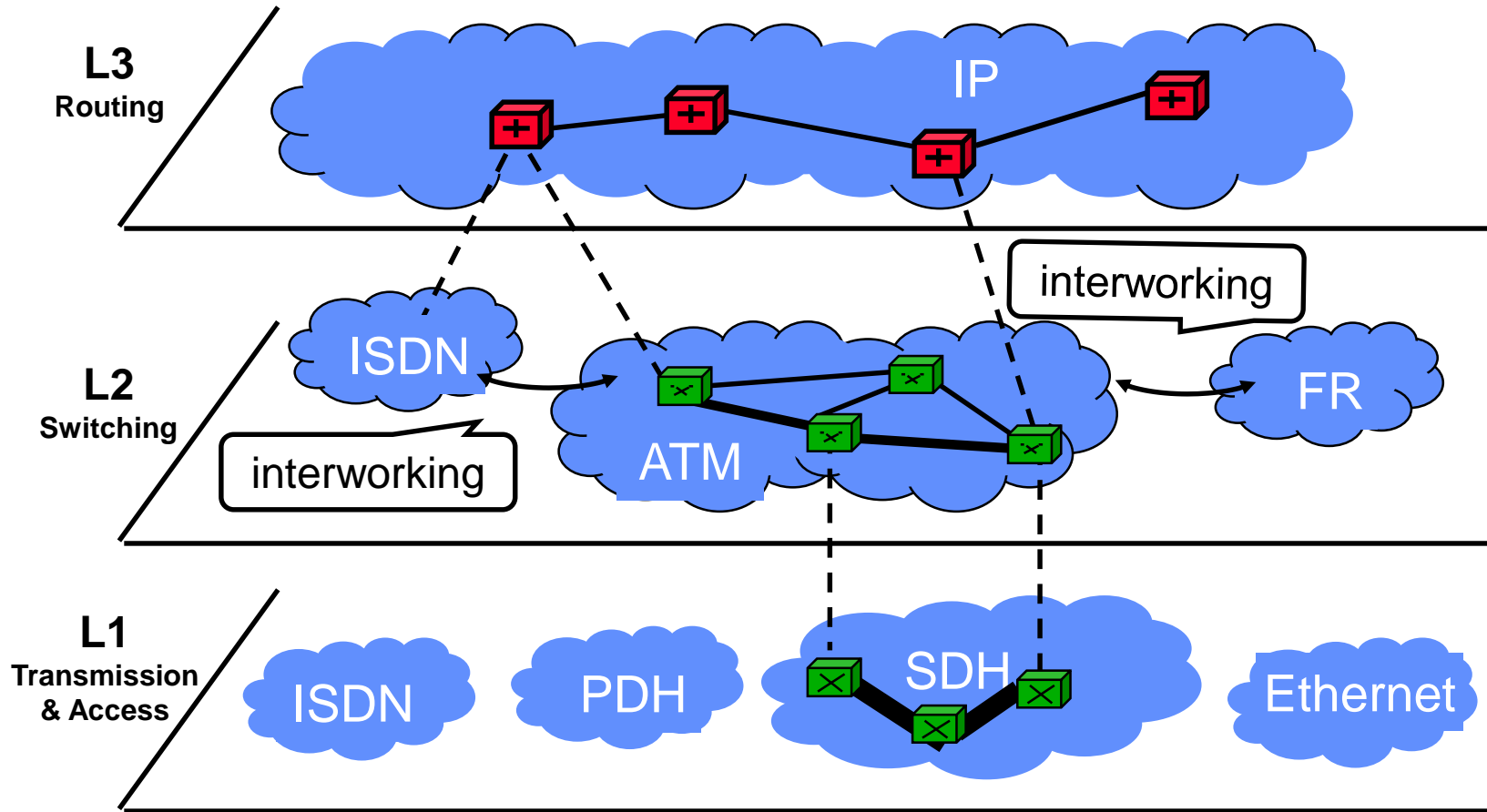
- Login attempt from UCLA to SRI (Stanford)
- L (ok)
- O (ok)
- G (crash)





Working groups

Layered Network Model



- The Internet Protocol (**IP**) defines an unreliable, connectionless, best-effort delivery mechanism:
 - specification of basic unit of transfer (packet)
 - IP software performs the routing function
 - IP includes the rules for unreliable, connectionless, best-effort delivery

IPv4 packet

0	4	8	16	19	24	31
vers	hlen	ToS	total length (bytes)			
identification			flags	fragment offset		
time to live		protocol	header checksum			
source IP address						
destination IP address						
IP options (if any)					padding	
user data						
...						

- Internet Registry (IR): IP address space allocation, protocol parameter assignment, DNS management
- Internet Corporation for Assigned Names and Numbers. Regional IRs:
 - ARIN (North America)
 - APNIC (Asia-Pacific)
 - LACNIC (Latin America, Caribbean)
 - RIPE NCC (Europe):
 - Local IRs (national):
 - ISPs:
 - » End-user

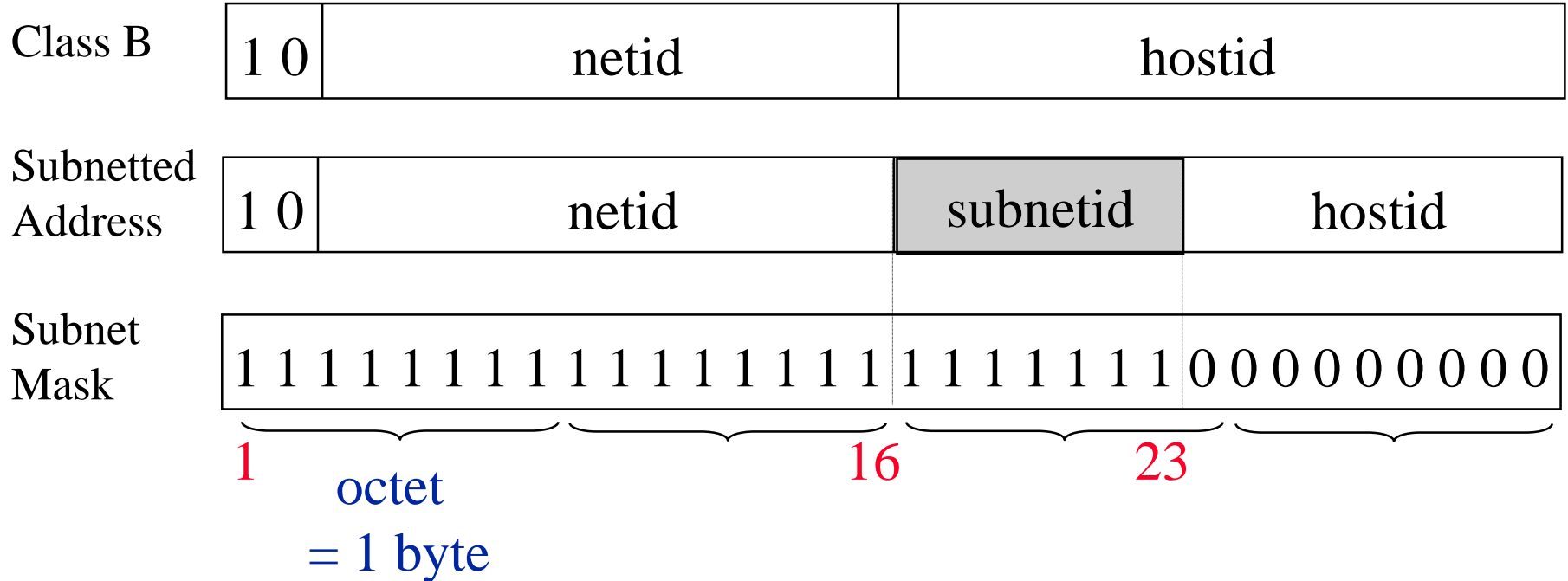
- Geographic (physical) versus logical (mobility)
- 32-bit (4-byte) addresses: $x.y.z.w$
- Netid and hostid
 - Subnet address
 - Classless Interdomain Routing (CIDR)

Class-based Addressing

	0	1	2	3	4	8	16	24	31																					
A	0	netid				hostid																								
B	1	0	netid								hostid																			
C	1	1	0	netid														hostid												
D	1	1	1	0	multicast address																									
E	1	1	1	1	0	reserved for future use																								

E.g. 129.82.6.25 belongs to the class B network 129.82.0.0
(129 = 1000 0001)

Subnet Addressing



Subnet address = IP address \otimes subnet mask

Subnet Mask: consists here of 23 one bits = 255.255.254.0

Subnet: $x.y.z.w/v$, where v is an integer smaller than 32

C:\Documents and Settings\TUD>ipconfig /?

USAGE:

```
ipconfig [/? | /all | /renew [adapter] | /release [adapter] |  
        /flushdns | /displaydns | /registerdns |  
        /showclassid adapter |  
        /setclassid adapter [classid] ]
```

where

adapter Connection name
 (wildcard characters * and ? allowed, see examples)

Options:

```
/?                      Display this help message  
/all                    Display full configuration information.  
/release                Release the IP address for the specified adapter.  
/renew                  Renew the IP address for the specified adapter.  
/flushdns               Purges the DNS Resolver cache.  
/registerdns            Refreshes all DHCP leases and re-registers DNS names  
/displaydns            Display the contents of the DNS Resolver Cache.  
/showclassid            Displays all the dhcp class IDs allowed for adapter.  
/setclassid            Modifies the dhcp class id.
```

The default is to display only the IP address, subnet mask and default gateway for each adapter bound to TCP/IP.

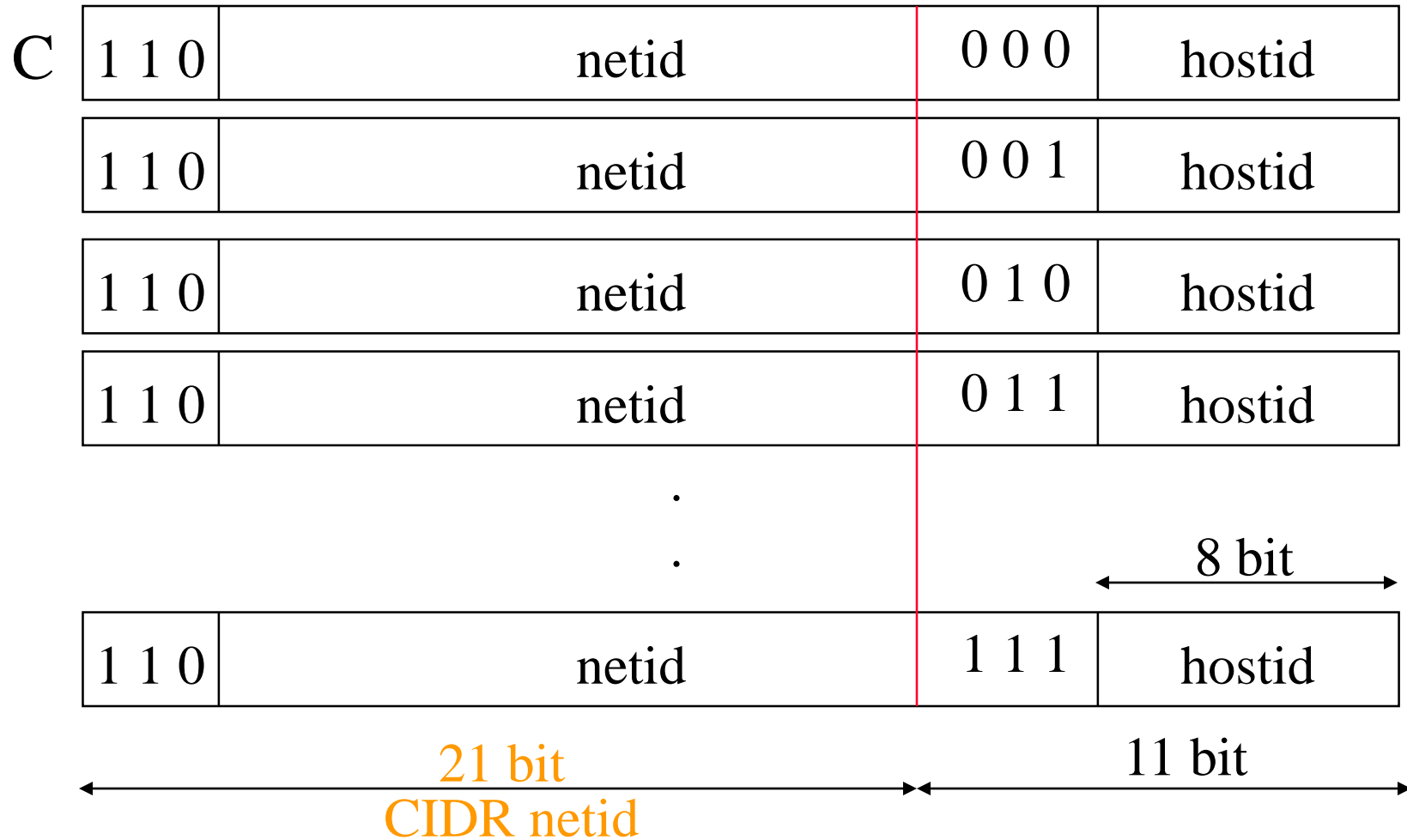
For Release and Renew, if no adapter name is specified, then the IP address leases for all adapters bound to TCP/IP will be released or renewed.

For Setclassid, if no ClassId is specified, then the ClassId is removed.

Examples:

```
> ipconfig                      ... Show information.  
> ipconfig /all                ... Show detailed information  
> ipconfig /renew               ... renew all adapters  
> ipconfig /renew EL*           ... renew any connection that has its  
                                 name starting with EL  
> ipconfig /release *Con*      ... release all matching connections,  
                                 eg. "Local Area Connection 1" or  
                                 "Local Area Connection 2"
```

C:\Documents and Settings\TUD>



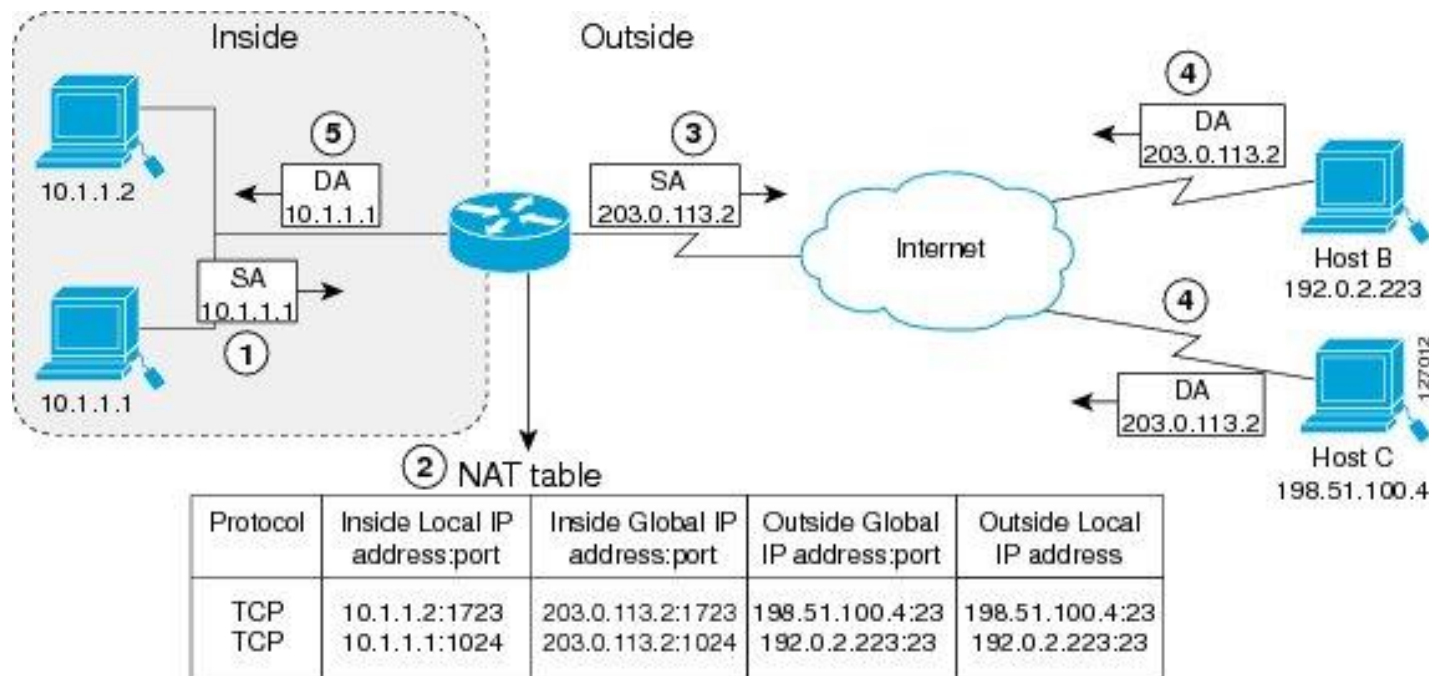
- Class-based addressing loosened into ranges using subnet masks
- Private networks
 - 10.0.0.0/8, 172.16.0.0/12, 192.168.0.0/16
- Link-local
 - 169.254.0.0/16 (The “the-Internet-is-down” range)
- Multicast
 - 224.0.0.0/4
 - 255.255.255.255/32 (local subnet only)

Network Address Translation

- Remaps private address outgoing flows to public address
- Hides internal network topology and addresses
- Conserves IP address usage
- Incoming connections handled through port mapping
 - Hence problematic with non-TCP or -UDP flows

Network Address Translation

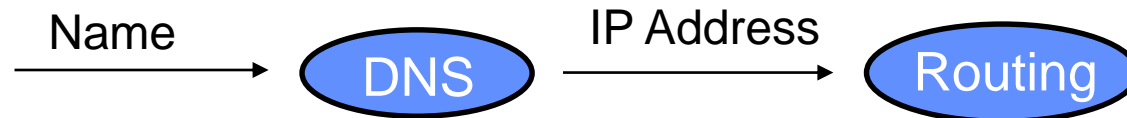
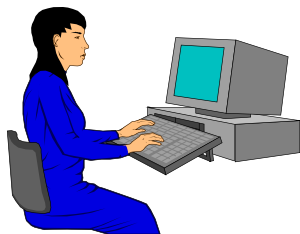
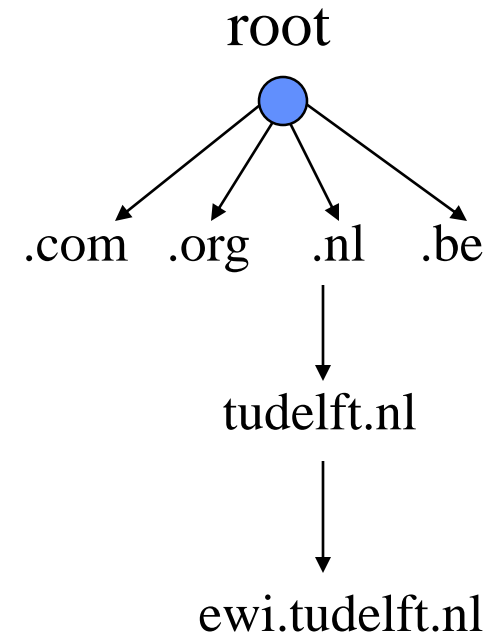
- Outgoing flows automatically added
- Incoming flows need prior existing rules



[http://www.cisco.com/c/en/us/td/docs/ios-xml/ios/ipaddr_nat/configuration/15-mt/nat-15-mt-book/iadnat-addr-consv.html]

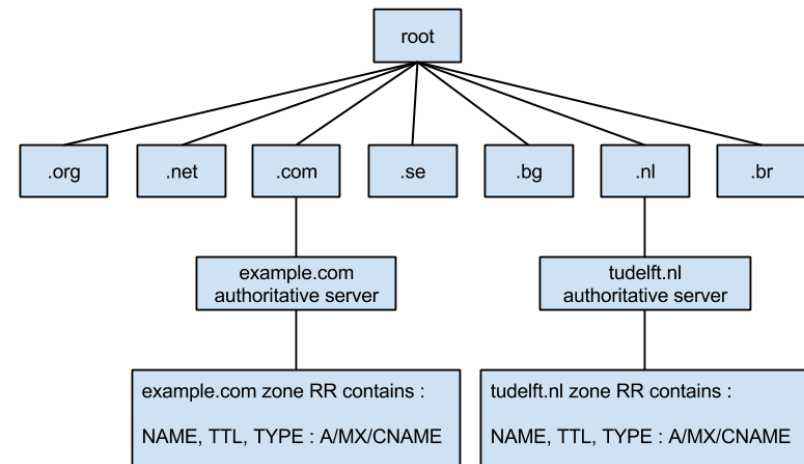
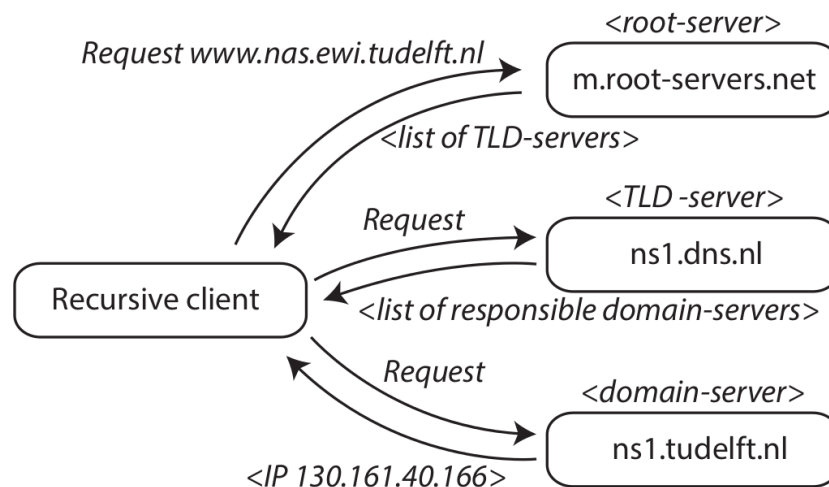
Domain Name System (DNS)

- Humans work best using the name of a host (e.g. ewi.tudelft.nl)
- IP Addresses used to identify a host, and for routing (e.g. 130.161.40.75)
- DNS
 - deals with mapping: Name \longleftrightarrow IP Address



Domain Name System (DNS)

- Hierarchically distributed database
 - Root Domain -> Top Level Domain -> Authoritative domain
- Clients use recursive servers
- High level of caching



C:\Documents and Settings\TUD>nslookup

Default Server: ns1.tudelft.nl

Address: 130.161.180.1

> ?

Commands: (identifiers are shown in uppercase, [] means optional)

NAME - print info about the host/domain NAME using default server

NAME1 NAME2 - as above, but use NAME2 as server

help or ? - print info on common commands

set OPTION - set an option

all - print options, current server and host

[no]debug - print debugging information

[no]d2 - print exhaustive debugging information

[no]defname - append domain name to each query

[no]recurse - ask for recursive answer to query

[no]search - use domain search list

[no]vc - always use a virtual circuit

domain=NAME - set default domain name to NAME

srchlist=N1[/N2/.../N6] - set domain to N1 and search list to N1,N2, etc.

root=NAME - set root server to NAME

retry=X - set number of retries to X

timeout=X - set initial time-out interval to X seconds

type=X - set query type (ex. A,ANY,CNAME,MX,NS,PTR,SOA,SRU)

querytype=X - same as type

class=X - set query class (ex. IN (Internet), ANY)

[no]mxfr - use MS fast zone transfer

ixfrver=X - current version to use in IXFR transfer request

server NAME - set default server to NAME, using current default server

lserver NAME - set default server to NAME, using initial server

finger [USER] - finger the optional NAME at the current default host

root - set current default server to the root

ls [opt] DOMAIN [> FILE] - list addresses in DOMAIN (optional: output to FILE)

-a - list canonical names and aliases

-d - list all records

-t TYPE - list records of the given type (e.g. A,CNAME,MX,NS,PTR etc.)

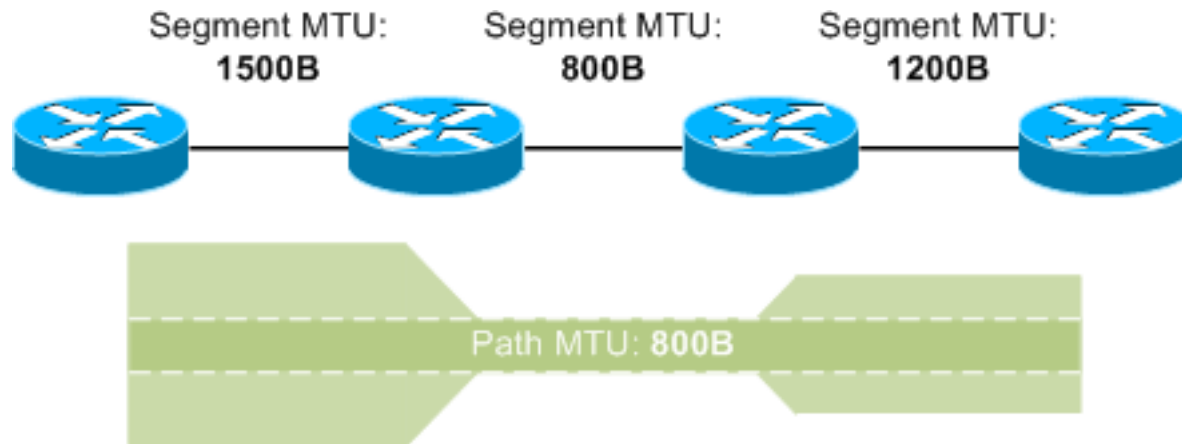
view FILE - sort an 'ls' output file and view it with pg

exit - exit the program

>

Maximum Transmission Unit

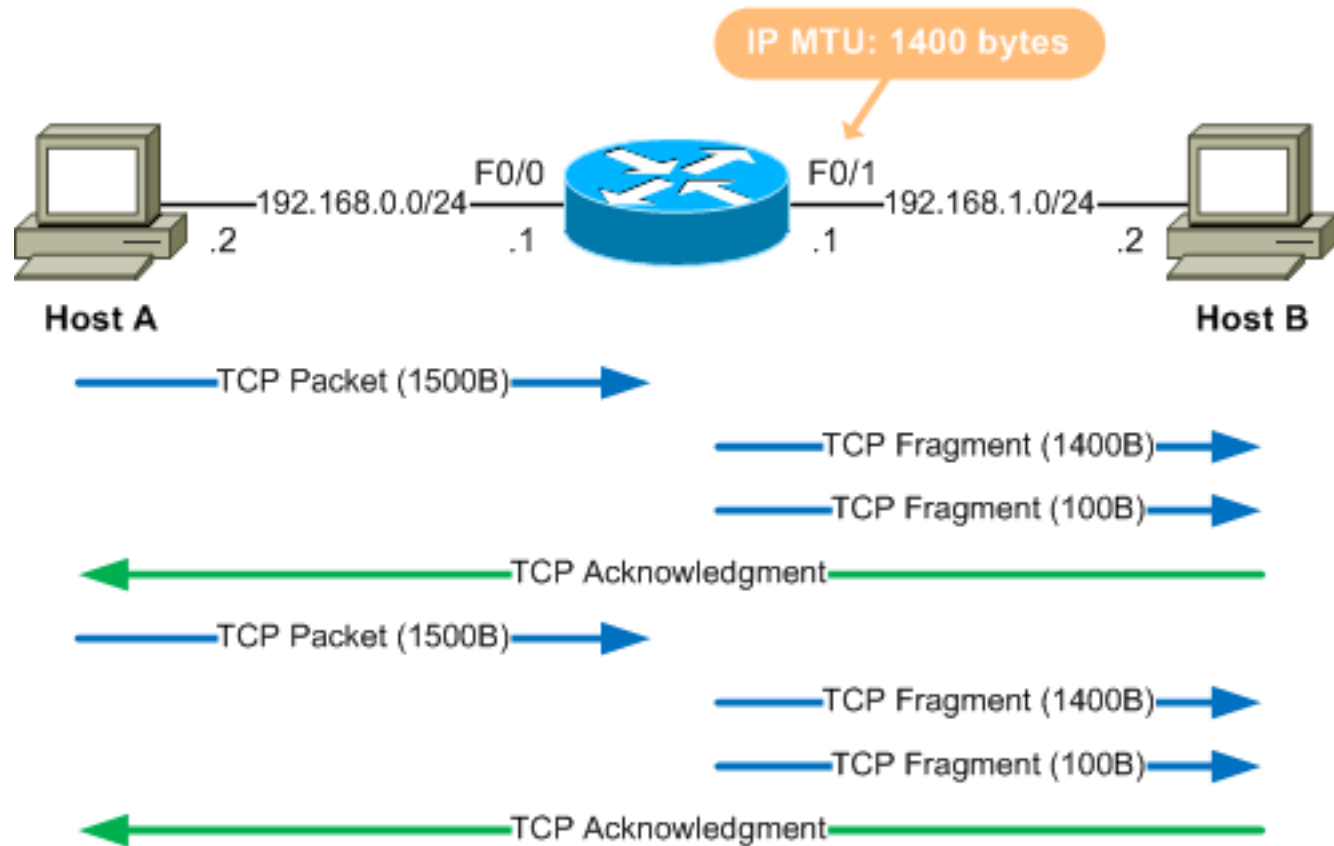
- Not all links support the same fragment size
- Connectionless and layered nature prevents learning



[<http://packetlife.net/blog/2008/aug/18/path-mtu-discovery/>]

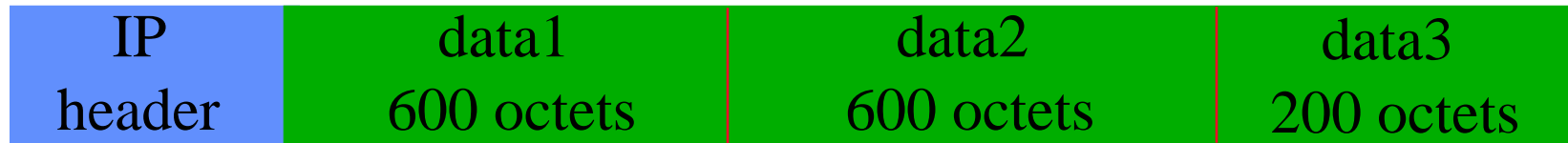
Fragmentation

- Instead, IPv4 fragments packets



[<http://packetlife.net/blog/2008/aug/18/path-mtu-discovery/>]

Fragmentation



Fragment1: offset 0



Fragment2: offset 600



Fragment3: offset 1200

- IPv6 is not a simple derivative of IPv4, but a definite improvement
- Four major changes over IPv4:
 1. *extended addressing capability:*
128 bits versus 32 bits in IPv4
 2. *a fixed format to all headers*
no option element, no header length field, but extension headers
 3. *no header checksum*
diminish cost of processing, other layers check & correct for errors
 4. *no hop-by-hop fragmentation*
unit of transmission = unit of control: use *path MTU discovery*
 5. *no Network Address Translation*

IPv6 Header Format



Header comparison IPv6-IPv4

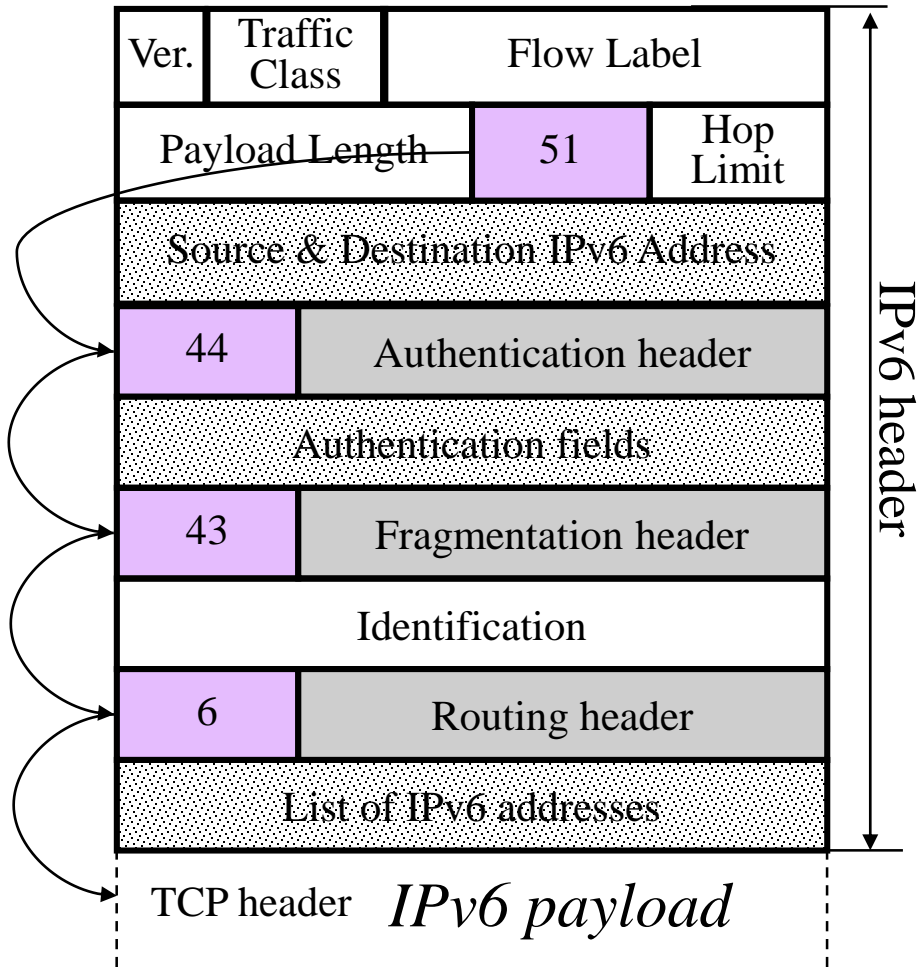
Ver.	Traffic Class	Flow Label		
Payload Length		Next Header	Hop Limit	
Source Address				
Destination Address				

Ver.	Hdr Len	Type of Service	Total Length	
Identification			Flg	Fragment Offset
Time to Live		Protocol	Header Checksum	
Source Address				
Destination Address				
Options...				

shaded fields have no equivalent in the other version

IPv6 header is twice as long (40 bytes) as IPv4 header without options (20 bytes)

Extension Header



- Text Representations:
 - $x_1:x_2:x_3:x_4:x_5:x_6:x_7:x_8$ where x_k hexadecimal value of 16-bit.
e.g. 1080:0:0:0:0:0:200C:417A
 - “zero” compressed notation: e.g. 1080::200C:417A
 - Mixed IPv4 and IPv6: $x:x:x:x:x:x:d.d.d.d$ where d decimal value of 8-bit IPv4 address
e.g. 0:0:0:0:0:0:13.1.68.3 = ::13.1.68.3
- Three types:
 - unicast, anycast and multicast (no broadcast)
 - multicast: FF: $x_2:x_3:x_4:x_5:x_6:x_7:x_8$ (1/256 of address space)
 - All types are assigned to interfaces, not nodes

No.	Time	Source	Destination	Protocol	Info
1	0.000000	fe80::cd94:9867:5e1ff02::1	ff33:de59	ICMPv6	Neighbor solicitation for fe80::7c6e:6ab4:f633:de59 from 00:24:d7:2f:c0:48
2	0.718421	fe80::390e:e68d:390ff02::1	ffbb:b426	ICMPv6	Neighbor solicitation for fe80::b5f2:fa7:f5bb:b426 from 00:21:5c:0a:f9:2d

```

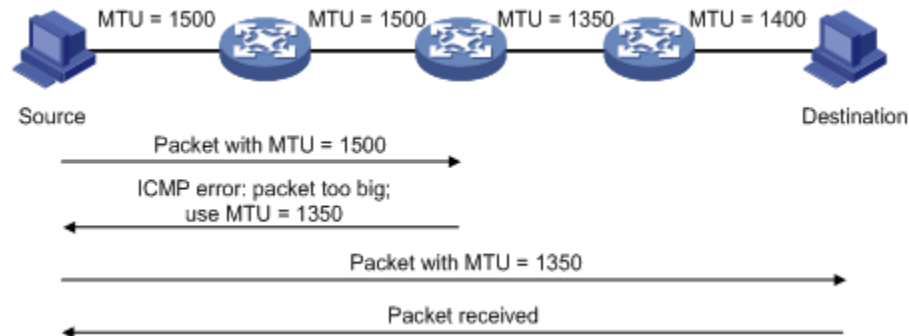
+ Frame 1: 86 bytes on wire (688 bits), 86 bytes captured (688 bits)
+ Ethernet II, Src: IntelCor_2f:c0:48 (00:24:d7:2f:c0:48), Dst: IPv6mcast_ff:33:de:59 (33:33:ff:33:de:59)
+ Internet Protocol Version 6, Src: fe80::cd94:9867:5e10:af74 (fe80::cd94:9867:5e10:af74), Dst: ff02::1:ff33:de59 (ff02::1:ff33:de59)
  + 0110 .... = Version: 6
  + .... 0000 0000 .... .... .... .... = Traffic class: 0x00000000
    .... .... .... 0000 0000 0000 0000 0000 = FlowLabel: 0x00000000
    Payload length: 32
    Next header: ICMPv6 (0x3a)
    Hop limit: 255
    Source: fe80::cd94:9867:5e10:af74 (fe80::cd94:9867:5e10:af74)
    Destination: ff02::1:ff33:de59 (ff02::1:ff33:de59)
+ Internet Control Message Protocol v6

```

0010	00000000	00000000	00000000	00100000	00111010	11111111	11111110	10000000
0018	00000000	00000000	00000000	00000000	00000000	00000000	11001101	10010100
0020	10011000	01100111	01011110	00010000	10101111	01110100	11111111	00000010	gA...	...
0028	00000000	00000000	00000000	00000000	00000000	00000000	00000000	00000000
0030	00000000	00000001	11111111	00110011	11011110	01011001	10000111	00000000	...3.Y...	...
0038	11010111	01000001	00000000	00000000	00000000	00000000	11111110	10000000	A.....	...
0040	00000000	00000000	00000000	00000000	00000000	00000000	01111100	01101110 n	...
0048	01101010	10110100	11110110	00110011	11011110	01011001	00000001	00000001	j...3.Y...	...

Path MTU Discovery

- IPv6 does not fragment like IPv4
- Relies on ICMP to find MTU
- Creates blackholes when blocked



[<http://bitsanddoctets.blogspot.nl/2012/05/mtu-and-mss-why-should-we-care.html>]

- Internet control message protocol (ICMP)
 - two modes: query (echo, timestamp, etc...) and error
 - Diagnosing error conditions sent back to the source when a router destroys a packet
 - Not to enhance reliability, but merely to provide feedback about network problems.

- Ping (via ICMP *query* mode ‘echo’)
- Ping -f -l <size in bytes> address
 - f: sets don’t fragment flag
 - l: sets payload size

- Traceroute (via ICMP *error* mode)
- Windows: tracert address
- Returns the path to the address
- Sends multiple UDP packets with increasing TTL
- Artefacts:
 - ICMP messages may be discarded
 - UDP packets might traverse multiple paths
 - Interface and no router addresses

- Give, in x.y.z.w. notation, the subnet mask that belongs to the address 130.140.150.160/20.
- Give three differences between the headers of IPv4 and IPv6 and explain the motivation behind these differences.

- Use *ping* to find RTTs,
- And *tracert* to find the paths to:
 1. Your gateway router
 2. Common destinations, such as
 1. Google website
 2. TU Delft website
 3. Major Dutch news site
 4. Major foreign (transatlantic) news site
- 3. Each other's laptop