

Tuan Nguen

**Question 1:**

a) The invariant that is maintained between the variables `bufin`, `bufout` and `bufcnt` is that if `bufin < bufout`, means that the array is wrapping around, so `bufcnt = bufin + NBUF - bufout`. Otherwise, `bufcnt = bufin - bufout`. We can use this to eliminate one of the variables. Calculate `bufcnt + bufout` and if it is larger than `NBUF`, we decrement the sum by `NBUF`. Otherwise, that is `bufin`.

```
void uart_handler(void) {
    if(UART_TXDRDY) {
        UART_TXDRDY = 0;
        if(bufcnt == 0)
            txidle = 1;
        else {
            UART_TXD = txbuf[bufout];
            bufcnt--;
            bufout = (bufout+1) % NBUF;
        }
    }
}

/* serial_putc - send output character */
void serial_putc(char ch) {
    while(bufcnt == NBUF) pause();

    intr_disable();
    if(txidle) {
        UART_TXD = ch;
        txidle = 0;
    }
    else {
        if(bufout+bufcnt > NBUF){
            txbuf[bufout+bufcnt-NBUF] = ch;
            bufcnt++;
        }
    }
    intr_enable();
}
```

b) We want to disable the interrupts before checking `txidle` because there could be a situation when the variable is set to 0, and we want it to remain that way, but if the handler is called, the `uart_handler` can change it back to 1.

c) The `bufcnt++` instruction can be implemented in assembly like this:

```

ldr r0, =bufcnt
ldr r1, [r0]
add r1, r1, #1
str r1, [r0]

```

If an interrupt happens between the second and the third line, then the effect of `bufcnt--` will be lost, because when the interrupt handler returns, it is an old value that is incremented.

d) The `wfi` command cannot see interrupts that have happened before it, whereas the `wfe` command can. So if an interrupt happens during the comparison between the `bufcnt` and `NBUF`, then the program would crash.

e)

```

/* serial_putc - send output character */
void serial_putc(char ch) {
    while(bufcnt == NBUF) pause();

    intr_disable();
    if(txidle) {
        intr_enable();
        UART_TXD = ch;
        txidle = 0;
    }
    else {
        if(bufout+bufcnt > NBUF){
            txbuf[bufin] = ch;
            intr_disable();
            bufcnt++;
            intr_enable();
            bufin = (bufin + 1) % NBUF;
        }
    }
}

```

### **Question 2:**

```

static const unsigned small[] = {
    0x2df0, 0x5fb0, 0x8af0
};

static unsigned row = 0;

void heartBigDelay(void) {
    unsigned n = 1050000
    while(n-- > 0){
        for(int p = 0; p < 3; p++){
            GPIO_OUT = heart[p]

```

```

    }
}

void heartSmallDelay(void) {
    unsigned n = 150000
    while(n-- > 0){
        for(int p = 0; p < 3; p++){
            GPIO_OUT = heart[p]
        }
    }
}

void smallDelay(void) {
    unsigned n = 150000
    while(n-- > 0){
        for(int p = 0; p < 3; p++){
            GPIO_OUT = small[p]
        }
    }
}

void timer1_handler(void) {
    if (TIMER1_COMPARE[0]) {
        heartBigDelay();
        smallDelay();
        heartSmallDelay();
        smallDelay();
        TIMER1_COMPARE[0] = 0;
    }
}

void init_timer(void) {
    TIMER1_STOP = 1;
    TIMER1_MODE = TIMER_Timer_Mode;
    TIMER1_BITMODE = TIMER_16Bit;
    TIMER1_PRESCALER = 4; // 1MHz = 16MHz / 2^4
    TIMER1_CLEAR = 1;
    TIMER1_CC[0] = 1000 * TICK;
    TIMER1_SHORTS = BIT(TIMER_COMPARE0_CLEAR);
    TIMER1_INTENSET = BIT(TIMER_INTEN_COMPARE0);
    TIMER1_START = 1;

    set_priority(TIMER1_IRQ, 3);
    enable_irq(TIMER1_IRQ);
}

```

```

void init(void) {
    GPIO_DIR = 0xffff0;
    GPIO_PINCNF[BUTTON_A] = 0;
    GPIO_PINCNF[BUTTON_B] = 0;
    GPIO_OUT = heart[0];

    init_timer();

    while (1) {
        pause();
    }
}

```

The drawbacks of the way that my program is designed(I guess there is less naïve way to do it) is that for every „picture” has to be written in its own separate function, and if we want several “pictures” to flash this way would be infeasible. The major drawback is that there is only one main program, so only one task can be executed at a time, and the other tasks have to be called at appropriate times, with their entire state stored in variables.

### **Question 3:**

\* I will explain my idea of what the program should be \*

The program will be similar to the one we have seen for the interrupt based prime number program. In the `rng_handler` will be checked why an interrupt has occurred, so we want to check if `RNG_VALRDY` is 1. If it is, then we test if the buffer is empty. If it is, then we get the number that has been generated and pass it to `randint()`, where we will take 4 random numbers, pad left each of them and then put the next one to the right of the previous one. If the buffer is not empty, get a number from the buffer and pass it to `randint()`. Decrease `bufcnt` and increase `bufout`. In the `randint()` function we check if the buffer is full, and if it is, then we pause, until at least one place in the buffer is freed. Disable interrupts, check if the program is idling, and depending on that 1) we put a randomly generated number into our number that we have to return; 2) fetch a number from the buffer and put that into our number that we have to return. Enable interrupts. For `roll()` the procedure would be the same, but instead of appending 4 8-bit numbers, we can calculate `mod6` of the decimal number, corresponding to the 8-bit binary.

### **Question 5:**

If the stack grew upwards, instead of downwards, the program would still be vulnerable. If there is a register `rd` that points to the start of the buffer, and there is an operation that involves `rd`, which there probably will be, we can write our code on the buffer. Then use that same register to return to our code and then execute it.