

# No-repeat Substring (hard)

#### We'll cover the following

- Problem Statement
- Try it yourself
- Solution
- Code
  - Time Complexity
  - Space Complexity

### **Problem Statement**

Given a string, find the **length of the longest substring** which has **no repeating characters**.

## Example 1:

Input: String="aabccbb"

Output: 3

Explanation: The longest substring without any repeating character

s is "abc".

### Example 2:

```
Imput: String="abbbb"
Output: 2
```



Explanation: The longest substring without any repeating character s is "ab".

#### Example 3:

```
Input: String="abccde"
Output: 3
Explanation: Longest substrings without any repeating characters ar
e "abc" & "cde".
```

# Try it yourself

Try solving this question here:

```
JS
                                            C++
               Python3
   Java
import java.util.*;
                                                                               class NoRepeatSubstring {
        public static int findLength(String str) {
                char pre = 0;
                int end = 0;
                int max = 0;
                List<Character> arr = new ArrayList<>();
                for (end = 0; end < str.length(); end++) {</pre>
                        arr.add(str.charAt(end));
                        if (arr.size() - 1 >= 0 && arr.get(arr.size() -1).equals(pre)
                                 max = Math.max(max, arr.size() - 1);
                                 arr.clear();
                                 arr.add(str.charAt(end));
                        pre = str.charAt(end);
                return max;
        }
}
```











This problem follows the **Sliding Window** pattern and we can use a similar dynamic sliding window strategy as discussed in Longest Substring with K Distinct Characters

(https://www.educative.io/collection/page/5668639101419520/567146485435596 We can use a **HashMap** to remember the last index of each character we have processed. Whenever we get a repeating character we will shrink our sliding window to ensure that we always have distinct characters in the sliding window.

### Code

Here is what our algorithm will look like:



```
import java.util,*;
 [>_] ⊖OUCQUIV⊖(/learn)
class NoRepeatSubstring {
 public static int findLength(String str) {
    int windowStart = 0, maxLength = 0;
    Map<Character, Integer> charIndexMap = new HashMap<>();
    // try to extend the range [windowStart, windowEnd]
    for (int windowEnd = 0; windowEnd < str.length(); windowEnd++) {
      char rightChar = str.charAt(windowEnd);
      // if the map already contains the 'rightChar', shrink the window from the beg:
      // we have only one occurrence of 'rightChar'
      if (charIndexMap.containsKey(rightChar)) {
        // this is tricky; in the current window, we will not have any 'rightChar' at
        // and if 'windowStart' is already ahead of the last index of 'rightChar', we
        windowStart = Math.max(windowStart, charIndexMap.get(rightChar) + 1);
      }
      charIndexMap.put(str.charAt(windowEnd), windowEnd); // insert the 'rightChar'
      maxLength = Math.max(maxLength, windowEnd - windowStart + 1); // remember the maxLength - windowStart + 1);
    }
    return maxLength;
 }
 public static void main(String[] args) {
    System.out.println("Length of the longest substring: " + NoRepeatSubstring.findLe
    System.out.println("Length of the longest substring: " + NoRepeatSubstring.findLe
    System.out.println("Length of the longest substring: " + NoRepeatSubstring.findLe
 }
}
```







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#### **Time Complexity**

The time complexity of the above algorithm will be O(N) where 'N' is the number of characters in the input string.

#### **Space Complexity**

The space complexity of the algorithm will be O(K) where K is the number of distinct characters in the input string. This also means K <= N, because in the worst case, the whole string might not have any repeating character so the entire string will be added to the **HashMap**. Having said that, since we

can expect a fixed set of characters in the input string (e.g., 26 for English letters), we can say that the algorithm runs in fixed space O(1); in this case,



we can use a fixed-size array instead of the **HashMap**.

