

Discrete Math (Honor) 2022-Fall Homework-5

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Problem 1. (12 Points)

Write the prenex normal form for each of the following formulas.

1. $(\forall x)(\forall y)((\exists z)P(x, y, z) \wedge ((\exists u)Q(x, u) \rightarrow (\exists v)Q(y, v)))$
2. $(\neg(\exists x)F(x) \vee (\forall y)G(y)) \wedge (F(a) \rightarrow (\forall z)H(z))$

Answer:

1. $(\forall x)(\forall y)((\exists z)P(x, y, z) \wedge (\neg(\exists u)Q(x, u) \vee (\exists v)Q(y, v)))$
 $(\forall x)(\forall y)(\exists z)(P(x, y, z) \wedge (\neg(\exists u)Q(x, u) \vee (\exists v)Q(y, v)))$
 $(\forall x)(\forall y)(\exists z)(P(x, y, z) \wedge ((\forall u)\neg Q(x, u) \vee (\exists v)Q(y, v)))$
 $(\forall x)(\forall y)(\exists z)(P(x, y, z) \wedge (\forall u)(\neg Q(x, u) \vee (\exists v)Q(y, v)))$
 $(\forall x)(\forall y)(\exists z)(P(x, y, z) \wedge (\forall u)(\exists v)(\neg Q(x, u) \vee Q(y, v)))$
 $(\forall x)(\forall y)(\exists z)(\forall u)(\exists v)(P(x, y, z) \wedge (\neg Q(x, u) \vee Q(y, v)))$
2. $(\neg(\exists x)F(x) \vee (\forall y)G(y)) \wedge (F(a) \rightarrow (\forall z)H(z))$
 $((\forall x)\neg F(x) \vee (\forall y)G(y)) \wedge (\neg F(a) \vee (\forall z)H(z))$
 $(\forall x)(\forall y)(\neg F(x) \vee G(y)) \wedge (\forall z)(\neg F(a) \vee H(z))$
 $(\forall x)(\forall y)((\neg F(x) \vee G(y)) \wedge (\forall z)(\neg F(a) \vee H(z)))$
 $(\forall x)(\forall y)(\forall z)((\neg F(x) \vee G(y)) \wedge (\neg F(a) \vee H(z)))$

Problem 2. (12 Points)

Write the Skolem normal form for each of the following formulae.

1. $(\forall x)(P(x) \rightarrow Q(x)) \rightarrow ((\exists x)P(x) \rightarrow (\exists x)Q(x))$
2. $(\forall x)(P(x) \rightarrow (\exists y)Q(x, y)) \vee (\forall z)R(z)$

Answer

1. Step 1:
 $(\forall x)(P(x) \rightarrow Q(x)) \rightarrow ((\exists x)P(x) \rightarrow (\exists x)Q(x))$
 $(\forall x)(\neg P(x) \vee Q(x)) \rightarrow (\neg(\exists x)P(x) \vee (\exists x)Q(x))$
 $(\forall x)(\neg P(x) \vee Q(x)) \rightarrow ((\forall x)\neg P(x) \vee (\exists x)Q(x))$
 $\neg(\forall x)(\neg P(x) \vee Q(x)) \vee ((\forall x)\neg P(x) \vee (\exists x)Q(x))$
 $(\exists x)\neg(\neg P(x) \vee Q(x)) \vee ((\forall x)\neg P(x) \vee (\exists x)Q(x))$
 $(\exists x)(P(x) \wedge \neg Q(x)) \vee ((\forall x)\neg P(x) \vee (\exists x)Q(x))$
 $(\exists x)(\forall y)(P(x) \wedge \neg Q(x)) \vee (\neg P(y) \vee (\exists x)Q(x))$
 $(\exists x)(\forall y)(\exists z)(P(x) \wedge \neg Q(x)) \vee (\neg P(y) \vee Q(z))$

Step 2:

$$(\forall y)(\exists z)(P(u) \wedge \neg Q(u)) \vee (\neg P(y) \vee Q(z))$$

$$(\forall y)(P(u) \wedge \neg Q(u)) \vee (\neg P(y) \vee Q(f(y)))$$

or $(\forall y)(P(a) \rightarrow Q(a)) \rightarrow (P(y) \rightarrow Q(f(y)))$

2. Step 1:

$$\begin{aligned}
& (\forall x)(P(x) \rightarrow (\exists y)Q(x, y)) \vee (\forall z)R(z) \\
& (\forall x)(\neg P(x) \vee (\exists y)Q(x, y)) \vee (\forall z)R(z) \\
& (\forall x)(\exists y)(\neg P(x) \vee Q(x, y)) \vee (\forall z)R(z) \\
& (\forall x)(\exists y)(\neg P(x) \vee Q(x, y) \vee (\forall z)R(z)) \\
& (\forall x)(\exists y)(\forall z)(\neg P(x) \vee Q(x, y) \vee R(z))
\end{aligned}$$

Step 2:

$$(\forall x)(\forall z)(\neg P(x) \vee Q(x, f(x)) \vee R(z))$$

Problem 3. (8 Points) Prove the following inference **by resolution**

$$(\forall x)(P(x) \vee Q(x)) \wedge (\forall x)(Q(x) \rightarrow \neg R(x)) \wedge (\forall x)R(x) \Rightarrow (\forall x)P(x)$$

Answer: Write the Skolem normal form of $\alpha \wedge \neg\beta$

$$\begin{aligned}
& (\forall x)(P(x) \vee Q(x)) \wedge (\forall x)(Q(x) \rightarrow \neg R(x)) \wedge (\forall x)R(x) \wedge \neg(\forall x)P(x) \\
& = (\forall x)(P(x) \vee Q(x)) \wedge (\forall x)(\neg Q(x) \vee \neg R(x)) \wedge (\forall x)R(x) \wedge (\exists x)\neg P(x) \\
& = (\forall x)(\forall y)(\forall z)(\exists v)((P(x) \vee Q(x)) \wedge (\neg Q(y) \vee \neg R(y)) \wedge R(z) \wedge \neg P(v))
\end{aligned}$$

The clause set is $S = \{P(x) \vee Q(x), \neg Q(y) \vee \neg R(y), R(z), \neg P(a)\}$

$$\left. \begin{array}{l} P(x) \vee Q(x) \\ \neg P(a) \end{array} \right\} \xrightarrow{\sigma=\{x/a\}} Q(a), \quad \left. \begin{array}{l} \neg Q(y) \vee \neg R(y) \\ Q(a) \end{array} \right\} \xrightarrow{\sigma=\{y/a\}} \neg R(a)$$

$$R(z) \xrightarrow{\sigma=\{z/a\}} R(a)$$

This gives us a contradiction $R(a) \wedge \neg R(a)$

Problem 4. (8 Points)

Formalize the following inference and prove it **by resolution**.

All SJTU students are smart. Bob is both an SJTU student and an NBA player. Therefore, some NBA player is smart.

Answer: We define following predicates a : “Bob” $P(x)$: “ x is an SJTU student.” $Q(x)$: “ x is smart.” $R(x)$: “ x is an NBA player.” We need to prove

$$(\forall x)(P(x) \rightarrow Q(x)) \wedge P(a) \wedge R(a) \Rightarrow (\exists y)(R(y) \wedge Q(y))$$

Write the Skolem normal form of $\alpha \wedge \neg\beta$

$$\begin{aligned}
& (\forall x)(P(x) \rightarrow Q(x)) \wedge P(a) \wedge R(a) \wedge \neg(\exists y)(R(y) \wedge Q(y)) \\
& = (\forall x)(\neg P(x) \vee Q(x)) \wedge P(a) \wedge R(a) \wedge (\forall y)(\neg R(y) \vee \neg Q(y)) \\
& = (\forall x)(\forall y)((\neg P(x) \vee Q(x)) \wedge (\neg R(y) \vee \neg Q(y)) \wedge P(a) \wedge R(a))
\end{aligned}$$

The clause set is $S = \{\neg P(x) \vee Q(x), \neg R(y) \vee \neg Q(y), P(a), R(a)\}$

$$\left. \begin{array}{l} \neg P(x) \vee Q(x) \\ P(a) \end{array} \right\} \xrightarrow{\sigma=\{x/a\}} \neg Q(a), \quad \left. \begin{array}{l} R(a) \\ \neg R(y) \vee \neg Q(y) \end{array} \right\} \xrightarrow{\sigma=\{y/a\}} \neg R(a)$$

This gives us a contradiction $R(a) \wedge \neg R(a)$